PY410 / 505 Computational Physics 1

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Next up : Linear algebra

- Covered in Chapter 4 of Garcia and Chapter 2 of Numerical Recipes
- Huge number of applications!
 - -Complex circuit diagrams
 - -Coupled oscillators
 - -General solution of fitting arbitrary curves
- We'll learn how to :
 - -Solve linear equations
 - -Compute eigenvalues
 - -Apply these to various applications

 You should all be familiar with the basics of linear algebra

Vectors
$$\mathbf{A}\mathbf{x} = \mathbf{b}$$

- -Matrices
- -Solving matrix equations
- -Gaussian (or Gauss-Jordan) elimination
 - <u>http://en.wikipedia.org/wiki/</u> <u>Gaussian_elimination</u>
- We'll go over the computational issues here

http://xkcd.com/184/

$$\begin{bmatrix} \cos 90^\circ & \sin 90^\circ \\ -\sin 90^\circ & \cos 90^\circ \end{bmatrix} \begin{bmatrix} \alpha_1 \\ \alpha_2 \end{bmatrix} = \underbrace{90^\circ}_{2^\circ} \underbrace{90^\circ}_{2^\circ}$$

• Recall Cramer's rule :

-http://en.wikipedia.org/wiki/Cramer's_rule

• If we have the determinant of a matrix A :

$$|\mathbf{A}| = \sum_{i=1}^{n} (-1)^{i+j} A_{ij} |\mathbf{R}_{ij}| ,$$

- where R_{ij} is the "residual matrix" or "minor" of A, removing row i and column j
- Then the inverse of the matrix is :

$$A_{ij}^{-1} = (-1)^{i+j} \frac{|\mathbf{R}_{ji}|}{|\mathbf{A}|}$$

• This is a recursive rule

- Computing this "brute force" way is okay for small n, but problematic for large n (>10)
- Scales as n factorial (n!)
- To see this, consider an expansion of the determinant :

$$|\mathbf{A}| = \sum_{P} (-)^{P} A_{1p_{1}} A_{2p_{2}} \dots A_{np_{n}},$$

- Here, P runs over the n! permutations of the indices
- 20! = 2.43x10¹⁸ yipes! (which is "yipes, factorial")
- OK, well, scratch that idea.
 What else ya got, Sal?

- Medium sized matrices (n~10-10³)
 - -Gaussian elimination, LU-decomposition, and Householder method
- Larger matrices $(n > 10^3)$
 - -Storage becomes a problem
 - -Cannot practically do this for arbitrary matrices
 - However, most matrices are "sparse" with lots of zeroes in practical applications
 - -Can, however, store and solve these fairly well

- Linear algebra is the raison d'être for numpy.
 Let's just use it.
- Most of the scipy algorithms for linear algebra are from LAPACK. (Linear Algebra Package)
 - <u>http://www.netlib.org/lapack/</u>
- Other options:
 - -BLAS (Basic Linear Algebra Subprograms)
 - <u>http://en.wikipedia.org/wiki/Basic_Linear_Algebra_Subprograms</u>
 - -LAPACK
 - -BOOST Basic Linear Algebra Library
 - <u>http://www.boost.org/doc/libs/1_54_0/libs/numeric/ublas/doc/</u> index.htm
 - -matlab (the "mat" in "matlab" stands for "matrix")
 - <u>http://www.mathworks.com/products/matlab/</u>

- Game plan:
 - -Use numpy software.
 - -Go over algorithms that are used internally
 - -Check into use cases

So if you recall the Gaussian Elimination, we have a matrix equation :

$$Ax = b$$
,

• Or, written out for the case of n=3:

$$\begin{pmatrix} a_{00} & a_{01} & a_{02} \\ a_{10} & a_{11} & a_{12} \\ a_{20} & a_{21} & a_{22} \end{pmatrix} \begin{pmatrix} x_0 \\ x_1 \\ x_2 \end{pmatrix} = \begin{pmatrix} b_0 \\ b_1 \\ b_2 \end{pmatrix} ,$$

$$\begin{pmatrix} a_{00}x_0 + a_{01}x_1 + a_{02}x_2 \\ a_{10}x_0 + a_{11}x_1 + a_{12}x_2 \\ a_{20}x_0 + a_{21}x_1 + a_{22}x_2 \end{pmatrix} = \begin{pmatrix} b_0 \\ b_1 \\ b_2 \end{pmatrix}$$

.

