

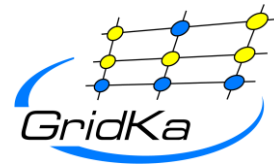
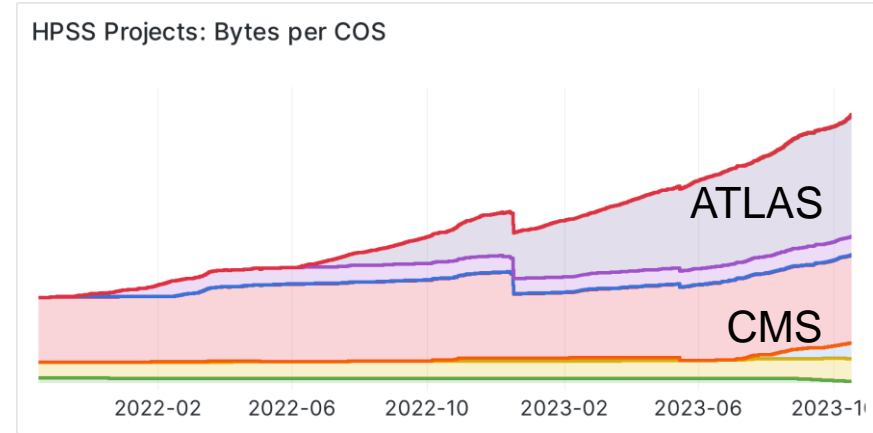
Exploring Technologies for HPSS Disk Caches

Dorin Lobontu, Preslav Konstaninov, Andreas Petzold, Doris Ressimann



HPSS at KIT

- Serves as tape system platform for Tier-1 and bwDataArchive
- ~100PB used today
- 4 tape libraries
- two sites ~8x J-E lengths apart



HPSS Tier-1 Integration

- Migration TSM → HPSS since 2020
 - Outside of dCache/xrootd
 - 20-40 streams (write+read for checksum)
- Writing to tape
 - dCache/xrootd pools call HPSS client locally
 - 40 streams per pool
- Reading from tape
 - dCache ENDIT provider + new HPSS-specific scheduling backend per VO
 - Backend sorts files per aggregate, triggers stage, copies files into pools
 - Plan to replace file-based provider-backend communication

Why another disk cache?

- Disk cache required for
 - Aggregation when writing
 - Transparent to client
 - Pack files into ~300GB aggregates
 - reduce number of tape marks
 - 380MB/s write speed per drive
 - Full Aggregate Recall (FAR)
 - Request for one file triggers recall of full aggregate
 - 400MB/s read speed per drive

Workload on Cache

■ Tier-1 writing to tape

- Write from client + read from client for checksum
- Writing to tape: read ~same as write from client



2:1 read:write

■ Tier-1 reading from tape

- Read from tape: write on cache one stream per drive
- Read from client: read from cache








1:1 read:write

HDD-based Setup

- 2 NetApp E5700 w/ 120 8TB drives each (~1.4PB usable)
 - Expect ~12GB/s per system with 70% read workload
- Observations
 - Never close to 24GB/s
 - System prioritizes writes → reads starved → writing to tape slow
 - DDP vs. RAID6 vs. RAID10 no big difference

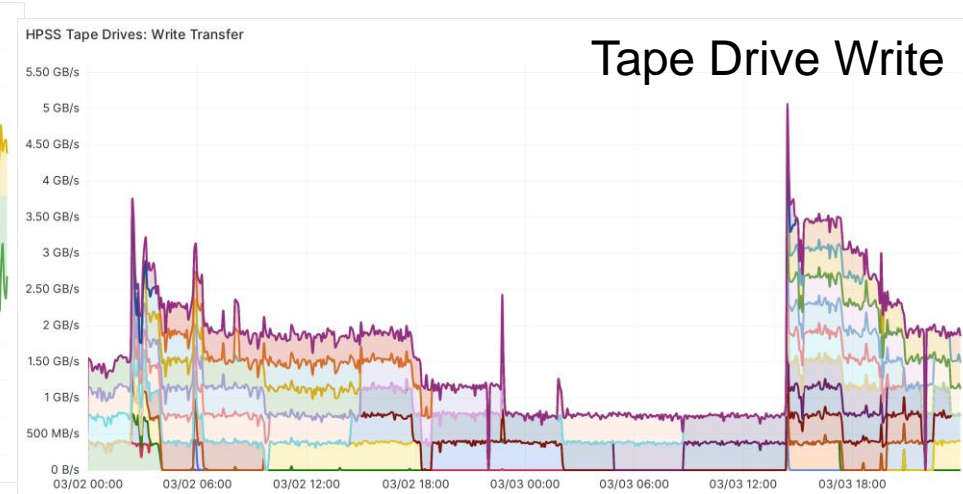


Workload on Cache

- Tier-1 writing to tape
 - Write from client + read from client for checksum
 - Writing to tape: read ~same as write from client  2:1 read:write
- Tier-1 reading from tape
 - Read from tape: write on cache one stream per drive
 - Read from client: read from cache  1:1 read:write
- Random I/O to/from clients  Random IO only
- Sequential I/O to/from tape drives  Random IO only
- Streams to tape drives need to be stable  more IOPS help

