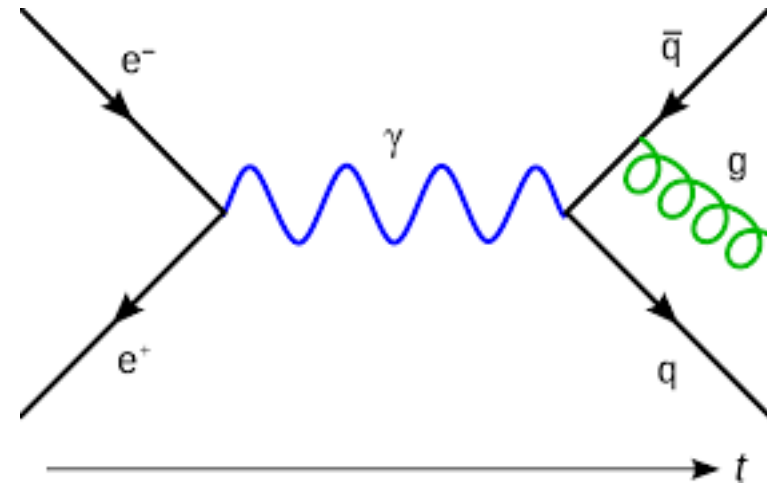
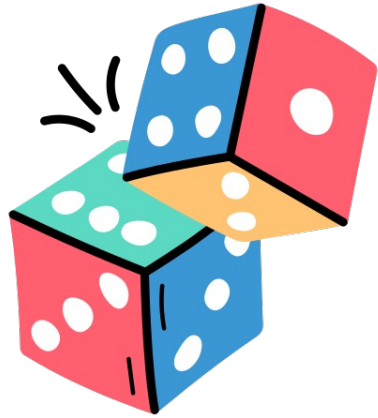


# Stochastic quantum simulations for scattering experiments



Oriel Kiss, Michele Grossi and Alessandro Roggero

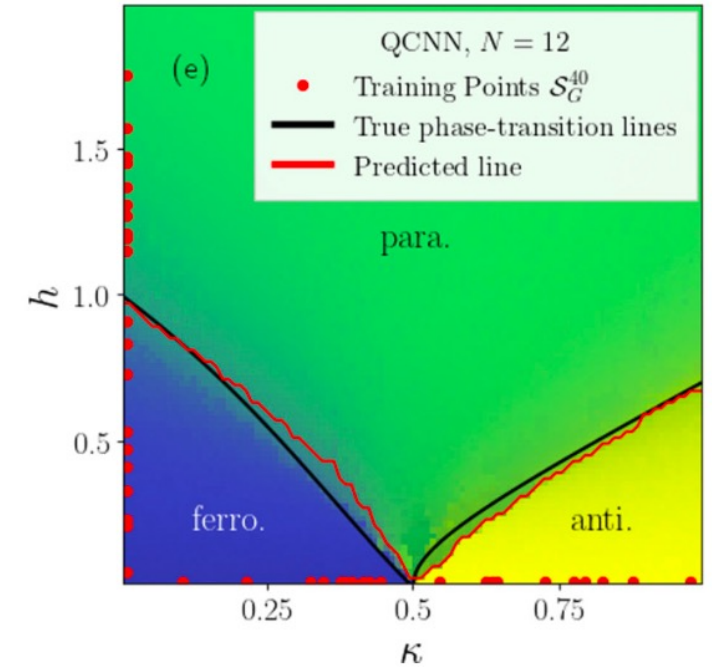
QTI Lecture, 5th April 2023

# Quantum Simulation

- Use quantum computers to simulate quantum physical systems via evolving the Hamiltonian in time (Schrödinger equation).
- Quantum many-body systems are challenging for classical computers due to the exponential scaling of the Hilbert space.
- Available and useful with only  $\sim 100$  (partially) error-corrected qubits on intermediate scale quantum devices (ISQ).
- We need better algorithms: smaller depth and lower overhead.

# Applications

- Spectrum: compute energies via quantum phase estimation
- Dynamical properties: response functions, density of states, correlators, spin dynamics, etc
- Quantum phase diagram: critical points



Monaco, Kiss, et al., Phys. Rev. B 107, L081105 (2023)  
Grossi, Kiss, et al, Phys. Rev. E 107, 024113 (2023)

# Outline:

1. Stochastic product formula (QDRIFT).
2. Implementation cost reduction with importance sampling.
3. Composite channels: how to reduce the implementation cost?
4. Numerical simulations: lattice effective field theory.
5. Response functions computations from scattering experiments.

# Stochastic product formula (QDRIFT)

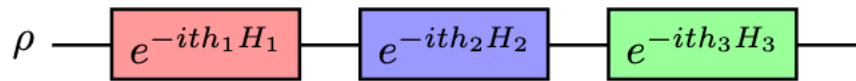
# Deterministic product formulas

**Problem:**  
scaling with  $L$

Hamiltonian  $H = h_1 H_1 + h_2 H_2 + h_3 H_3$

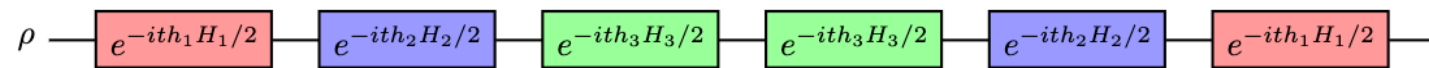
Time evolution  $e^{-iHt}$

First order: apply term sequentially



$$U_1(t) = \prod_{\gamma}^{\rightarrow} e^{-itH_{\gamma}} = U(t) + \mathcal{O}(t^2),$$

Second order



$$U_2(t) = \prod_{\gamma}^{\rightarrow} e^{-it/2H_{\gamma}} \prod_{\gamma}^{\leftarrow} e^{-it/2H_{\gamma}} = U(t) + \mathcal{O}(t^3),$$

# QDRIFT (random compiler)

We evolve terms in the Hamiltonian randomly.

