

Fast Pattern Matching in Quantum Circuits

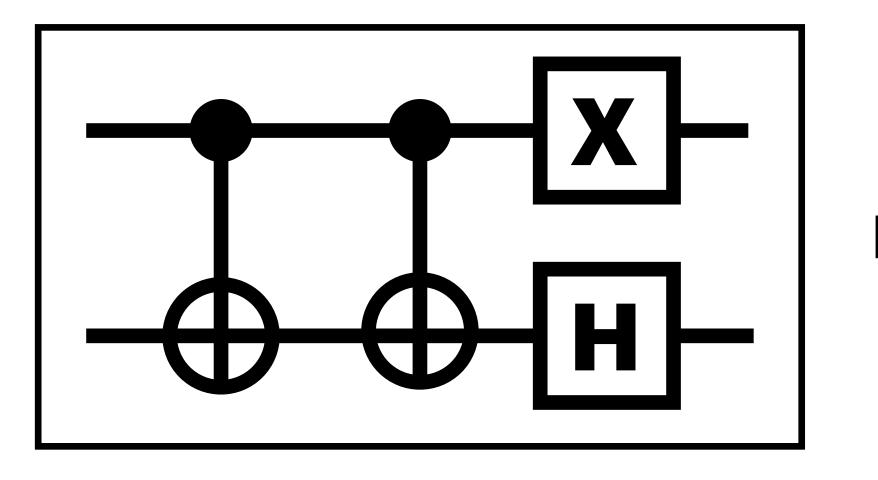
Luca Mondada & Pablo Andres-Martinez ICE-8 Conference, Santiago Spain, 1st June 2023

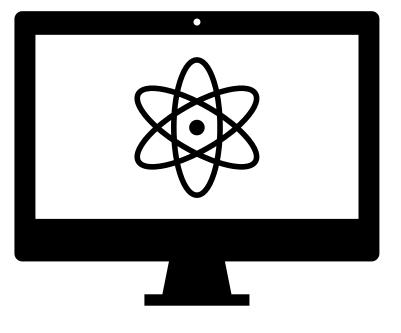
QUANTINUUM



<u>arXiv:2302.06717</u> [quant-ph]

Quantum Circuit compilation pipeline

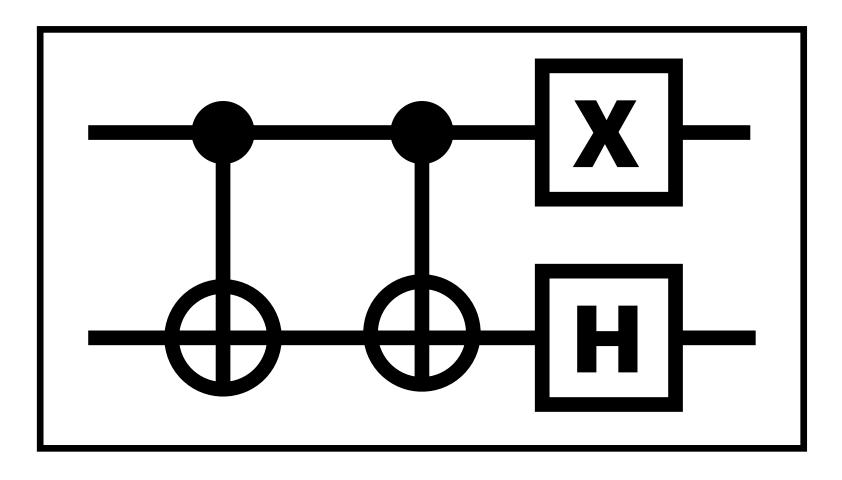




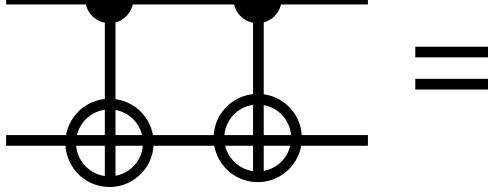
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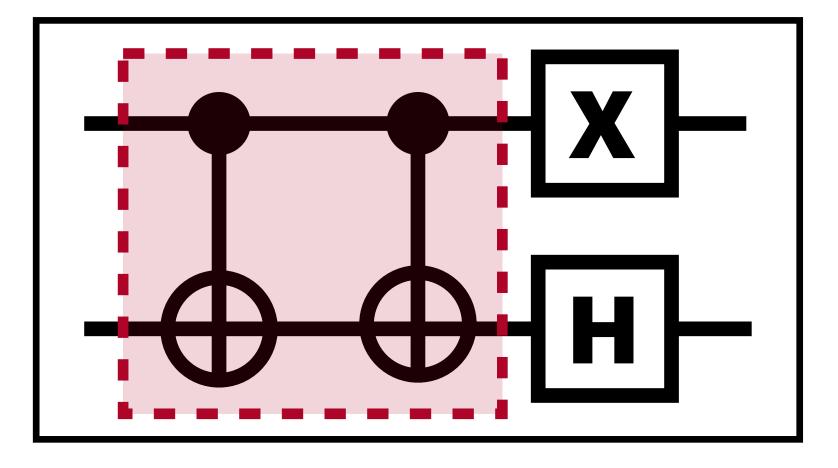
Quantum Circuit optimisation ... is about to get a lot harder

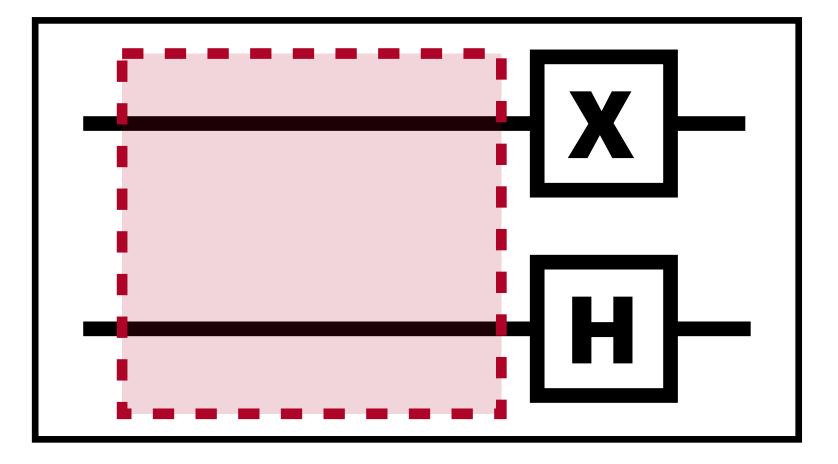
- Ever larger circuits 1.
- Ever larger instruction set 2.
- Every bit of optimisation will matter 3.