- We take linear combinations of the rows to convert the matrix into "reduced row echelon form":
 - –Multiply first equation by a(10)/a(00) and subtract from second:

 $a_{10}x_0 + a_{11}x_1 + a_{12}x_2 - (a_{10}/a_{00})(a_{00}x_0 + a_{01}x_1 + a_{02}x_2) = b_1 - (a_{10}/a_{00})b_0.$

-Then x0 is eliminated from the second equation:

 $(a_{11} - a_{10}a_{01}/a_{00})x_1 + (a_{12} - a_{10}a_{02}/a_{00})x_2 = b_1 - (a_{10}/a_{00})b_0 ,$

-which can be written as:

 $a_{11}'x_1 + a_{12}'x_2 = b_1' \; .$

-Repeat until you run out of rows

-Example:

$$\begin{bmatrix} 1 & 3 & 1 & 9 \\ 1 & 1 & -1 & 1 \\ 3 & 11 & 5 & 35 \end{bmatrix} \rightarrow \begin{bmatrix} 1 & 3 & 1 & 9 \\ 0 & -2 & -2 & -8 \\ 0 & 2 & 2 & 8 \end{bmatrix} \rightarrow \begin{bmatrix} 1 & 3 & 1 & 9 \\ 0 & -2 & -2 & -8 \\ 0 & 0 & 0 & 0 \end{bmatrix} \rightarrow \begin{bmatrix} 1 & 0 & -2 & -3 \\ 0 & 1 & 1 & 4 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

• Can also think about this in terms of an "augmented" matrix where you add the vector b as another column:

| (a_{00}) | a_{01} | a_{02} | b_0 |
|------------|----------|----------|--|
| a_{10} | a_{11} | a_{12} | $\begin{pmatrix} b_1 \\ b_2 \end{pmatrix}$ |
| a_{20} | a_{21} | a_{22} | b_2 |

• You eliminate each row iteratively, reducing the dimension by one each time:

$$\begin{pmatrix} a_{00} & a_{01} & a_{02} & b_0 \\ \frac{a_{10} - a_{00} a_{10}}{a_{00}} & a_{11} - \frac{a_{01}a_{10}}{a_{00}} & a_{12} - \frac{a_{02}a_{10}}{a_{00}} & b_1 - b_0 a_{10}/a_{00} \\ a_{20} - \frac{a_{00}a_{20}}{a_{00}} & a_{21} - \frac{a_{01}a_{20}}{a_{00}} & a_{22} - \frac{a_{02}a_{20}}{a_{00}} & b_2 - b_0 a_{20}/a_{00} \end{pmatrix}$$

$$= \begin{pmatrix} a'_{00} & a'_{01} & a'_{02} & b'_{0} \\ 0 & a'_{11} & a'_{12} & b'_{1} \\ 0 & a'_{21} & a'_{22} & b'_{2} \end{pmatrix}$$

$$\begin{pmatrix} a'_{00} & a'_{01} & a'_{02} & b'_{0} \\ 0 & a'_{11} & a'_{12} & b'_{1} \\ 0 & a'_{21} - a'_{11}a'_{21}/a'_{11} & a'_{22} - a'_{12}a'_{21}/a'_{11} & b'_{2} - b'_{1}a'_{21}/a'_{11} \end{pmatrix}$$

$$= \begin{pmatrix} a''_{00} & a''_{01} & a''_{02} & b''_{0} \\ 0 & a''_{11} & a''_{12} & b''_{1} \\ 0 & 0 & a'''_{22} & b''_{2} \end{pmatrix}$$

- To solve for the actual equation, we then use "back substitution", starting at the last row
 - -For instance

$$x_2 = \frac{b_2''}{a_{22}''}$$
.

-Then we use the second equation

$$x_1 = \frac{1}{a_{11}''} \left[b_1'' - a_{12}'' x_2 \right] \; .$$

-And finally

$$x_0 = \frac{1}{a_{00}^{\prime\prime}} \left[b_0^{\prime\prime} - a_{01}^{\prime\prime} x_1 - a_{02}^{\prime\prime} x_2 \right] \; .$$

- Gaussian elimination is is $O(n^3)$ operations -n² matrix elements, and O(n) row operations on each
- Back-substitution is O(n²) operations
- So the total is $O(n^3) + O(n^2) \sim O(n^3)$ for large n
- Much better than $O(n!) >>O(n^3)$
- Easy to extend to arbitrarily large n's
 - But, as we mentioned earlier, storage becomes a problem

- You can see one problem already
- If any of the rows have a zero as the first element, then you divide by zero and get an exception



- So, need to make sure this doesn't happen by "partial pivoting" :
 - -Before performing the ith row operation, search for the element a_{ki} (k=i,...n-1) with the largest magnitude
 - -If k != i, interchange rows i and k of the augmented matrix, and interchange x_i <---> x_k
 - -Perform as usual
- Will not fail for small a_i's, also stable to roundoff errors
- Note : in "full" pivoting you swap rows AND columns