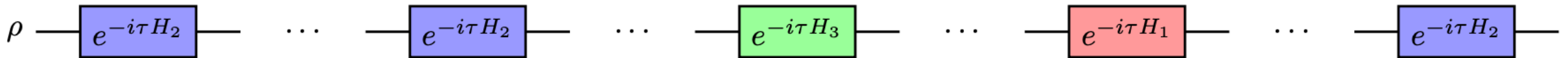


$$\begin{aligned} &h_1 H_1 \\ &h_2 H_2 \\ &h_3 H_3 \end{aligned}$$

$$p(x) = \frac{h_x}{\lambda}$$

$$\tau = \frac{t\lambda}{N}$$

$$\left( \begin{aligned} \lambda &= \sum_{\ell=1}^L h_{\ell}, \tau = \frac{\lambda t}{N} \\ H &= \sum_{j=1}^L h_j H_j \end{aligned} \right)$$



**Results:** if  $\tau = \frac{t\lambda}{N}$ , qDrift matches Trotter at first order.

$$N \geq \frac{2t^2\lambda^2}{\epsilon} \quad \text{ensures} \quad \|(\mathbb{E}V)^N - U\| \leq \frac{\epsilon}{2}.$$

Campbell, PRL **123**, 070503 (2019)

Improvement from Chen, PRX Quantum 2, 040305

# How is this working?

Do a Taylor expansion!

QDRIFT

$$\mathcal{E}(t, N)[\rho] = \sum_{j=1}^L p_j e^{-i\tau H_j} \rho e^{i\tau H_j} = \sum_{j=1}^L p_j (\mathbb{1} + i\tau [H_j, \rho]) + \mathcal{O}(\tau^2) = \mathbb{1} + \frac{it}{N} [H, \rho] + \mathcal{O}\left(\left(\frac{t\lambda}{N}\right)^2\right)$$

EXACT

$$\mathcal{U}(t/N)[\rho] = e^{-iHt/N} \rho e^{iHt/N} = \mathbb{1} + \frac{it}{N} [H, \rho] + \mathcal{O}\left(\left(\frac{t\lambda}{N}\right)^2\right)$$

ERROR

$$\|\mathcal{U}(t) - \mathcal{E}^N(t, N)\|_{\diamond} = \|\mathcal{U}^N(t/N) - \mathcal{E}^N(t, N)\|_{\diamond} \leq N \|\mathcal{U}(t/N) - \mathcal{E}(t, N)\|_{\diamond} \leq \mathcal{O}\left(\left(\frac{t\lambda}{N}\right)^2\right)$$

Sub-additivity

$$V_j(t) = \prod_{k=1}^N e^{-i\tau_{j_k} H_{j_k}}. \quad (6)$$

The qDrift channel is then built as the arithmetic average of the  $M$  experiments

$$\mathcal{E}(t; N, M)[\rho] = \frac{1}{M} \sum_m^M [V_{j_m} \rho V_{j_m}^\dagger], \quad (7)$$



# Implementation cost reduction with importance sampling

# Importance sampling for QDRIFT

**Idea?** Sample from an alternative probability distribution  $q(j)$ .

**Why?** Computational cost reduction.

**How?** Considering the implementation cost on hardware.

Importance sampling

$$\mathbb{E}_p[f(x)] = \sum_x q(x) \frac{p(x)}{q(x)} f(x) \equiv \mathbb{E}_q[\omega(x) f(x)],$$

Popular in Monte Carlo for variance reduction.

$$\begin{aligned} \mathcal{E}_q(t)[\rho] &= \sum_j q(j) e^{-i\tau_j H_j} \rho e^{i\tau_j H_j} \\ &\equiv \sum_j q(j) e^{\tau_j \mathcal{L}_j}(\rho) \\ &= \left( 1 + \sum_j q(j) \tau_j \mathcal{L}_j + \sum_{n=2}^{\infty} \sum_j q(j) \tau_j^n \mathcal{L}_j^n \right) (\rho). \end{aligned} \quad (14)$$

We match Trotterization at first order if:

$$\tau_j = \frac{th_j}{Nq_j}.$$

# What can we show for the IS QDRIFT?

Kiss et al, arXiv:2212.05952

**Bias error bound**

$$\|\mathcal{U}(t) - \mathcal{E}_q(t; N, 1)\|_{\diamond} \leq \frac{t^2 \lambda^2}{N} (1 + \mathbb{E}_p[\omega(j)]).$$

**Original QDRIFT:**

$$\begin{aligned} \mathbb{E}_p[\omega(j)] &\geq \mathbb{E}_p\left[\frac{p(j)}{p(j)}\right] \\ &= \mathbb{E}_p[1] = 1 \end{aligned}$$

**Concentration bound**

$$\begin{aligned} &\Pr \left[ \left\| \mathcal{E}_q(t; N, M) - \mathbb{E}_q \left[ \prod_{k=1}^M e^{-i\tau_k H_k} \right] \right\| \geq \epsilon/2 \right] \\ &\leq 2^{n+1} \exp \left\{ -\frac{NM\epsilon^2}{11t^2\lambda^2(1 + \max_k \omega(k))^2} \right\}. \end{aligned}$$

**Price to pay:**  
Increase in the number of samples  $N$ .

**Efficient parallelization because of concentration**

To be  $\epsilon$  close with probability  $(1 - \delta)$ :

$$NM = 11 \frac{t^2 \lambda^2}{\epsilon^2} \left( 1 + \max_k \omega(k) \right)^2 (n + 1) \log \left( \frac{2}{\delta} \right)$$

# Computational cost reduction

$$\lambda_c = \sum_l \frac{h_l}{C_l},$$

$$q_c(j) = \frac{h_j}{C_j \lambda_c}.$$

implementation cost of the generator  $H_j$ .

**Theorem:**

$$N_{q_c} \mathbb{E}_{q_c}[C] \leq N_p \mathbb{E}_p[C].$$

- Number of two-qubit native gates (CNOT).
- Connectivity (length of the Pauli string).
- Non locality.
- On error-corrected devices: minimize the number of T gates.

