

Numerical Methods for Solving Large Linear Systems

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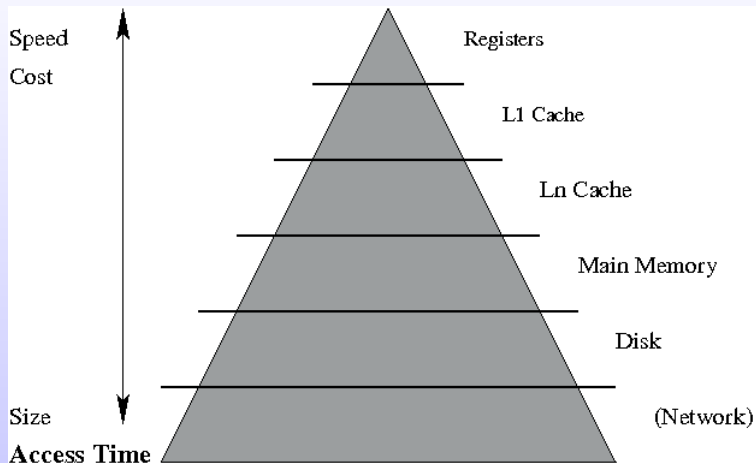
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- 1 High Performance Computing
- 2 Sparse Matrices
- 3 Linear Systems
- 4 Error Analysis
- 5 Eigensolvers





Cache misses. Data movement. Storage by row/columns. Data locality.

Linear Algebra operations decomposed into simpler operations.



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$$\begin{array}{ll} \text{BLAS-1:} & y := y + \alpha x \quad x, y \in \mathbb{R}^n \\ & \text{dot} := \alpha + x^T y \end{array}$$

$$\begin{array}{ll} \text{BLAS-2:} & y := y + Ax \quad L \in \mathbb{R}^{n \times n}, x, y \in \mathbb{R}^n \\ & y := A^{-1}x \quad L \in \mathbb{R}^{n \times n} \wedge \text{triangular} \end{array}$$

$$\begin{array}{ll} \text{BLAS-3:} & C := C + AB \quad A, B, C \in \mathbb{R}^{n \times n} \\ & C := L^{-1}B \quad L \in \mathbb{R}^{n \times n} \wedge \text{triangular} \end{array}$$

Linear Algebra operations decomposed into simpler operations.

BLAS-1:	$y := y + \alpha x$ $dot := \alpha + x^T y$	$x, y \in \mathbb{R}^n$
BLAS-2:	$y := y + Ax$ $y := A^{-1}x$	$L \in \mathbb{R}^{n \times n}, x, y \in \mathbb{R}^n$ $L \in \mathbb{R}^{n \times n} \wedge \text{triangular}$
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BLAS	#FLOPS	Mem. refs.	Ratio	Proc. use
Level 1	$2n$	$3n$	2/3	low
Level 2	$2n^2$	n^2	2	medium-low
Level 3	$2n^3$	$4n^2$	$n/2$	very high

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- Practical definition in terms of cost:
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A matrix is sparse when it has enough zeros
that pays off to exploit them (Wilkinson)

Objectives:

- Storage space.
- Accessing, inserting matrix elements.
- Matrix operations and fill in.

Sparse Matrices — Structured Matrices — Dense Matrices

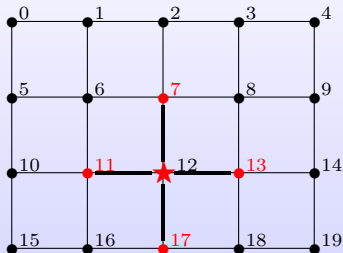
$$\begin{bmatrix} \star & & & & \star \\ & \star & \star & & \star \\ & \star & \star & & \star \\ & & & \star & \star \\ & \star & & \star & \\ \star & & & & \star \\ & & \star & \star & \star \end{bmatrix} \quad \begin{bmatrix} \star & \star & & & \\ \star & \star & \star & & \\ & \star & \star & \star & \\ & & \star & \star & \star \\ & & & \star & \star & \star \\ & & & & \star & \star & \star \\ & & & & & \star & \star \end{bmatrix}$$

- Structured matrices:
bidiagonal, tridiagonal, banded, blocked, ...

Small number of non-zeros (NNZ), but **known** structure!