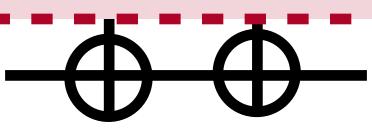
$-\mathbf{H}\mathbf{H} = ----$







Problem: too many rewriting rules!



Digression String pattern matching

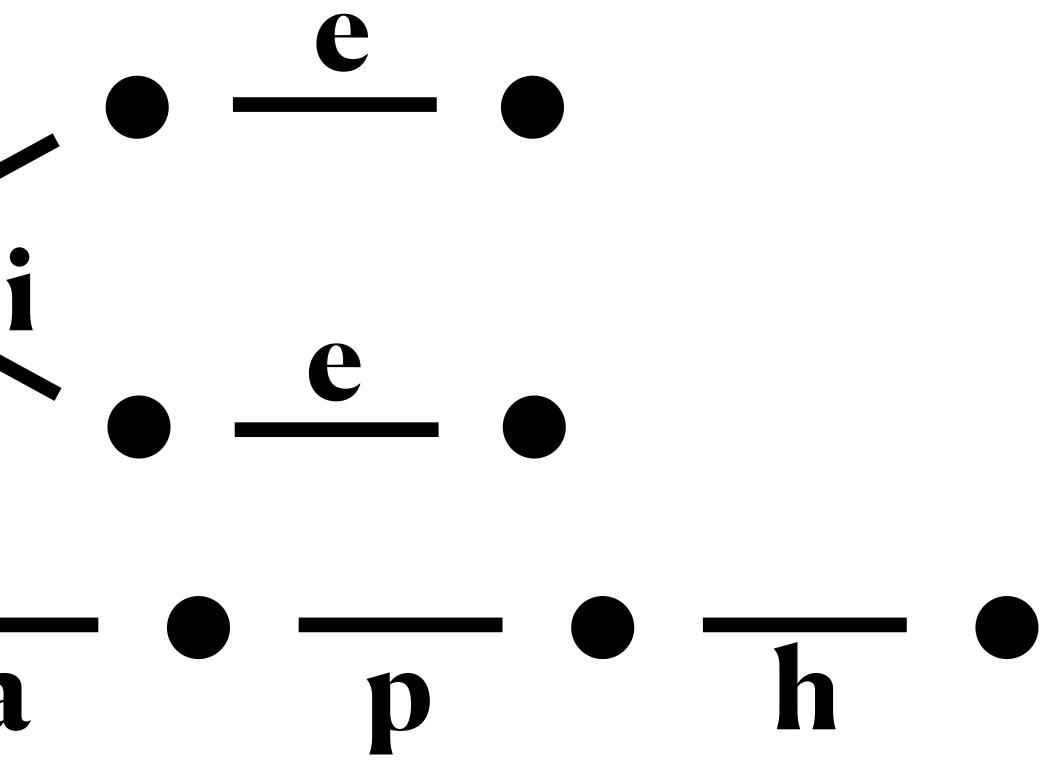
Tries: storing string patterns

$\frac{r}{t}$ $\frac{r}{t}$

