## **High Performance Graph** Segmentation for ATLAS **GNN Track Reconstruction**

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## Track Reconstruction & the GNN4ITk Pipeline



## Tracking in ATLAS HL-LHC Inner Tracker (ITk)

#### [ATL-PHYS-PUB-2021-024]



- Track finding requires associating each hit to a track candidate
- Number of hits per  $pp \rightarrow t\bar{t}$  event: 311,000 +/- 35,000
- · Number of particles per  $pp \rightarrow t\bar{t}$  event: 16,000 +/- 1,700
- Innermost pixel layer 25x100  $\mu m^2$ , all other pixel layers 50x50  $\mu m^2$
- · Strip layers are at millimeter resolutions
- We focus on Athena simulation in the following slides

# ATL-PHYS-PUB-2021-024]



UNSG-2021-57

## GNN4ITk Pipeline



- Pipeline receives clusters = collections of energy deposits on silicon. These are associated with 3D spacepoints, to be used as nodes for stage 1 onwards
- Out of stage 3 we obtain a set of track candidates, each is an unordered set of spacepoints
- For processing in Athena track fitting chain, we associate these back to the original clusters, and order in increasing distance from beamspot origin

## Track Reconstruction Performance



- Tracking efficiency compared with current combinatorial kalman filter (CKF) technique
- Behaviour across  $\mathbf{\eta}$  and  $\mathbf{p}_{\tau}$  similar to CKF good sanity check!

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## HL-LHC Offline & Online Track Reconstruction Needs

•

	LHC Run 3	HL-LHC
L0 trigger accept	100 kHz	1 MHz
Event Filter accept	1 kHz	10 kHz
Event size	1.5 MB	4.6 MB

•	Event filter	(high leve	l trigger)	contains tracking
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- CPU-based Fast Tracking: 23.2 HS06s/event (around 1 second per single-core CPU), small drop in track efficiency: 1-2% on average, 5% for pT in [1,1.5]GeV
  - GPU-based GNN4ITk pipeline: First two steps run in 400-600 ms/event. But final step has previously taken around **42 seconds** to run

$\langle \mu \rangle$	Tracking	Release	Byte Stream	Cluster	Space	Si Track	Ambiguity	Total
			Decoding	Finding	Points	Finding	Resolution	ITk
140	default	21.9	2.2	6.4	3.5	31.6	43.4	87.1
	fast			6.1	1.0	13.4	-	22.7
200	default	21.0	2.2	8.3	4.9	66.1	64.1	146.6
	fast	21.9	5.2	8.1	1.2	23.2	-	35.7

CPU-based Fast tracking vs Default tracking timing (HS06 x s) [ATLAS-TDR-029-ADD-1]

Stage	Pipeline			
	Metric Learning	Module Map		
1. Graph Construction	505 ms	69 ms		
2. Edge Scoring	108 ms	323 ms		

**CNN4ITk pipeline execution times of first two stages, per event [ATL-PHYS-PUB-2024-018]** 

## Graph Segmentation



## Stage 3: Graph Segmentation

The task:

- Given a collection of hits in detector and directed edges (a hypothesis that the two hits were created successively by the same particle), each with a score...
- Produce a set of track candidates (sets of nodes)
- Some nodes may belong to background tracks, some may be noise, some may be shared between target tracks



## Stage 3: Graph Segmentation



- There are many ways to solve this problem
- In the case where we do not allow shared nodes, this is a classic *graph partitioning* problem



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Would like to produce: [1, 2, 3, 4, 5, 6] [7, 2, 8, 9, 10] [11, 12, 13, 14]

### Recall that we have the GNN predicted edge scores:











- The simplest idea is (weakly) Connected Components
- All nodes belonging to the same component after removing low-scoring edges, are assigned to a track
- That is, the nodes must be reachable via an undirected path (hence "weakly")



- This is an extremely fast algorithm with the Scipy implementation, requiring **only a few milliseconds for O(1000) components,** on a single thread
- Each node and edge visited only once, giving good scaling performance: O(N+E)
- Iterative Depth-first Search (DFS) means predictable and reliable memory behaviour (as opposed to recursive DFS)



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- Iterative Depth-first Search (DFS) means predictable and reliable memory behaviour (as opposed to recursive DFS)
- However note that (as in example), it can merge tracks into a single candidate if the score threshold is too low
- We can raise the threshold, but then we are just as likely to *split* a track in multiple candidates
- Both lead to low track finding efficiency

All connected components (of size > 3) are assigned to a track



## Walking through the Hit Graph



## Walkthrough Algorithm



- To avoid merging tracks, we want out candidates to be "chain-like", with maximum one incoming and one outgoing edge
- The walkthrough algorithm is inspired by the way traditional track finding is performed

• It performs well, but is much slower than CC



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## Walkthrough Algorithm



- Repeat steps 4 and 5 until all nodes have been used, or have no neighbours in the graph
- Clearly, as this is sequential, some tracks will arbitrarily be built that split later tracks





Precisely, the walkthrough requires the following sub-algorithms:

- 1. Remove Cycles
- 2. Filter Graph
- 3. Extract Chains
- 4. Topological Sort
- 5. Build Paths









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- Remove Cycles: Ensure the directed graph is acyclic by directing edges in increasing R, scaling as ~ O(E)
- 2. Filter Graph
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- 5. Build Paths



## Graph with number of nodes N, number of edges E

Precisely, the walkthrough requires the following sub-algorithms:

- 1. Remove Cycles
- 2. Filter Graph: Remove low-scoring edges  $\sim O(E)$
- 3. Extract Chains
- 4. Topological Sort
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21

Precisely, the walkthrough requires the following sub-algorithms:

- 1. Remove Cycles
- 2. Filter Graph
- 3. Extract Chains: Apply connected components, and any chain-like components are instantly submitted as candidates ~ O(N+E)
- 4. Topological Sort
- 5. Build Paths



# Precisely, the walkthrough requires the following sub-algorithms:

- 1. Remove Cycles
- 2. Filter Graph
- 3. Extract Chains
- 4. Topological Sort: Order node indices such that earlier nodes can visit subsequent nodes if a path exists between the two  $\sim O(N+E)$
- 5. Build Paths



Implemented as a Breadth-first Search (visit each edge only once): Kahn's Algorithm\*

\* A. B. Kahn. 1962. Topological sorting of large networks. Commun. ACM 5, 11 (Nov. 1962), 558-562. https://doi.org/10.1145/368996.369025

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Precisely, the walkthrough requires the following sub-algorithms:

- 1. Remove Cycles
- 2. Filter Graph
- 3. Extract Chains
- 4. Topological Sort
- 5. Build Paths: For each starting node, choose longest path - this is a DAG source-to-sink algorithm. Currently no heuristic to choose between equally-long paths ~ best case O(N), worst case O(N!)



Precisely, the walkthrough requires the following sub-algorithms:

- 5. Build Paths: To trade off between best case O(N) and worst case O(N!), we do a sort of beam search:
  - a. The highest scoring neighbour is always considered for the longest path (regardless of edge score)
  - b. Any neighbours with edge score > 0.6 are considered for the longest path

Reduces combinatorics significantly, while avoiding splitting tracks (every node will have at least one outgoing edge)



Precisely, the walkthrough requires the following sub-algorithms:

- 1. Remove Cycles
- 2. Filter Graph
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- 5. Build Paths

To produce public physics results, most of these steps used NetworkX implementation. This took *O(minutes)* per event - we need something faster!



## Optimisations to Segmentation



## FastWalkthrough

A variety of improvements can be made:

- 1. Remove Cycles: Unchanged
- Filter Graph: Use Pytorch Geometric (PyG) graph representation
- 3. **Extract Chains**: Use Scipy CC and PyG scatter\_max across components
- 4. **Topological Sort**: Pure python+numba implementation
- 5. **Build Paths**: Pre-process graphs to get max+min edges (with scatter operations), then pure python+numba implementation

(Timing is per-event)



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Chain-like: all hits must have in-degree and out-degree less or equal to 1



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- The preprocessing of junctions in FastWalkthrough can be extended to any heuristics For example, first remove chain-like components as usual
- Then remove all "junctions" (nodes with more than one outgoing or incoming edge)
- Then re-run connected components



Connected Components

- One could apply many kinds of choices
- E.g. only remove "X-junctions", but allow "Y-junctions"



Junction NOT removed (in-degree ≠ out-degree)



- The best physics performance actually came simply from choosing the highest-scoring incoming and outgoing edges
- Then apply connected components, as usual
- This is Connected Components + Junction Removal Version 3 (CC+JR v3)



## Comparison of Segmentation Approaches

We are able to reduce the running time of this stage (compared with that used to produce public physics results "CTD23") by 3 orders of magnitude, and *increase* the physics performance!

Stage	Efficiency (Relative Difference, %)	Running Time (ms)
CTD23 Walkthrough		42,000
FastWalkthrough	+0.53	120
CC	-1.33	6.0
CC+JR	+0.93	40

Per-event execution times and relative difference in integrated physics efficiencies of the various graph segmentation techniques available in Stage 3 (graph segmentation). Differences are calculated relative to the baseline CTD23 Walkthrough as (eff - CTD)/CTD. The score cut on CC set to 0.01, with the minimum and additive thresholds of walkthroughs set to 0.1 and 0.6 respectively. The running times are evaluated on a single CPU core (AMD EPYC 7763). CTD23 Walkthrough is the same as that used in IDTR-2023-06.

#### [ATL-PHYS-PUB-2024-018]









## Next Steps

- Small tweaks required to run on GPU (Pytorch, PyG support GPU out-of-the-box, numba requires some massaging) - expect to reduce existing CC+JR version closer to 10ms
- Integration with production systems Athena and ACTS
- C++ implementation needed for integration, and expected to be somewhat faster than PyG+numba
- Current FastWalkthrough and CC+JR still involve duplication (e.g. running across number of edges E with CC ~O(E) then JR ~O(E))
- Fusing these steps will reduce again by some factor
- Expect to obtain an algorithm of >1KHz on single GPU with native CUDA