

# Characterisation of qutrit universal gates

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# Outline

- How to estimate the performance of quantum gates?
- Why is randomised benchmarking chosen?
- Universal qutrit randomised benchmarking.

## Definition

- Three-level system:  $\text{span}(|0\rangle, |1\rangle, |2\rangle)$ .

## Advantages with respect to qubits

- Larger Hilbert space.
- More natural: avoid the truncation of a native higher level system.

## Qutrit experiments [3, 6, 14, 7, 4, 15, 5, 11, 2, 9]

- Superconductors.
- Ion traps.
- Photons (interferometer).
- More.

# Gates and groups

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 $g \in G$

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Physical implementation or gate

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 $[g] = [U_g]$

Math. description: channel  
 $\rho \mapsto \mathcal{E}(U_g \rho U_g^\dagger)$

For instance,  $\mathcal{E}$  (read calligraphic capital E  $\neq$  epsilon) a depolarising channel.

# Estimating performance = characterising

## Why?

- Estimate the performance to improve the gates.
- Know how many operations we can perform before the computation is purely random.

## Why is characterising not easy?

- Stochastic nature of quantum operations.
- Tomography (reconstruction) is expensive.

## Fidelity between states

- How similar are two states  $\rho$  and  $\sigma$ .
- Defined by  $F(\rho, \sigma) := \text{Tr} (\sqrt{\sqrt{\rho}\sigma\sqrt{\rho}})^2$ .

## Average gate fidelity of a channel $\mathcal{E}$

- How close is the output of a noisy implementation  $\rho \mapsto \mathcal{E}(U\rho U^\dagger)$  to the  $\rho \mapsto U\rho U^\dagger$  on **average**.
- The average is computed over a set of “pure” states.
- Defined as follows. Let  $X$  be the set of pure states with probability measure  $\mu$ .

$$\bar{F} = \int d\mu F(U\rho U^\dagger, \mathcal{E}(\rho)) = \int d\mu F(\rho, U^\dagger \mathcal{E}(\rho) U).$$

# Introduction to Randomised Benchmarking

## Summary of randomised benchmarking

- Goal of randomised benchmarking: estimate  $\bar{F}$ .
- RB requires: a gate set  $\mathcal{G}$ , preparation of the state  $|0\rangle$ , a measurement  $\langle 0|$ .

## Group properties

- Composition of two gates is another gate from the gate set.
- Each gate has an inverse operation.

Group elements  
 $g \in G$

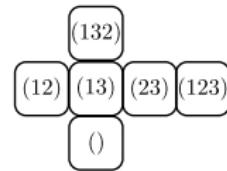
Physical implementation or gate

$$[g] = [U_g]$$

## Algorithm [10]

- ① For  $k$  in set of numbers of gates to compose (circuit depth) (commonly  $k \in \{0, \dots, 200\}$ ).
- ② For  $N$  in number of repetitions per  $k$  ( $N$  commonly  $\sim 20$ ).
  - ① Prepare  $|0\rangle$ .
  - ② Randomly sample  $k$  gates and apply them to  $|0\rangle$ .
  - ③ Generate the inversion gate (for the sequence) and apply to the state from the previous step.
  - ④ Measure  $\langle 0|$ .

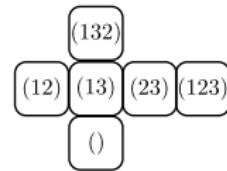
# Algorithm: toy example $S_3$



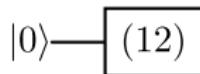
Prepare the state  $|0\rangle$

$|0\rangle$

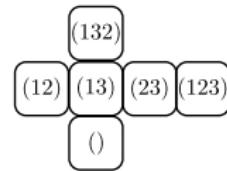
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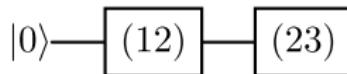
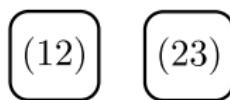
Roll the dice and apply the corresponding gate



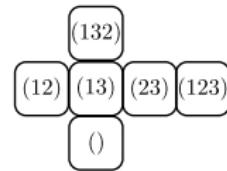
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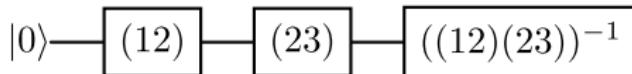
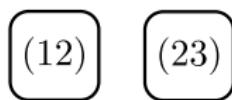
Roll a second time and apply the corresponding gate



