

Conference on Computing in High Energy and Nuclear Physics

19-25 October 2024

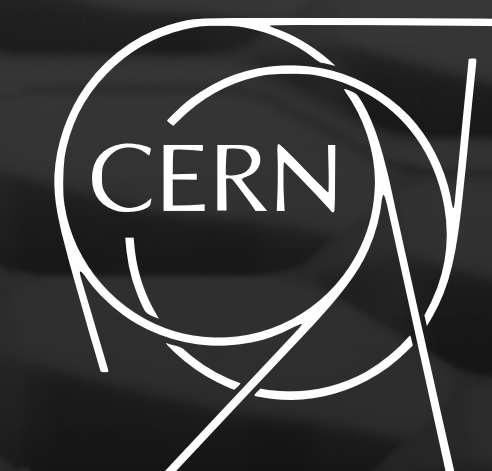
Kraków, Poland



Evolution of the CERN Tape Archive scheduling system



Dr. Jaroslav Guenther (CERN)
on behalf of the CERN IT CTA Team





Scheduling Evolution in CERN Tape Systems

~1990 SHIFT

manual scheduling



~1998 STK libraries

robotic scheduling, namespace → CERN IT
repack introduced !

next ~22 years CASTOR

centrally deployed scheduling



~2020 CTA

distributed multi-threaded scheduling
using ObjectStore technology



~ 2024

preparation for **Run 4**

CTA adopted by **wider community** !
new challenges ahead



Scheduling DB Stores Transient Metadata



Disk buffer
XrootD SSI / gRPC
WorkFlow Engine

request

CTA Frontend
XrootD SSI / gRPC
Scheduler thread

data

CTA Tape Server
tape drive daemons
Scheduler thread



Scheduler DB
Ceph ObjectStore

Catalogue DB
Oracle
PostgreSQL

transient
metadata

permanent
metadata
tape file
namespace



Scheduler DB Implemented as ObjectStore

Architecture

- **motivated by performance for archival and retrieval queueing operations**
- **multi-threaded interface to Ceph**
- **protobuf serialised objects in key/value store**
(archive, retrieve requests and queues)
- **code design ensures**
 - **high performance**
 - **scaling**
 - **reliability**
(despite > storage round trips than DB)
- **delivered very well for Run 3**

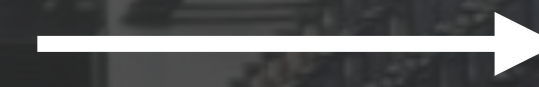




Development Cost

- **complex distributed transaction management system**
- **high maintainability cost**
 - additional dependency
 - requires extensive learning effort
- **scheduler logic is tightly coupled to ObjectStore implementation**
 - lack of indexes is a serious constraint on implementation of scheduler algorithms (e.g. cancelling requests)
 - as we scaled up to ~200 tape drives, global locking caused scheduler contention
- **workflows beyond original design proven difficult**

but



design allows multiple backends

and

code complexity reduction possible

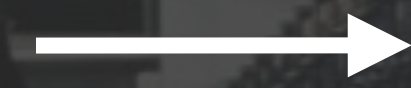
and

we have off-the-shelf solution = Relational DB



Operational Cost

- **ObjectStore requires additional technology expertise**
 - sites which do not run Ceph
 - FTS scheduling using Postgres
- **complex backend object structure**
 - **object introspection, forensics and cleanup is difficult**
 - "schema" updates difficult to manage
- **high priority fixes still required several times a year**
(e.g. for object deletion, empty shard handling, infinite loops, global locking issues, repack exhausting resources, object size handling, etc.)



