

Zero-Copy Merge with RNTuples

EUGENIO MARINELLI

Supervisors:

Jakob Blomer

Javier Lopez Gomez

Introduction

- ▶ Datasets in HEP are often produced in a parallel manner
 - ▶ partial results that are merged into a single file at the end of the data processing workflow
- ▶ **Merging** the content of many RNTuples requires significant resource consumption
- ▶ Modern file systems allow for zero-copy cloning of data blocks for a file
 - ▶ potentially **high impact** for the merging of HEP data (no physical copying required)

Zero-Copy Merge Project

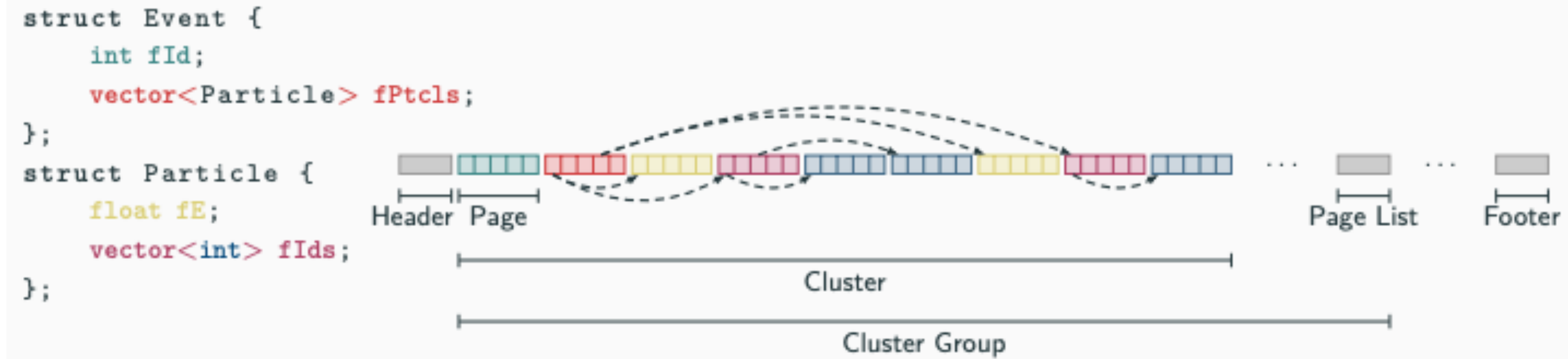
- ▶ In this work:
 - ▶ analyze:
 - ▶ the zero-copy functionality in modern file systems
 - ▶ the requirements to use it with ROOT RNTuple
 - ▶ ROOT RNTuple is a new I/O subsystem (still in development)
 - ▶ implement a first version of Zero-Copy Merge on ROOT RNTuples
 - ▶ comparison with alternative implementations

RNTuple Merging Operation

- ▶ RNTuple file:

- ▶ Metadata

- ▶ Events data



- ▶ We want one resulting RNTuple with:

- ▶ Metadata updated

- ▶ Same data content of all the source RNTuples

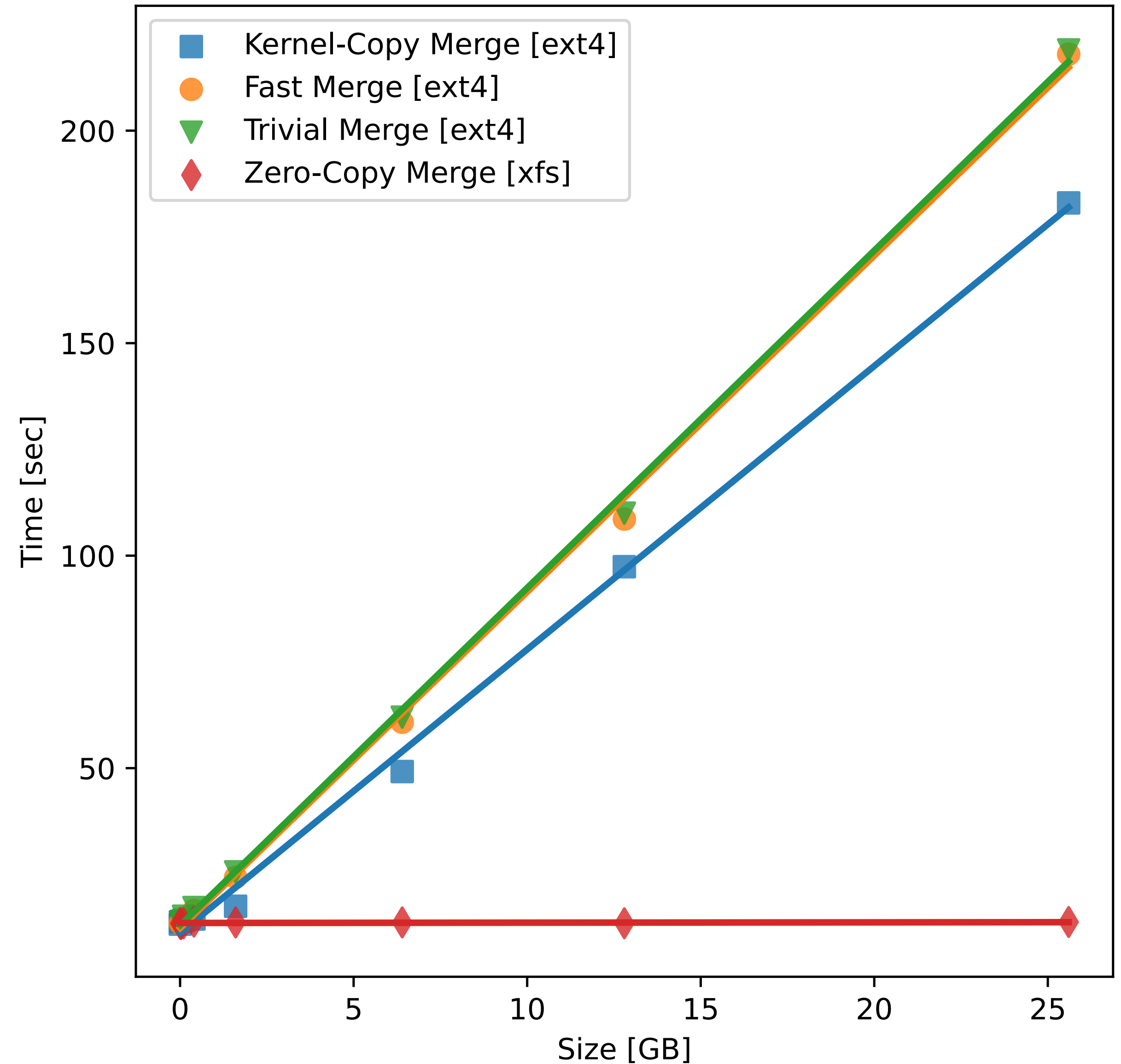
- ▶ Note: we want to **clone** the content without copying it (zero-copy)

Evaluation

- ▶ Zero-Copy merge compared with:
 - ▶ Trivial merge
 - ▶ read and write data using standard RNTuple API
 - ▶ Fast merge
 - ▶ read from source files and write on destination file (user space)
 - ▶ Kernel merge
 - ▶ file copy in kernel space

