Dealing with dynamic and mixed workloads

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HTCondor WS, 2021, Sep 20

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INFN-T1, Current Status (HTC-CE/HTC, 2020)

- $(6+1) \times \text{HTC-CE}$, $1 \times \text{CM}$, $875 \times \text{WN}$, (40000 Cpus, $\sim 435 \text{ KHS06})$
- $1 \times SN$ for Remote Submission (from local UI, auth via FS_REMOTE)
- HTC 8.8.15, HTC-CE, 3.4.1
- We plan upgrading to HTC 9.0.5, HTC-CE 5.1 in the next days

From last year talk about LSF→HTC migration

Need (or wish) to improve our fairshare setup to consider different HS06 of WNs.
 Useful when job distribution of one AcctGroup is not homogeneus through the nodes.

This talk is mostly about what has been done about the above bullet. A new policy is in use from Early July 2021.

Usergroups, workloads, pledges and shares

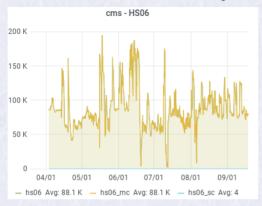
- ~ 50 User groups: 24 Grid VOs, ~ 25 local.
- LHC VOs are the 4 major players (total pledge sum up to $412\,\mathrm{KHS}06$)
- each group has GROUP_QUOTA_DYNAMIC_<group>= x_i and $\sum x_i = 1.0$

Multicore and singlecore

Note: in the following, $sc \rightarrow singlecore$, mc or $mcore \rightarrow multicore$. $mc \rightarrow RequestCpus = 8$ in job.ad

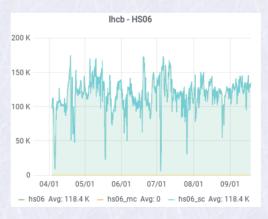
- Multicore jobs are $\sim 42\%$ (6 months average)
- 2,4,8 core (mostly 8; other size used by non-LHC groups)
- CMS → mc only, ATLAS → (quickly) variable mix
- ALICE → steady mix (from May '21), LHCb → sc only

LHC VOs jobs, last 6 months









Multicore provisioning

Initially: DEFRAG and recipe from the htcondor wiki

Later: also added a set of statically dedicated nodes.

Problems

- 1. Unused slots during mcore shortage, or overpledge of a group over another one
- 2. un-even opportunities: sc jobs tend to start sooner than mc of the same group
- 3. mcore of some groups last much longer than others (days vs hours)
- 4. fairshare makes no distinction on sc and mc and is unaware of different core power
- 5. whenever a mc job ends (and claim expired) one sc is likely to be the "next one" in queue and claim one of 8 just freed cores. Thus a new defrag time is needed.

Ideas for an alternative setup

- 1. Adding a grace time: when a mc job ends, only accept another mc for some time.
- 2. Order groups having pending jobs by expected HS06 usage current HS06 usage (also consider size of pending jobs). First groups are "rich", last are "poor".
- 3. have (a subset of) machines imposing an upper limit to the rich, and accepting jobs from the poor group: mc or sc, depending on which has more pending.

Implementation

- 1. A STARTD cronjob defines a boolean machine classad MC_GRACE defined as True if 8 Cpus are free for less than a few minutes (currently: 8 mins).
- 2. This requires centrally collecting data for all running and pending jobs (see below)
- 3. Another STARTD cronjob (JOB_CTL) to set a few more custom machine classads

Collecting data Every 17 minutes we run:

```
condor_q -global -all -cons 'Member(JobStatus,{1,2}) && (JobUniverse != 7)' \
-af JobStatus 'split((RemoteHost ?: "u@PEND.t1"),"@.")[1]' \
'split(AcctGroup,".")[0]' 'time() - (JobStartDate ?: time())' \
'CpusProvisioned ?: min({RequestCpus ?: 1,8})' \
'((int(MATCH_t1_wn_hs06 ?: 400) + 0.0)/(MATCH_TotalSlotCpus ?: 40))'

# MATCH_t1_wn_hs06 = node power as machine classad, inherited by the job
# MATCH_TotalSlotCpus = same as NUM_CPUS, inherited by the job
# NOTE: We already collect elsewhere these data for monitoring
```

The output is worked on and appended to an auxiliary shares_Error.log file:

# TO	VO	PLEDGE	DC	SC	MC	DHS	DPHS	PSC	PMC
1627101243	cms	87100	-542	0	4888	-3321	-29403	0	12784
1627101243	atlas	105300	-236	1608	4720	-489	-32021	4	2572
1627101243	alice	71400	1150	4170	1432	12594	-8785	242	410

Final result goes to a shared file:

```
      condor_q -glob ... -af ...
      sharectl.py → /shared/sharectl.txt

      AcctGroup_hs06.txt
      sharectl.py → /shared/sharectl.txt
```

The sharectl file

Since CMS has less than expected and ALICE has more, we reduce by 8 the maximum number of allowed ALICE cores per machine. We do this according to the number of pending CMS jobs. We set this target into the shared file, which looks like this:

```
~$ cat /shared/sharectl.txt
cn-609-05-06 alice 56 cms 8
cn-610-02-03 alice 56 cms 8
cn-608-06-06 alice 56 cms 8
...
```

The JOB_CTL cronjob. Run by STARTD every 13 mins. It sets the following classads:

```
cn-609-05-06 ~]# condor_status -comp -af:ln t1_CurrentJobs t1_TargetGroups t1_Targetcores
t1_CurrentJobs = alice:64:lhcb:3:atlas:5
t1_TargetGroups = { "alice","cms" }
t1_Targetcores = { 56,8 } # { 0,0 } means no target
```

It checks for its hostname into sharectl.txt and sets t1_Target* accordingly

Now to the START expression (just a little bit cumbersome)

```
cn-609-05-06 ~]# ccv StartJobs
True && (!t1_overheat) && (t1_mc_grace) && t1_sharectl

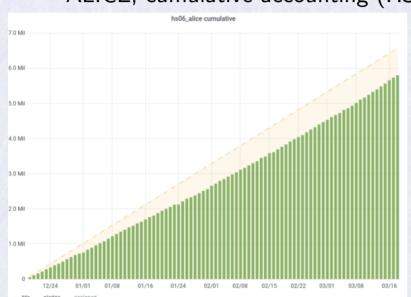
#Prevent singlecore when MC_GRACE is True
cn-610-05-06 ~]# ccv t1_mc_grace
( (TARGET.RequestCpus > 1) || ((TARGET.RequestCpus == 1) && !(MC_GRACE ?: False)) )

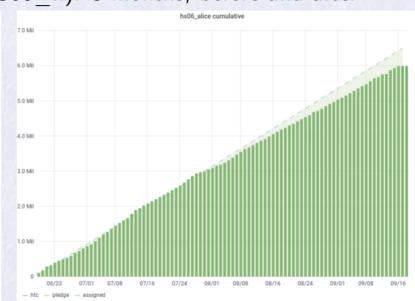
cn-610-05-06 ~]# ccv t1_sharectl
( (t1_Targetcores[0] =?= 0) || \
    (( split(AcctGroup,".")[0] =?= t1_TargetGroups[1] && RequestCpus =?= t1_Targetcores[1] ) || \
    ( AcctGroup =?= t1_TargetGroups[0] && \
    (t1_Targetcores[0] ?: 0) > int(split(t1_CurrentJobs ?: "none:0",":")[1]))))
```

```
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How does it work?

Started on July 2021, some stop (kernel upgrade and reboot). Steady from August. ALICE, cumulative accounting (HS06 h). 3 months, before and after



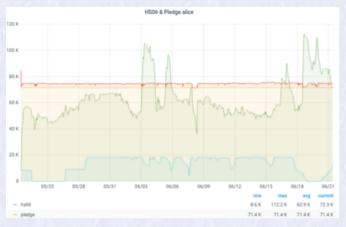


December 17 to March 16

June 16 to September 16

dashed line is target value.

ALICE, ATLAS, 30 days, old setup vs new









LHCb, CMS, 30 days, old setup vs new



Observations

- This setup is active on a subset $(\sim 50\%)$ of all computing power.
- The DEFRAG daemon was stopped a few days after
- The Errors used to compute sharectl.log are in HS06 units.

Possible improvements

- The control policy only takes the latest value for $e_g(t) = \text{current} \text{expected quota}$. We expect better results by averaging it with past values.
- The same should hold for pending jobs, to ensure that we consider groups having a regular submission rate