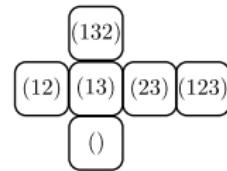
# Algorithm: toy example $S_3$



Apply the **inversion gate**



## Algorithm: toy example $S_3$

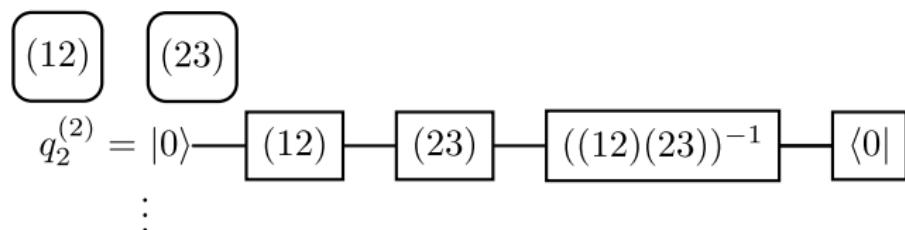
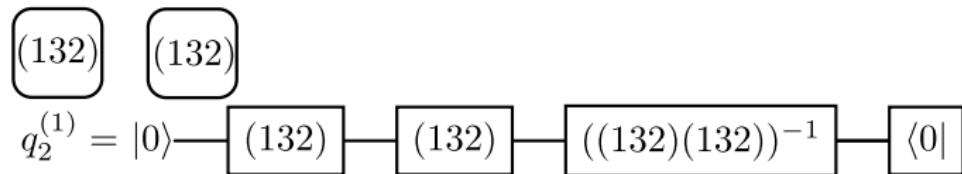


Measure with respect to the ground state  $\langle 0|$

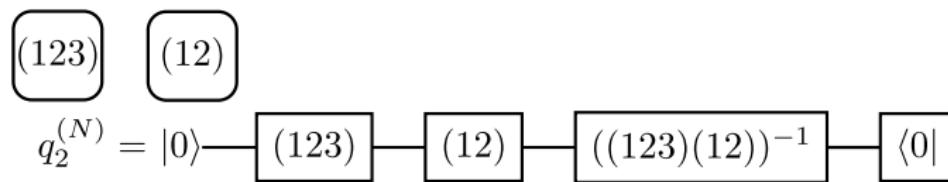
$$q_2^{(2)} = |0\rangle \xrightarrow{(12)} \xrightarrow{(23)} \xrightarrow{((12)(23))^{-1}} \langle 0|$$

## Algorithm: toy example $S_3$

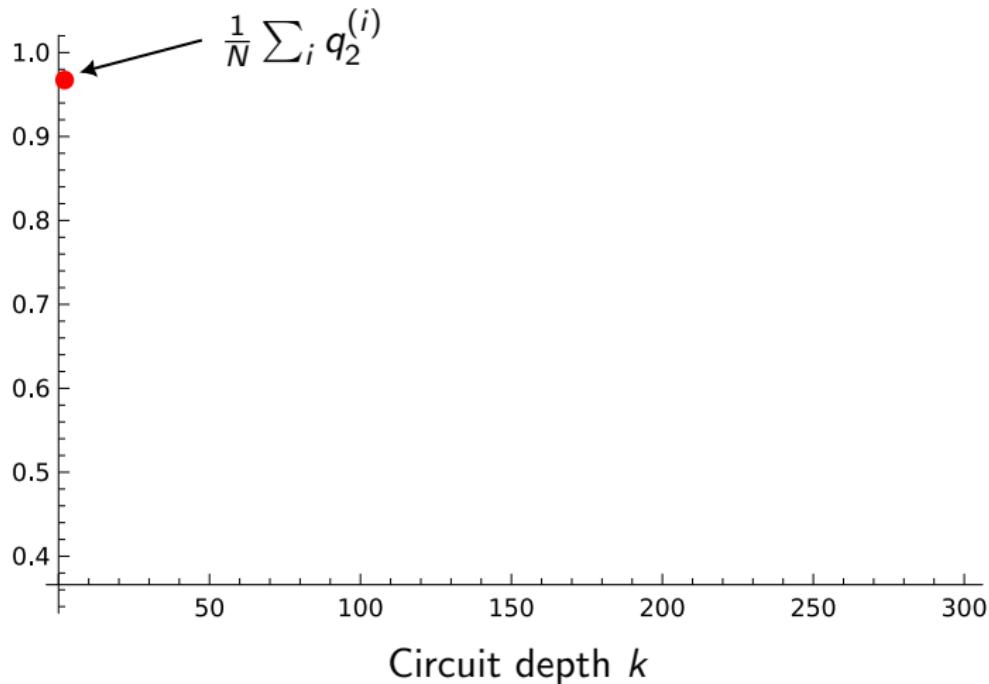
We repeat the process above  $N$  times: sample  $N$  sequences of 2 gates