- Another variation is Gauss-Jordan elimination
 - -Does not require the backsubstitution step
 - Replaces A by the inverse "in place", avoiding memory copy
- So, we have another augmented equation (again in n=3):

$$\begin{pmatrix} a_{00} & a_{01} & a_{02} \\ a_{10} & a_{11} & a_{12} \\ a_{20} & a_{21} & a_{22} \end{pmatrix} \begin{pmatrix} x_0 & a_{00}^{-1} & a_{01}^{-1} & a_{02}^{-1} \\ x_1 & a_{10}^{-1} & a_{11}^{-1} & a_{12}^{-1} \\ x_2 & a_{20}^{-1} & a_{21}^{-1} & a_{22}^{-1} \end{pmatrix} = \begin{pmatrix} b_0 & 1 & 0 & 0 \\ b_1 & 0 & 1 & 0 \\ b_2 & 0 & 0 & 1 \end{pmatrix}$$

• This is a way of stating :

$$Ax = b$$
, and $AA^{-1} = 1$

- Algorithm :
 - -Start with zeroth row, divide by a(00) and subtract
 - -Subtract all zeroth column entries (as in Gaussian elimination)
 - -For row i=1,...n-1, divide by diagonal element a(ii) and eliminate elements in column i other than the diagonal by subtracting a(ij) times the ith row elements from the jth row
 - Gaussian elimination : subtracts only those below diagonal
 - Gauss-Jordan : subtracts ALL elements where i != j
 - -Simultaneously perform on the augmented (b 1) matrix on the RHS

• In the end we get :

$$\begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix} \begin{pmatrix} \mathsf{b0} & a_{00}^{-1} & a_{01}^{-1} & a_{02}^{-1} \\ \mathsf{b1} & a_{10}^{-1} & a_{11}^{-1} & a_{12}^{-1} \\ \mathsf{b2} & a_{20}^{-1} & a_{21}^{-1} & a_{22}^{-1} \end{pmatrix} = \begin{pmatrix} x_0 & a_{00}^{-1} & a_{01}^{-1} & a_{02}^{-1} \\ x_1 & a_{10}^{-1} & a_{11}^{-1} & a_{12}^{-1} \\ x_2 & a_{20}^{-1} & a_{21}^{-1} & a_{22}^{-1} \end{pmatrix}$$

- A is replaced by a unit matix
- b is reduced to input vector x
- Unit matrix on RHS is replace by inverse A⁻¹
- Can also implement pivoting in this algorithm to make sure we have numeric stability

- Another variation is "LU decomposition" ("lower-upper")
- Reduce A to an L*U product: Ax = b,

$$\begin{pmatrix} a_{00} & a_{01} & a_{02} \\ a_{10} & a_{11} & a_{12} \\ a_{20} & a_{21} & a_{22} \end{pmatrix} = \begin{pmatrix} \alpha_{00} & 0 & 0 \\ \alpha_{10} & \alpha_{11} & 0 \\ \alpha_{20} & \alpha_{21} & \alpha_{22} \end{pmatrix} \begin{pmatrix} \beta_{00} & \beta_{01} & \beta_{02} \\ 0 & \beta_{11} & \beta_{12} \\ 0 & 0 & \beta_{22} \end{pmatrix}$$

- Then solve the problem in two steps :
 - -Solve $\mathbf{L}\mathbf{y} = \mathbf{b}$ with forward substitution:

$$y_0 = \frac{b_0}{\alpha_{00}}$$
, $y_i = \frac{1}{\alpha_{ii}} \left[b_i - \sum_{j=0}^{i-1} \alpha_{ij} y_j \right]$, $i = 1, 2, \dots, n-1$.

-Solve $\mathbf{U}\mathbf{x} = \mathbf{y}$, since $\mathbf{A}\mathbf{x} = \mathbf{L}(\mathbf{U}\mathbf{x}) = \mathbf{L}\mathbf{y} = \mathbf{b}$, with backward substitution :

$$x_{n-1} = \frac{y_{n-1}}{\beta_{n-1,n-1}}, \qquad x_i = \frac{1}{\beta_{ii}} \left[y_i - \sum_{j=i+1}^{n-1} \beta_{ij} x_j \right], \qquad i = n-2, \dots, 0.$$

