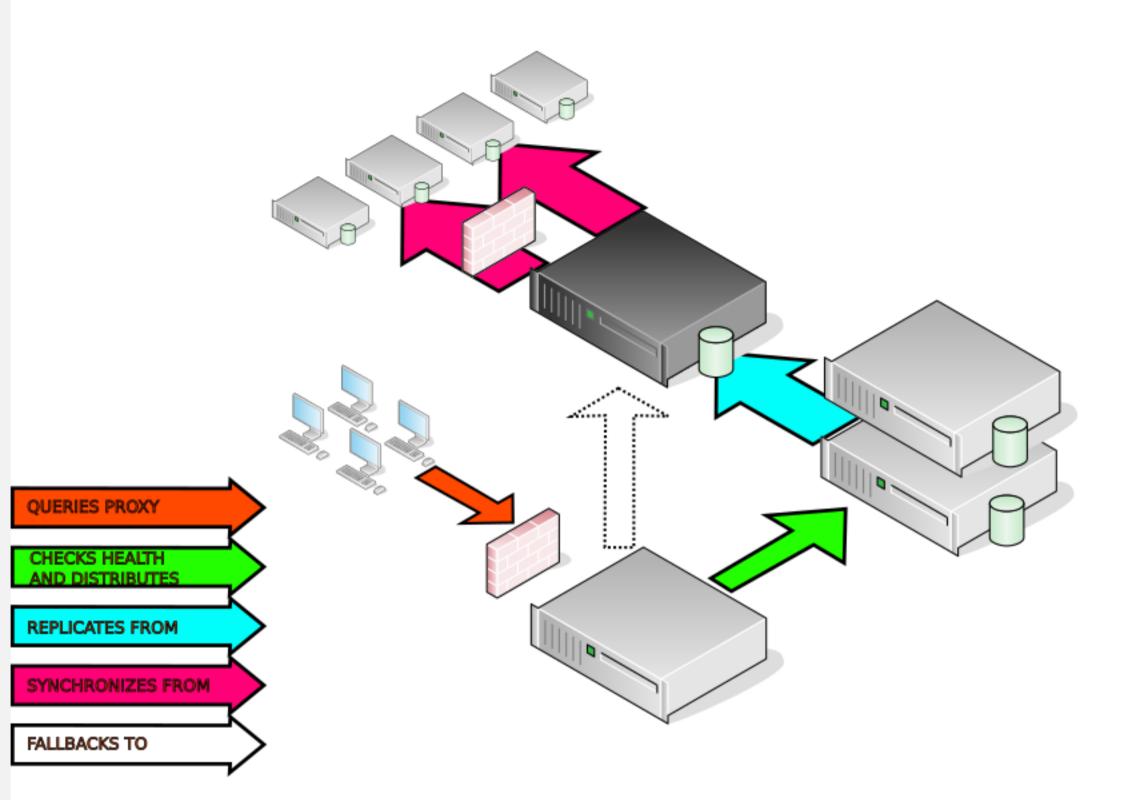
EGI's top BDII



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The EGI Foundation

Agenda

- Available options to run a top BDII
- Implementation design
- Current status and challenges
- Plan

Infra at EGI

- Managing collaborations tools
- < 30 VMs, mostly Debian
- One sysadmin

How to run BDII? (pick one)

Running CentOS and outdated software

or

Running BDII on a modern debian

or

Trying something else?

Problem: running on CentOS

- Does not fit in our infra (Debian)
- Hence, everything must be special-cased: monitoring, logging, automation, security, update cycles, ...
- Backporting is a pain
- Personal distaste & experiences
- Too much effort for me
- What were the reasons for which EGI has been asked to run this service?

Problem: using BDII on debian

- BDII is made of a <u>many</u> (integrated) pieces
- Each of those will have to be reviewed and ported accordingly
- In the long run better than having CentOS, in my opinion
- But still too much effort for me
- What were the reasons for which EGI has been asked to run this service?

More problems

- BDII Code = Favorite Italian dish?
- Most of it is not needed for a top-BDII.
- Py3 migration + text file processing = bad feelings (*)
- Not operating a site: de facto not an user, never touched a real BDII.
- Dilemma: fixing bugs in unused functionality?