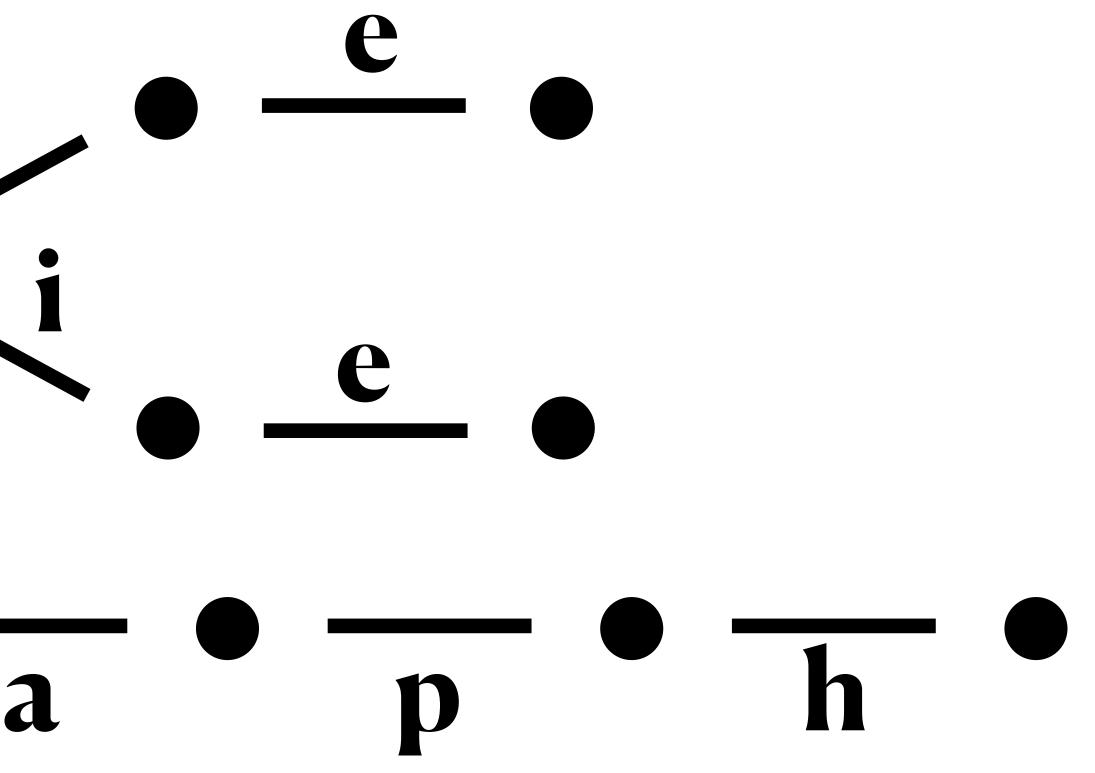
e •

Tries: storing string patterns e r 2 r 2

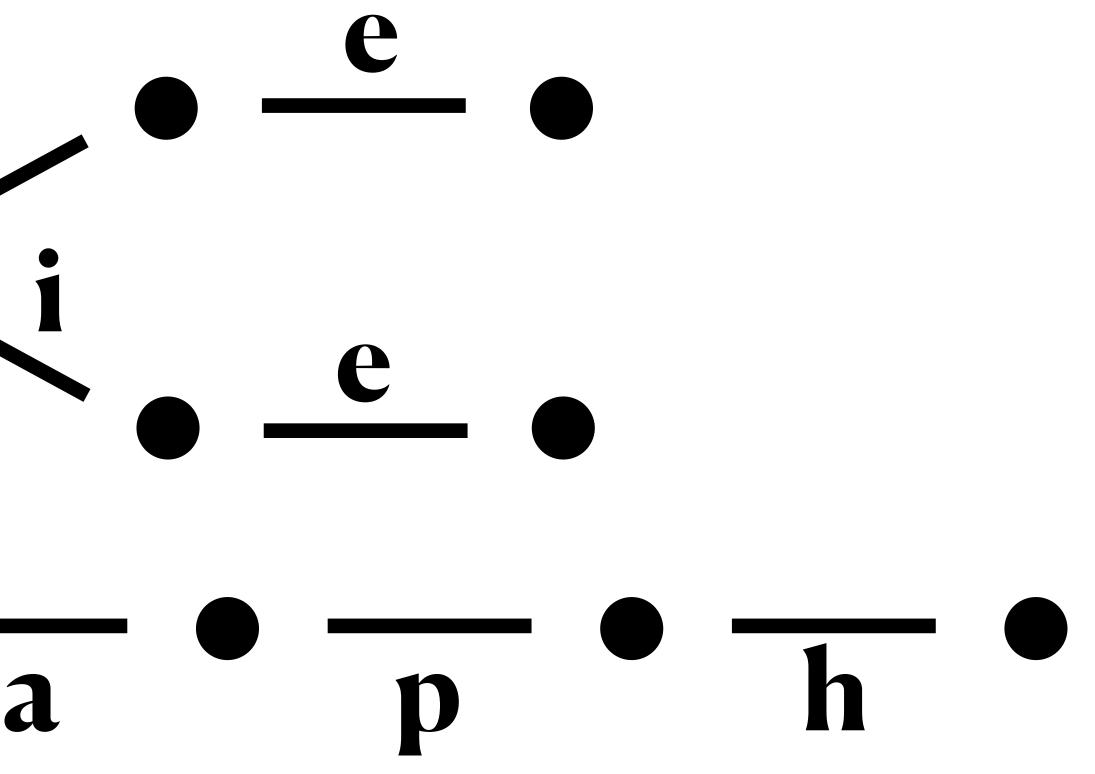
Tries: storing string patterns e r 0 r



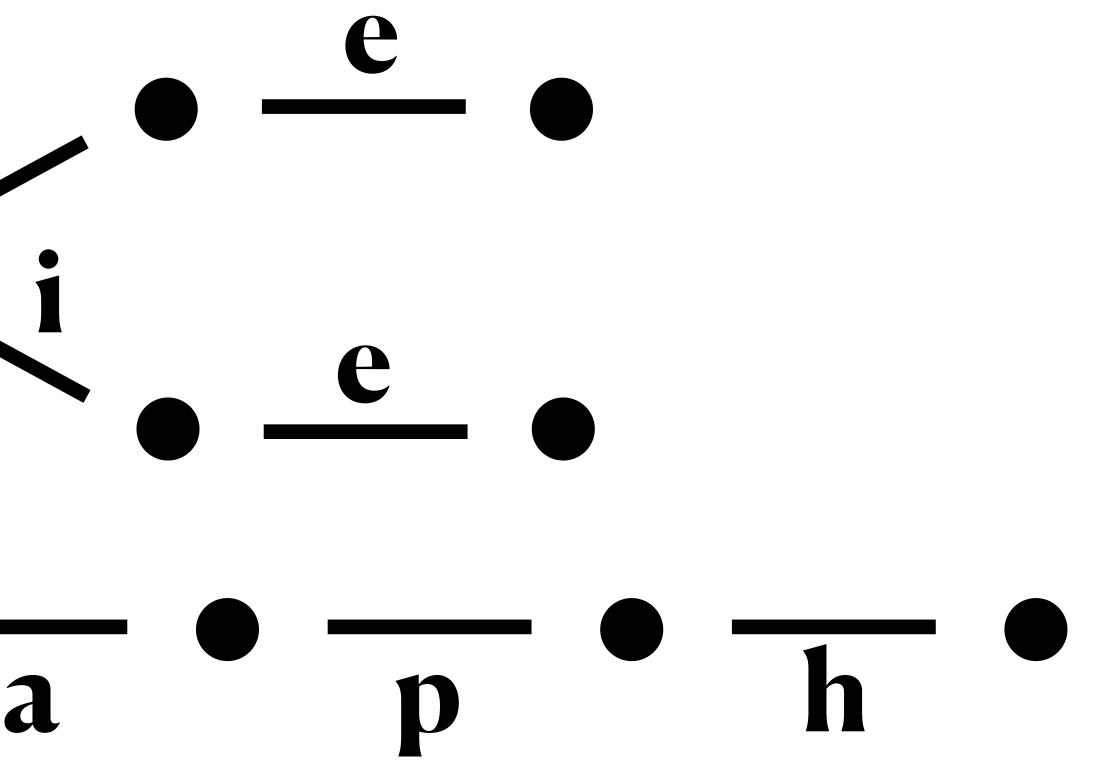
$\frac{r}{g} \cdot \frac{r}{r} \cdot \frac{e}{r}$



