Quantum Computing Quantum Monte Carlo

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Outline

- Background & motivation
- Quantum Computing QMC
- Sign problem
- Summary & outlook

Background & motivation

 Quantum Computing QMC Sign problem Summary & outlook

Background

Given molecular Hamiltonian $H=\sum_i H_i$, find the ground state of H

Talk on Thursday, we discuss how to solve physical problems using quantum embedding

• Challenges: understanding the problem (interpretation, partition ,ansatz design)

Results: Qubit reduction What about the circuit depth?

The other aspect: given the Hamiltonian, design quantum algorithms

Minimization problem: min θ $E\big(\vec{\theta}\big)$ with $E\big(\vec{\theta}\big) = \big\langle \psi\big(\vec{\theta}\big)|H|\psi\big(\vec{\theta}\big)\big\rangle, |\psi(\vec{\theta})\big\rangle = U\big(\vec{\theta}\big)|\psi_\text{ref}\rangle$ with VQE

Background

VQE approaches, in principle, have several drawbacks: Limited expressivity:

- Due to the presence of various noises, the circuit depths are restricted. Thus, their expressiveness is limited.
- Barren plateau¹: Once the ansatz is too expressive (i.e., reaches unitary 2-design), the gradients of the parameters vanish exponentially.

Challenge of converge and training issue:

- The training of the VQE algorithm is hard³ in general.
- A super-polynomial number of local minimums is presented near the global minimum in the VQE loss landscape, avoiding convergence to the global minimum ².

Intercorrelated

Can we further enhance the performance of NISQ devices (shallow circuits)? Benefit from advanced classical approaches

1 McClean et al, Nat. Commun. 9, 4812 (2018). 2 Anschuetz and Kiani, Nat. Commun. 13, 7760 (2022). 3 Bittel and Kliesch PhysRevLett.127.120502 (2021).

Imaginary time evolution

$$
|\psi(0)\rangle = \sum_{i} \alpha_{i} |E_{i}\rangle
$$

\n
$$
|\psi(\tau)\rangle = \frac{e^{-H\tau}|\psi(0)\rangle}{\sqrt{\langle\psi(0)|e^{-2H\tau}|\psi(0)\rangle}}
$$

\n
$$
|\psi(\tau \to \infty)\rangle = |E_{0}\rangle
$$

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$$
\omega_{\text{challenge: non-unitarity.}}
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\omega_{\text{challenge: non-unitarity.}}
$$

Review on Quantum Monte Carlo

Key idea: stochastically propagating the *walkers* when implementing $e^{-\Delta \tau H}$

Walkers: states in a chosen (complete orthonormal) basis set.

Key challenge: *sign problem* causes exponential sample complexity of the walkers

$$
\langle O \rangle = \text{Tr}(e^{-\tau H}O)/Z
$$

\n
$$
Z = \text{Tr}(e^{-\tau H}) = \int d^{Nd}r \langle r|e^{-\tau H}|r\rangle
$$

\n
$$
= \int dr_0 dr_1 ... dr_M \langle r_0|e^{-\delta \tau H}|r_M\rangle \cdots \langle r_2|e^{-\delta \tau H}|r_1\rangle \langle r_1|e^{-\delta \tau H}|r_0\rangle
$$

Insert complete basis

Projector QMC

Path integral Monte Carlo (QMC): configurational space propagation.

Full Configuration interaction Quantum Monte Carlo (FCIQMC): Fock space propagation.

The sign problem causes great statistical fluctuation.

Current progress of QC-QMC

Key idea of Huggins et al.'s work

Introduce quantum ansatz

Original wave function vs. Importance sampled wave function

Current progress of QC-QMC

Exponential challenge¹ is encountered even in sign-problem-free cases.

The overlap in the denominator could be exponentially small.

the number of samples for estimating the quantity could grow exponentially.

Mazzola and Carleo, Arxiv 2205.09203 (2022).

QC-QMC: Overview

Quantum Computing-Full Configuration Interaction QMC (QC-FCIQMC) algorithm.

- Improve both the VQE and classical QMC methods:
	- \triangleright Provide systematic improvements over the given VQE-prepared state
	- \triangleright A general way for mitigating the sign problem

The walkers are entangled states (opposed to QMC)

- \triangleright Given the VQE-prepared state $|\psi(\vec{\theta})\rangle = U(\vec{\theta})|\psi_{ref}\rangle$, the basis is chosen to be $\{|\phi_i\rangle = U(\vec{\theta})|i\rangle$, $\forall i\}$, where each $|i\rangle$ is a classical product state.
- \triangleright Each state (walker) in the basis is orthonormal since $\langle \phi_i | \phi_j \rangle = \delta_{ij}$.

