# ANALOG CONTROL OF THE DIAMOND QUANTUM PROCESSOR

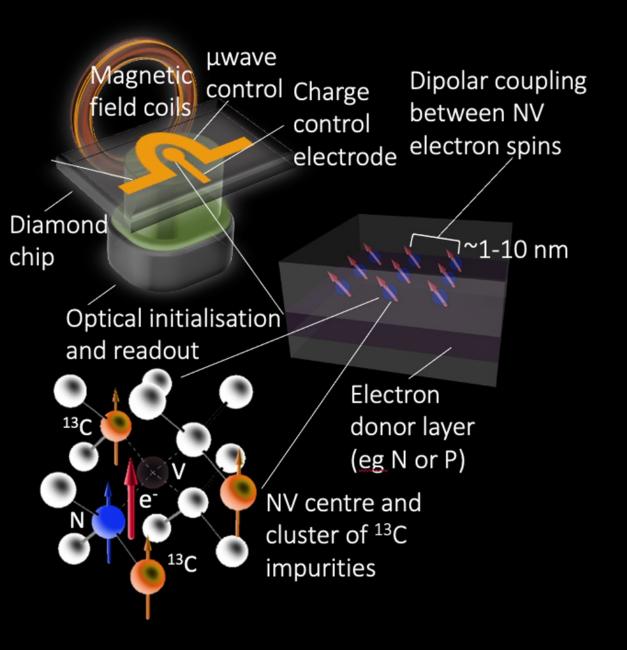
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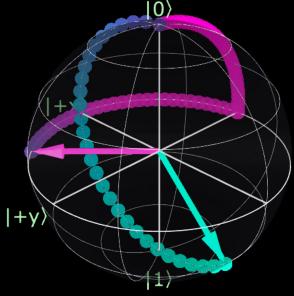
## DIAMOND QUANTUM COMPUTING

- Uses the NV-centre, an optical defect.
- Long coherence time at room temperature.
- Controlled via MW and RF pulses.
- Quantum Brilliance is in the process of designing a fully scalable array of NV centres.
- Currently able to easily fabricate diamond samples with many NV centres.
- Currently, computation using diamond is only done via the gate paradigm- perhaps there are alternatives...



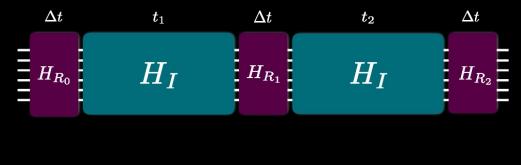
## ANALOG QUANTUM COMPUTING

- Sequence of gates can be thought of a single rotation about an arbitrary axis.
- No gates- rotate straight to end result.
- Continuous driving, rather than discrete pulses.



- Potential error robustness
- AQC compilation is difficult.
- The challenge is an opportunity.

- ParityQC and D-Wave- quantum annealing
- Not the same as CVQC encoding.
- Better for diamond: DAQC- fusing gates and annealing together.