(*: possibly unwarrented)

New design: goals

- Implementation simplicity
- Cost/resource effectiveness
- High integration, low maintenance
- Long term operations
- Having fun and learning something

Non-goals

- Maintaining and evolving the BDII infrastructure

Different choices

- Does <u>only</u> one thing: aggregation of remote BDIIs
- Per site(s) synchronizations (<u>small step</u>)
 - rather than a global aggregation (big step)
- Trade peak memory usage for more CPU (in theory)
- Native libldap2 calls
 - avoiding LDIF processing and encoding issues(*)
- Using newer OpenLDAP mmap db engine and on-line configuration (OLC)
 - An in-memory database rather than an in-memory filesystem
- (*: in reality, things are more complex)

Design pros

- Concurrency flattens resource usage (vs spiky batch behavior)
- Finer grained monitoring and metrics is dead simple to add on
- Luajit, libldap2 and systemd are here to stay...
 - ... for a long time
 - ... without breaking changes*
- (*: excepting systemd, maybe...)

Design cons

- Need to keep track of entries per site (for deletions)
- Complicates site removal (to avoid orphaned entries)
- LMDB engine is not a perfect fit for this workload*

(*: mitigation planned, cf. issues slides and thereafter)

Other bonus

- Read-scalability

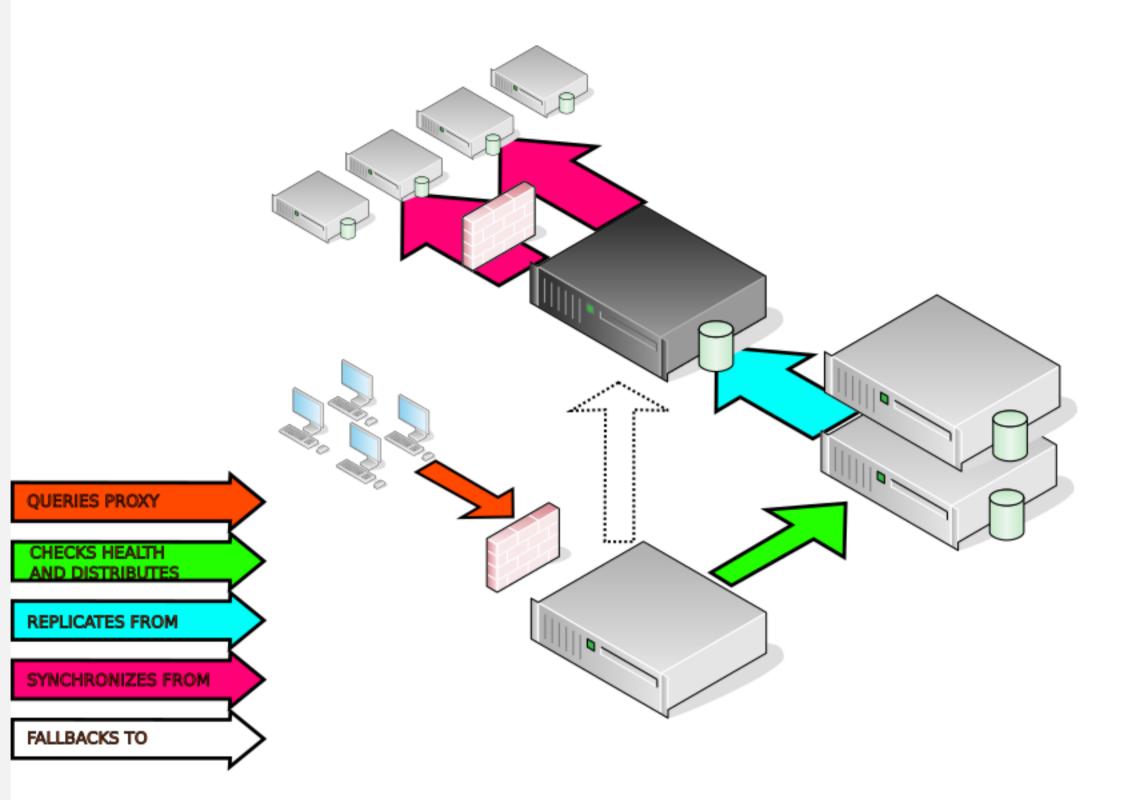
Comparison

- (talk here)