QC-FCIQMC: Implementation

Key idea: effectively realize the diagonal and off-diagonal terms of the ITE operator in a probabilistic way $e^{-\Delta \tau (H-E_s)} \approx \mathbb{I} - \Delta \tau (H-E_s)$

Given $|\phi_0\rangle = U(\bar{\theta})|0\rangle|0\rangle \equiv |\psi_{\text{ref}}\rangle$. We may initialize the walker set to be $\{|\phi_0\rangle\}$, and repeat

- **Spawn**: For each walker $|\phi_i\rangle$ in the current walker set \rightarrow walkers $|\phi_j\rangle$ with probability $p(j|i)=|-\varDelta\tau\langle\phi_j|H|\phi_i\rangle|/\sum_j|-\varDelta\tau\langle\phi_j|H|\phi_i\rangle|$, $\forall j\neq i$ adjusting the sign of the corresponding walker
- **Death/Cloning**: For each $|\phi_i\rangle$ in the old generation, if $p(i) = \Delta \tau(\langle \phi_i|H|\phi_i\rangle E_s) > 0$, then kill the walker with probability $p(i)$; Otherwise, clone with $-p(i)$
- **Annihilation**: Sum up the new and old generation of walkers

QC-FCIQMC: Implementation

Key techniques: For walker $|\phi_i\rangle$, sample the $|\phi_j\rangle$ s according to $|H_{ji}|^2$, $\forall j$. Denote $\Pi_j = |j\rangle\langle j|$, and let $H = \sum_k h_k P_k$. We have

$$
H_{ji}|^2 = \sum_{kk'} h_k h_{k'} p_{kk'}^i(j), \text{ s.t. } \sum_j p_{kk'}^i(j) \le 1,
$$

where $p_{kk'}^i(j) = \text{Re}\langle i|U^\dagger P_k U \Pi_j U^\dagger P_{k'} U|i\rangle.$

Figure 1: Circuit samples $p_{kk'}^i(j)$

Figure 2: Circuit estimates H_{ji} .

Sign problem

An example:

Consider the partition function of the thermal state in the form of path integrals,

$$
Z = \text{Tr}\left(e^{-\beta H}\right) = \sum_{n=0}^{\infty} \frac{(-\beta)^n}{n!} \text{Tr}\left(H^n\right)
$$

$$
= \sum_{n=0}^{\infty} \sum_{i_1, \dots, i_n} \frac{(-\beta)^n}{n!} \langle i_1 | H | i_2 \rangle \langle i_2 | H | i_3 \rangle \cdots \langle i_n | H | i_1 \rangle
$$

$$
= \sum_{n=0}^{\infty} \sum_{i_1, \dots, i_n} p(i_1, \dots, i_n) = \sum_c p(c)
$$

Each path integral forms a closed path

The sign problem:

- Presents whenever $p(c)$ contains negative terms. E.g, fermions, negative weights arise from the Pauli exclusion principle under particle exchange.
- Basis dependent.

Troyer and Wiese, PhysRevLett.94.170201 (2005)

Origin of the sign problem

For arbitrary observable A, the standard way for evaluation is to sample w.r.t. the bosonic system with $|p(c)|$ and $s(c) = sign(p(c))$, we have

$$
\langle A \rangle = \frac{\sum_{c} A(c) s(c) |p(c)| / \sum_{c} |p(c)|}{\sum_{c} s(c) |p(c)| / \sum_{c} |p(c)|} = \frac{\langle A \rangle'}{\langle s \rangle'},
$$

The average sign decays exponentially with the free energy difference Δf ,

$$
\langle s \rangle = \frac{z}{z'} = \exp(-\beta N \Delta f)
$$

Fermionic and corresponding bosonic system

$$
z = \sum_{c} p(c), z' = \sum_{c} |p(c)|
$$

Thus, the relative error grows exponentially:

$$
\frac{\Delta s}{\langle s \rangle} = \frac{\sqrt{\langle s^2 \rangle - \langle s \rangle^2 / M}}{\langle s \rangle} = \frac{\sqrt{1 - \langle s \rangle^2}}{\sqrt{M} \langle s \rangle} \sim \frac{e^{\beta N \Delta f}}{M}
$$

 M is the number of samples.

Stoquasticity and mitigation of the sign problem

Definition (Stoquastic Hamiltonian^{*a*})

^aBravyi et al. 2006.

For any Hamiltonian H, s.t. $H_{ij} \leq 0, \forall i \neq j$. These types of Hamiltonian shall not present sign problem, *i.e.*, all terms in the path integrals are positive.

Proof idea: Denote $G = \alpha I - H$ with $\alpha = \max$ $\mathop{ax} H_{ii}$, $Z = Tr(e^{-\beta H}) = e^{-\beta \alpha} Tr(e^{\beta G})$ $(G_{ij} \ge 0)$ Thus, stoquastic Hamiltonians are sign-problem-free

Mitigation: The sign problem is basis dependent. A universal approach for suppressing the sign problem is by similarity transformation of H, s.t, to approach **stoquastic Hamiltonians** $H(U) = U^{\dagger} H U$

Remark: By expanding the wave function in the VQE-unitary-rotated basis set, *our method effectively implements the similarity transformation with* $U = U(\vec{\theta})$ prepared by VQE

Measure of the sign problem

Relative deviation from the stoquastic Hamiltonians

Definition (Non-Stoquastic Indicator)

Denoting bosonic form of H as \tilde{H} :

 $\tilde{H}_{ij} = H_{ij}$ if $i = j$ or $i \neq j$ and $H_{ij} < 0$; $\tilde{H}_{ij} = -H_{ij}$ if $i \neq j$ and $H_{ij} > 0$.