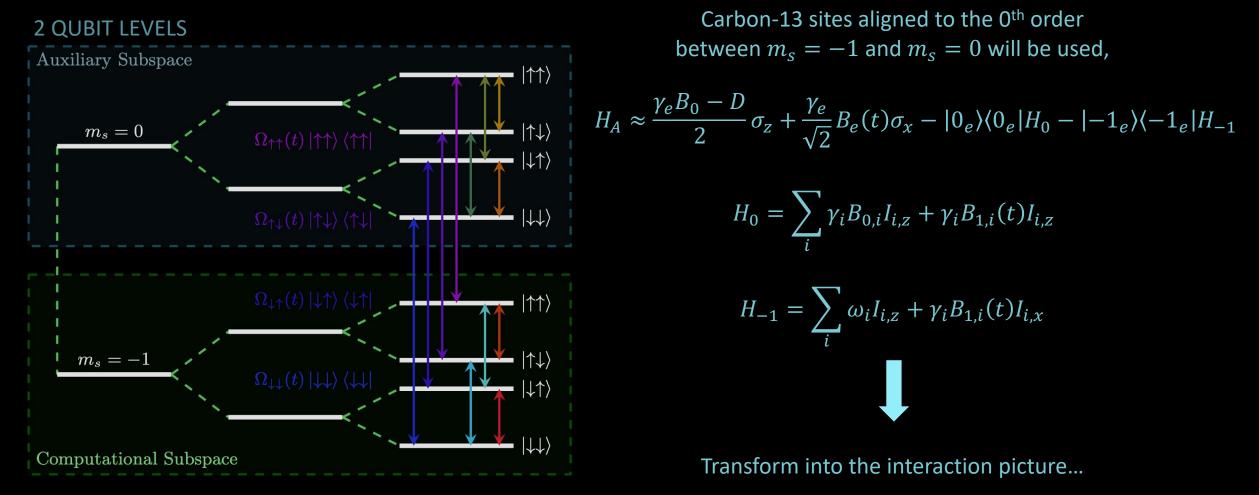


$\Delta t$	$t_1 - 3\Delta t/2$	$\Delta t$	$t_2 - 3\Delta t/2$	$\Delta t$	
 $H_{R_0}$		$H_{R_1}$		$H_{R_2}$	
+	$H_{I}$	+	$H_{I}$	+	
$H_I$	1	$H_I$	1	$H_{I}$	

- Expand quantum annealing work.
- Begin to address analog compilation challenges.
- Implement DAQC on diamond.

## HAMILTONIAN

$$H_{S} = D\left(S_{z}^{2} - \frac{2}{3}\right) + \gamma_{e}\left(S_{z}B_{0} + S_{x}B_{1}(t)\right) + \sum_{i}\vec{S}\cdot\vec{A}_{i,z}\cdot\vec{I}_{i} - \gamma_{i}\left(I_{i,z}B_{0} + I_{i,x}B_{1}(t)\right)$$



## THE ISING MODEL

 $H_{C} = H_{S}(t) + |0_{e}\rangle\langle 0_{e}|H_{0} + |-1_{e}\rangle\langle -1_{e}|H_{-1}$  $H_{S}(t) = \frac{\gamma_{e}}{\sqrt{2}}B_{1}(t)\sum_{\alpha} (\sigma_{x}\cos\Lambda_{\kappa_{\alpha}}t - \sigma_{y}\sin\Lambda_{\kappa_{\alpha}}t) \left(\prod_{i} m_{\kappa_{\alpha i}}I_{i,z} + \frac{1}{2}\right)$ 

Collect like terms to write each transition as a linear combination of nuclear spin operators multiplied by applied control fields...

$$H_{Ising} = S_x \left( \sum_{i < j} \Omega_{i,j}(t) I_{i,z} I_{j,z} + \sum_i \Omega_i(t) I_{i,z} \right) + \sum_i \Xi_i(t) I_{i,x} + \Xi_i(t) I_{i,y}$$

Two qubit example:  $\kappa_{\alpha} = \{\uparrow\uparrow, \uparrow\downarrow, \downarrow\uparrow, \downarrow\downarrow\}$  $H_{S}(t) = \frac{\gamma_{e}}{\sqrt{2}} B_{1}(t) \left\{ \frac{1}{4} (I_{1z}I_{2z} + I_{1z} + I_{2z} + 1) \right\}$  $\times \left(\sigma_x \cos \Lambda_{\uparrow\uparrow} t - \sigma_y \sin \Lambda_{\uparrow\uparrow} t\right) \\ + \frac{1}{4} \left(-I_{1z} I_{2z} + I_{1z} - I_{2z} + 1\right)$  $\times \left(\sigma_x \cos \Lambda_{\uparrow\downarrow} t - \sigma_y \sin \Lambda_{\uparrow\downarrow} t\right) \\ + \frac{1}{4} \left(-I_{1z} I_{2z} - I_{1z} + I_{2z} + 1\right)$  $\times \left(\sigma_x \cos \Lambda_{\downarrow\uparrow} t - \sigma_y \sin \Lambda_{\downarrow\uparrow} t\right)$  $+ \frac{1}{4}(I_{1z}I_{2z} - I_{1z} - I_{2z} + 1)$  $\times \sigma_x \cos \Lambda_{\downarrow\downarrow} t - \sigma_v \sin \Lambda_{\downarrow\downarrow} t$ 

Gather  $\Lambda$ s together in linear combination to produce  $\Omega_{1,2}$ ,  $\Omega_1$ ,  $\Omega_2$ ,  $\Xi_1$  and  $\Xi_2$ .