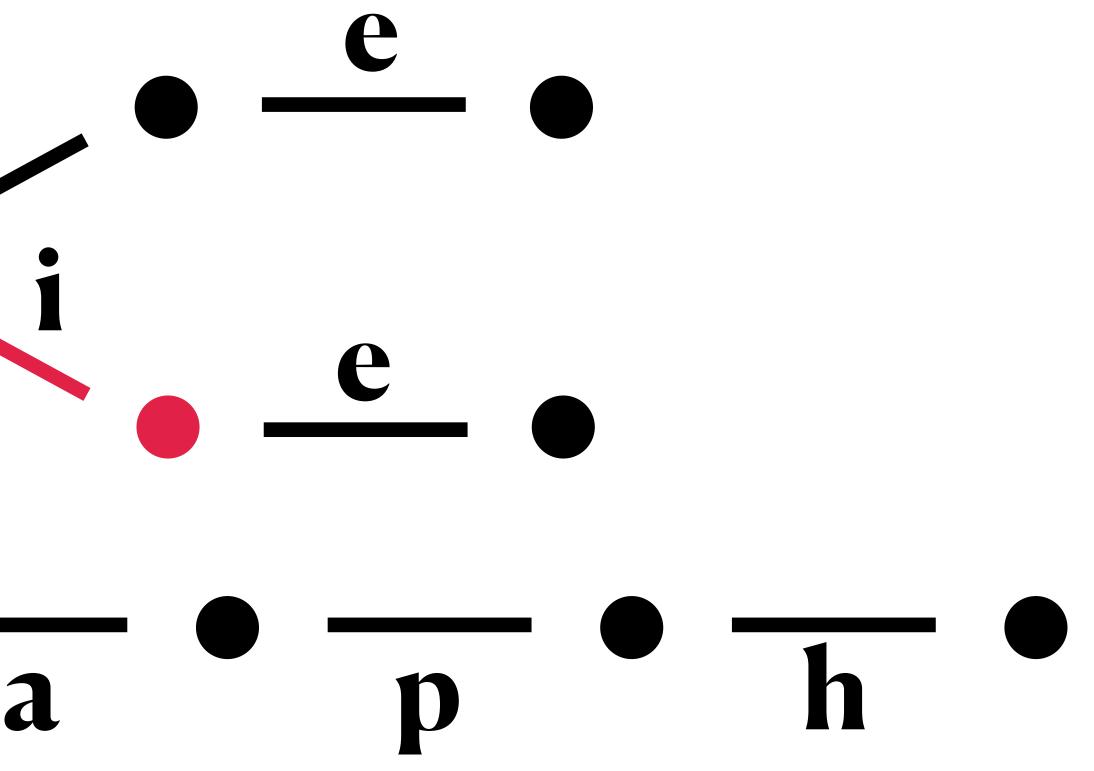
$\frac{r}{g} \cdot \frac{r}{r} \cdot \frac{e}{r}$