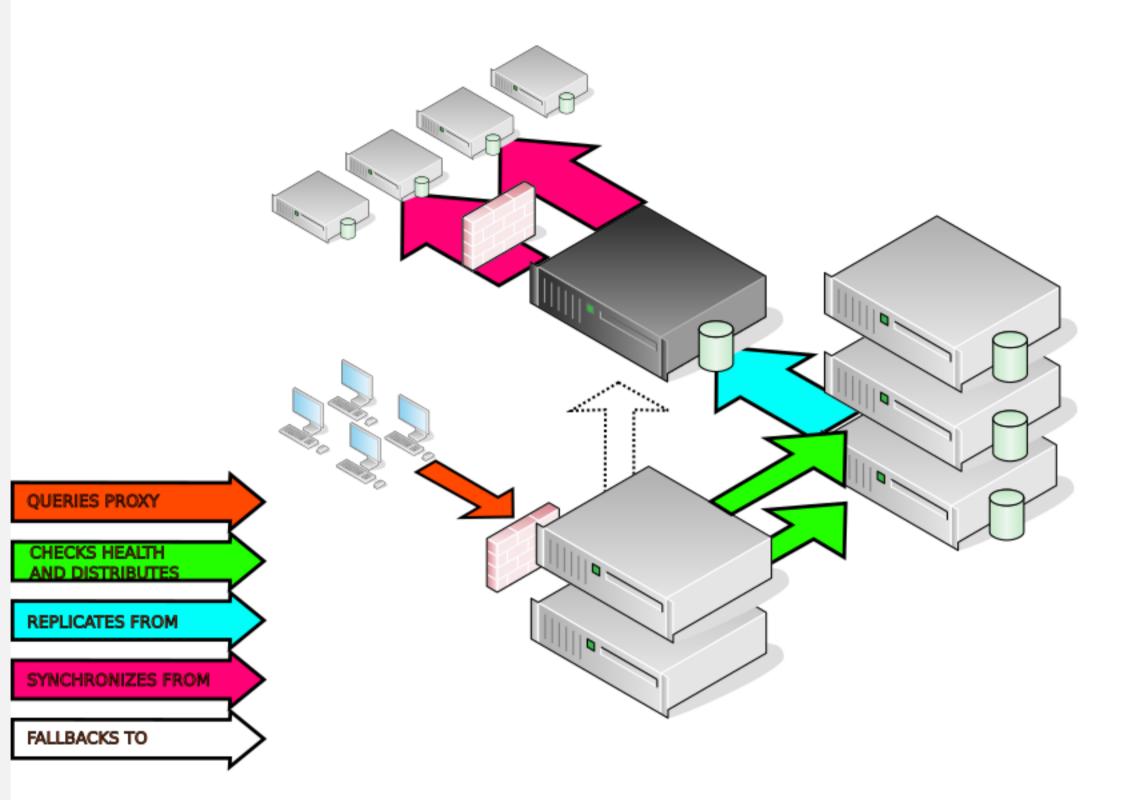
Status

- POC implementation (2nd half August)
- Public instance w/ basic optimizations (1st week of September)
- Replicated (September)
 - to provide availability and read scalability
- Just above 1 kLOC of Lua, deployed with Ansible