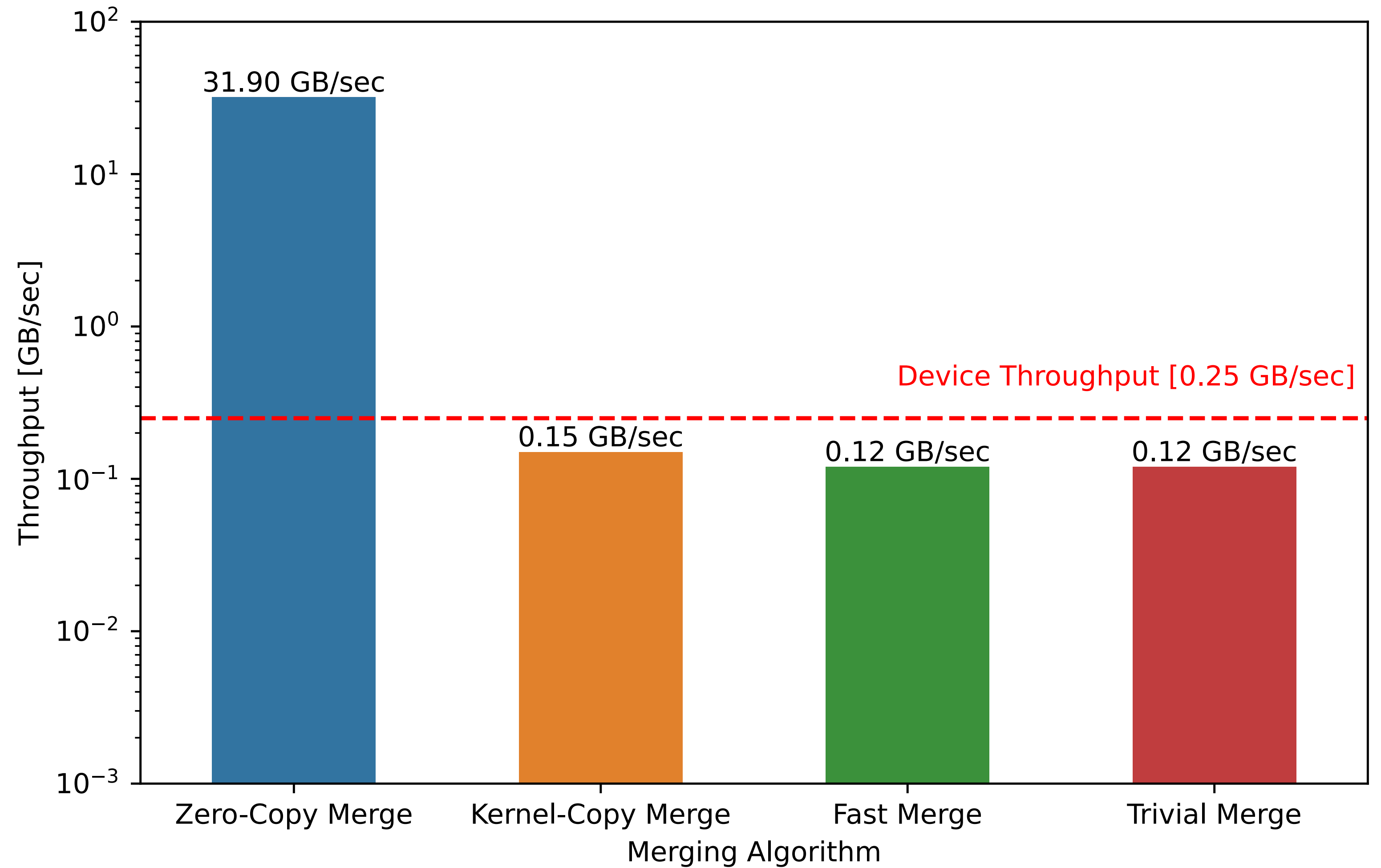
Evaluation: Execution Time

- ▶ Zero copy merge nearly constant
 - ▶ metadata updates is the only contribution
- ▶ Other merge grows linearly
- ▶ Test data is uncompressed



Evaluation: Throughput

- ▶ We achieve higher throughput than the device limit



Conclusion

- ▶ This first Zero-Copy merge implementation enables the RNTuple merge in almost constant time
- ▶ These results have significant impact on ROOT's object store support
 - ▶ HEP data in object stores (such as Intel DAOS) is by nature prepared to exploit zero-copy merging

THANK YOU!