$\frac{r}{g} \cdot \frac{r}{r} \cdot \frac{e}{r}$

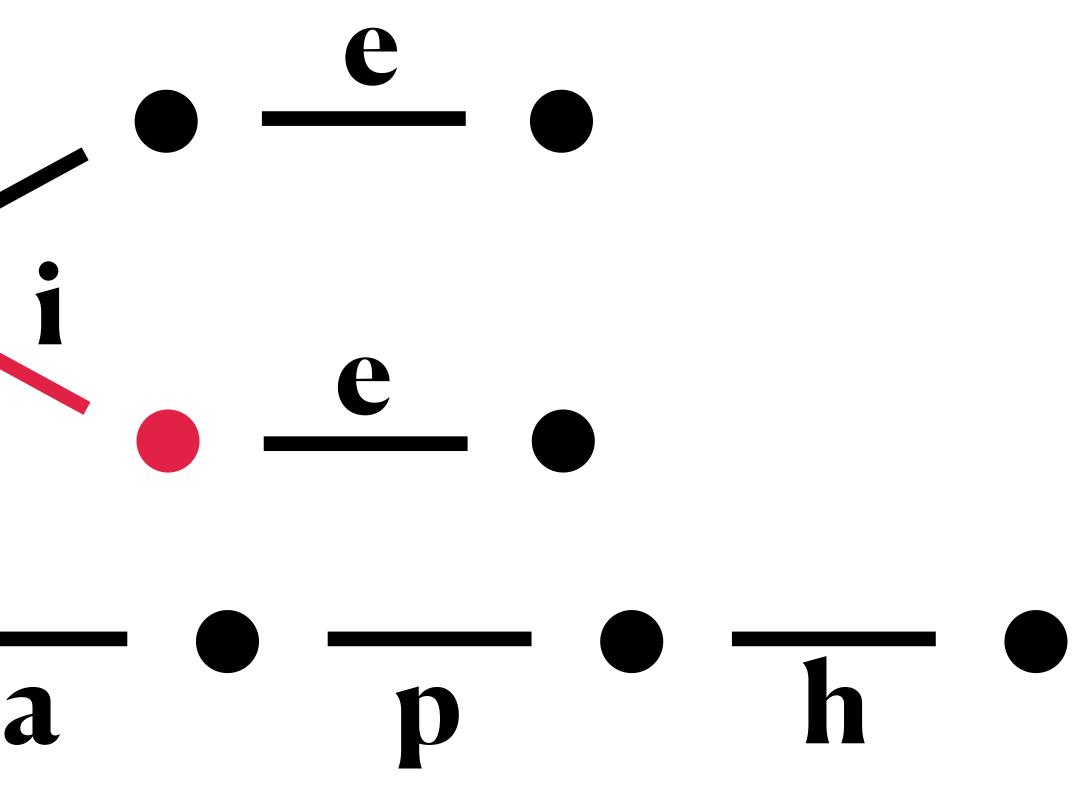


r r r g r r



Tries are deterministic finite state machines

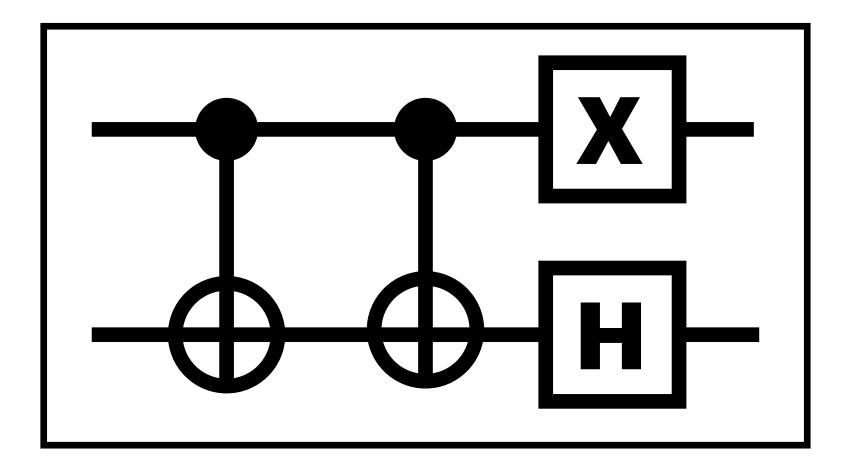
$\frac{r}{g} \cdot \frac{r}{r} \cdot \frac{e}{r}$



Tries also work for circuits!

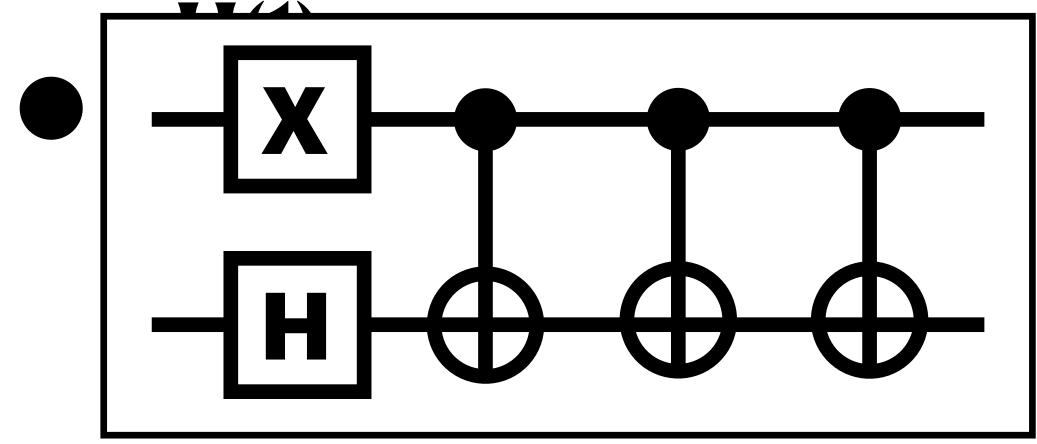
CX(0,1) • CX(0,1) • (0)

H(1)