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 Factoring the matrix A = LU can be done with "Crout's Algorithm" :

$$-Set alpha(ii) = 1$$
 for i=0,...n-1

-for each j = 0,...n-1 :
for i = 0,...j
compute
$$\beta_{ij} = a_{ij} - \sum_{k=0}^{i-1} \alpha_{ik} \beta_{kj}$$

for j = j+1,...n-1
compute
$$\alpha_{ij} = \frac{1}{\beta_{jj}} \left[a_{ij} - \sum_{k=0}^{j-1} \alpha_{ik} \beta_{jk} \right]$$

• Replaces A by LU "in place":

$$\begin{pmatrix} a_{00} & a_{01} & a_{02} \\ a_{10} & a_{11} & a_{12} \\ a_{20} & a_{21} & a_{22} \end{pmatrix} \longrightarrow \begin{pmatrix} \beta_{00} & \beta_{01} & \beta_{02} \\ \alpha_{10} & \beta_{11} & \beta_{12} \\ \alpha_{20} & \alpha_{21} & \beta_{22} \end{pmatrix}$$

The RH matrix is NOT a matrix! It is a storage unit in the computer!

- Simple special case : Tridiagonal matrices
- If you have a problem such as :

$$\mathbf{M} = \begin{pmatrix} 2-c & -1 & 0 & \cdots & 0 & 0 \\ -1 & 2-c & -1 & \cdots & 0 & 0 \\ 0 & -1 & 2-c & \cdots & 0 & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & 0 & \cdots & 2-c & -1 \\ 0 & 0 & 0 & \cdots & -1 & 2-c \end{pmatrix},$$

• This matrix is sparse!

-Can solve in O(n) operations

• Equations are:

$$M_i^- u_{i-1} + M_i^0 u_i + M_i^+ u_{i+1} = b_i$$
, $i = 1, ..., n-1$.

• Can show that (recursively):

$$\alpha_{i-1} = -\frac{M_i^-}{M_i^0 + \alpha_i M_i^+} , \qquad \beta_{i-1} = \frac{b_i - \beta_i M_i^+}{M_i^0 + \alpha_i M_i^+} .$$

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Can start from the right boundary value for i=n-2,...0

 $u_n = \alpha_{n-1}u_{n-1} + \beta_{n-1}$ if $\alpha_{n-1} = 0$, $\beta_{n-1} = u_n$,

and then solve from the left boundary value for i=1,...n-1

$$u_{i+1} = \alpha_i u_i + \beta_i \; ,$$

So we "sweep twice" and obtain O(n) operations

- Examples:
 - -Polynomial fits (again)
 - -Circuit diagrams
 - -Boundary value problems

• This is a matrix equation, so we define the "design matrix" :

$$A_{ij} = \frac{Y_j(x_i)}{\sigma_i}$$

$$\mathbf{A} = \begin{bmatrix} Y_1(x_1)/\sigma_1 & Y_2(x_1)/\sigma_1 & \dots \\ Y_1(x_2)/\sigma_2 & Y_2(x_2)/\sigma_2 & \dots \\ \dots & \dots & \dots & \dots \end{bmatrix}$$

• Then our chi2 minimization becomes :

$$(\mathbf{A}^T \mathbf{A})\vec{a} = \mathbf{A}^T \vec{b}$$

• SO :

$$\vec{a} = (\mathbf{A}^T \mathbf{A})^{-1} \mathbf{A}^T \vec{b}$$

• If we define the "correlation matrix" :

$$\mathbf{C} = (\mathbf{A}^T \mathbf{A})^{-1}$$

• Then the uncertainty on a_j is :

$$\sigma_{a_j} = \sqrt{C_{jj}}$$

• As a first example, let's look at polynomial fits

$$y = \sum_{k=0}^{m-1} a_k x^k$$

- Slight generalization of the linear fit we did previously
- General solution is to minimize the chi2 :

$$\chi^2(\vec{a}) \equiv \sum_{i=0}^{n-1} \left(\frac{y_i - y(x; \vec{a})}{\sigma_i} \right)^2$$

• In this case :

$$\chi^{2}(\vec{a}) = \sum_{i=0}^{n-1} \left(\frac{y_{i} - \sum_{j=0}^{M} a_{j} x^{j}}{\sigma_{i}} \right)^{2}$$

• Our design matrix is therefore :