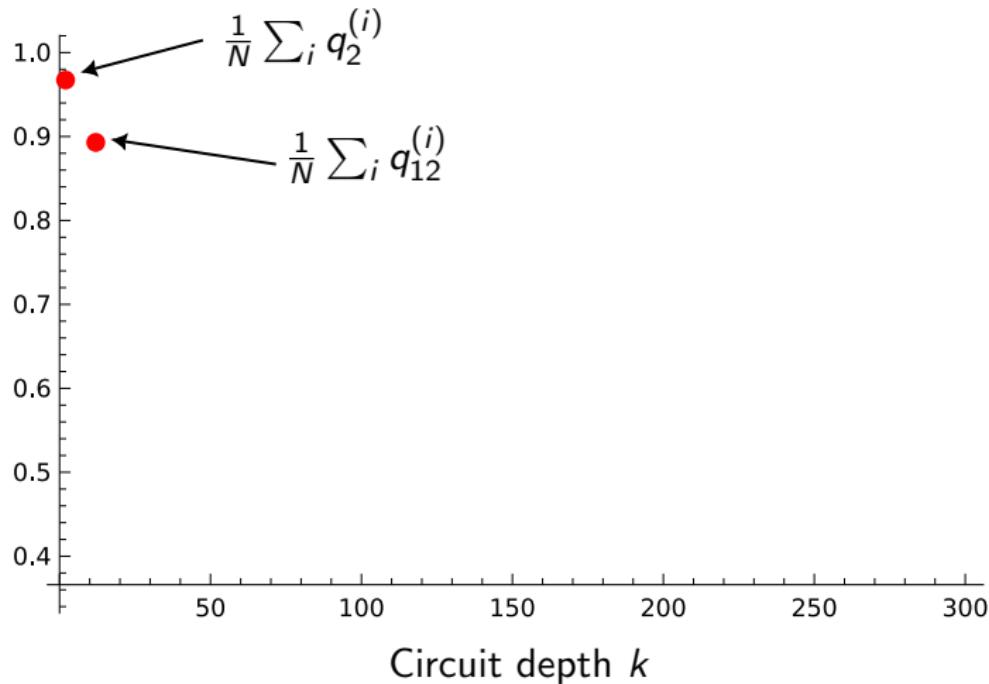
⋮



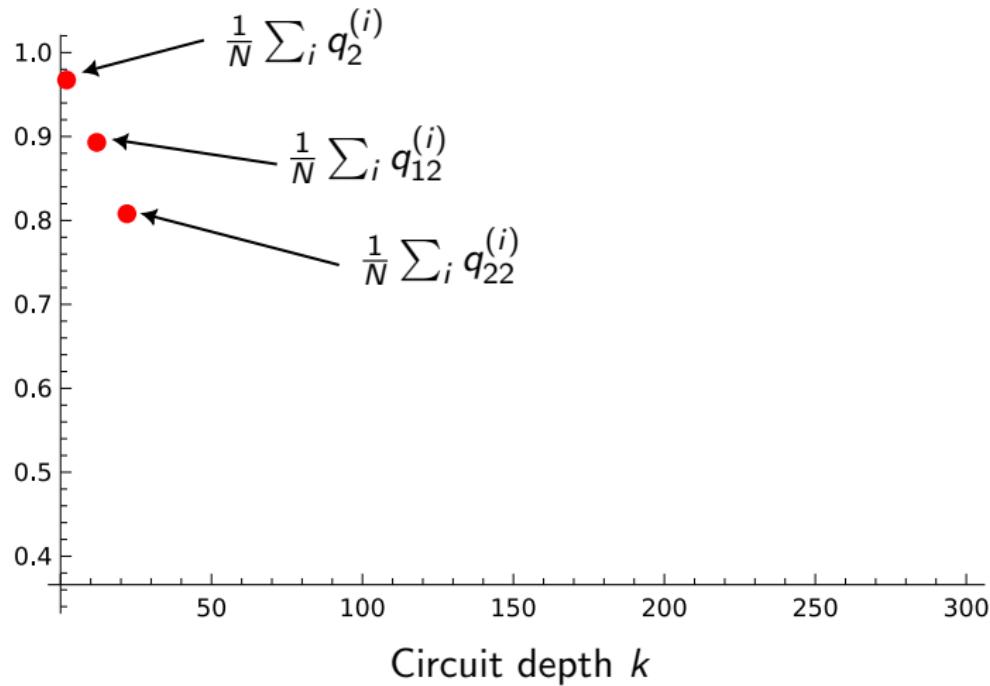
# Survival probability curve; Clifford



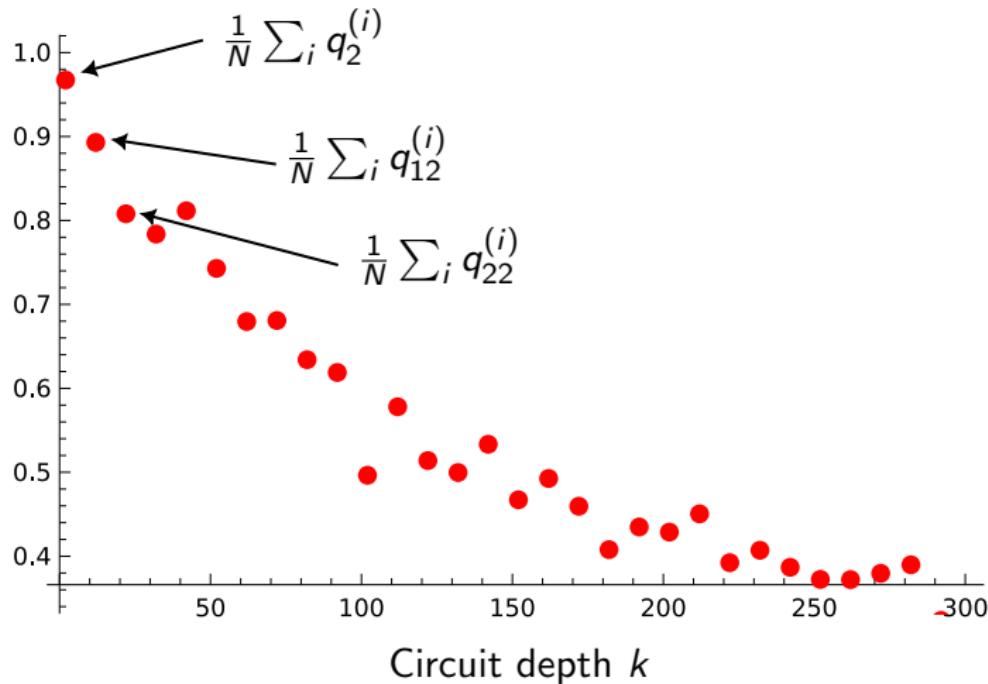
