

EPFL

MARVEL



NATIONAL CENTRE OF COMPETENCE IN RESEARCH



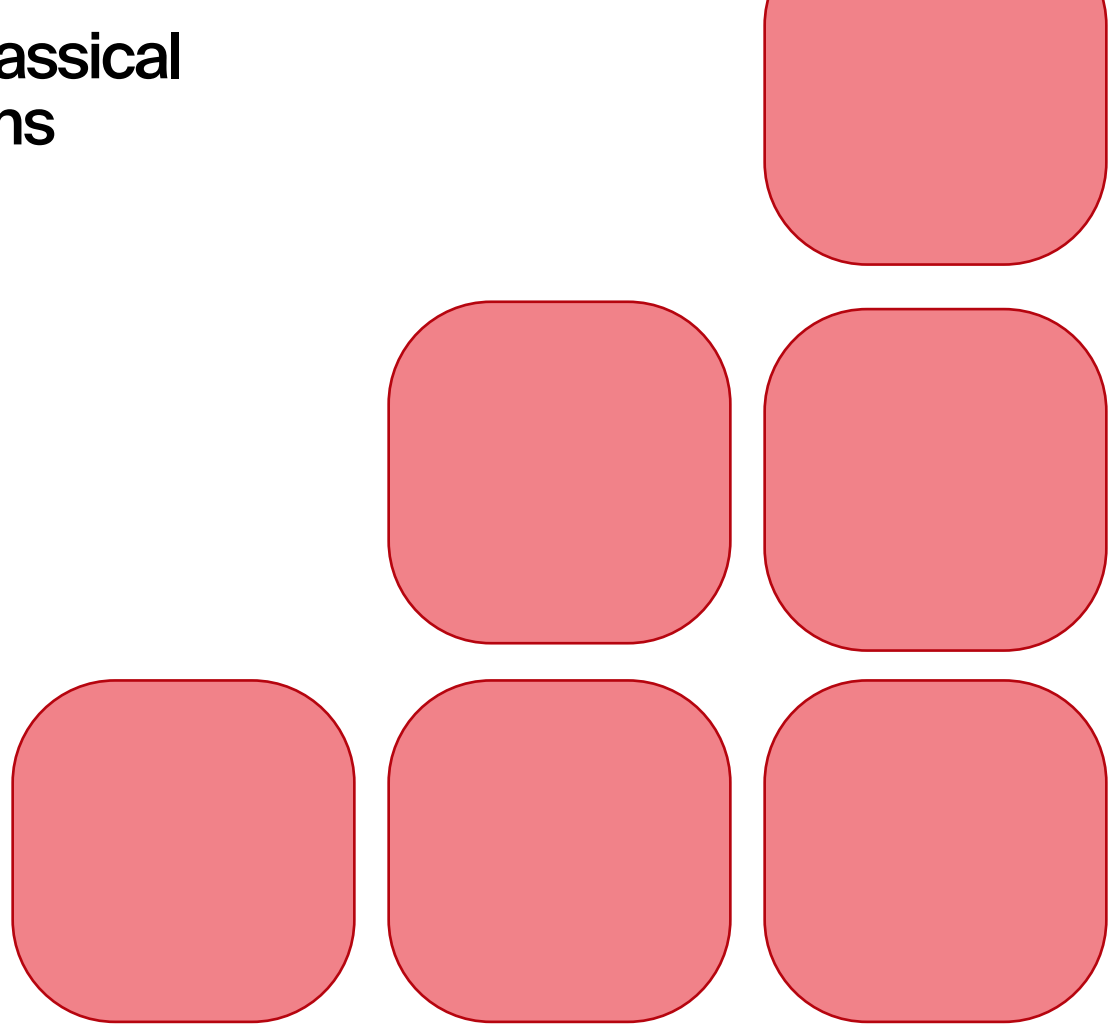
Computational
Quantum
Science
Laboratory