## Monte Carlo iterations

$$q(j) \approx p(j), \text{ while regularising } th_j / (Nq(j)) \approx k\pi.$$

# Summary of QDRIFT

- Stochastic channel: good for non-uniform distribution of the coefficients, e.g. chemistry
- Concentration:  $N$  can be kept small.
- Importance sampling, concentration with NM
- Higher order: qSWIFT = qDRIFT + correction terms

PRL **123**, 070503 (2019)

PRX Quantum **2**, 040305 (2021)

Kiss et al, arXiv:2212.05952

arXiv:2302.14811

**Better error scaling**  $d_{\diamond}(\mathcal{U}, \mathcal{E}^{(K)}) \in \mathcal{O}\left(\left(\frac{(\lambda t)^2}{N}\right)^K\right).$

# **Composite channels: how to reduce the implementation cost?**

# Composite channels

**Given:** a partition of the Hamiltonian.

$$\mathbf{H} = \mathbf{A} + \mathbf{B}$$

**Idea:** simulate each terms with a different channel (Trotter & QDRIFT).

**Time evolution = Trotter + QDRIFT**

## Advantages:

- Greater than the sum of the parts.
- Take advantage of the specific properties of each sub-system.
- Cost and error reduction.

$$\left\| \mathcal{I}_H(t) - \tilde{\mathcal{U}}_H(t/r)^r \right\|_{\diamond} \leq \frac{t^2}{r} \left( \sum_{i < j} a_i a_j \| [A_i, A_j] \|_{\infty} + \sum_{ij} a_i b_j \| [A_i, B_j] \|_{\infty} + \frac{\lambda_B^2 (1 + \mathbb{E}_p[\omega(j)])}{N_B} \right)$$

Hagan and Wiebe, arXiv: 2206.06409 (2022)

# How to choose the partition?

1. **Perturbation theory:**  $H = A + \beta B, \beta \ll 1$ . Example: SU(4) lattice gauge theory.
2. **Cost oriented:**  $A$  contains the easiest terms while  $B$  the most expensive ones.  
Example: lattice effective field theory (Hubbard like) potential VS hopping.
3. **Connectivity oriented:**  $H = \sum A_i$  where  $A_i$  can be simulated with minimal swap overhead.

## 4. T gate efficient:

$$|th_j - k_j\pi| < \epsilon$$

$$k_j \in \mathbb{Q}$$

$$\begin{aligned} tH &= \sum_j th_j H_j = \sum_j (k_j\pi) H_j + \sum_j (th_j - k_j\pi) H_j \\ &= A + B. \end{aligned}$$



# Numerical simulations: lattice effective field theory

# Scalable model for a nuclei

Pionless EFT on a lattice with  $M \times M$  sites and  $A$  nucleons from  $N_f$  different species.

$$H = -t \sum_{f=1}^{N_f} \sum_{\langle i,j \rangle} c_{i,f}^\dagger c_{j,f} + 2dtA$$

Kinetic energy (hopping)

$$+ U \sum_{i=1}^{N_f} \sum_{f < f'} n_{i,f} n_{i,f'} + V \sum_{f < f' < f''} \sum_{i=1}^{N_f} n_{i,f} n_{i,f'} n_{i,f''}$$

Two and three-body on-site interaction

$$+ U \sum_{f=1}^{N_f} n_{1,f} + V \sum_{f < f'}^{N_f} n_{1,f} n_{1,f'}$$

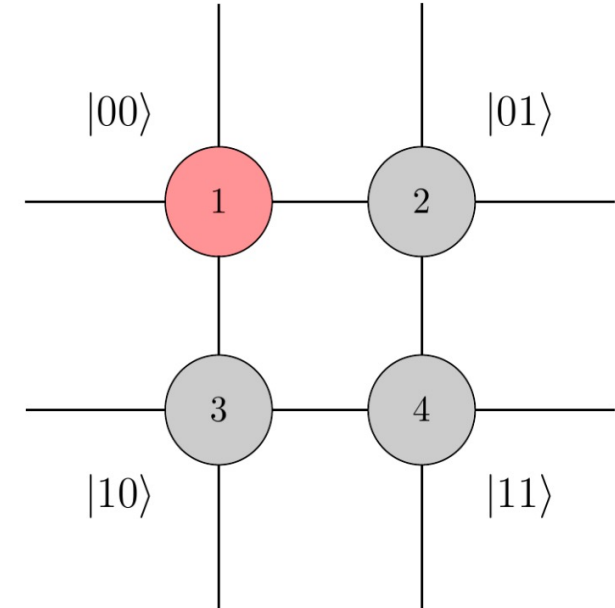
Potential of a frozen nucleon

Roggero et al., Phys. Rev. D **101**, 074038 (2020)

## Mapping to qubits:

First quantisation

needs  $\log_2(M^2)N_f$  qubits.



Second quantisation (via Jordan-Wigner):

More natural but needs  $M^2 N_f$  qubits.

# Cost of the different implementations

Assuming a 1423 connectivity

generator	cost	generator	cost
$X_k$	0.1	$Z_1 Z_2$	6
$Z_k$	0.1	$Z_3 Z_4$	6
$Z_1 Z_4$	2	$Z_1 Z_2 Z_3 Z_4$	6
$Z_2 Z_4$	2	$Z_1 Z_3 Z_4$	8
$Z_2 Z_3$	2	$Z_1 Z_2 Z_3$	8
$Z_1 Z_2 Z_4$	4	$Z_1 Z_3$	10
$Z_2 Z_3 Z_4$	4		

Table 1: Implementation cost for the different generators appearing in the two considered Hamiltonians.