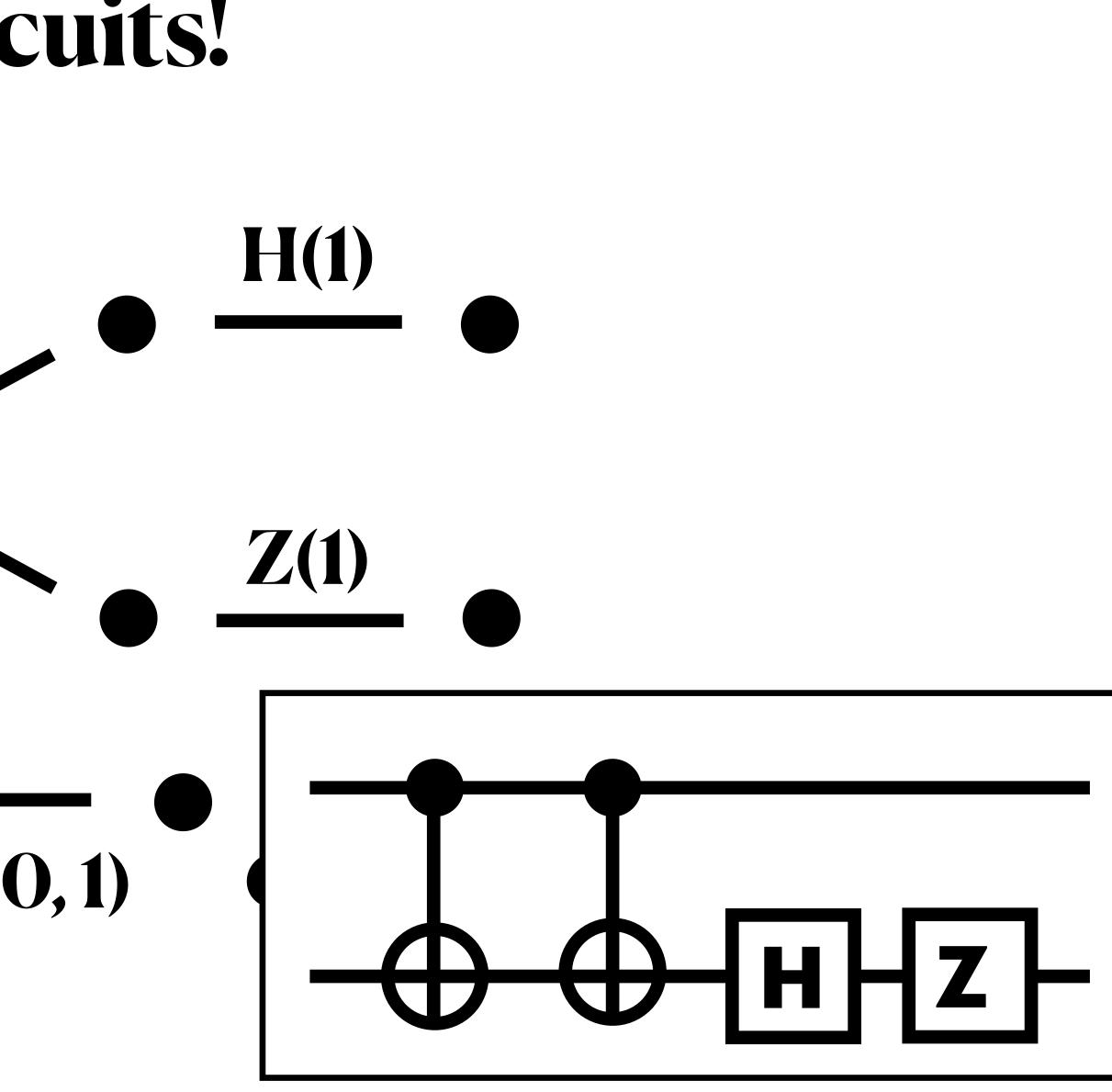


Tries also work for circuits!

X(0) CX(0,1) CX(0,1) Η **H(1) X(0) CX(0,1) CX(0,1) CX(0,1)**

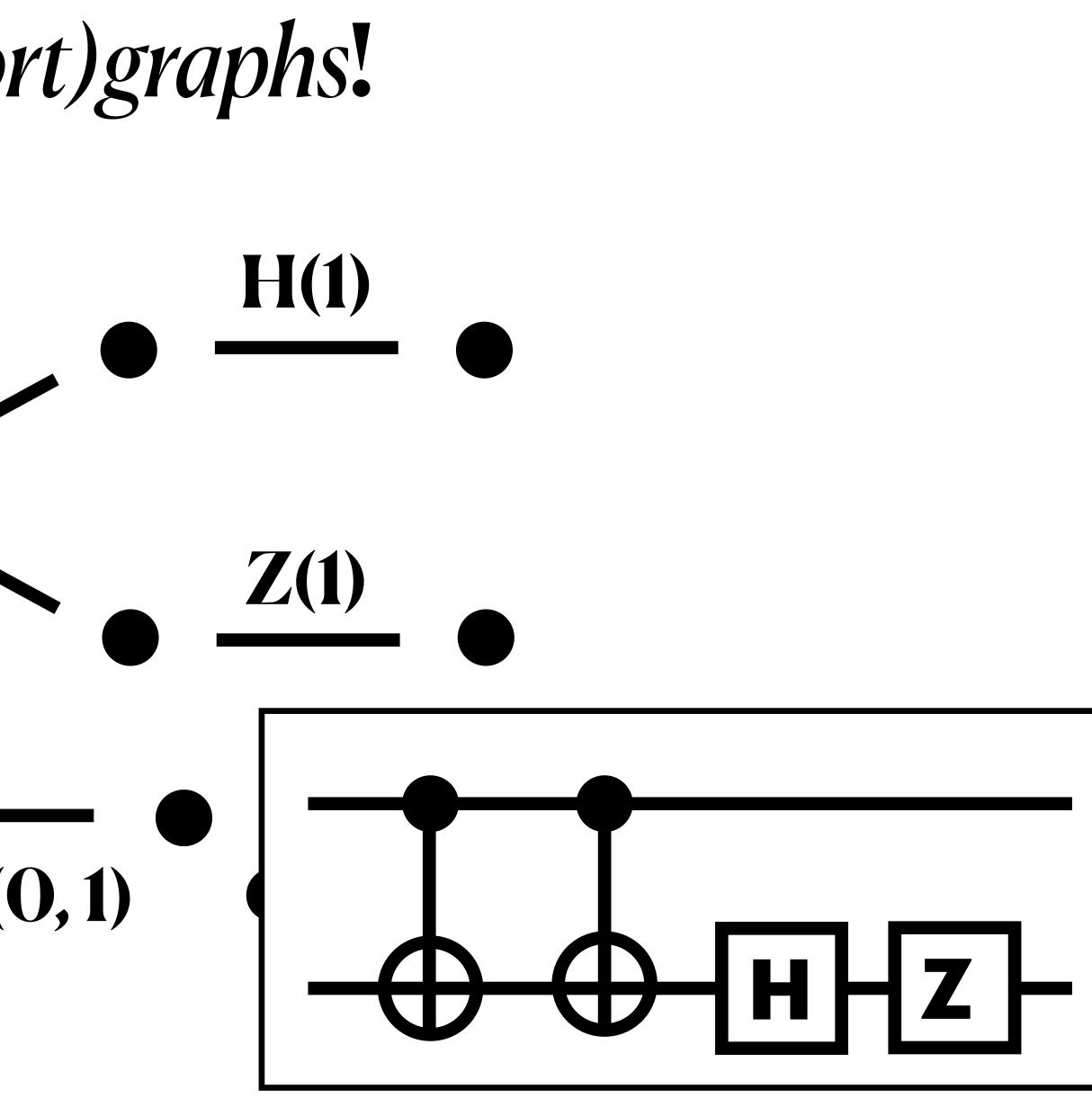


Tries also work for circuits! **X(0) CX(0,1) CX(0, 1) H(1)** H(1) **X(0) CX(0,1)**