Overhead-constrained circuit knitting for variational quantum dynamics

QTM 2023 – CERN

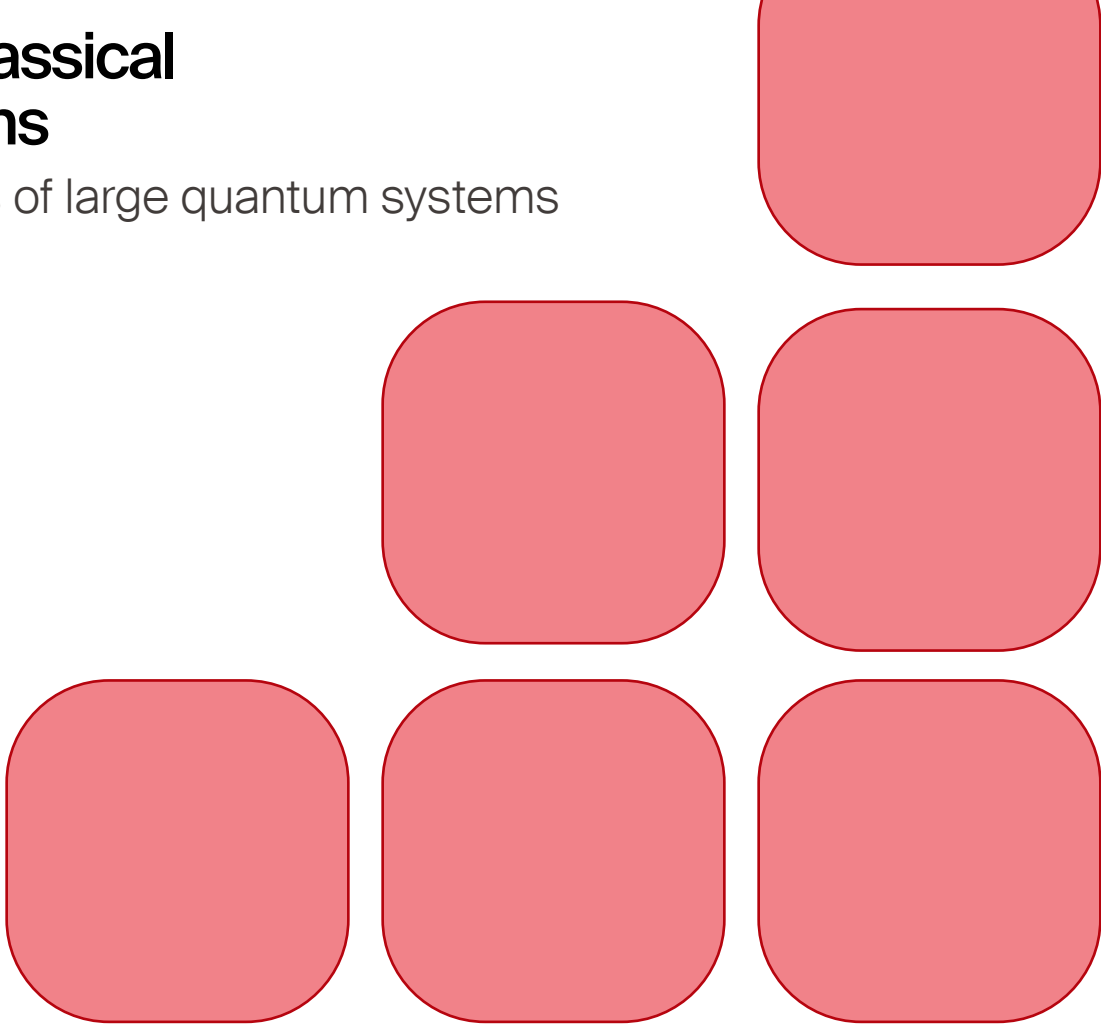
Gian Gentinetta

EPFL Hybrid quantum-classical
quantum simulations



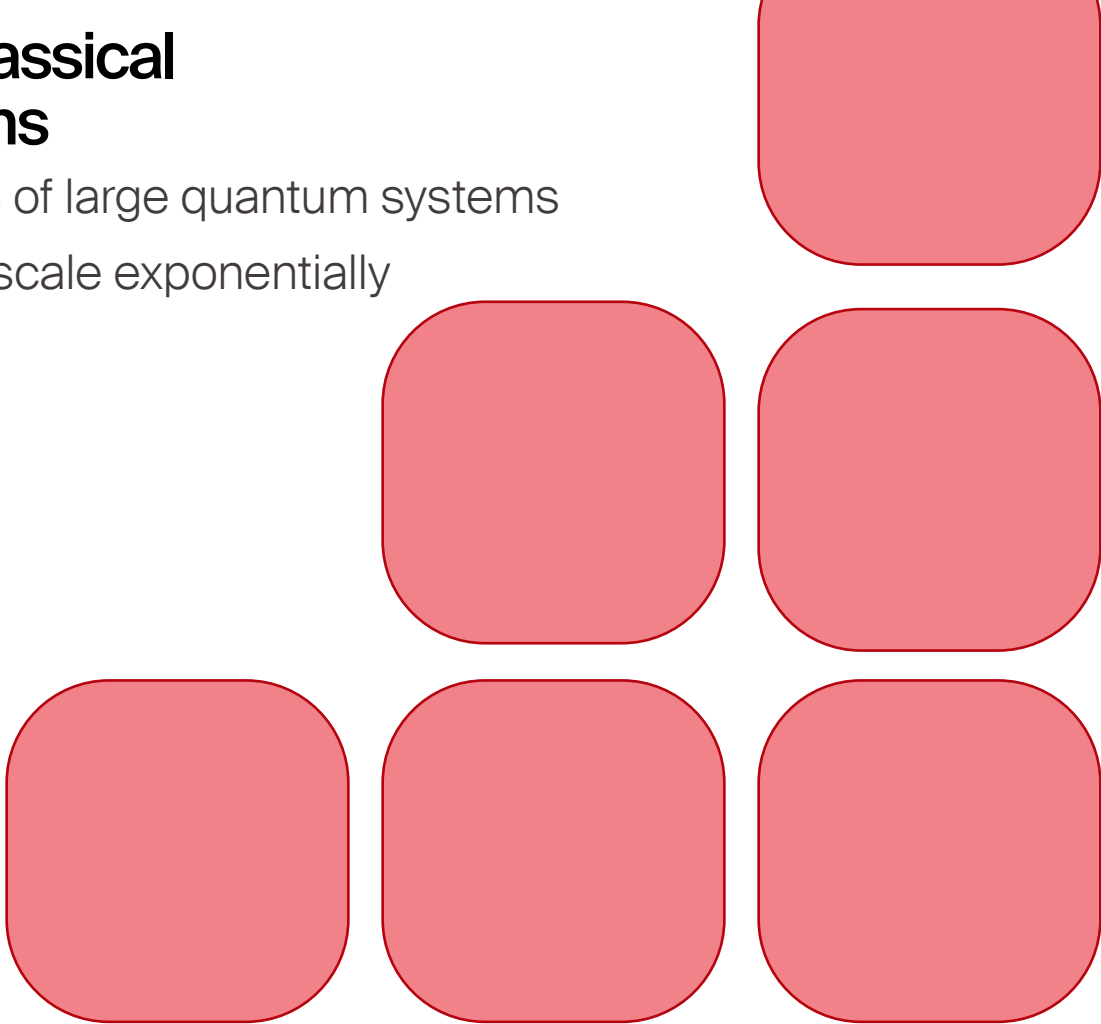
EPFL Hybrid quantum-classical quantum simulations

- Goal: Simulate dynamics of large quantum systems



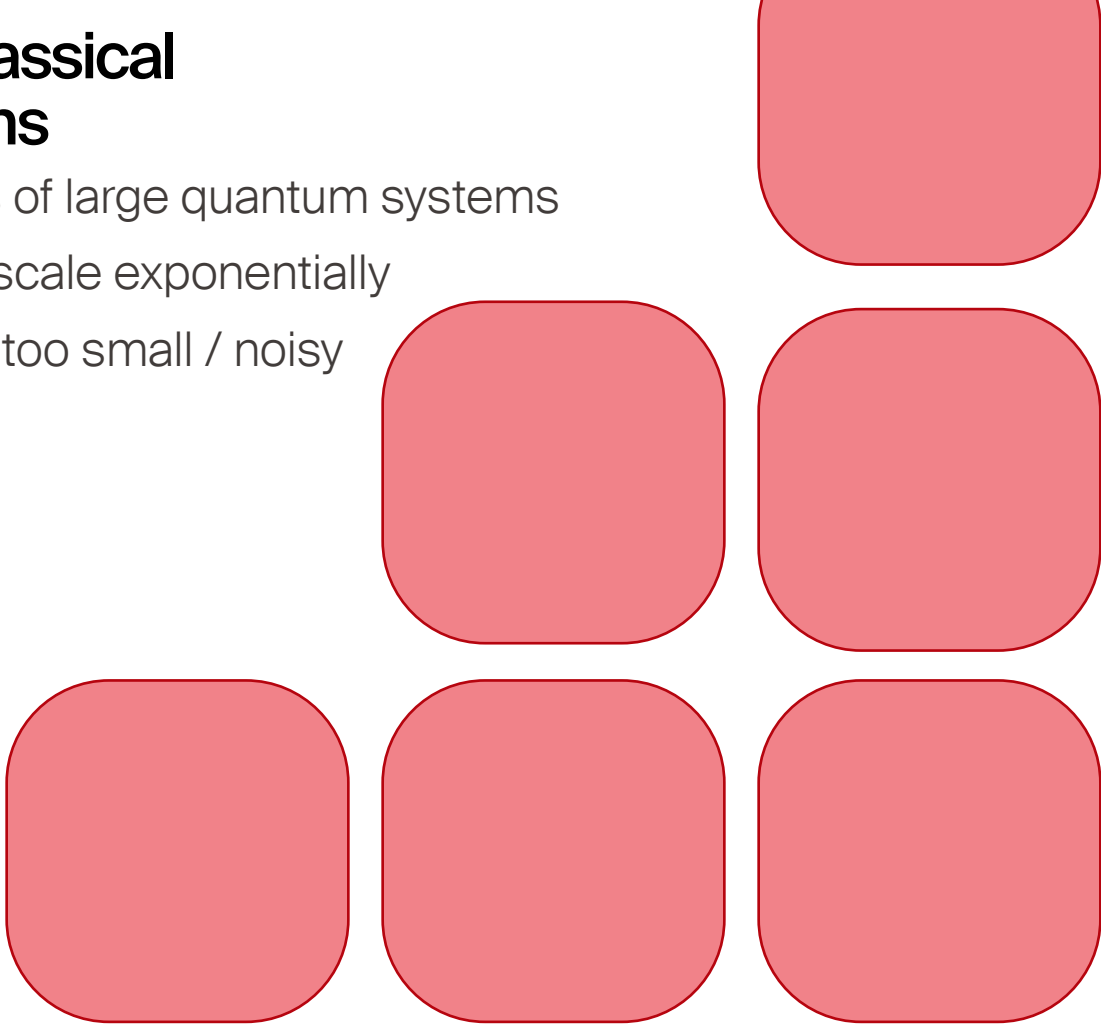
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- Goal: Simulate dynamics of large quantum systems
 - Classical algorithms scale exponentially