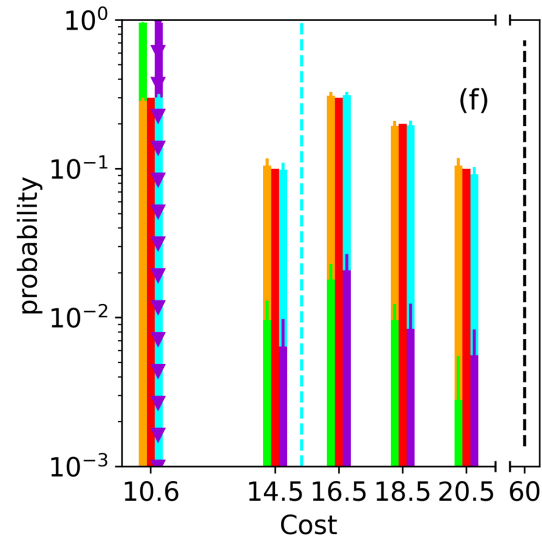
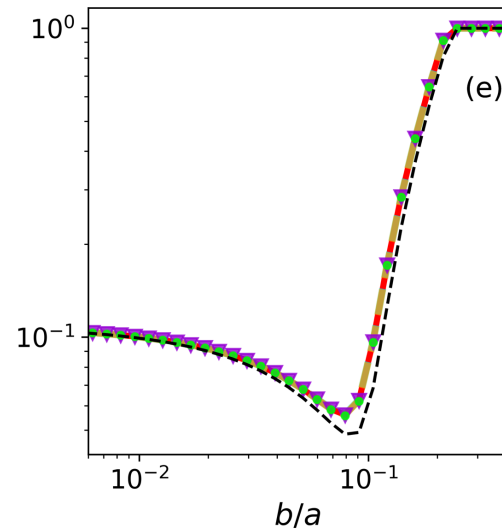
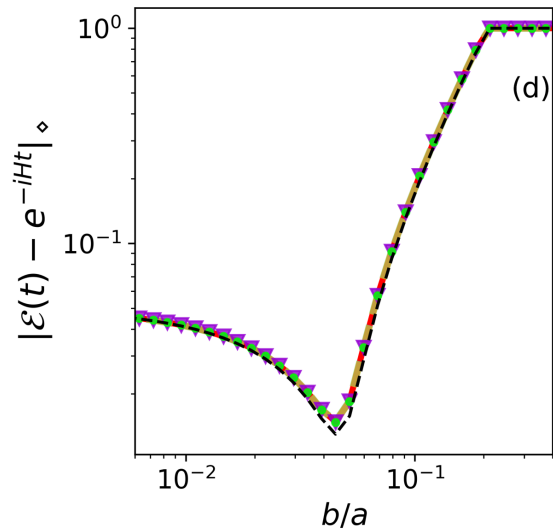
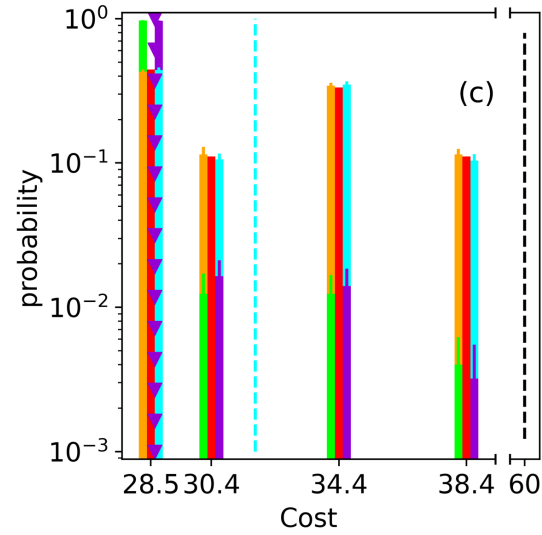
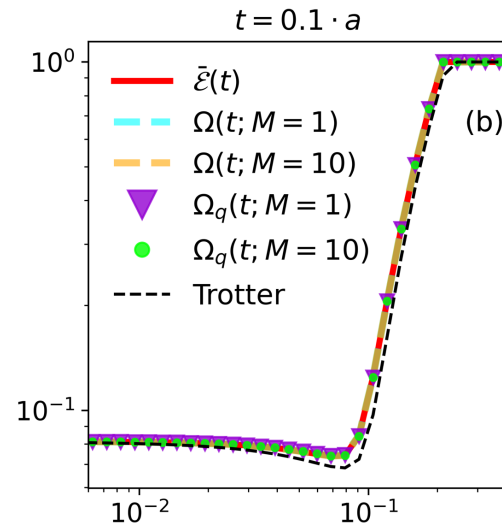
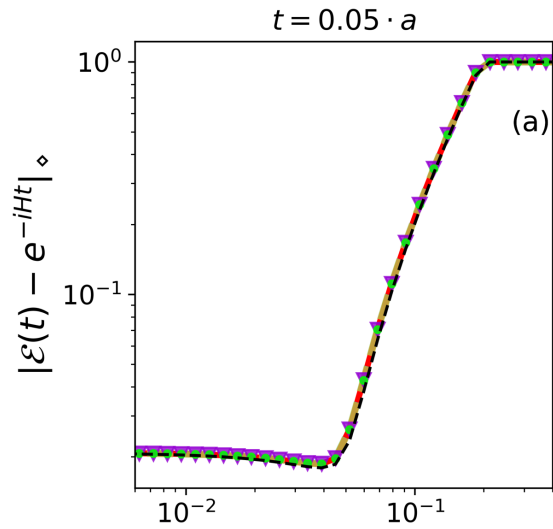
Two regimes (**A** with Trotter, **B** with QDRIFT):

1. **A** describes the bulk ( $U = -4V$ ), while **B** is the linear interpolation towards **realistic coefficients**.
2. **A** contains the **easy** terms and **B** the **expensive** ones.

$$A^{(1)} = \sum_{k=1}^4 X_k + Z_1 + Z_1 Z_4 + Z_2 Z_3 + Z_2 Z_4 + Z_1 Z_2 Z_4 \quad (63)$$

$$B^{(1)} = \sum_{k=2}^4 Z_k + Z_1 Z_2 + Z_1 Z_3 + Z_3 Z_4 + Z_1 Z_2 Z_3 + Z_2 Z_3 Z_4 + Z_1 Z_3 Z_4 + Z_1 Z_2 Z_3 Z_4. \quad (64)$$

# Numerical Simulation (N=1)

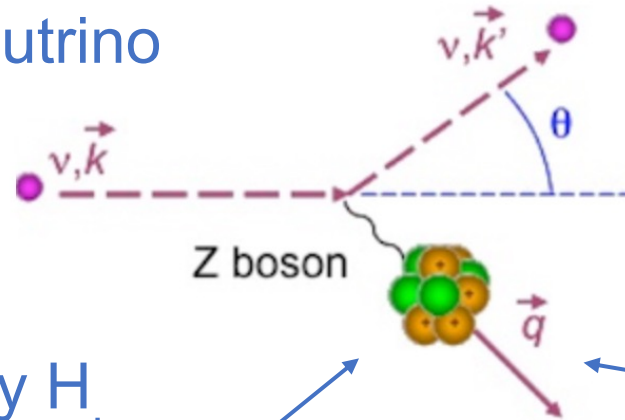


- Cost reduction in the B evolution of factor 10.
- QDRIFT and IS QDRIFT match Trotterization for all values of  $\beta$ .