 $A_{ij} = x_i^j / \sigma_i$

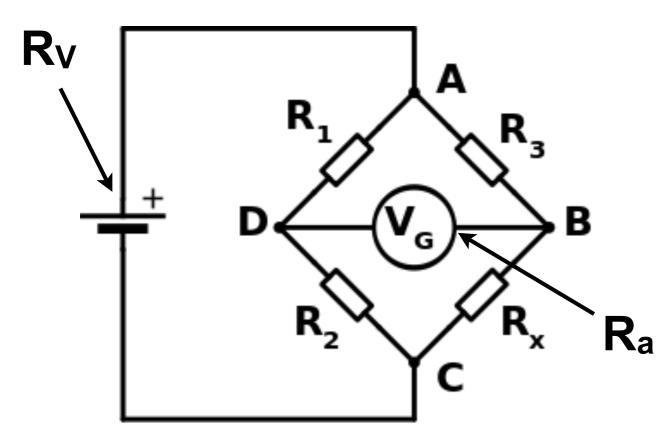
 Caveat : This oftentimes is ill-formed, so don't go too crazy here. Typically we do quadratic, cubic, quartic, but above that it strains credibility.

Polynomial fits

Can finally generalize our formalism to arbitrary functional fits

• Example : quadratic fit for our CO2 data!

- Example :
 - -Kirchoff's Law for a Wheatstone bridge :
 - -http://en.wikipedia.org/wiki/Wheatstone_bridge



Solve for Rx given R1,R2,R2, i and V

• This is a matrix equation: $\mathbf{Ri} = \mathbf{v}$,

$$\begin{pmatrix} R_1 + R_v & R_2 & R_v \\ R_1 + R_a & -R_a & -R_3 \\ R_x + R_a & -R_2 - R_x - R_a & R_x \end{pmatrix} \begin{pmatrix} i_1 \\ i_2 \\ i_3 \end{pmatrix} = \begin{pmatrix} v \\ 0 \\ 0 \end{pmatrix}$$

• Try to "zero" the potential at V_G:

$$V_G = \left(\frac{R_x}{R_3 + R_x} - \frac{R_2}{R_1 + R_2}\right) V_s$$

• Example : boundary value problems

-Consider:

$$\frac{\mathrm{d}^2 u}{\mathrm{d} x^2} = -\frac{\pi^2}{4}(u+1) \; ,$$

-with Dirichlet boundary conditions u(0) = 0 and u(1)=1-We can discretize this :

$$\frac{2u_i - u_{i+1} - u_{i-1}}{h^2} = \frac{\pi^2}{4}(u_i + 1), \qquad i = 1, \dots, N-1$$

-This is therefore a sparse matrix equation with $c = h^2$ pi^2/4:

$$\mathbf{M} = \begin{pmatrix} 2-c & -1 & 0 & \cdots & 0 & 0 \\ -1 & 2-c & -1 & \cdots & 0 & 0 \\ 0 & -1 & 2-c & \cdots & 0 & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & 0 & \cdots & 2-c & -1 \\ 0 & 0 & 0 & \cdots & -1 & 2-c \end{pmatrix}, \qquad \mathbf{x} = \begin{pmatrix} u_1 \\ u_2 \\ u_3 \\ \vdots \\ u_{N-2} \\ u_{N-1} \end{pmatrix}, \qquad \mathbf{b} = \begin{pmatrix} u_0 + c \\ c \\ \vdots \\ c \\ u_N + c \end{pmatrix},$$

Hands on!

- Today : Eigenvalues and eigenvectors, more hands on
- See Chapter 11 of Numerical Recipes
- In this case, even they recommend using packaged software for eigenvalues and eigenvectors, but let's get the general gist here

 Example : normal modes of a harmonic oscillator between n objects

$$\mathcal{L} = \frac{1}{2} \sum_{j,k}^{n} M_{jk} \dot{q}_j \dot{q}_k - \frac{1}{2} \sum_{j,k}^{n} A_{jk} q_j q_k ,$$

• "Normal mode" is the mode of the system where all of the coordinates oscillate with some frequency omega:

$$q_j(t) = x_j e^{i\omega t}$$
, $\sum_{j=1}^n \left[A_{jk} - M_{jk} \omega^2 \right] x_j = 0$, $(\mathbf{A} - \mathbf{M} \omega^2) \mathbf{x} = 0$.

 Homogeneous matrix equation has solutions for omega for which the determinant is zero:

$$det \left| \mathbf{A} - \mathbf{M} \omega^2 \right| = 0.$$

Eigenvalues and Eigenvectors

- This is an instance of a general class of eigenvalue and eigenvector problems
- So, if A is an nxn matrix, x is a column vector (the "right" eigenvector), lambda is a number (the "eigenvalue"), and :

$$Ax = \lambda x$$
,

• Then the equation is satisfied when :

$$\mathbf{A} - \lambda \mathbf{1} = \sum_{k=0}^{n} a_k \lambda^k = 0 \; ,$$

- This is an n-th degree polynomial in lambda that depends on the matrix elements
- N-degree polynomial ==> n different eigenvalues, so we need to solve for them