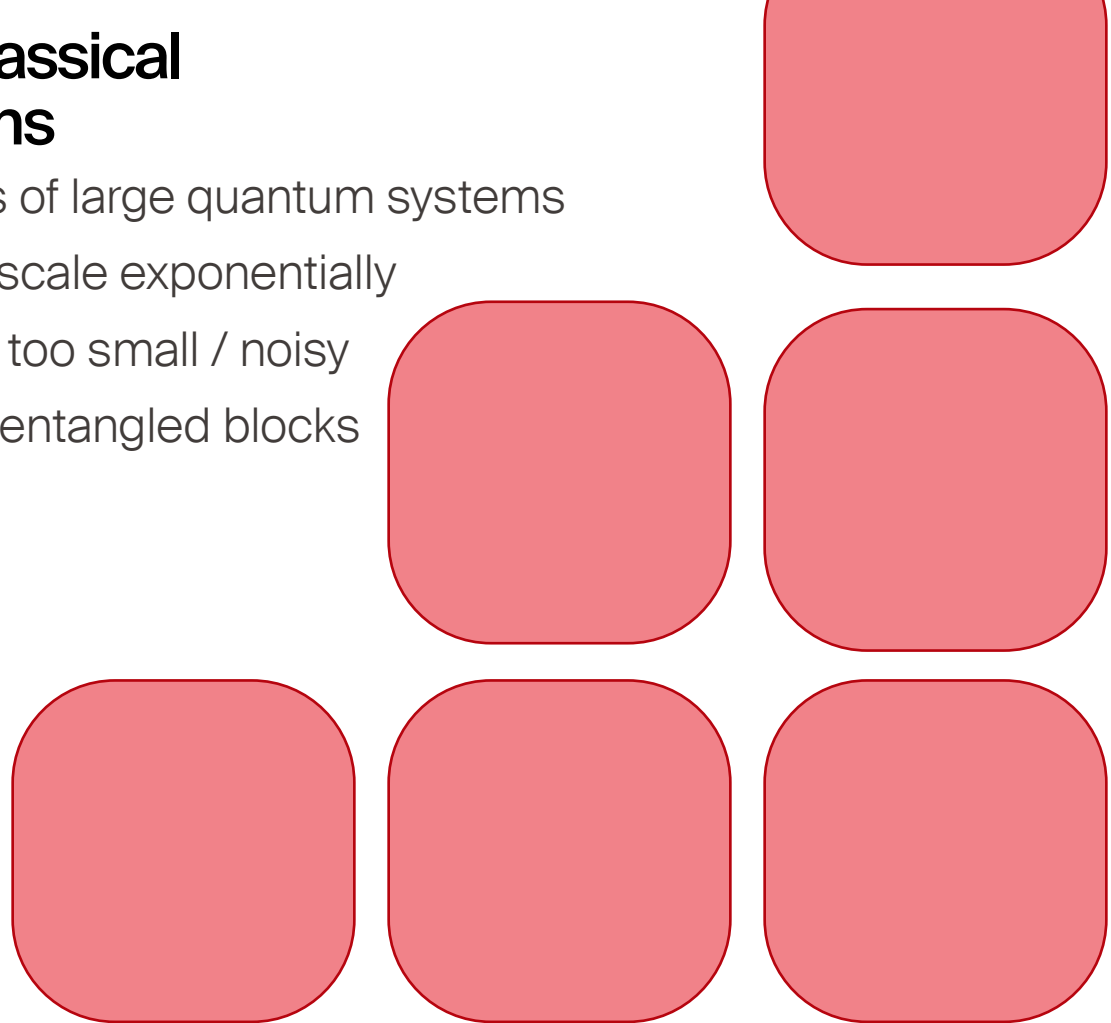
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- Goal: Simulate dynamics of large quantum systems
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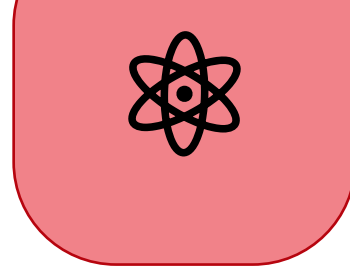
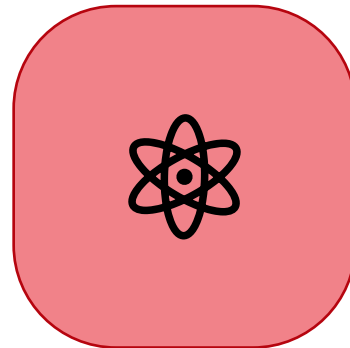
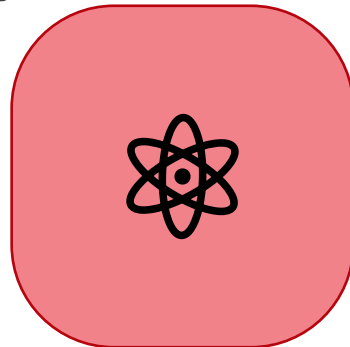
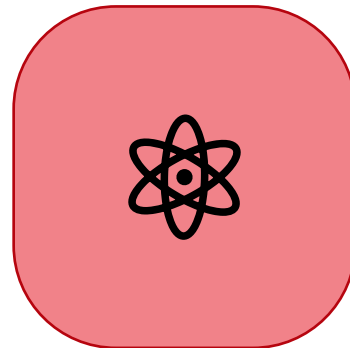
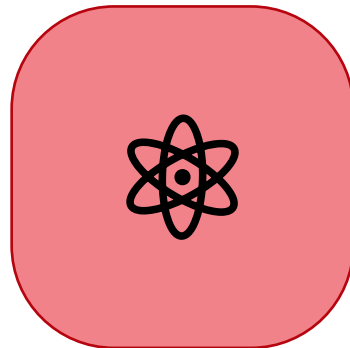
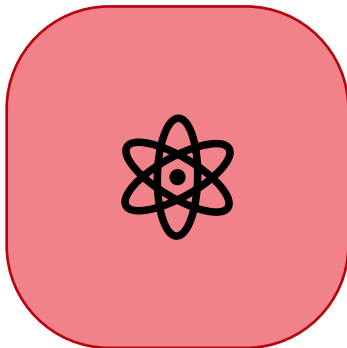
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- Goal: Simulate dynamics of large quantum systems
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- Split system into weakly entangled blocks