# Response functions computations from scattering experiments

# Physical Model

Incoming neutrino



Nuclei described by  $H$ ,  
in its ground state.

momentum transfer  
described by  $\hat{O}$   
(external probing).

$$H = 4.5 \cdot \mathbb{1} - 2 \sum_{i=1}^4 X_i$$

$$+ 1.75 \left( \sum_{i < j < k} Z_i Z_j Z_k + Z_1 Z_4 + Z_2 Z_3 \right)$$

$$\hat{O}(\vec{q}_k) = \sum_{f=1}^{N_f} \rho_f(\vec{q}_i) = \sum_{f=1}^{N_f} e_f \sum_i e^{i\vec{q}_k \cdot \vec{r}_i} n_{i,f},$$

# Linear response function

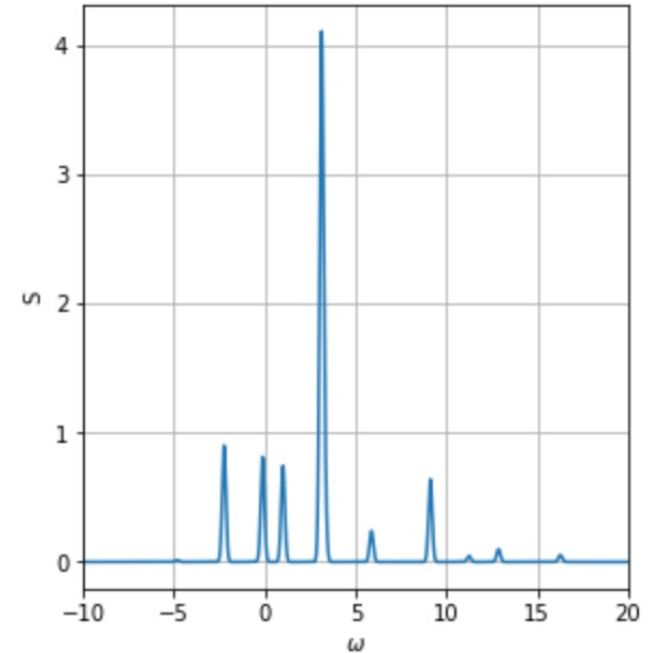
$$\begin{aligned} S(\omega, \vec{q}) &= \langle \Psi_0 | \hat{O}(\vec{q})^\dagger \delta(\omega - (E_0 - E_f)) \hat{O}(\vec{q}) | \Psi_0 \rangle \\ &= \sum_f \left| \langle \Psi_0 | \hat{O}(\vec{q}) | f \rangle \right|^2 \delta(\omega - (E_0 - E_f)), \end{aligned}$$

Instead: Integral transform with Gaussian Kernel + truncation

$$\Phi_N^\chi(\nu) = \frac{1}{\chi \|H\|} \sum_{n=-N}^N g_n^\chi(\nu) \langle \Psi_0 | \hat{O}(\vec{q})^\dagger e^{-in\delta t H} \hat{O}(\vec{q}) | \Psi_0 \rangle,$$

Groudstate with the Variational Quantum Eigensolver

Challenging, since it requires the full spectrum



**Hamiltonian moments**  
**Easy with quantum computers!**

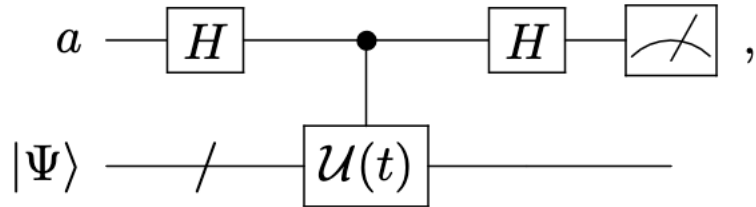
arXiv:2211.00790

# Estimating expectation values

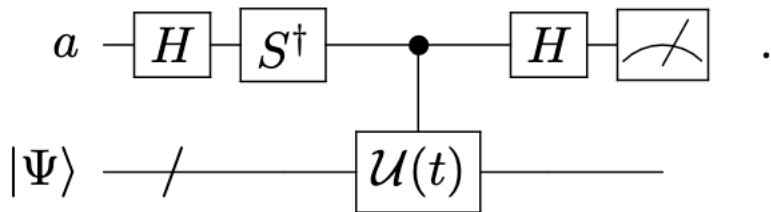
$$g(t) = \langle \Psi_0 | \hat{O}(\vec{q})^\dagger e^{-itH} \hat{O}(\vec{q}) | \Psi_0 \rangle,$$