Eigenvalues and Eigenvectors

• Example :

$$\mathbf{A} = \left(\begin{array}{cc} 1 & 1\\ 1 & 2 \end{array}\right)$$

• The eigenvalue equation and eigenvalues are :

$$\mathbf{A} = \begin{pmatrix} 1-\lambda & 1\\ 1 & 2-\lambda \end{pmatrix}$$
$$(1-\lambda)(2-\lambda)-1 = \lambda^2 - 3\lambda + 1 = 0 \quad \Rightarrow \qquad \lambda_0 = \frac{3+\sqrt{5}}{2}, \quad \lambda_1 = \frac{3-\sqrt{5}}{2}$$

• We solve for the eigenvectors :

$$\begin{pmatrix} 1 & 1 \\ 2 & 1 \end{pmatrix} \begin{pmatrix} x_0 \\ x_1 \end{pmatrix} = (1 \pm \sqrt{2}) \begin{pmatrix} x_0 \\ x_1 \end{pmatrix} \qquad \mathbf{x}^{(0)} = \begin{pmatrix} \sqrt{\frac{1}{3}} \\ \sqrt{\frac{2}{3}} \end{pmatrix}, \qquad \mathbf{x}^{(1)} = \begin{pmatrix} \sqrt{\frac{1}{3}} \\ -\sqrt{\frac{2}{3}} \end{pmatrix}.$$

Eigenvalues and Eigenvectors

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• Also consider :

$$\mathbf{A} = \begin{pmatrix} 1 & 1 \\ -1 & 1 \end{pmatrix}$$

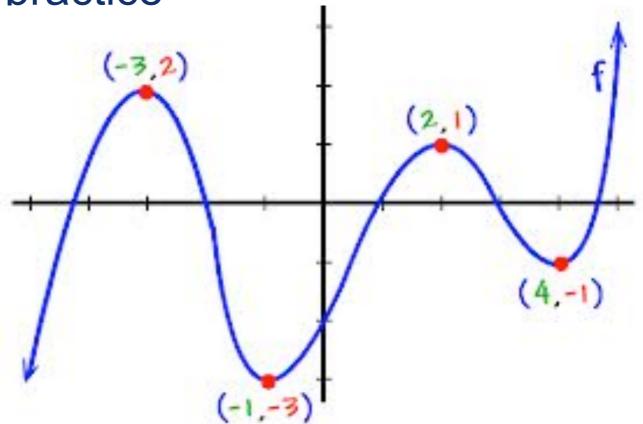
• So we have:

$$(1 - \lambda)^2 + 1 = \lambda^2 - 2\lambda + 2 = 0$$

• And thus the eigenvalues are complex :

$$\lambda_0 = 1 + i$$
, $\lambda_1 = 1 - i$, where $i = \sqrt{-1}$.

- "Hey Sal! I know what we can do!" you say. "We can just find the roots of the characteristic polynomial with our rootfinding methods!"
- "Excellent idea in principle, my clever students," I say. "But unfortunately, it's hard in practice"
 - -Need to know roughly where the roots are ahead of time
 - -You don't get the eigenvectors this way either



• So, we'll look at ways to do both!

• First, define "left" eigenvectors:

$$\mathbf{y}\mathbf{A} = \lambda \mathbf{y}$$
,

- In this case, y is a row vector.
- This implies : $\mathbf{A}^{\mathrm{T}}\mathbf{y} = \lambda \mathbf{y} \;,$

• Since :
$$|\mathbf{A}^{\mathrm{T}} - \lambda \mathbf{1}| = |\mathbf{A} - \lambda \mathbf{1}|$$
,

then the left and right eigenvalues are the same

- However the left and right eigenvectors are not necessarily the same
- Left eigenvector 0, though, is orthogonal to right eigenvector 1, etc

• If we consider a complex matrix A, we can define the Hermitian conjugate :

$$\left(\mathbf{A}^{\dagger}\right)_{ij} = a_{ji}^{*}$$
,

- -This is also referred to as the "conjugate transpose" -If the a_{ij} are real, this is just the transpose
- Define a "normal" matrix if the conjugate transpose commutes with the matrix:

$$\mathbf{N} \cdot \mathbf{N}^{\dagger} = \mathbf{N}^{\dagger} \cdot \mathbf{N}$$
.