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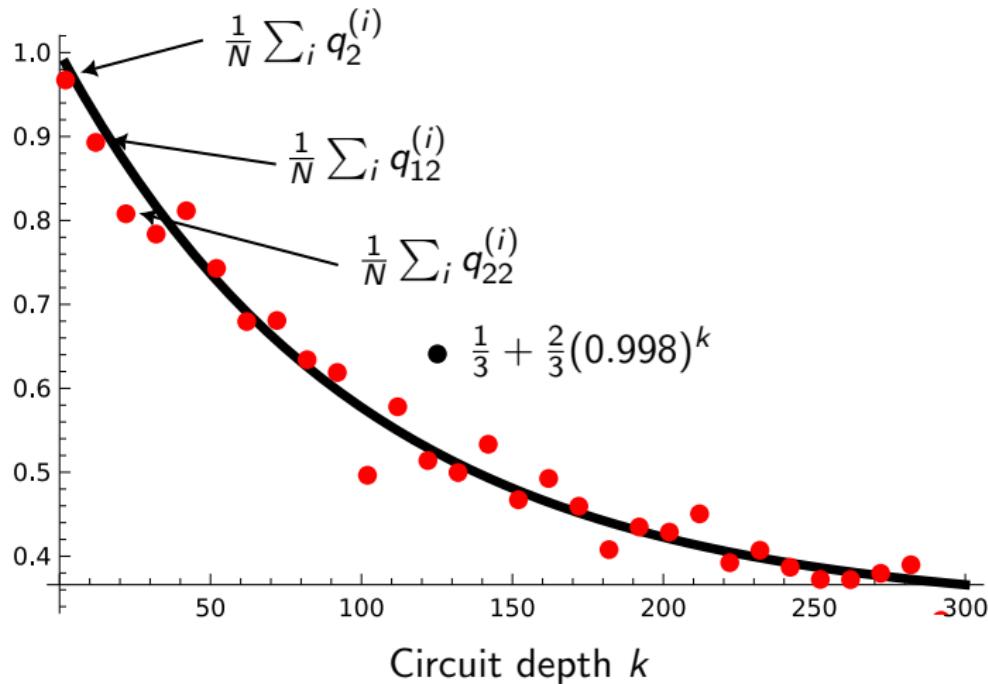
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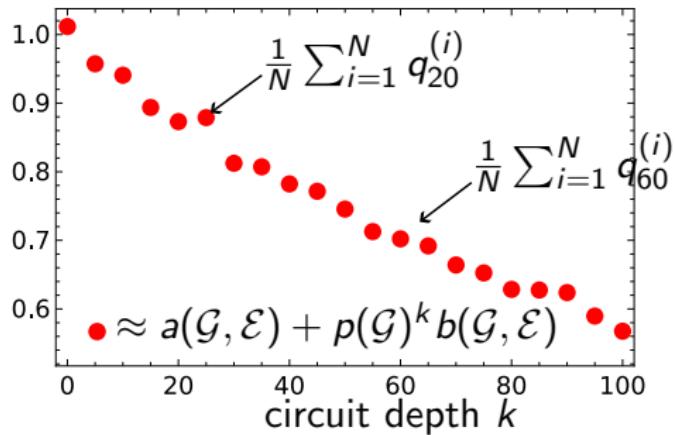