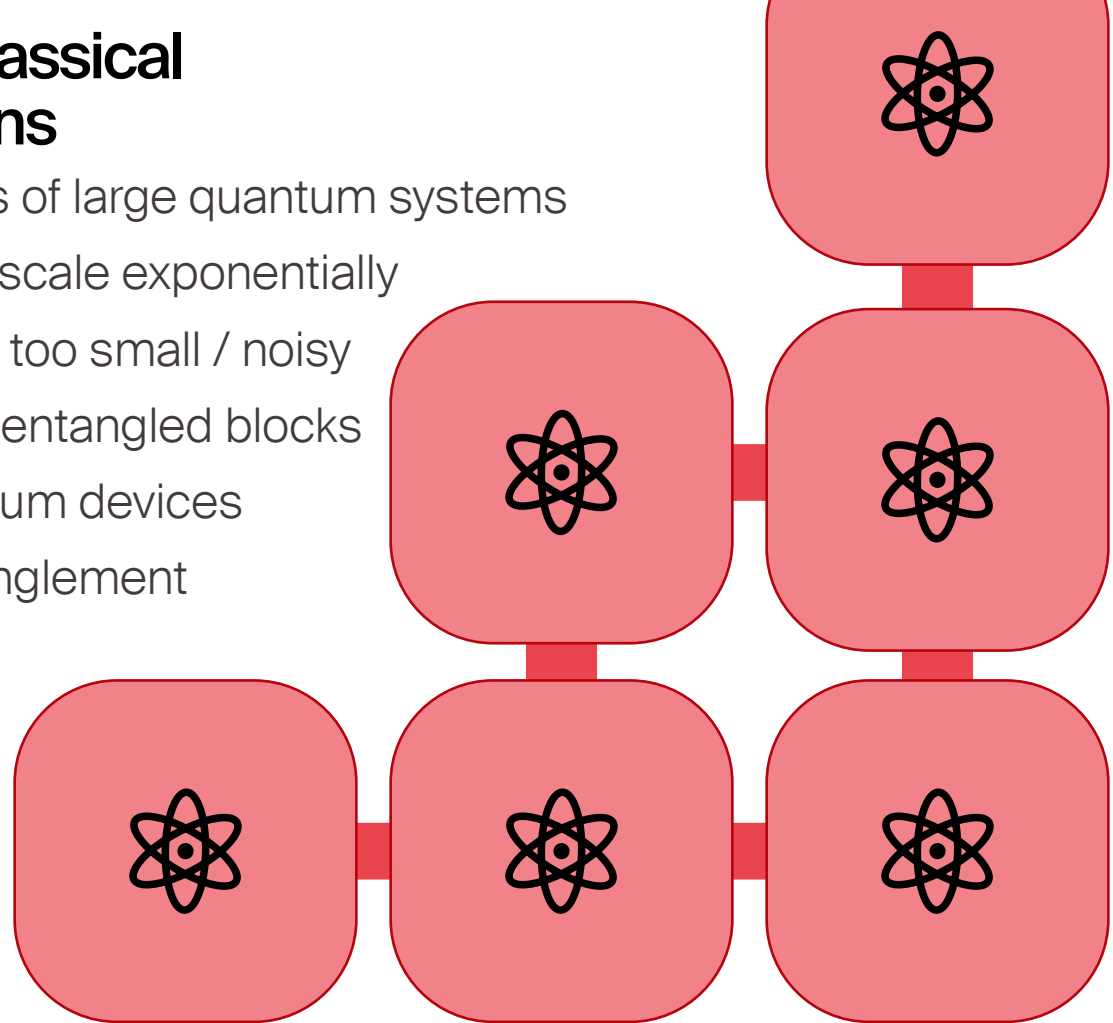
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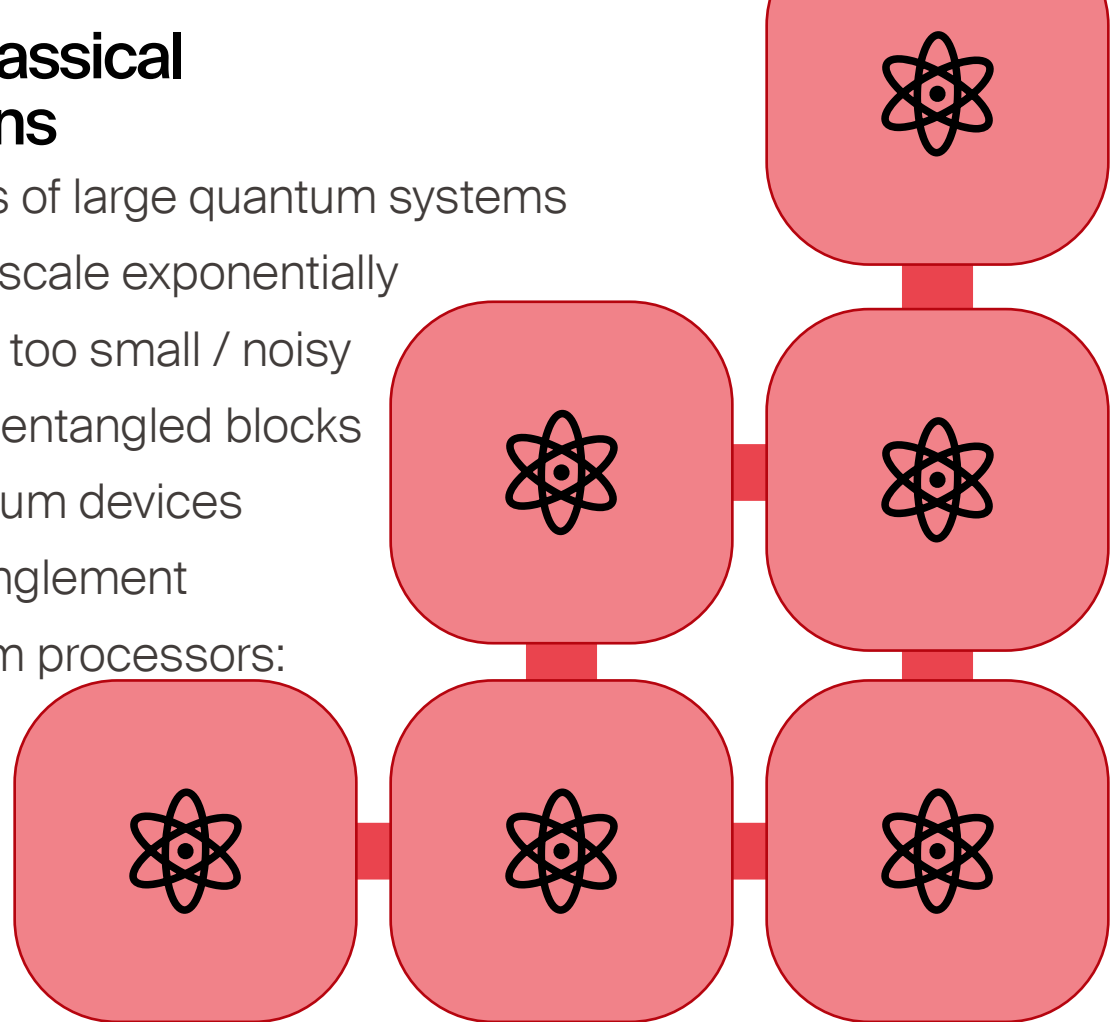
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- Next generation quantum processors:



- Simulate dynamics of spins in a transverse field Ising model

$$H = \sum_{\langle ij \rangle} J_{ij} Z_i Z_j + \sum_i X_i$$

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Strongly entangled spins

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Strongly entangled spins

Weakly entangled blocks

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Strongly entangled spins

Weakly entangled blocks

⇒ Cut system into blocks

- Time evolution with fixed circuit depth

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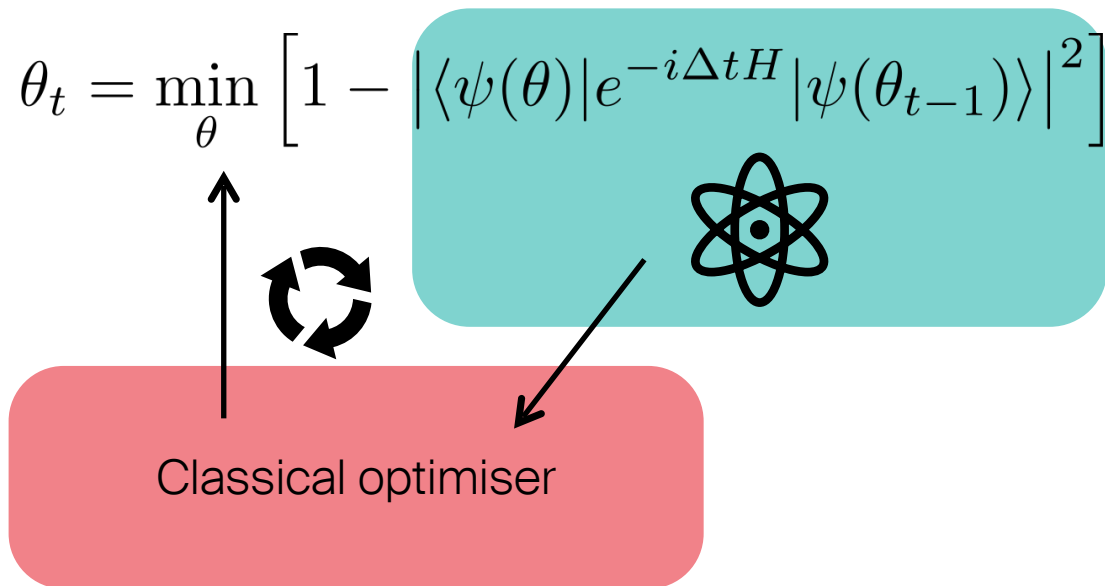
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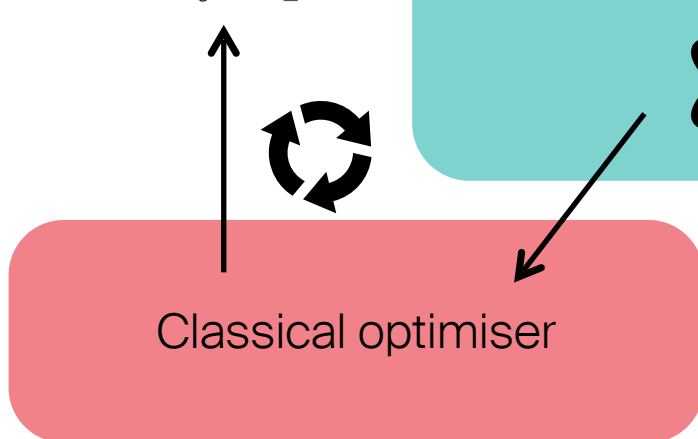
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⇒ Only requires measuring fidelities