## Hadamard test (phase kickback)

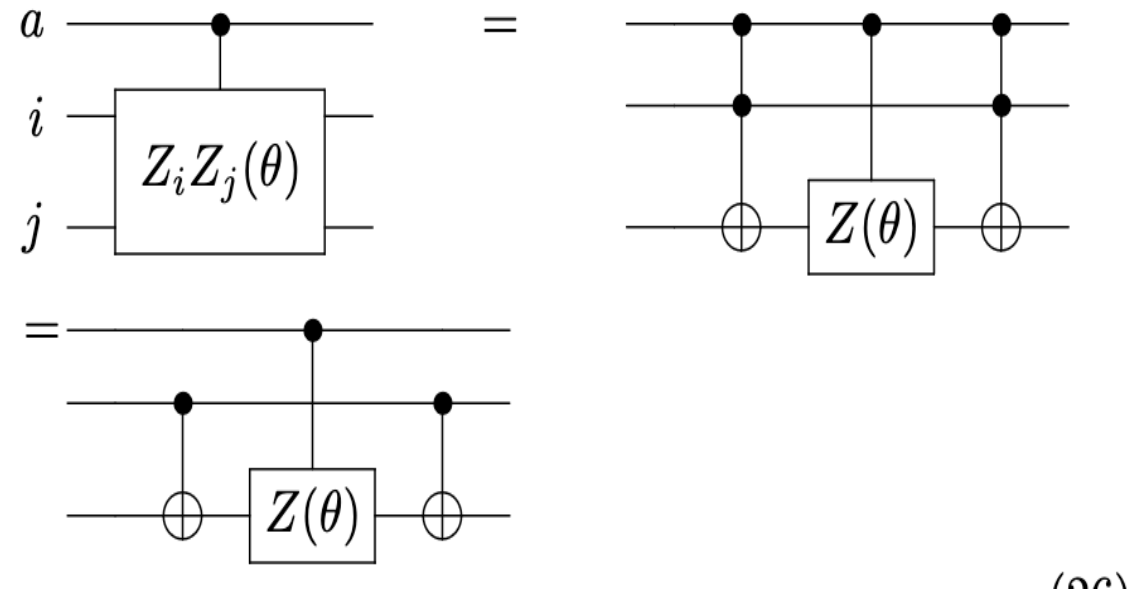
Real



Imag.



In practice, control operations are expensive!





# Control Reversal Gates

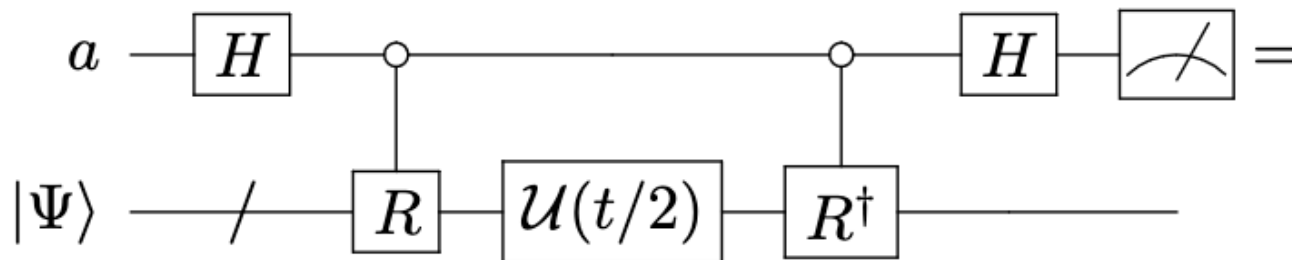
## How to avoid the control operation?

R anti commutes with H:  $\{H,R\} = 0$   
 Use R to toggle the flow of time

$$R \exp\{-iHt\} R^\dagger = R \sum_n \frac{(-itH)^n}{n!} R^\dagger =$$

$$\sum_n \frac{(itH)^n}{n!} R R^\dagger = \exp\{itH\},$$

### Hadamard test with control reversal gates



### Caveats:

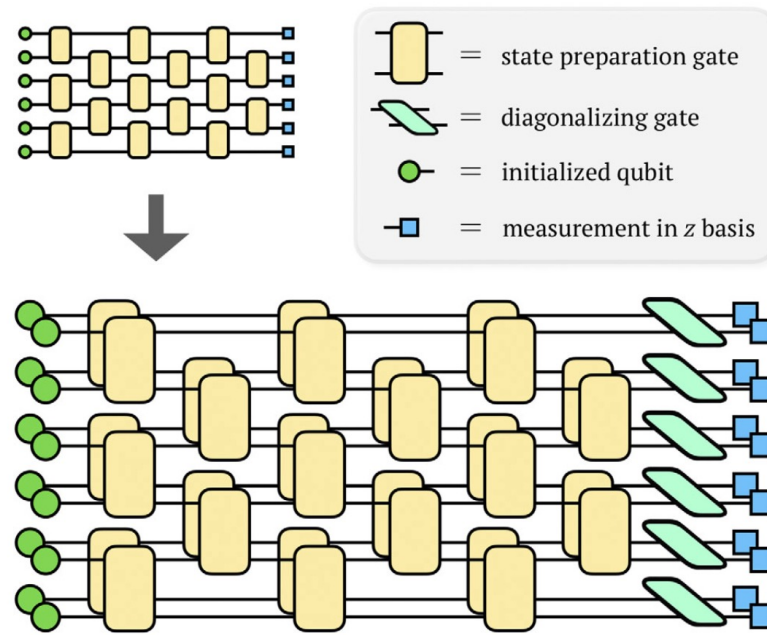
- we may need to split the Hamiltonian, and/or insert the CRG between each Trotter steps!
- **We were able to avoid this here!**
- Require **even order** product formula.

# Error mitigation: Virtual Distillation

Uses multiple copies to suppress the non-dominating components.

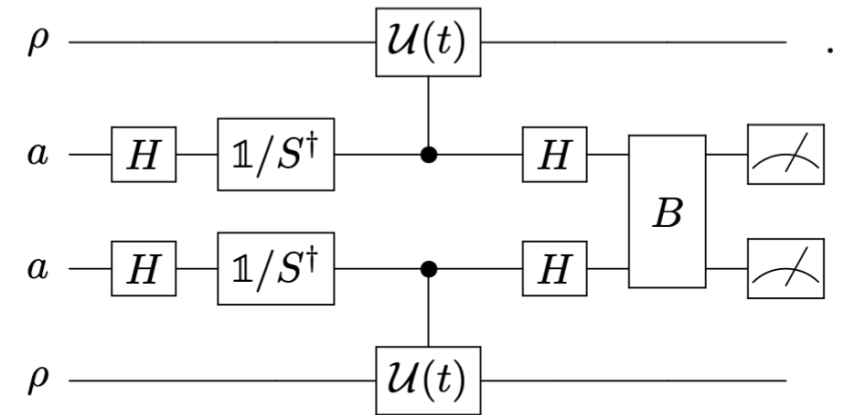
$$\rho^M = \left( \sum_i p_i |i\rangle\langle i| \right)^M = \sum_i p_i^M |i\rangle\langle i|,$$

$$\langle O \rangle_{\text{VD}} \equiv \frac{\text{Tr}(O \rho^M)}{\text{Tr}(\rho^M)}.$$



- requires connectivity between every copy.
- The diagonalising gate can be challenging to implement in general.