# Survival probability curve; Clifford



# Data-analysis

- Survival probability curve.



- It can be proven with some representation theory results that:

$$\bar{F}(\mathcal{G}) = \frac{3p(\mathcal{G}) - 1}{2}. \quad (1)$$

- Recall that  $\bar{F}$  is the average fidelity over a gate set.

# Representation theory

## PL representation [8]

- Given an orthonormal basis for the set of  $3 \times 3$  matrices:  $\{W_i\}$ .
- $\Gamma(\mathcal{E})_{i,j} := 3^{-1/2} \text{Tr}(W_i^\dagger \mathcal{E}(W_j))$ .
- $\Gamma$  is a representation.

## Irreps

For a unitary irrep  $\gamma$ ,  $\Gamma \cong \gamma \otimes \bar{\gamma}$ , where  $\bar{\cdot}$  denotes complex conjugate.

## Average gate fidelity [12]

For a given channel  $\mathcal{E}$ ,

$$\bar{F}(\mathcal{E}, \mathbb{I}) = \frac{3 \text{Tr}(\Gamma) + 9}{36}. \quad (2)$$

If the noise  $\mathcal{E}$  is the same for each gate-set member  $\bar{F}(\mathcal{E}, \mathbb{I}) = \bar{F}(\mathcal{G})$ .

# Survival probability theoretical curve

It can be shown that:

$$\sum_{\mathbf{g} \in G^k} \langle 0 | \tilde{U}_{\text{inv}(\mathbf{g})} \tilde{U}_{g_1} \cdots \tilde{U}_{g_k} | 0 \rangle = \langle\langle 0 | \mathcal{T}(\mathcal{E})^k | 0 \rangle\rangle, \quad (3)$$

where  $\tilde{U} := \rho \mapsto \mathcal{E}(U\rho U^\dagger)$  and

$$\mathcal{T}(\mathcal{E}) := \sum_{g \in G} \Gamma(U_g)^\dagger \Gamma(\mathcal{E}) \Gamma(U_g). \quad (4)$$

The matrix  $\mathcal{T}$  is known as the **twirl** of  $\mathcal{E}$ .

# Pros and cons for randomised benchmarking

## Pros

- Resistant against state preparation and measurement (SPAM) errors.
- Scales(\*) with the number of qutrits.
- Simple data analysis.