- Need to measure $|\langle \psi(\theta) | e^{-i\Delta t H} | \psi(\theta_{t-1}) \rangle|^2$ across blocks

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Global quantum channel

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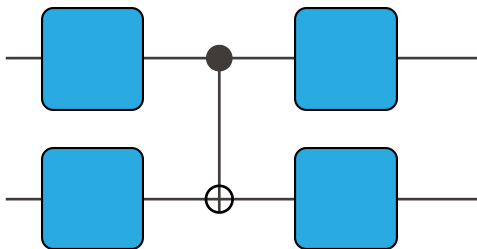
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Cost: Sampling overhead $\omega = \left(\sum_k |\alpha_k| \right)^2$

Choosing gates for circuit knitting

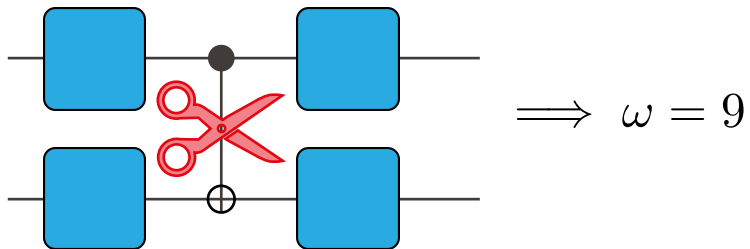
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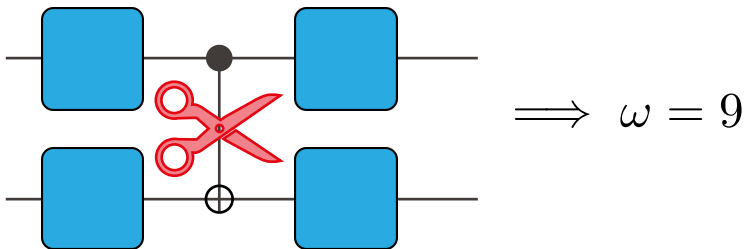


Choosing gates for circuit knitting

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Cut multiple gates \implies Multiply overhead

2 CNOTs $\implies \omega = 81$

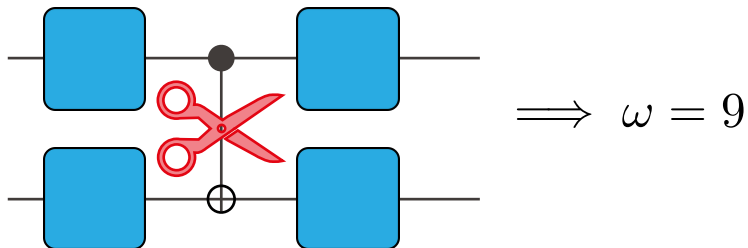
3 CNOTs $\implies \omega = 243$

...

L CNOTs $\implies \omega = 9^L$

Choosing gates for circuit knitting

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- CNOT not ideal, want to control overhead

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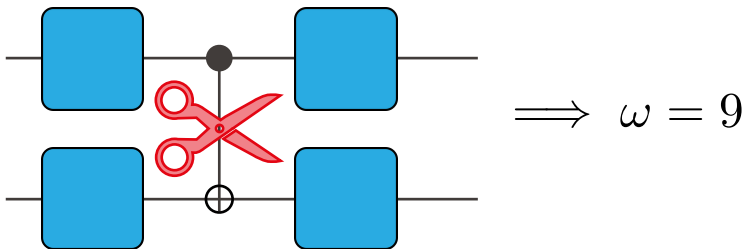
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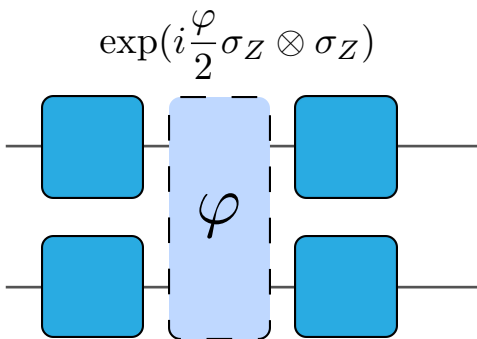
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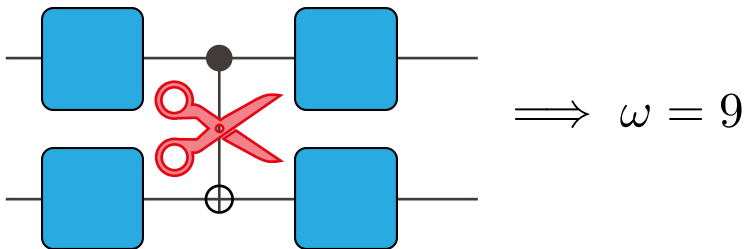
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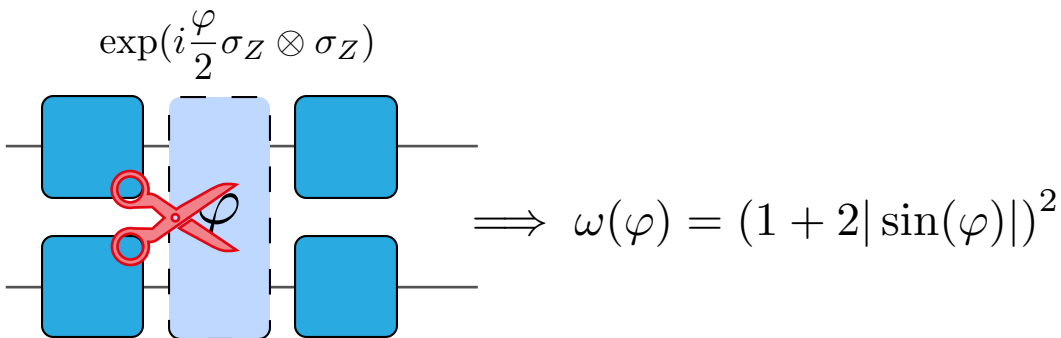
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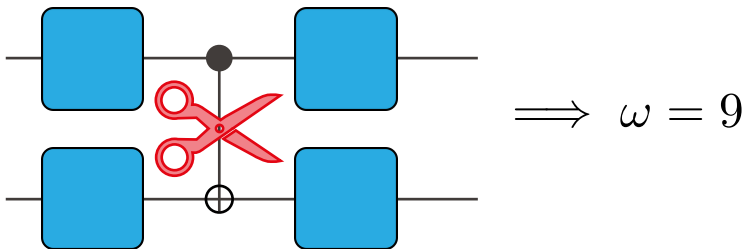
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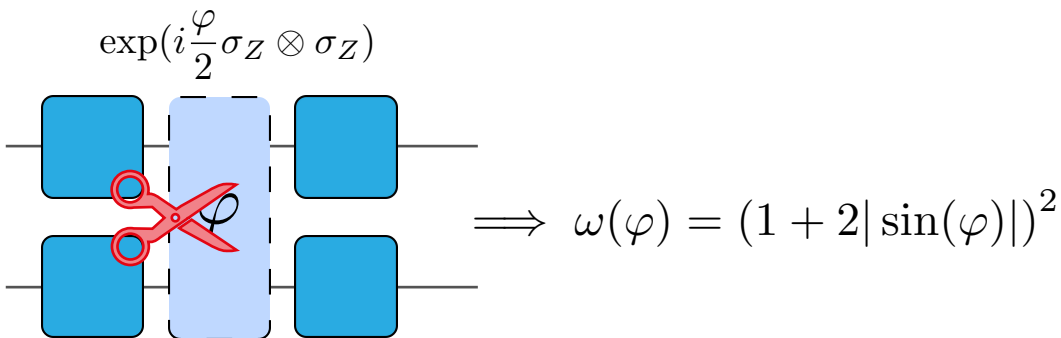
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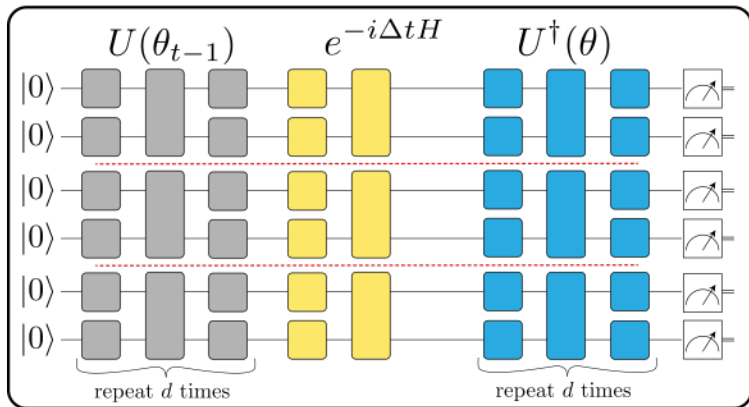
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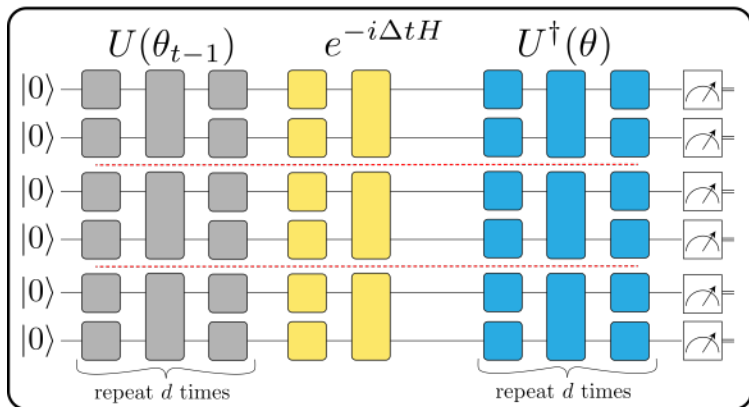
Small angles