**we only need to connect the ancilla.**



Huggins, et al., Phys. Rev. X 11, 041036 (2021)

# Error mitigation: Self-verification

Assuming depolarising noise

$$\rho_1 = (1 - p)|\Psi_1\rangle\langle\Psi_1| + \frac{p}{2^n}\mathbb{1}. \quad (38)$$

The expectation value of a Pauli operator  $P$  can then be computed as

$$\text{Tr}(P\rho_1) = (1 - p)\langle\Psi_1|P|\Psi_1\rangle, \quad (39)$$

**We could correct if we knew  $p$ !**

**Idea: compute  $p$  using a known circuit.**

We **prepare**, then **un-prepare** the targeted state (two Trotter steps).

$$|\Psi_2\rangle = \overline{A^\dagger AB} |0\rangle,$$

$$\text{Tr}(Q\rho_2) = (1 - p)\langle\Psi_2|Q|\Psi_2\rangle + p\frac{\text{Tr}(Q)}{2^n},$$

And use this value for the correction

$$\langle\Psi|\mathcal{U}(2t)|\Psi\rangle_{\text{SV}} = \frac{\langle 0|B^\dagger\mathcal{U}(2t)B|0\rangle}{\langle 0|B^\dagger\mathcal{U}(t)\mathcal{U}(-t)B|0\rangle},$$

O'Brien, PRX Quantum 2, 020317 (2021)

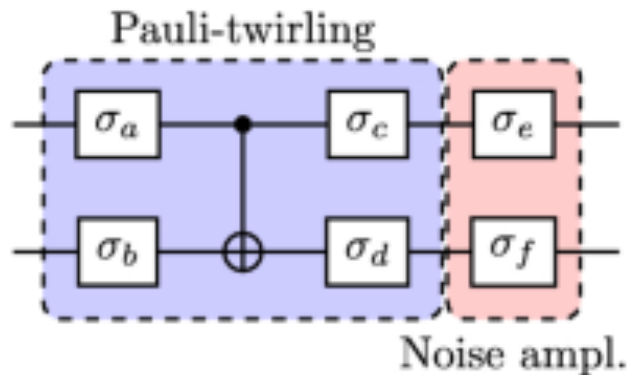
# How to make the noise more depolarising?

## Use Pauli Twirling!

Turn a noisy operator into a Pauli channel, via gate conjugation.

$$\mathcal{T}_W(\overline{M}) = \frac{1}{|W|} \sum_{w \in W} \overline{w M w^\dagger}.$$

In practice: sample!



(a) Pauli-twirling and noise amplification.

$\sigma_a$	$\mathbb{1}$	$\mathbb{1}$	$\mathbb{1}$	$\mathbb{1}$	$\sigma_x$	$\sigma_x$	$\sigma_x$	$\sigma_x$	$\sigma_y$	$\sigma_y$	$\sigma_y$	$\sigma_y$	$\sigma_z$	$\sigma_z$	$\sigma_z$	$\sigma_z$
$\sigma_b$	$\mathbb{1}$	$\sigma_x$	$\sigma_y$	$\sigma_z$	$\mathbb{1}$	$\sigma_x$	$\sigma_y$	$\sigma_z$	$\mathbb{1}$	$\sigma_x$	$\sigma_y$	$\sigma_z$	$\mathbb{1}$	$\sigma_x$	$\sigma_y$	$\sigma_z$
$\sigma_c$	$\mathbb{1}$	$\mathbb{1}$	$\sigma_z$	$\sigma_z$	$\sigma_x$	$\sigma_x$	$\sigma_y$	$\sigma_y$	$\sigma_y$	$\sigma_y$	$\sigma_x$	$\sigma_x$	$\sigma_z$	$\sigma_z$	$\mathbb{1}$	$\mathbb{1}$
$\sigma_d$	$\mathbb{1}$	$\sigma_x$	$\sigma_y$	$\sigma_z$	$\sigma_x$	$\mathbb{1}$	$\sigma_z$	$\sigma_y$	$\sigma_x$	$\mathbb{1}$	$\sigma_z$	$\sigma_y$	$\mathbb{1}$	$\sigma_x$	$\sigma_y$	$\sigma_z$

(b) Valid combinations for Pauli-twirling of the CX gate.