## Cons

- Info provided by  $\bar{F}$  is sometimes incomplete.
- SPAM independence only under special cases.
- Independent sampling of gates is difficult/unrealistic.
- Limited (explicit case) to the Clifford case.

## Normaliser

Let  $N$  and  $G$  groups such that  $N \leq G$ .  $G$  **normalises**  $N$  if  $\forall g \in G, h \in N$ ,  $ghg^{-1} \in N$ .

## Pauli group for qutrits

$X$  and  $Z$  matrices and phases  $\omega := \exp(2\pi i/3)$ .

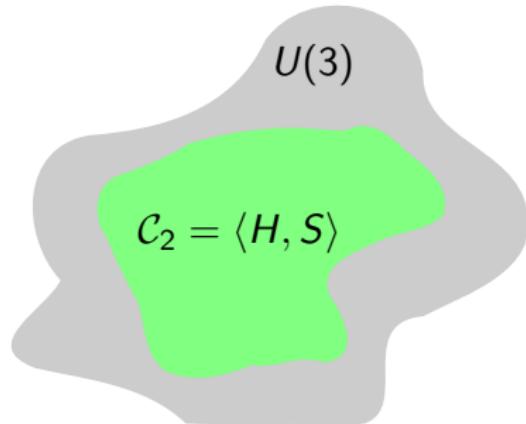
## Characteristics

- Normaliser of the Pauli group.
- Can be efficiently simulated on a classical computer.
- Cannot efficiently approximate an arbitrary unitary.

# Clifford is not universal

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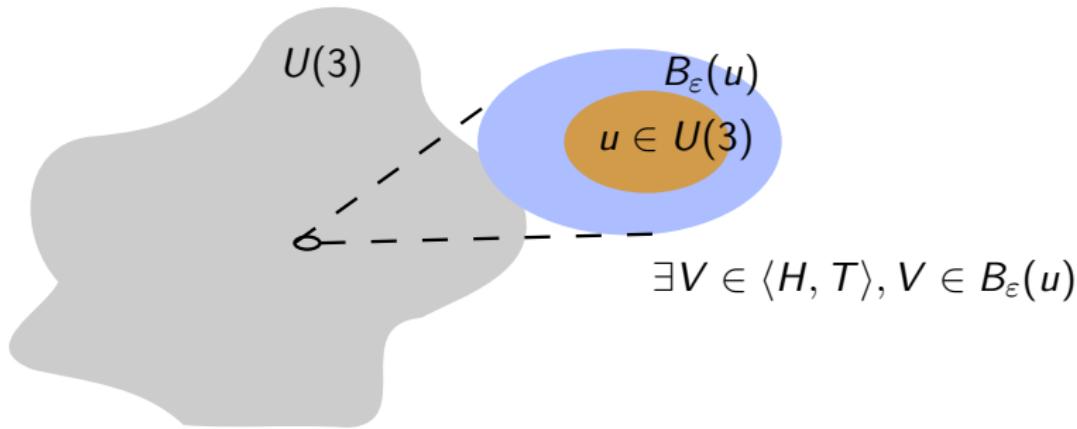


$$H := 3^{-1/2} \begin{bmatrix} 1 & 1 & 1 \\ 1 & \omega & \omega^2 \\ 1 & \omega^2 & \omega \end{bmatrix}, S := \begin{bmatrix} 1 & & \\ & 1 & \\ & & \omega^2 \end{bmatrix}. \quad (5)$$

# Universal gates

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- Gates with up to three qutrits.
- Finite set of gates approximate up to arbitrary accuracy any evolution.



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## Efficient

- Polylog with respect to the inverse of the accuracy  $\varepsilon$ .
- Meaning, increasing one digit of precision in accuracy increases polynomially.

## Solovay-Kitaev theorem [13]

- Efficient computation of the circuit.
- Efficient circuit depth.

# T gates

## T gate

- A T gate is any gate that is a generator of a universal gate set which is non-Clifford.
- Conjugating Pauli matrices by a T gate produces a Clifford gate.
- For instance,

$$T = \begin{bmatrix} 1 & & \\ & \exp(2\pi i/9) & \\ & & \exp(2\pi i/9)^8 \end{bmatrix}; \quad T^9 = \mathbb{I}. \quad (6)$$

# The HDG: characterisation of T (my contribution)

HDG stands for Hyperdihedral group.

## Motivation

- Generalise dihedral benchmarking [1].
- Fewer group elements than the Clifford group.
- Smallest number of parameters in the expression for  $\bar{F}$ .