\Rightarrow small overhead

BPA

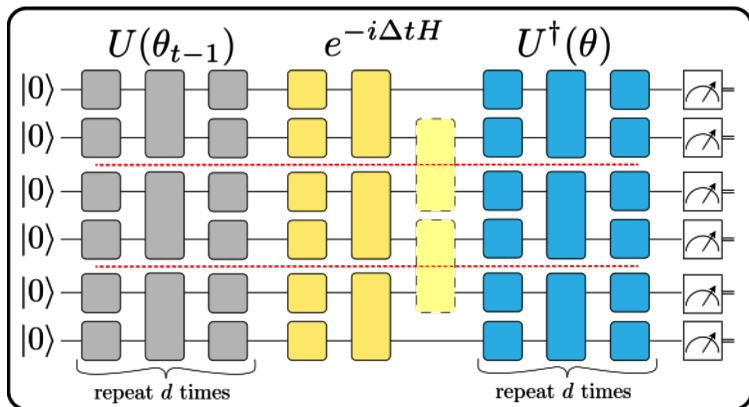


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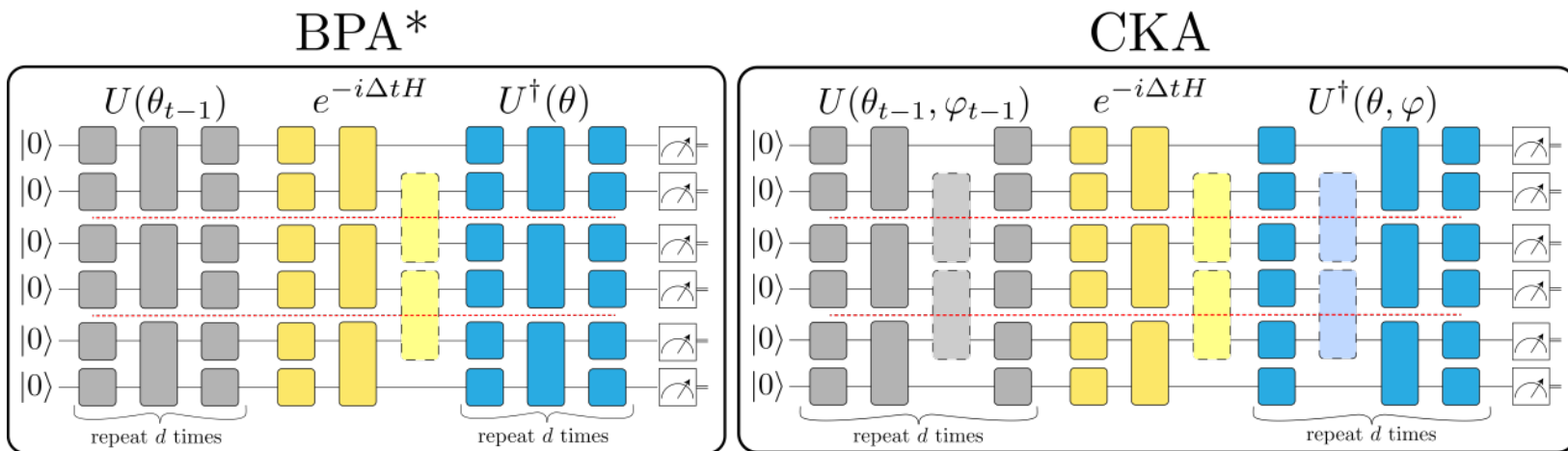


- Block product approximation (BPA): No entangling gates

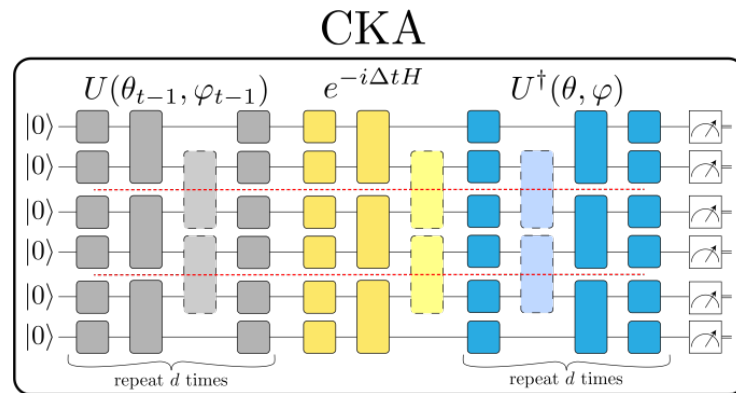
BPA*



- Block product approximation (BPA): No entangling gates
- BPA*: Only Trotter overhead $\omega_{\Delta t} = (1 + 2|\sin(2J\Delta t)|)^{2L}$