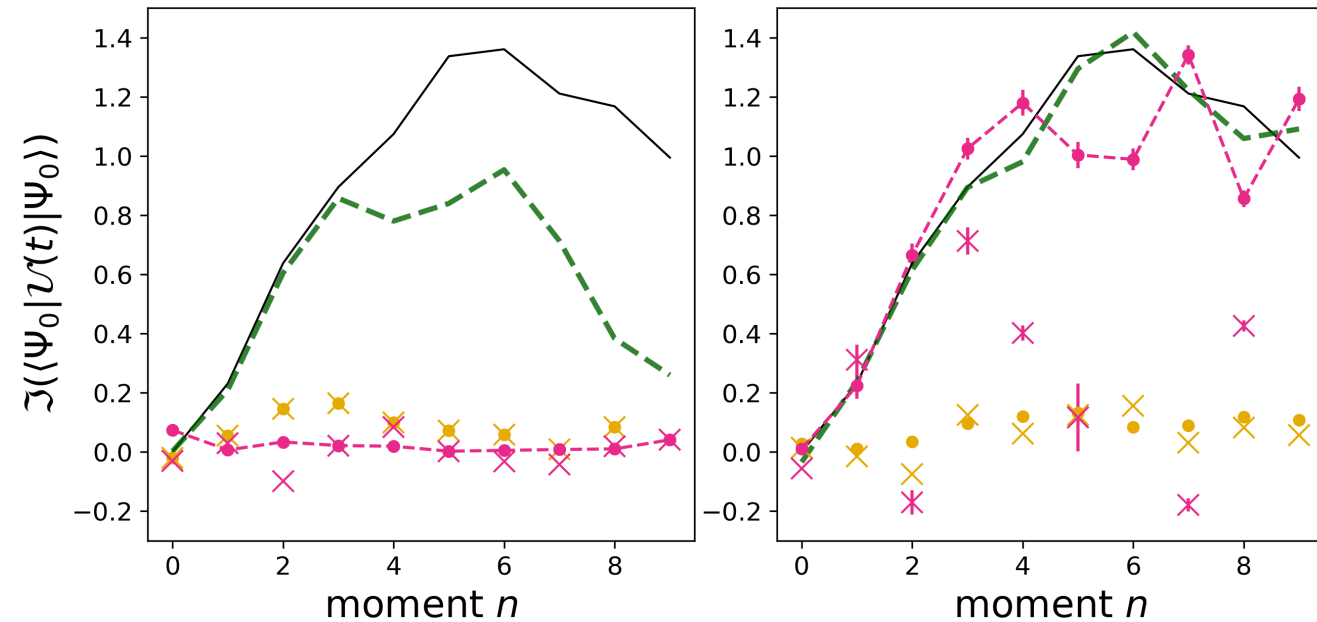
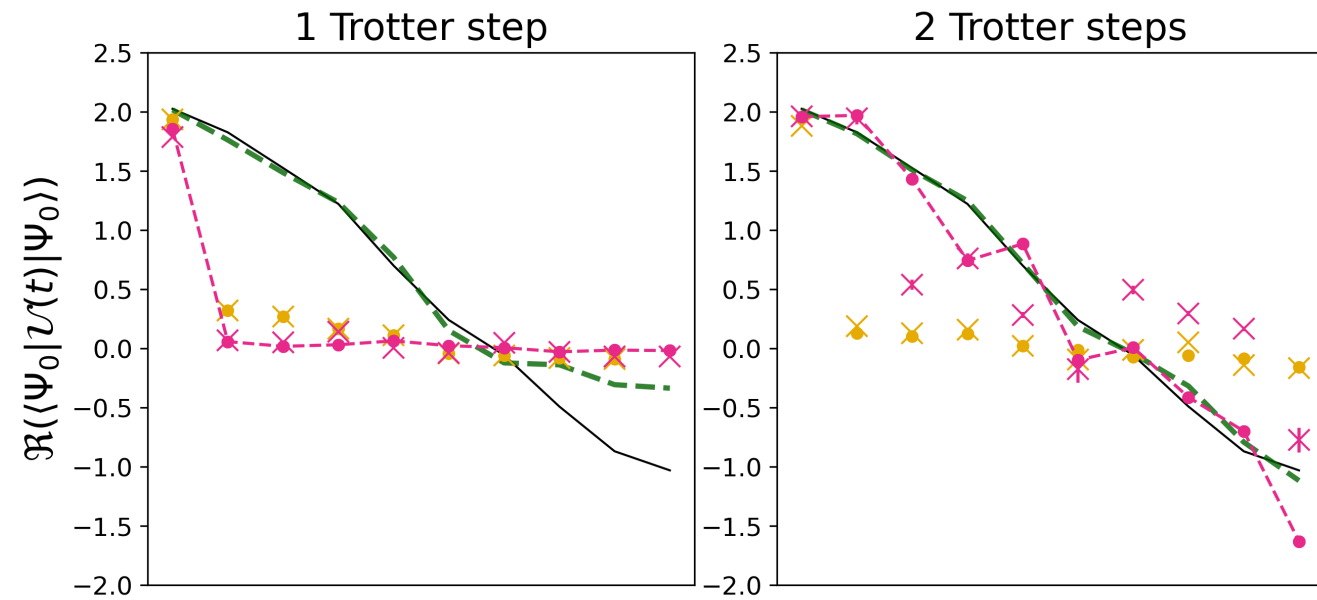
Fuchs, et al., *Eur. Phys. J. Plus* **135**, 353 (2020)

# Results on superconducting quantum hardware

	linear connectivity	full connectivity
one Trotter step	38	27
two Trotter steps	68	49

Number of CNOTs

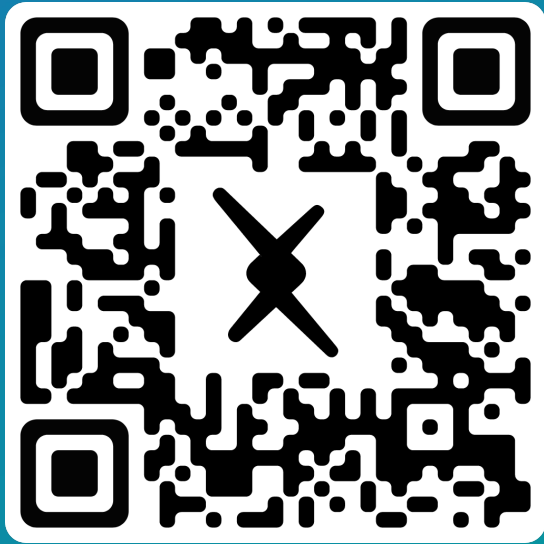
- VD does not improve the results (noise too important)
- Pauli twirling alone is not helpful.
- Self Verification works well in conjunction with Pauli Twirling.



# Conclusions

- **QDRIFT** generates random product formulas whose size **does not depend** on the number of terms and concentrates fast (small variance).
- **Importance sampling** can be used to **reduce the actual implementation cost** on hardware (guaranteed cost reduction with the same accuracy).
- **Composite channels** use different techniques to simulate different parts of the Hamiltonian (use the best one in each case!)
- **Response functions** can be estimated from expectation values!
- **Self-verification** and **twirling** are working well together.

# Thank you for your attention!



Michele Grossi



Alessandro Roggero

Important sampled QDRIFT  
arXiv:2212.05952



QUANTUM  
TECHNOLOGY  
INITIATIVE

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