SSDs

- Added 2 Dell ME5024 + Extension Enclosures with 2x 48 3.84TB SSDs
 - ~250TB usable space
- Better latencies → much improved tape write rates
- Limited throughput of ME5024 controllers → isn't there something better?



Cache Requirements

- Low latencies → Flash/NVMe
- High throughput → AFA+NVMe or NVMe in server
- Storage redundancy → AFA or other NVMe-optimized RAID

- Big vendor AFA way too expensive
- NVMe-optimized RAID solutions
 - GRAID SupremeRAID: GPU-accelerated RAID
 - Xinnor xiRAID: software RAID

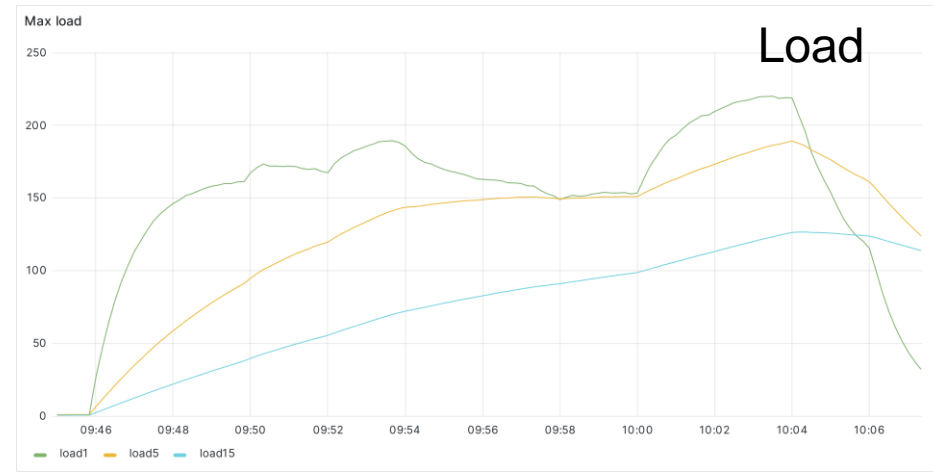
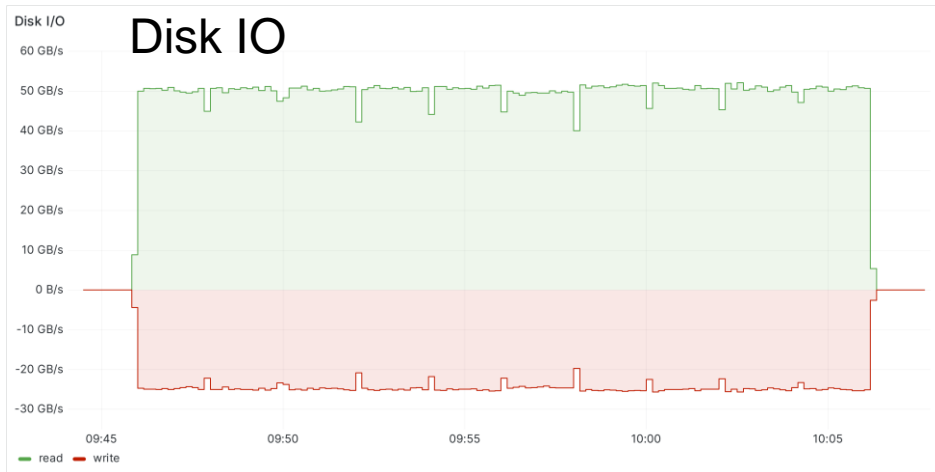
Setup

- 2U Supermicro AS-2015CS-TNR
 - Single AMD EPYC 9554P 64-Core 3.1GHz
 - 512GB RAM
 - 10x 30TB Micron 9400 NVMe devices (7GB/s)
 - 4x 100Gbit/s Ethernet
-
- xiRAID licensed per device used in RAID6
 - ➔ NVMe name spaces have to be licensed too :-)
 - Single xiRAID6 with several regular LVs on top
 - ~240TB usable space
 - LVs needed due to HPSS IO connection limits per disk device