Architecture diagram



Possible redundance



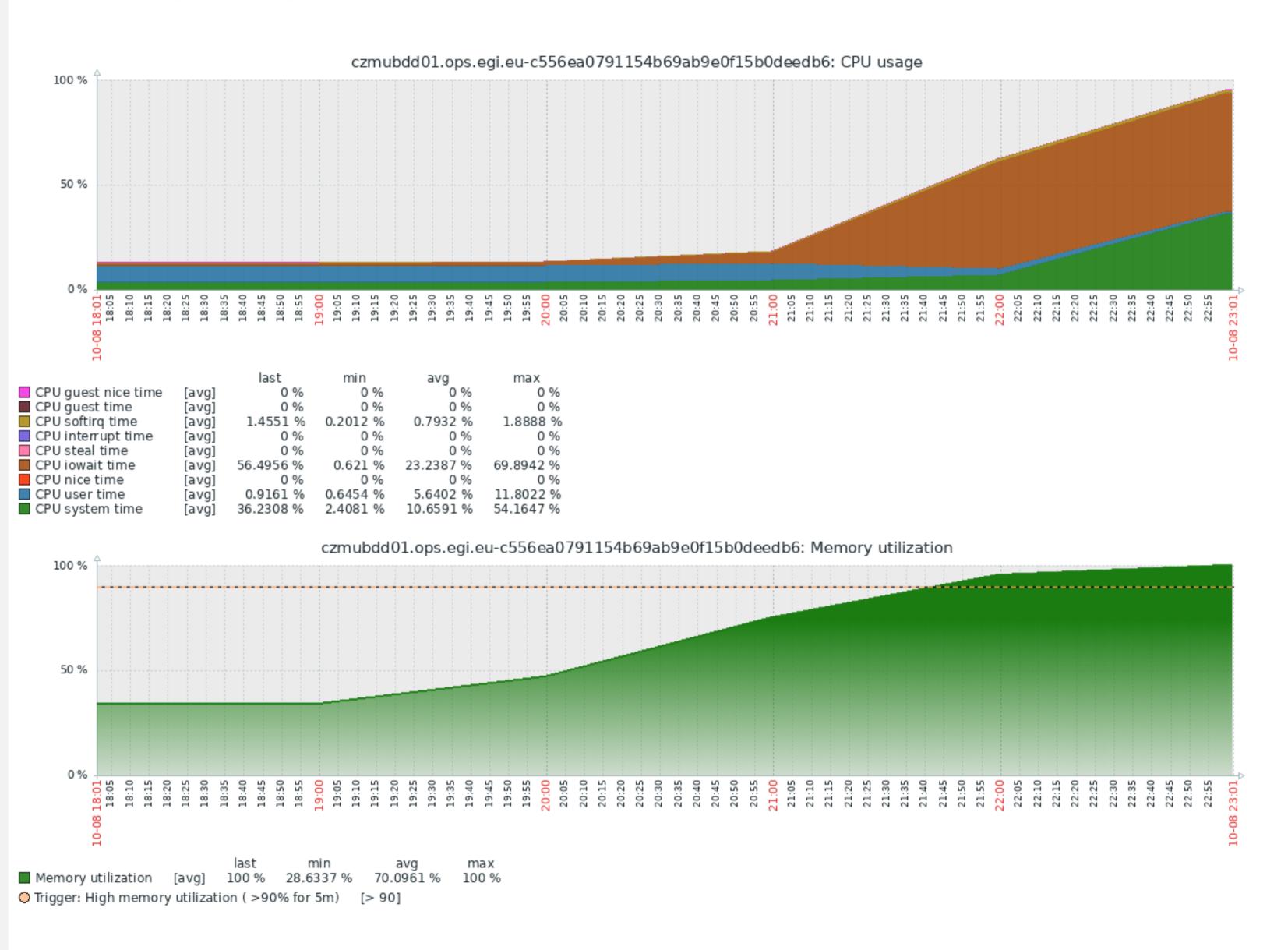
Lightweight

- Few dependencies: OpenLDAP, Luajit, systemd, xmlstarlet, sh
- HAProxy
- Master node has 2 CPU, 4GB RAM 20GB storage.

