Multithreaded Data Quality Assessment in AthenaMT (+ consequences)

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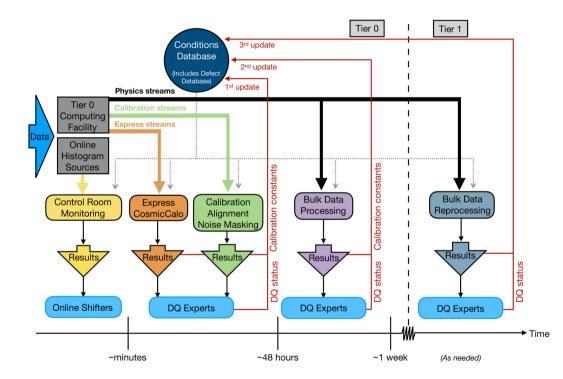
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Overview

- ATLAS offline data quality monitoring (DQM) chain is used to sign off data for good run lists and to provide feedback to optimize operation of the detector
- Histograms produced in Athena reconstruction are a major input to the system
 - runs both offline and online
- LHC shutdowns + migration to multithreaded framework provide opportunity to update system



Introduction

- Performance of detectors & calibrations is monitored with histograms created in the reconstruction workflow
 - O(10%) of data promptly processed in the "express" data stream, then full "Main" data 48 hours after end of run after calibration loop
 - detector monitoring algorithms need access to RAW data or quantities only available during reco (not stored in AOD) – monitoring algorithms must coexist with reco
- Reconstruction has been migrated to multithreaded AthenaMT in order to improve required memory/CPU core ratio
 - earlier effort: multiprocess (copy-on-write) "AthenaMP"
 - DQ a major cause of large memory use of AthenaMP each process gets its own copy of histograms
 - histograms need to be shared between threads to achieve good DQ memory scaling

Strategy

- Share a new jointly-developed core framework with High Level Trigger monitoring
 - able to leverage performance improvements, bug fixes, feature implementations from both
- MT-safe histogramming is tricky (and requires cooperation between all units): centralize hard parts
 - THistSvc design does not provide sufficient MT safety in histogram creation & management operations (MT-safe *filling* is "simple" part)
 - avoid global ROOT locks as much as we can
 - approach: programmers no longer touch raw histograms; all operations are handed in core libraries
 - AthenaMT works fairly hard to make the other parts of algorithms reentrant (execute methods are const, etc.)

Programmer Interface

- Histograms defined in Athena job Python configuration
 - extensive API to simplify tasks, e.g. for defining arrays of histograms for different detector regions
 - TH*, TProfile, TGraph, TEfficiency, TTree supported: can also support other backends although not used at present
 - histograms defined by variable names to be plotted
- Histograms filled in C++ event loop by providing variable names & data to histogramming tool
 - actual fill of histograms occurs in centralized code that determines what histograms can be filled given provided data

std::vector<int> detstatevec(xAOD::EventInfo::nDets+1); std::vector<int> detstatevec_idx(xAOD::EventInfo::nDets+1); std::iota(detstatevec_idx.begin(), detstatevec_idx.end(), 0);

auto detstates = Collection("detstates", detstatevec); auto detstates_idx = Collection("detstates_idx", detstatevec_idx);

// fill vectors

```
detstatevec[xAOD::EventInfo::nDets] = worststate;
fill(group, detstates, detstates_idx);
```

Critical Path: Lookup

- User code provides a list of what amounts to (variable name, variable data) pairs to framework in the fill() call
- Framework uses variable names to resolve which histograms need to be filled on each call
 - variable name scoped within a "group" (= Athena tool instance); can have many groups for a single monitoring algorithm
 - slightly pathological case: "a bunch of histograms for each of > 1000 muon detector chambers, distinguished by the name of the variables we call fill() with"
 - patterns quite different in offline and in trigger code
- User code usually calls fill() with a very small set of possible arguments, and set of histograms doesn't change once job starts: get enormous speedup by caching the lookups
- Also provide methods for users to provide vector data instead of scalar data: only do expensive lookup once per event instead of inside tight inner loops
 - "cutmask" variables can be provided for histogram definitions to only plot a subset of a vector
- Performance checked with profiler runs with regularity
 - can be quite slow if naively used, but rarely a problem once code is optimized