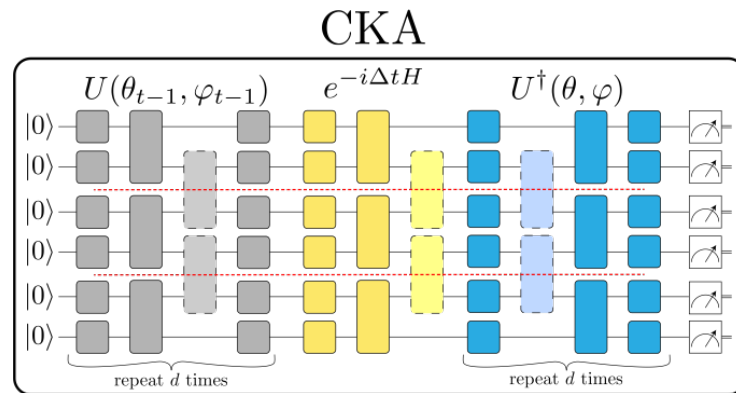
- Block product approximation (BPA): No entangling gates
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- Circuit knitting ansatz (CKA): Introduce entangling gates between blocks



- Overhead is multiplicative

$$\omega(\varphi) = \omega_{\Delta t} \cdot \left(\prod_{i=1}^M (1 + 2|\sin(\varphi_i)|) \right)^2$$

- Need to constrain the overhead!

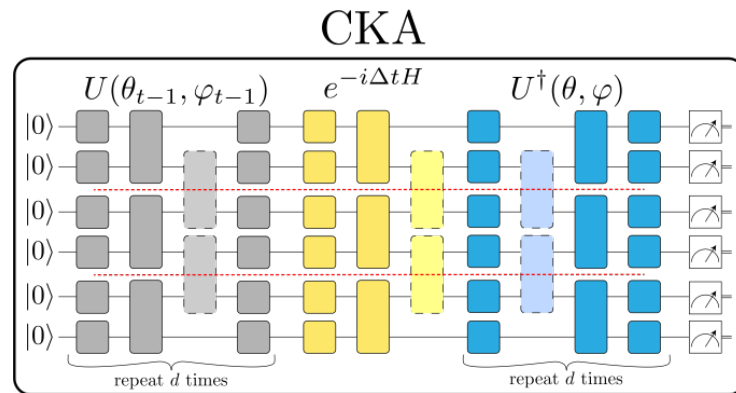


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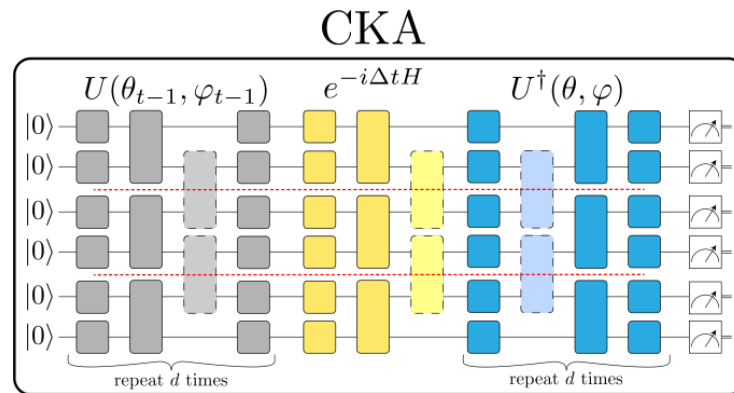
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s.t. $\omega(\varphi) \leq \tau$