The NSI is defined as

$$
S(H)=\frac{\text{Tr}[e^{-\beta \tilde{H}}]-\text{Tr}[e^{-\beta H}]}{\text{Tr}[e^{-\beta H}]}.
$$

Bosonic form: stoquastic Hamiltonians

Remark: The NSI is computationally non-trivial.

Measure of the sign problem

Theorem (Non-Stoquastic Indicator upper bound)

$Define$

 $(H_{-})_{ij} = H_{ij}$ if if $i = j$ or $i \neq j$ and $H_{ij} < 0$; wanted $(H_+)_{ij} = H_{ij}$ if $i \neq j$ and $H_{ij} > 0$. Unwanted

The NSI is upper-bounded by

 $S(H) \leq 2e^{\beta ||(\alpha - H_{-})||_{L_1}} \sinh(\beta ||H_{+}||_{L_1}).$

Here, $||M||_{L_1} := \sum_{i,j} |M_{ij}|$ is the matrix norm for matrix M.

Remark:

- ≻ Mitigate the sign problem needs minimize H_+ an H_- simultaneously.
- \triangleright When H_+ vanishes completely, $S(H) = 0$ as expected for stoquastic Hamiltonians.
- ➢ NSI does not sufficiently guarantee the performance

Proof sketch

1. Expand the NSI into path integral form

$$
S(H) = \text{Tr}(e^{-\beta \tilde{H}}) - \text{Tr}(e^{-\beta H})
$$

= $\text{Tr}(e^{-\beta(\alpha I - \tilde{G})}) - \text{Tr}(e^{-\beta(\alpha I - G)})$
= $e^{-\beta \alpha} \left(\sum_{k=0}^{\infty} \frac{\beta^k}{k!} \sum_{\{|\phi\rangle\}} (|\langle \phi_0 | G | \phi_1 \rangle \langle \phi_1 | G | \phi_2 \rangle \cdots \langle \phi_{k-1} | G | \phi_0 \rangle| - \langle \phi_0 | G | \phi_1 \rangle \langle \phi_1 | G | \phi_2 \rangle \cdots \langle \phi_{k-1} | G | \phi_0 \rangle) \right)$

2. Define $G_+ = \alpha - H_+$ and $G_- = \alpha - H_-$. The above subtraction is supposed to be -2 -fold of sum of all negative path integrals. all possible negative terms can be enumerated by combination of terms in G_+ and G_- :

$$
=2\sum_{k=0}^{\infty}\frac{(\beta)^{2k+1}}{(2k+1)!}\left(\binom{2k+1}{1}\|G_{-}\|_{L_{1}}\|G_{+}^{2k}\|_{L_{1}}+\binom{2k+1}{3}\|G_{-}^{3}\|_{L_{1}}\|G_{+}^{2k-2}\|_{L_{1}}+\cdots+\binom{2k+1}{2k+1}\|G_{-}^{2k+1}\|_{L_{1}}\right)\\+2\sum_{k=0}^{\infty}\frac{(\beta)^{2k}}{(2k)!}\left(\binom{2k}{1}\|G_{-}\|_{L_{1}}\|G_{+}^{2k-1}\|_{L_{1}}+\binom{2k}{3}\|G_{-}^{3}\|_{L_{1}}\|G_{+}^{2k-3}\|_{L_{1}}+\cdots+\binom{2k}{2k-1}\|G_{-}^{2k-1}\|_{L_{1}}\|G_{+}\|_{L_{1}}\right)
$$

Group by odd and even terms

- 1. By property of L1 norm such that $||A^m||_{L_1} \leq ||A||_{L_1}^m$, we can relax each term above by $||G^a_ L_1$ || G_+^b $L_1 \leq ||G_-||_{L_1}^a ||G_+||_{L_1}^b$.
- 2. Finally, by the binomial theo $\frac{1}{2}((b+a)^n (b-a)^n) = \sum_{i=1,odd}^n {n \choose i} a^x b^{n-x}$, we arrive at our results.

Verification by numerics

Verification by numerics

Results vs classical FCIQMC algorithms.

Distribution of the walkers become more concentrated to the initial walker $|\phi_0\rangle$ as it get closer to the ground state

• fewer walkers

Summary & outlook

Summary

- Introduce the QC-FCIQMC algorithm that improves the ability of shallow-depth quantum circuits and each state (walker) could be prepared by shallow depth circuit available by NISQ devices.
- Systematic suppression of the sign problem is achieved if the basis (set of walkers) is refined.
- NSI to characterize severity of the sign problem (not sufficient for performance guarantee)
- Limitation: our method demands a huge number of measurements for sampling the walkers that could be challenging for NISQ devices.

Possible future works

- The upper bound on NSI can serve as a low-cost loss function for easing the sign problem for classical QMC
- Explore other kinds of unitary construction for mitigating the sign problem, one with performance guarantee
- Performance analysis under certain noise channels
- Compatibility with classical shadows to reduce measurement costs