Critical Path: Histogram Filling

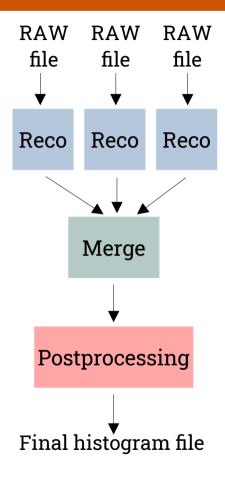
- Biggest issue: minimize memory motion of data values
 - avoid intermediate representations of data when possible
 - provide an interface to fill with arbitrary functions of elements of a collection (as long as collection supports operator[])
- As always, lots of subtle places for race conditions to creep in
 - e.g. rebooking of time-dependent histograms in multiple threads
 - THistSvc is missing an atomic "check if histogram exists, if not book" operation
- Actual filling protected behind per-histogram locks
 - Even running nothing but monitoring algorithms on an AOD, at 40 threads, DQ framework locks show no significant contention (much bigger effects in event I/O, etc.)

Histogram Framework Status

- At this point the framework is quite mature
 - still adding a few features as requested, e.g. "rolling last-N luminosity block" plots or prettier autogeneration of histogram titles
- Almost all Run 2 monitoring algorithms have been migrated to the new framework
- Strategy of isolating histogram-handling code in central, carefully-managed code has paid off
 - framework bugs can be fixed for all systems at once
 - clear best practices, things aren't reinvented by every system
- DQ CI & other tests fairly frequently catch breakage in other domains (unintentional AOD output changes, subtle problems in multithreaded I/O ...)
- Some small reduction in functionality largely addressed by next part of presentation...

Histogram Postprocessing

- Histogram production has several phases
 - an accumulation step with commuting & reversible operations (usually addition to a bin of an array) in a single reconstruction job
 - merging of results from multiple jobs
 - optional "postprocessing" (e.g. make a new plot showing mean residuals)
- Similar to map-reduce (except that often reduce is "trivial" histogram addition)
- Postprocessed histograms are in general not possible to merge between jobs coherently
 - e.g. efficiency plots really need to keep numerator and denominator separate until final plot making, which is why ROOT now has TEfficiency
- In past, used to allow arbitrary C++ postprocessing inside Athena (run at termination of a job)
 - this isn't compatible with the new monitoring architecture (user algorithms have no access to the histogram data) so is now forbidden
 - introduce a new system to handle postprocessing



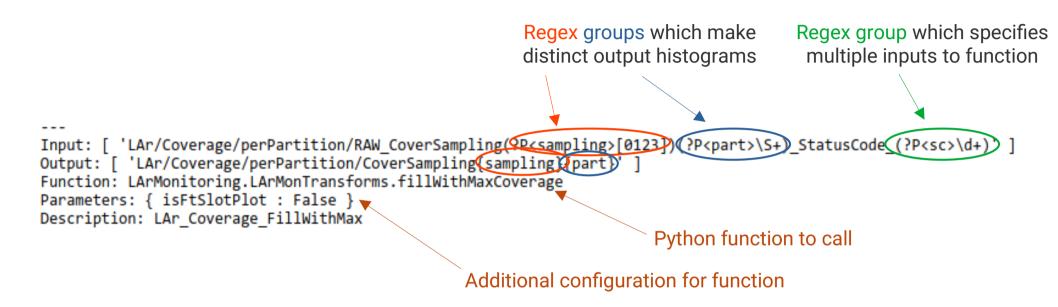
Postprocessing Engine

- Introduce generic framework for operations on histograms: histgrinder
 - complete factorization of histogram processing logic from access: handled via I/O plugins (ROOT file, ATLAS online histogramming system, Athena THistSvc, ...)
 - framework does not require any specific histogram technology
 - processing algorithms written in Python (but can use cppyy for speed)
 - pattern matching to simplify processing of similar histograms (e.g. different detector layers/regions)
 - for online operations: accepts histograms as they arrive & updates outputs
- ATLAS implementations:
 - offline: ROOT file \rightarrow ROOT file
 - online: distributed online histogram (OH) system \rightarrow OH
 - Athena piggyback for online: THistSvc \rightarrow THistSvc in parallel to the reco job

Identical Python postprocessing code in multiple environments



Example Histgrinder Configuration



Same YAML configuration & user code used for offline and online applications; only difference is I/O plugins

Summary

- New histogram production framework introduced with multithreaded AthenaMT upgrade
- Necessitated parallel deployment of a new histogram postprocessing framework
- Both deployed and used for Run 2 data reprocessing validation & cosmics data taking