Tries also work for (port)graphs! **X(0) CX(0,1) CX(0, 1) H(1)** H(1) **X(0) CX(0,1)**

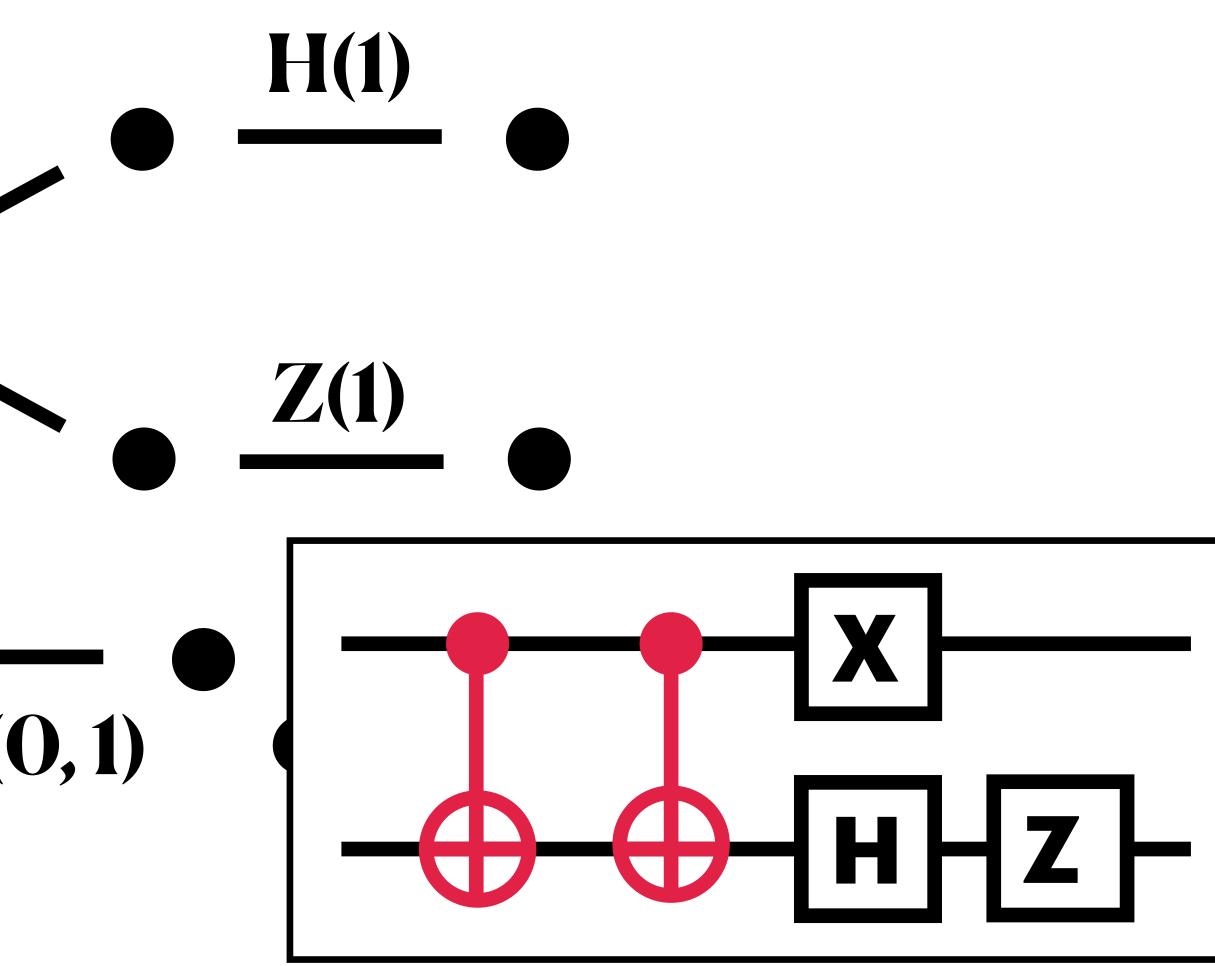




Tries also work for (port)graphs!