## Construction

- Every product of  $X$  and  $T$ .

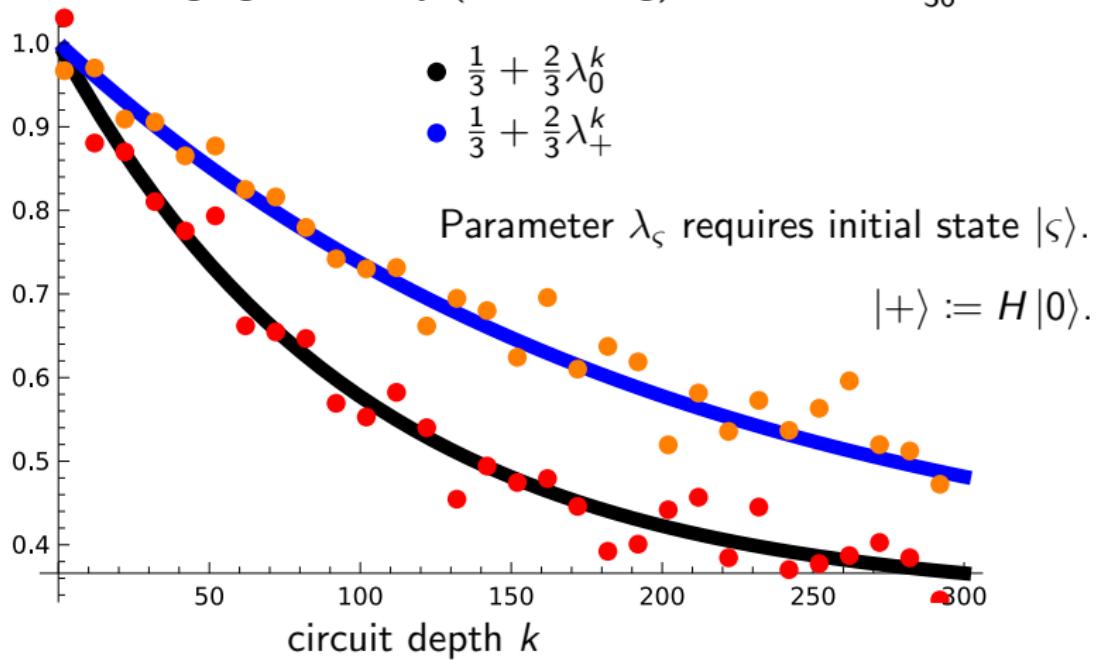
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$$X := \begin{bmatrix} 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \end{bmatrix}, T = \begin{bmatrix} 1 & & \\ & \exp(2\pi/3) & \\ & & \exp(2\pi/3)^8 \end{bmatrix}. \quad (7)$$

- $X^x T^\alpha T'^\beta$ ,  $x \in \{0, 1, 2\}$  and  $\alpha, \beta \in \{0, \dots, 8\}$ .
- $T' := XTX^\dagger$ .

## Resulting scheme

The average gate fidelity (for the hdg) is  $\bar{F} = \frac{9+3(1+2\lambda_0+6\lambda_+)}{36}$ .



# Resulting scheme

## Order

- Has order 81 (Clifford has order 216) [16].

## Satisfies imposed conditions

- Fewer gates than Clifford.
- Only two parameters:  $\lambda_0$  and  $\lambda_+$ .

## Outlook

- The HDG is the group  $C_3 \ltimes C_9 \times C_9$ .
- (Semidirect product structure) efficient sampling.
- Allows the use of induced representation results.
- Can be extended to prime level systems.

# Conclusions

- Estimating the performance of quantum gates using the average gate fidelity.
- Randomised benchmarking is the Academy and industry standard for the performance quantification of quantum gates.
- I presented the generalisation, mainly done by the introduction of a group HDG, of the randomised benchmarking scheme to characterise universal qutrits gates.



## Collaborators

- Dr. Hubert de Guise (Lakehead University).
- Dr. Barry C. Sanders (University of Calgary).

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# Why two parameters?

- $\Gamma(\text{HDG}) = \Sigma_I + \Sigma_0 + \Sigma_0^* + \Sigma_+ + \Sigma_+^*$ .
- Each irrep has associated a parameters in the average gate fidelity.
- For high-fidelity gates, the phase is negligible.