Note that normal matrices have the same left and right eigenvector sets

• A Hermitian matrix (or self-adjoint matrix) is defined when

$\mathbf{H}^{\dagger}=\mathbf{H}\;,$

 When the elements are real, this is called a "symmetric" matrix:

$\mathbf{S}^{\mathrm{T}} = \mathbf{S}$,

- These special matrices have :
 - -n real eigenvalues
 - -eigenvectors are orthogonal
 - -the set of eigenvectors is "complete" (spans the n-dim space)

- Computation of eigenvalues/eigenvectors for symmetric or Hermitian matrices depends on a nice property
- If we consider the matrices formed by the right and left eigenvectors X and Y:

$$\mathbf{X} = \begin{pmatrix} \sqrt{\frac{1}{3}} & \sqrt{\frac{1}{3}} \\ \sqrt{\frac{2}{3}} & -\sqrt{\frac{2}{3}} \end{pmatrix}, \qquad \mathbf{Y} = \begin{pmatrix} \sqrt{\frac{2}{3}} & \sqrt{\frac{1}{3}} \\ \sqrt{\frac{2}{3}} & -\sqrt{\frac{1}{3}} \end{pmatrix} = \frac{2\sqrt{2}}{3} \mathbf{X}^{-1}$$

- In general : $~~{
 m Y} \propto {
 m X}^{-1}$
- Can choose the normalization appropriately so that

$$\mathbf{Y} = \mathbf{X}^{-1}$$

• How does this help us?

$$\mathbf{Y} \cdot \mathbf{A} \cdot \mathbf{X} = \mathbf{X}^{-1} \cdot \mathbf{A} \cdot \mathbf{X}$$
$$= \mathbf{X}^{-1} \cdot \lambda \cdot \mathbf{X}$$
$$= \lambda$$

• Can exploit this fact to compute the eigenvalues!

• This is a similarity transformation!

$$\mathbf{A} \rightarrow \mathbf{Z}^{-1} \cdot \mathbf{A} \cdot \mathbf{Z}$$

 The interesting bit is that symmetry transformations leave the eigenvalues unchanged:

$$det \left| \mathbf{Z}^{-1} \cdot \mathbf{A} \cdot \mathbf{Z} - \lambda \mathbf{1} \right| = det \left| \mathbf{Z}^{-1} \cdot (\mathbf{A} - \lambda \mathbf{1}) \cdot \mathbf{Z} \right|$$
$$= det \left| \mathbf{Z} \right| det \left| \mathbf{A} - \lambda \mathbf{1} \right| det \left| \mathbf{Z}^{-1} \right|$$
$$= det \left| \mathbf{A} - \lambda \mathbf{1} \right|$$

• So, we can solve this easier problem instead

- Let X be the matrix of eigenvectors
 - If the matrix A is real and symmetric, then the eigenvectors are real and orthonormal, and :

$$\mathbf{X}^{-1} = \mathbf{X}^{\mathrm{T}}$$

– If the matrix A is is Hermitian, the matrix of eigenvectors is unitary :

$$\mathbf{X}^{-1} = \mathbf{X}^{\dagger}$$

Generalized eigenvalue problem : solve

 $\mathbf{A} \cdot \mathbf{x} = \lambda \mathbf{B} \cdot \mathbf{x}$

- Strategy : Successively apply similarity transformations until the matrix is "almost" diagonal
 - -Either diagonal, block diagonal, or tridiagonal and also easy to solve
- Nonsymmetric matrices will not, in general, have similarity matrices with real components
 - -So, cannot use real similarity matrix
 - -But! "Almost" can : will reduce to block-diagonal with two-by-two blocks replacing the complex eigenvalues

-Think
$$\begin{pmatrix} \cos\theta & \sin\theta \\ -\sin\theta & \cos\theta \end{pmatrix}$$

 Can then repeatedly apply similarity transformations until we can solve the problem "easily":

$$\begin{array}{cccc} \mathbf{A} & \rightarrow & \mathbf{P}_1^{-1} \cdot \mathbf{A} \cdot \mathbf{P}_1 & \rightarrow & \mathbf{P}_2^{-1} \cdot \mathbf{P}_1^{-1} \cdot \mathbf{A} \cdot \mathbf{P}_1 \cdot \mathbf{P}_2 \\ \\ \rightarrow & \mathbf{P}_3^{-1} \cdot \mathbf{P}_2^{-1} \cdot \mathbf{P}_1^{-1} \cdot \mathbf{A} \cdot \mathbf{P}_1 \cdot \mathbf{P}_2 \cdot \mathbf{P}_3 & \rightarrow & \text{etc.} \end{array}$$