Sparsity:

interactions between particles, components, neighbors, degrees of freedom.



The finer the discretization, the higher the sparsity.

- 50% of NNZ \rightarrow NOT a sparse matrix.
- 10% of NNZ: if $n = 100.000$, then 10.000 interactions per row. Still not very sparse.
- Large sparse matrices: $\text{NNZ} \ll 1\%$.

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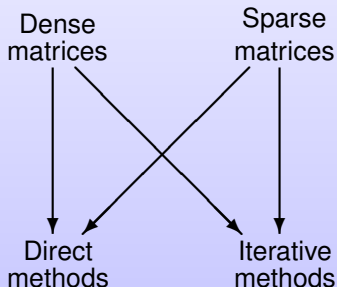
Iterative methods: Gauss-Seidel, Conjugate Gradient, GMRES,



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Iterative methods: Gauss-Seidel, Conjugate Gradient, GMRES,

Dense vs. Sparse (vs. Structured)



Direct

- Accuracy: fixed (cond num)
- Matrix-matrix operations
- Cost: $O(n^3)$, predictable
- Factorization re-use: multiple right-hand sides
- Fill-in, reordering

Iterative

- Variable accuracy
- Matrix-vector operations
- Cost not known: $k O(n^2)$
Convergence (spectrum)
Preconditioning
Stopping criteria
- Single/few right-hand sides
- Exploit sparsity

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BLAS, LAPACK, FLAME & PETSc
HSL, MUMPS, UMFPACK, ...

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PETSc, Trilinos, ...

Google: “Linear Algebra Software”

Survey of freely available libraries



30k, 1 core

Memory:	7GB
LU fact:	1640 secs (28m)
$Ax = b$:	1.9 secs (1 rhs)
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100k — extrapolating:

Memory	80GB
# of 8-cores	8 with 10-12GB each
LU fact	20 minutes

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$$Ax = b$$

- Acquisition and representation errors:

$$A \rightarrow \hat{A} = A + \delta A \quad b \rightarrow \hat{b} = b + \delta b$$

$$(A + \delta A)\hat{x} = b + \delta b$$

- $\hat{x} = x + \delta x$
- $\frac{\|\delta x\|}{\|x\|} = \mu(A) \frac{\|\delta A\|/\|A\| + \|\delta b\|/\|b\|}{1 - \mu(A)\|\delta A\|/\|A\|}$
- $\mu(A) = \|A\|\|A^{-1}\|$ is the **condition number** of A

Sensitivity to perturbations. Independent of the solution method.
Well vs. ill conditioned problems.



$f : X \rightarrow Y$ \hat{f} is an implementation of f

- Question: $|f - \hat{f}|$?

Exact arithmetic

$$x \rightarrow f(x)$$

Floating point arithmetic

$$\begin{aligned} x &\rightarrow \hat{f}(x) \\ (\hat{x} &\rightarrow \hat{f}(\hat{x})) \end{aligned}$$

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- **Forward stability:** $\forall x \ \|f(x) - \hat{f}(x)\|$ is small.
- Let \bar{x} be such that $\hat{f}(x) = f(\bar{x})$. Exact sol. to a different probl.
Backward stability: $\forall x \ \exists \bar{x} \ . \ \|x - \bar{x}\|$ is small.

Factorizations are backward stable.

Iterative methods \rightarrow convergence & convergence rate.



$$AV = V\Lambda$$



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- Three stages:
 - 1) Reduction to tridiagonal form
 - 2) **Tridiagonal eigensolver** ($TZ = Z\Lambda$)
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- Reduction: $O(n^3)$, perfectly stable, destroys sparsity.
- Backtransformation: matrix-matrix multiplication, $O(n^3)$, perfectly stable.
- Tridiagonal eigensolvers: MR³, QR, D&C, ...

Cost: $O(n^2)$ — $O(n^3)$

Accuracy: $\|Z^T Z - I\| \leq c n \epsilon \quad \wedge \quad \|TZ - Z\Lambda\| \leq c n \epsilon \|T\|$



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- Eigenvalues AND(?) eigenvectors? How many? Accuracy?
- LAPACK, PMR3, ScaLAPACK. Sparse solver: ARPACK.



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- Knowledge from applications.
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Thank you!

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