- No total ordering of the nodes
- No qubits
- No fixed node degrees

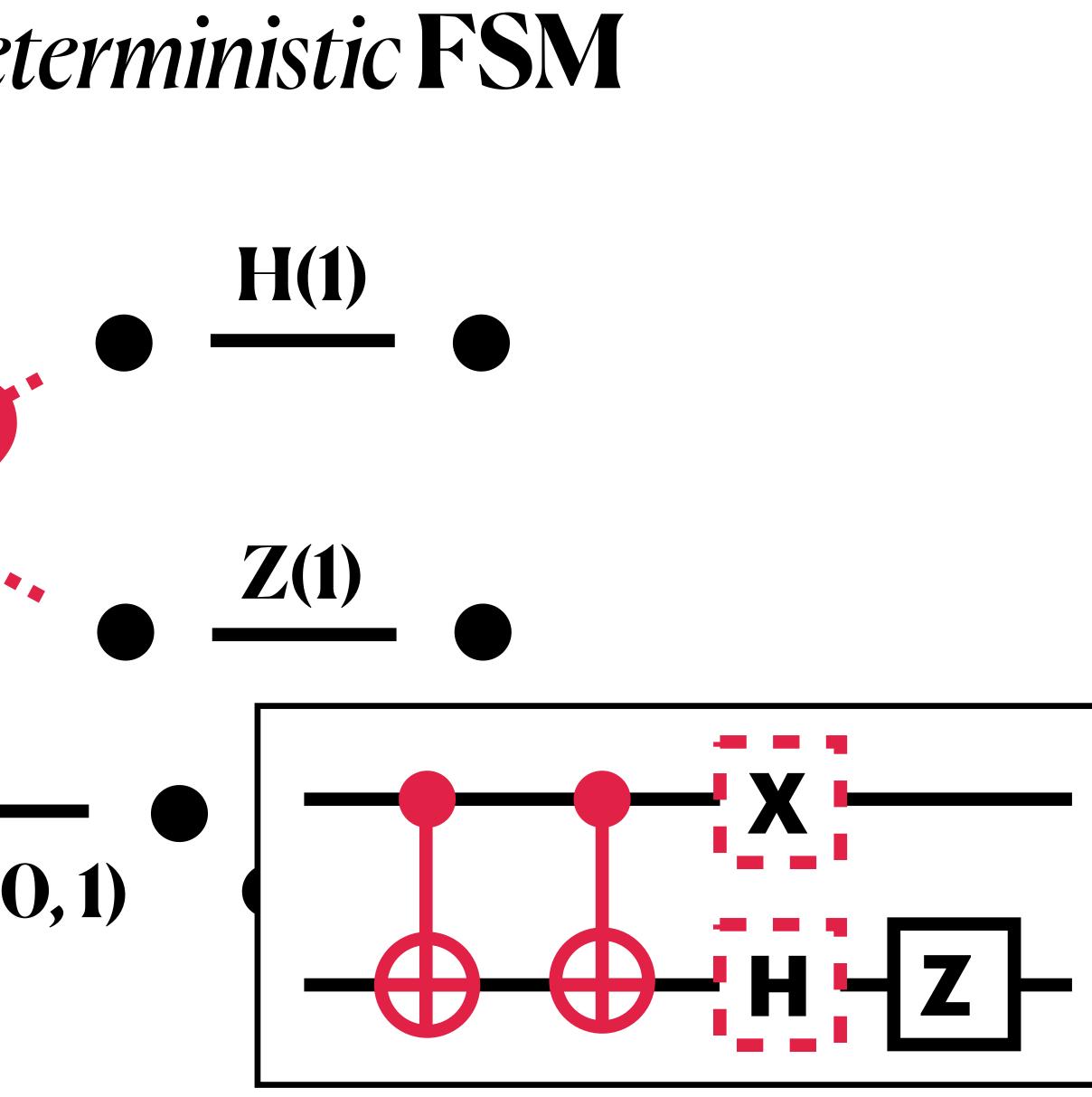
$\begin{array}{c} CX(0,1) \\ CX(0,1) \\ H(1) \\ H(1) \\ X(0) \\ CX(0,1) \end{array}$





Graph Tries are non-deterministic FSM

$\begin{array}{c} CX(0,1) \\ H(1) \\ H(1) \\ CX(0) \\ CX(0,1) \\ H(1) \\ CX(0) \\ CX(0,1) \end{array}$





Bound the number of non-deterministic states!

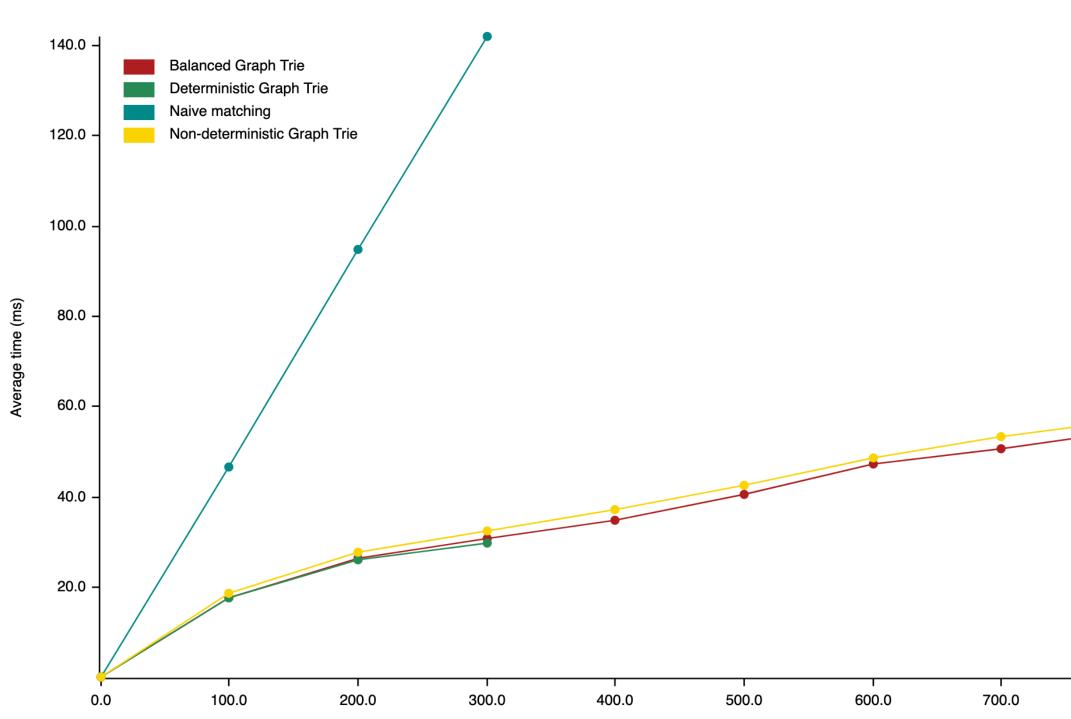
from root to leaf in the graph trie will have at most Q nondeterministic states.

Bound independent of input circuit size Graph trie depth bound by $Q \cdot d$

Thm. If every pattern circuit has at most Q qubits, then any path



All code available at github.com/lmondada/portmatching



Many Patterns Matching: Comparison

Input Size (Elements)

Please reach out! luca@mondada.net