#### TWO-CHANNEL DAQC

As per the Rodriguez formula, single qubit gates can be written as,

$$U_{R_{\phi}} = e^{iR_{\phi}\hat{n}\cdot\overline{\sigma}}$$

All entanglement can be done with an Ising coupling gate.

$$U_{I_{zz}} = e^{\overrightarrow{I_z} \overleftarrow{\Theta} \overrightarrow{I_z}}$$

Where,

$$\vec{I}_{z} = (I_{1z}, I_{2z}, \dots, I_{nz})$$
$$\overleftrightarrow{\Theta} = \begin{bmatrix} \Theta_{11} & \cdots & \Theta_{1n} \\ \vdots & \ddots & \vdots \\ \Theta_{n1} & \cdots & \Theta_{nn} \end{bmatrix}$$

Also, as the **Ising model commutes with itself** all entangling gates *and* single qubit gates can be performed in a single step-provided the hardware is capable of rotations about any arbitrary axis.

Diamond computers can do this.

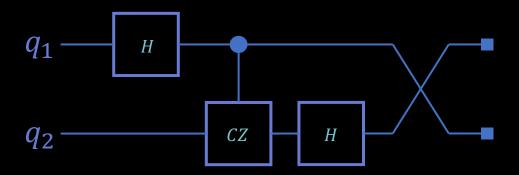
	1 QUBIT GATE	2 QUBIT GATE	1 QUBIT GATE	•••
Qubit 1	$R(\phi_1, \hat{n})$	$R_{zz}(\Theta_{12})$	$R(\phi_1, \hat{n})$	
Qubit 2	$R(\phi_2, \hat{n})$	$R_{zz}(\Theta_{12})$	$R(\phi_2, \hat{n})$	

*New way to control diamond quantum processors: Continuous driving of superposition of electron spin states.* 

#### **EXAMPLE: ANALOG QFT**

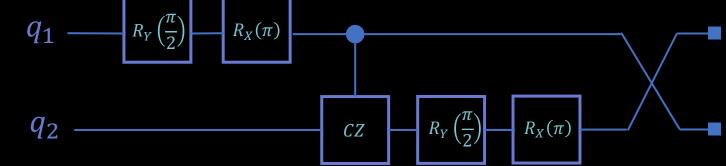
 $U(t) = e^{S_x \left(\sum_{i < j} \Omega_{i,j}(t) I_{i,z} I_{j,z} + \sum_i \Omega_i(t) I_{i,z}\right) + \sum_i \Xi_{i,x}(t) I_{i,x} + \sum_i \Xi_{i,y}(t) I_{i,y}}$  $U(t) = e^{S_x \left(\Omega_{ZZ,12}(t)I_{1,Z}I_{2,Z} + \Omega_{Z,1}(t)I_{1,Z} + \Omega_{Z,2}(t)I_{2,Z} + \Lambda_0\right) + \Xi_{1,X}(t)I_{1,X} + \Xi_{2,X}(t)I_{2,X} + \Xi_{1,Y}(t)I_{1,Y} + \Xi_{2,Y}(t)I_{2,Y}}$  $F_4 = \frac{1}{2} \begin{vmatrix} 1 & 1 & 1 & 1 \\ 1 & i & -1 & -i \\ 1 & -1 & 1 & -1 \\ 1 & -i & -1 & i \end{vmatrix}$ 

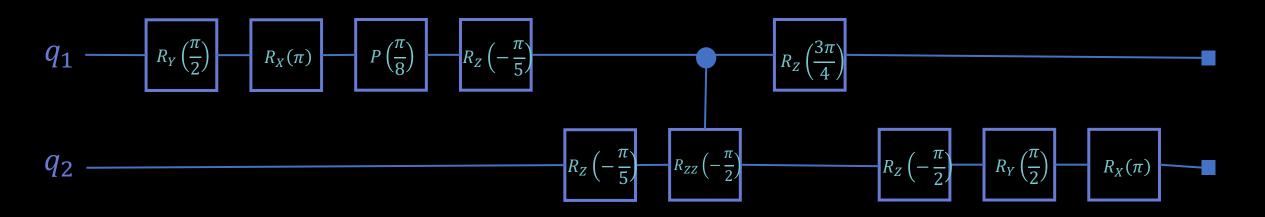
## 2-QUBIT QFT



Gate-based QFT compiled to native diamond gates. Circuit depth increases.