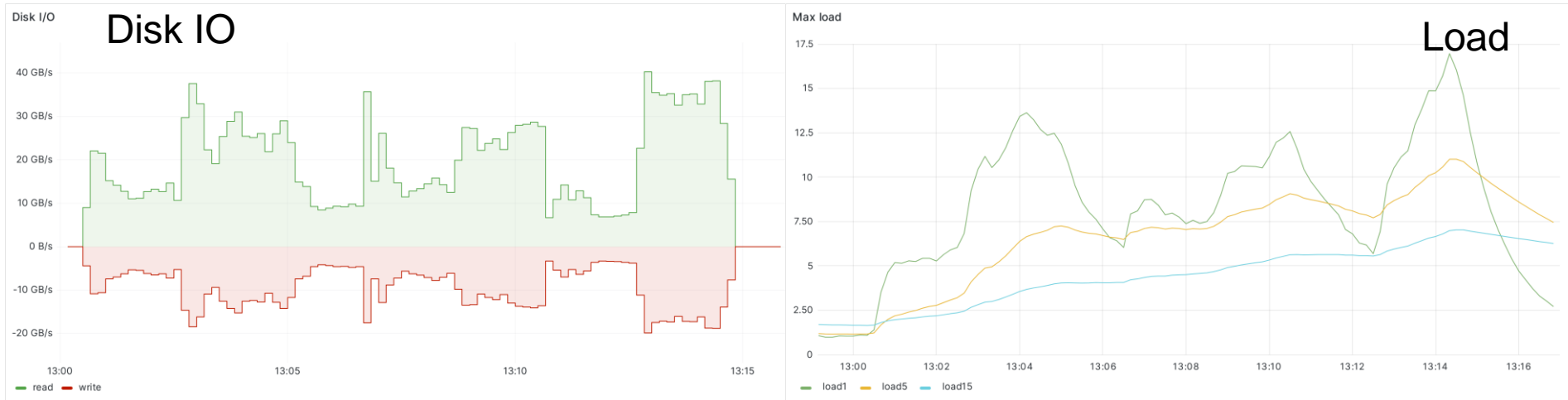
Benchmarks 1

- fio benchmarks with different block sizes/file sizes/number of clients always with 2:1 read:write ratio
- Monitor throughput, CPU load



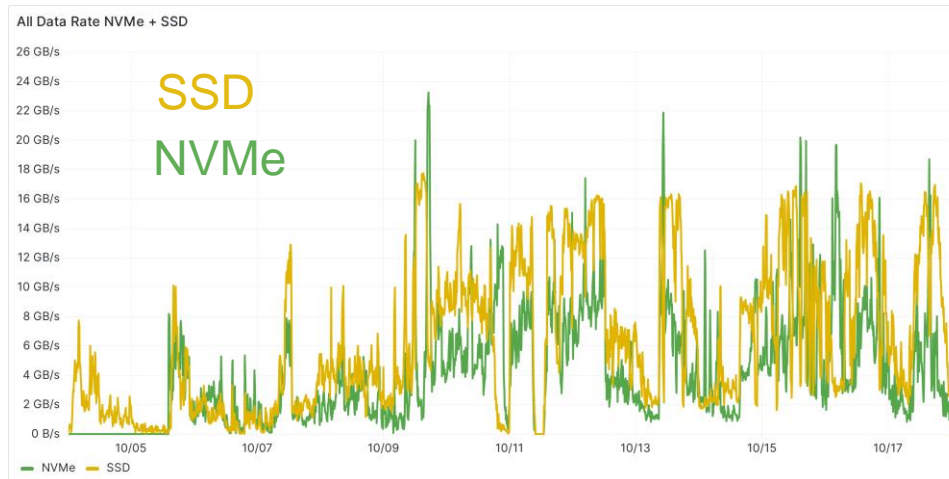
Benchmarks 2

- Same benchmark runs, but limit xiraid to 64 threads



In production

- Since two weeks, next to SSD systems
- No interference between xiRAID and HPSS mover process at current workload level
- SSDs still receive larger share of IO → HPSS tuning required



Summary and Outlook

- Workload on tape system cache requires lots of IOPS
 - Many disks or flash
- Servers with large local NVMe storage powerful and cost effective
 - Excellent latencies and throuput
- Redundancy requirement
 - Proprietary RAID solutions or simple mdraid RAID1?
- Would like to test PCIe connected storage enclosures
 - Decouple NVMe devices from servers