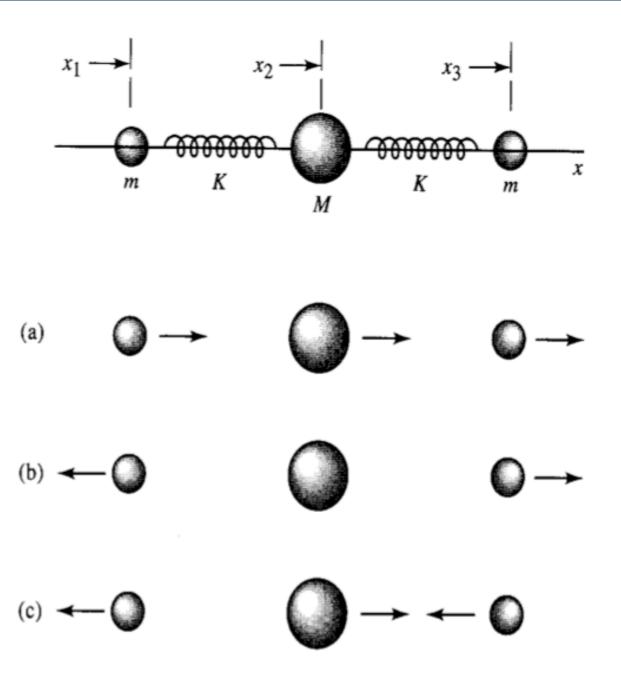
 If we end up in completely diagonal form, then the eigenvectors are just the columns of the total transformation:

$$\mathbf{X}_R = \mathbf{P}_1 \cdot \mathbf{P}_2 \cdot \mathbf{P}_3 \cdot \ldots$$

 Sometimes just want eigenvalues, in which case can stop when we reduce to triangular form
 Then eigenvalues are the diagonal elements!

- One strategy we will look at :
 - Reduce to tridiagonal form by the Householder algorithm (Chapter 11 Section 2 of Numerical Recipes)
 - Solve the tridiagonal matrix problem using a QL algorithm with implicit shifts (Chapter 11 Section 3 of Numerical Recipes)
 - Q = orthogonal
 - L = lower triangular matrix
- These are basically "uninteresting" to go through the gory details, because they're mostly just math tricks
- Instead, let's just focus on an application

- Consider a linear triatomic molecule
 - -Fowles and Cassiday, Chapter 11, Section 4
 - -Taylor, Chapter 11, Section 6
- Approximate by masses on springs



• Lagrangian for the general case is :

$$L = \frac{1}{2}m_1\dot{x}_1^2 + \frac{1}{2}m_2\dot{x}_2^2 + \frac{1}{2}m_3\dot{x}_3^2 - \frac{1}{2}\left[k_{12}(x_1 - x_2)^2 + k_{23}(x_2 - x_3)^2\right]$$

The equations of motion are

$$m_1 \ddot{x}_1 = \frac{d}{dt} \left(\frac{\partial L}{\partial \dot{x}_1} \right) = \frac{\partial L}{\partial x_1} = -k_{12} x_1 + k_{12} x_2$$
$$m_2 \ddot{x}_2 = \frac{d}{dt} \left(\frac{\partial L}{\partial \dot{x}_2} \right) = \frac{\partial L}{\partial x_2} = k_{12} x_1 - (k_{12} + k_{23}) x_2 + k_{23} x_3$$
$$m_3 \ddot{x}_3 = \frac{d}{dt} \left(\frac{\partial L}{\partial \dot{x}_3} \right) = \frac{\partial L}{\partial x_3} = -k_{23} x_3 + k_{23} x_2$$

• For our specific problem, we have

$$m_1 = m_3, m_2 = 2m_1, k_{12} = k_{23} = K$$

• Can be written as a matrix equation:

$$M\ddot{q} = -Kq$$

• with normal modes of the form :

 $\mathbf{q}(t) = \mathbf{a}\cos(\omega t - \delta)$

• The normal mode frequencies are eigenvalues of the generalized eigenvalue equation:

$$\mathbf{K}\mathbf{a} = \omega^2 \mathbf{M}\mathbf{a}$$

• The eigenvalues are therefore:

$$\omega_1^2 = 0 \ , \qquad \omega_2^2 = \frac{K}{m} \ , \qquad \omega_3^2 = \frac{K}{m} \left(1 + \frac{2m}{M}\right)$$

• And the eigenvectors are :

$$\mathbf{a}_{1} = a_{11} \begin{pmatrix} 1\\1\\1 \end{pmatrix} , \qquad \mathbf{a}_{2} = a_{12} \begin{pmatrix} 1\\0\\-1 \end{pmatrix} , \qquad \mathbf{a}_{3} = a_{13} \begin{pmatrix} 1\\-\frac{2m}{M}\\1 \end{pmatrix} ,$$

• "triatomic" python notebook