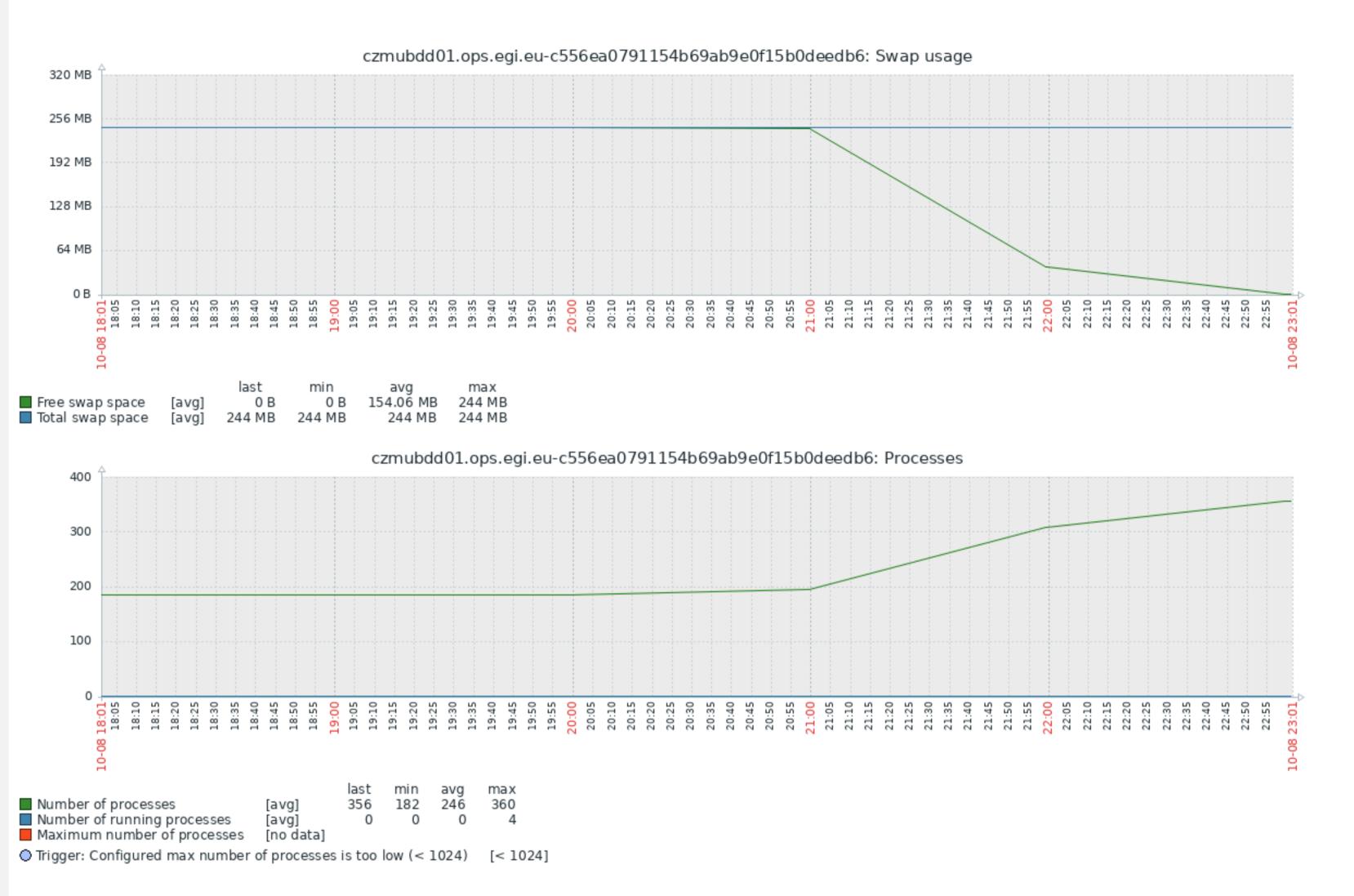
Issues

- Surprising mmap behaviour under Linux (for a non-MySQL guy)
- An interesting catastrophic failure phenomenon (see next slide)
- New software, unknown bugs, still a bit crude
- Rare LDAP errors (yet to be investigated)
- Improve and polish
- ENOTIME

Issues (cont.)