Splitting up the CZ gate gives an Ising coupling gate.

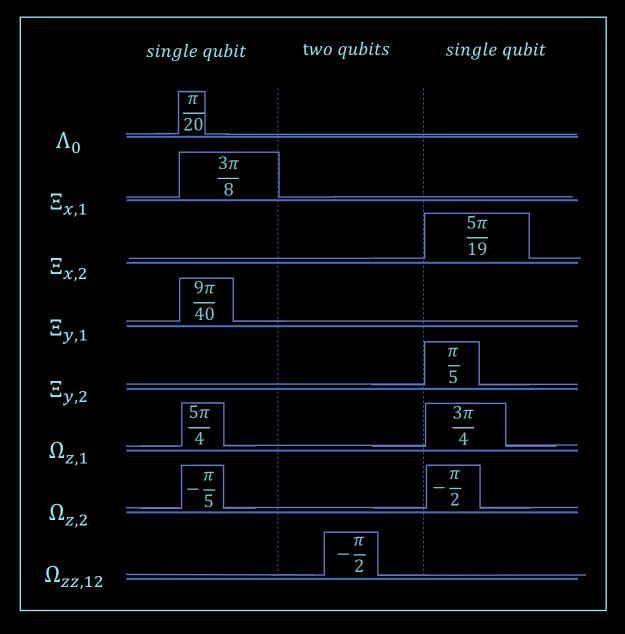




## **REDUCE DEPTH**

Invoking the Rodriguez formula, this can be compacted back into 3 gates.

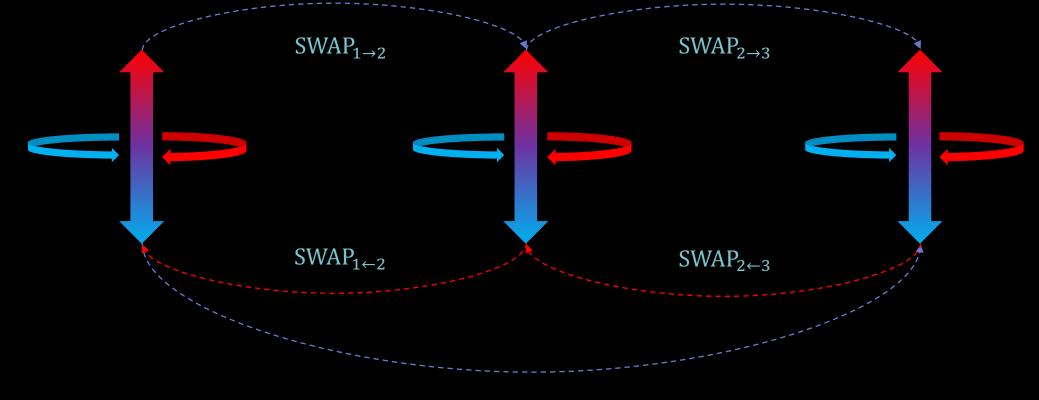
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Where each of the components of  $\hat{n}$ becomes our pulse components.

## CONNECTIVITY

 $e^{i(\Theta_{12}I_{1z}I_{2z}+\Theta_{23}I_{2z}I_{3z}+\Theta_{13}I_{1z}I_{3z})} = e^{i(\Theta_{12}I_{1z}I_{2z}+\Theta_{23}I_{2z}I_{3z})}.SWAP_{2\to 3}.e^{i(\Theta_{13}I_{1z}I_{3z})}.SWAP_{3\to 2}$ 



 $\Theta_{13}I_{1z}I_{3z}$ 

## FUTURE WORK

- See how this formulations performs error-wise compared to gate model implementations of the Fourier transform.
- See how number of SWAP gates scales with size and limited connectivity- and if this outweighs potential benefits.
- Expand past the hybrid model, looking to full universal AQC.
- Investigate what this means for optimal control.
- Find systematic way to compile AQC pulses for diamond.

## THANK YOU

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