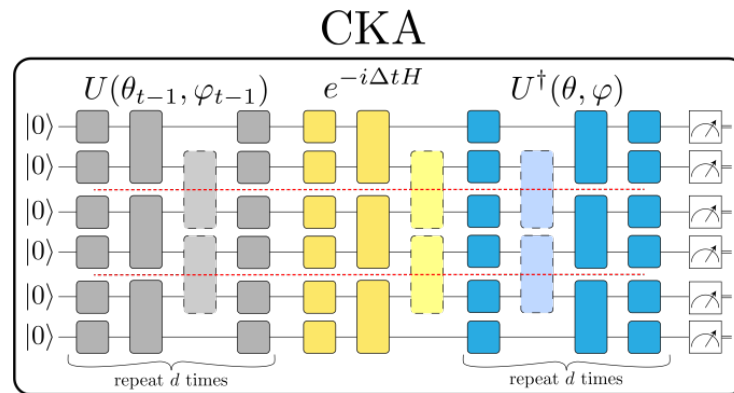
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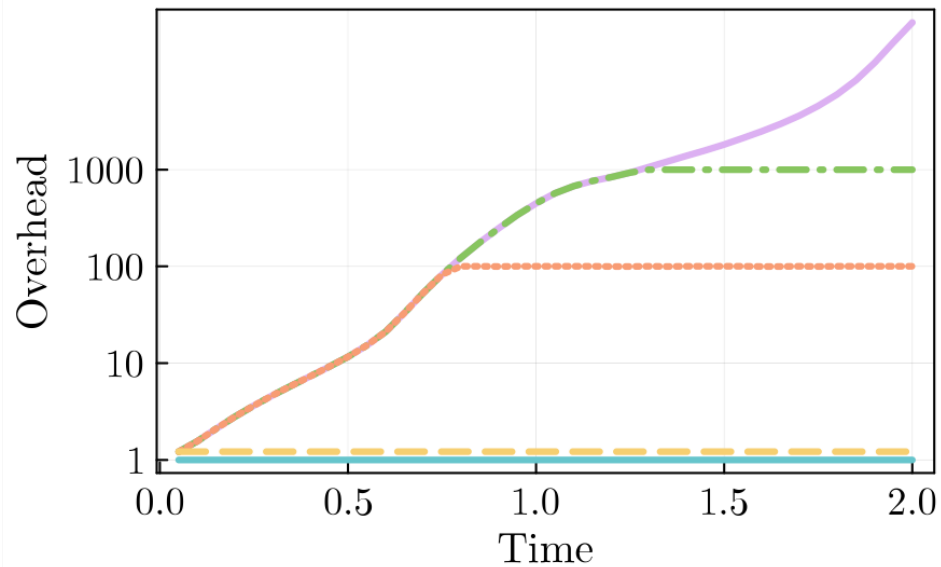
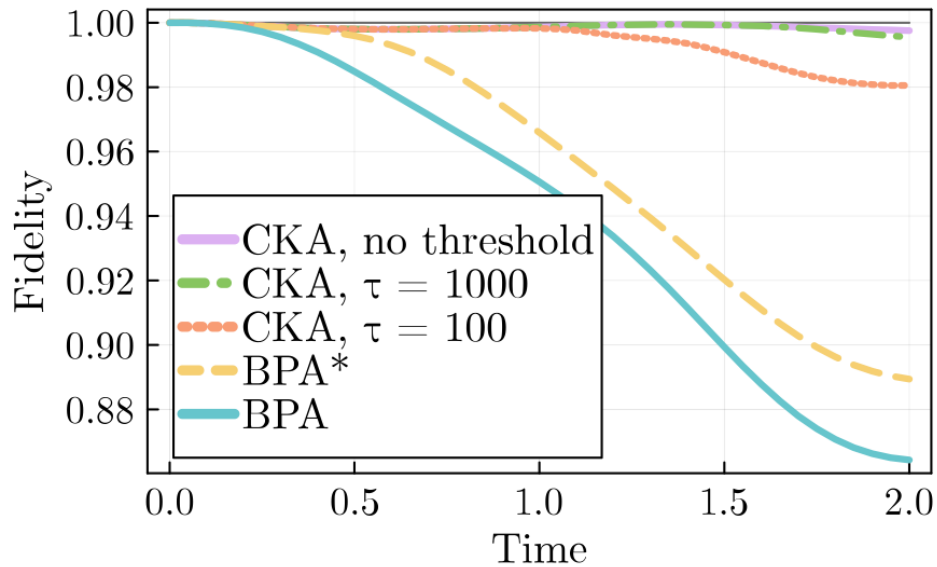
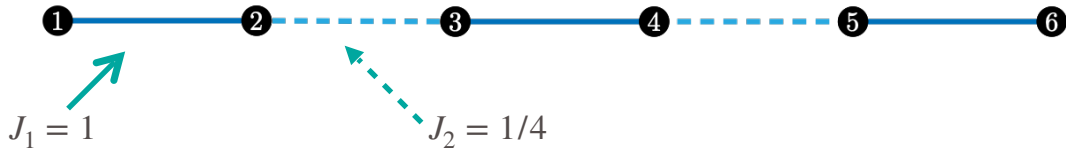
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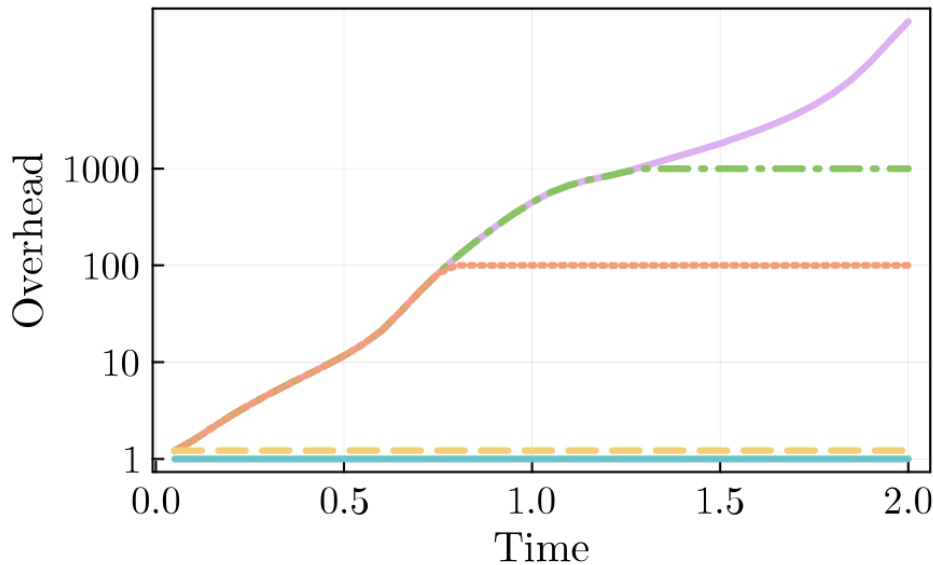
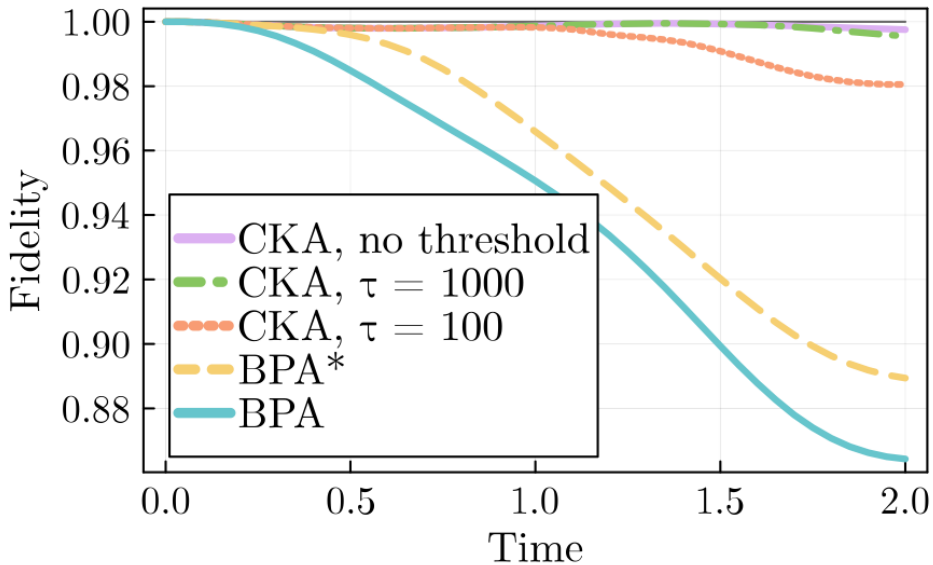
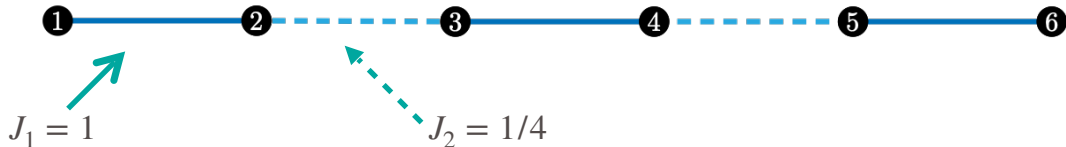
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Higher threshold $\tau \implies$ higher expressibility

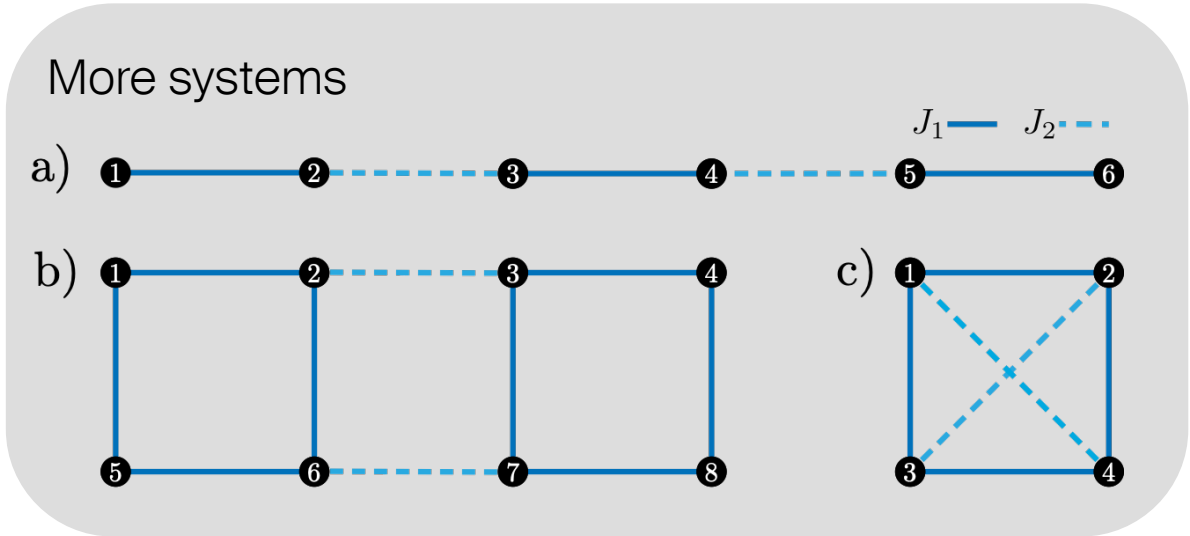




Higher threshold \implies Higher fidelity

Shot noise

Entanglement entropy



GG, Friederike Metz, Giuseppe Carleo (2023),
 arXiv:2309.07857

Conclusion and Outlook

- Simulate dynamics of quantum systems across multiple devices
- Control total budget of shots using overhead threshold

- Can we find cheaper decompositions for the circuit cutting scheme?
- Hardware implementation?
- Application to fermionic systems?

Thank you!



More on the arXiv : [GG, Friederike Metz, Giuseppe Carleo \(2023\)](#)