Issues (cont.)



```
Issues (cont.)
- DB growth caused by lmdb's MVCC
- Affects both changelogs and databases
- Begins to swap when core size > some % of RAM (80-90?)
  ... Then the fun starts :D

    Possible mitigation?

    smaller transactions

    less concurrency

  - adaptive scheduling: f(load), f(nconn), other?

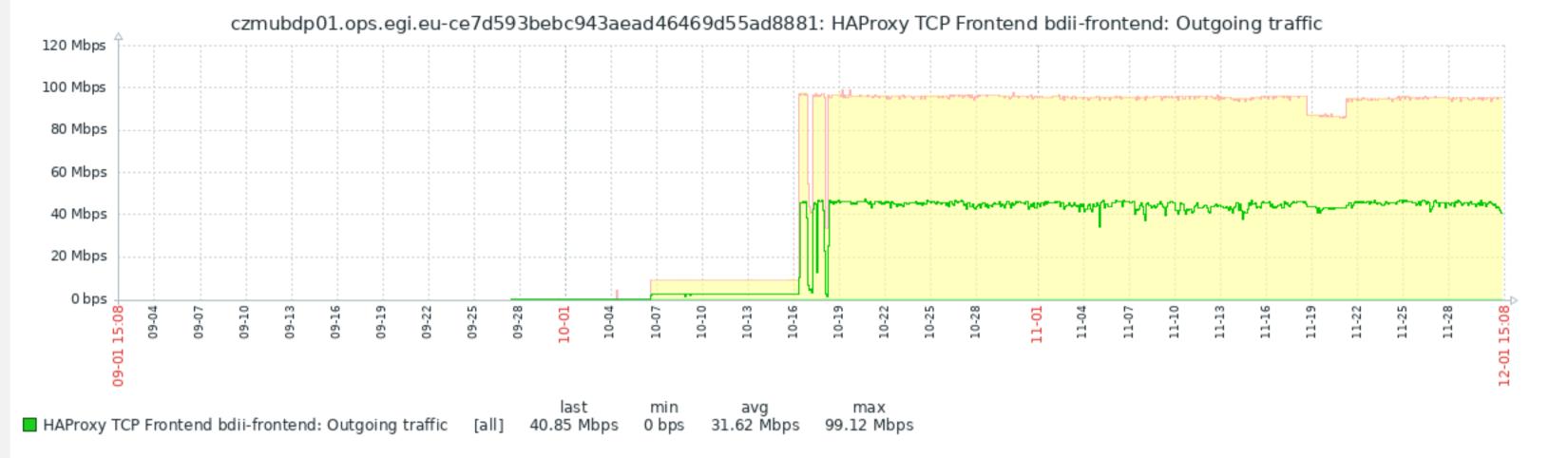
    compaction cycle

  - more RAM/$$ (*)
  (*: affording to be bothered less often, not a fix per se)
```

Issues (cont.)

- Even greater DB growth on replica
- However no swapping observed: working set always in memory.
- Hypothesis (speculation):
 - write txn are much bigger due to the way the replication stream is committed
 - pages are less fragmented (# of writer = 1)
- Unkown: how will it behave under load/clients?
- Mitigations:
 - compaction cycle
 - timelimit (?)
 - -> real observation under load is needed!

Graph: proxy outbound



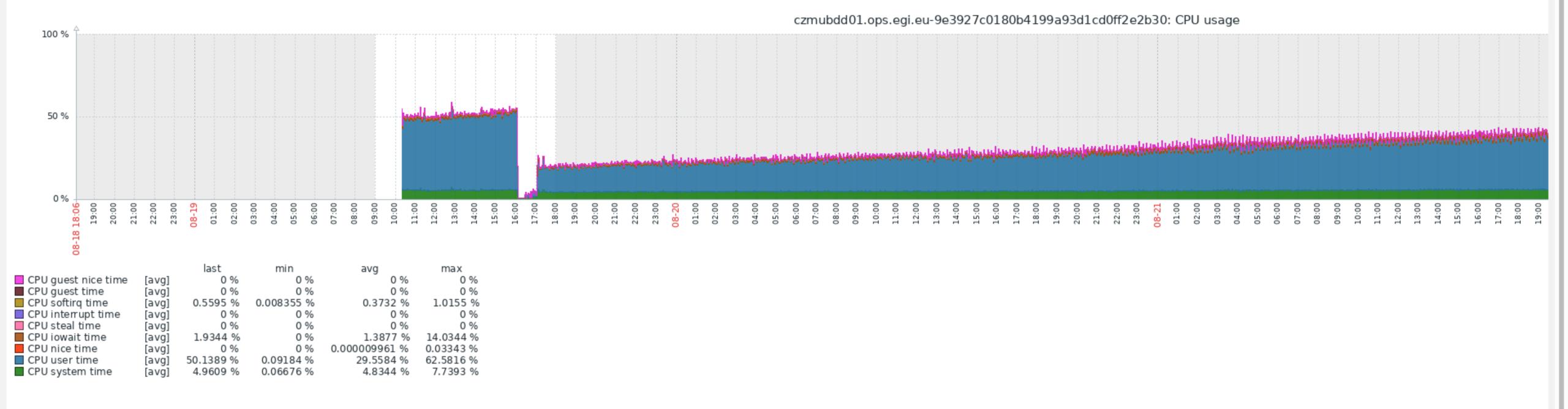
last 3 months

Graph: disk usage



master, last 3 months

Graph: no-optimization



Plan (ideally)

- Implementing concurrency mitigations
- Better systemd integration
- Metric extractions, grafana dashboard, ...