we can consolidate on common technologies

and

have Relational DB rows for operators

and

simpler operational tools



Advantages of Relational DB as the Solution

New Scheduler DB

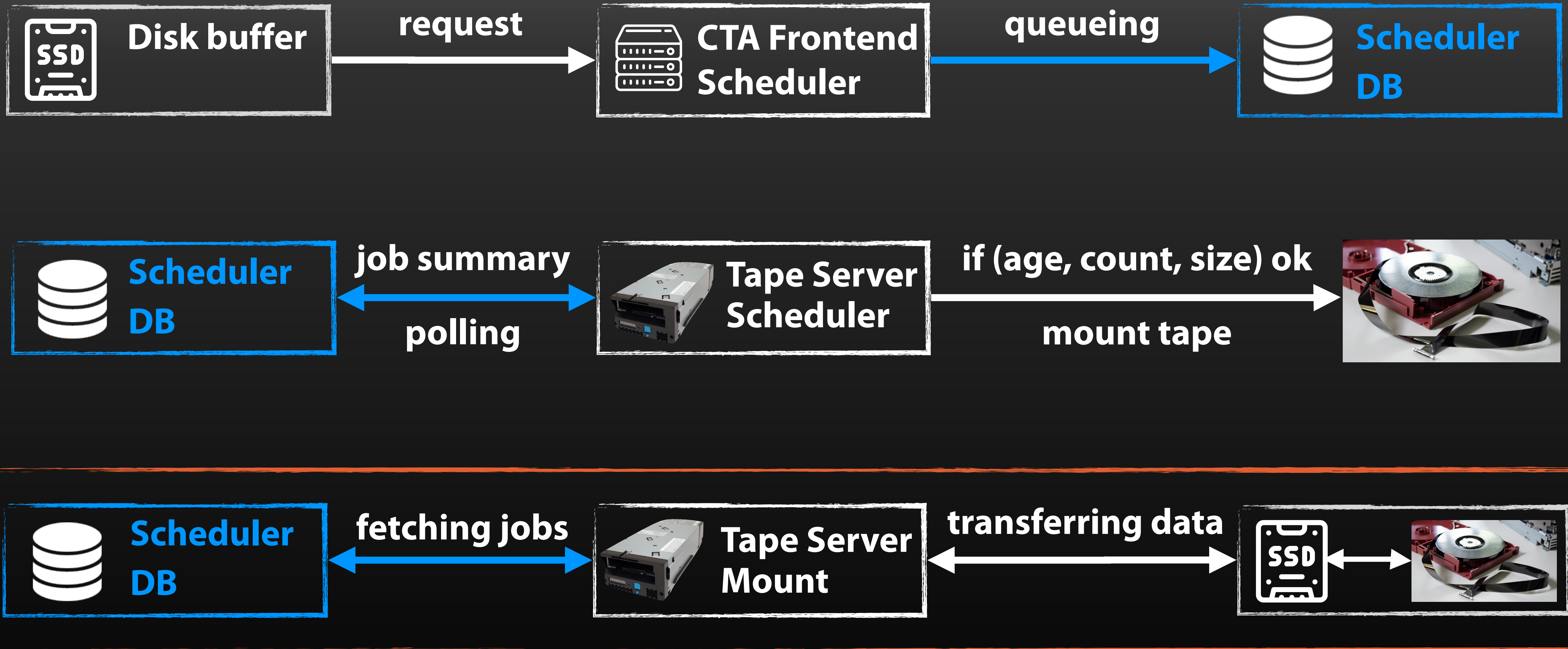
- **lower code complexity**
 - ▶ **no overhead of transactional management code**
MVCC, indexing (+sync) "for free" from DB
- **straightforward schema updates**
- **multi-index queues → more flexibility to improve and extend Scheduling algorithms**
- **extensible to multiple database backends**
(PostgreSQL DB, Oracle DB)
- **no global scheduling lock**

Overview of Scheduling Workflow

1/2



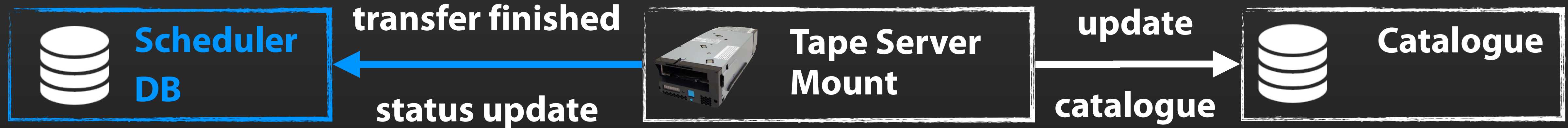
time



blue = scheduler DB tasks

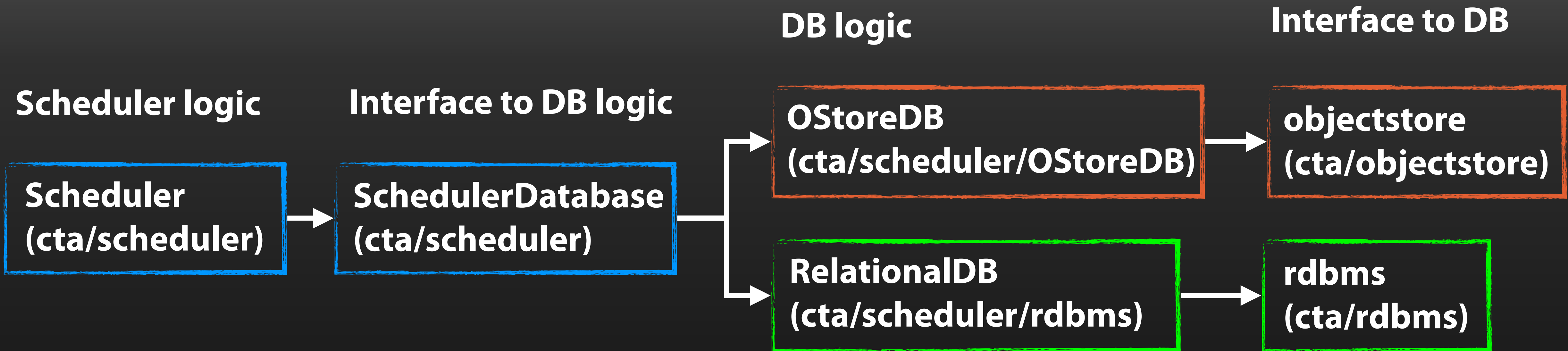


time



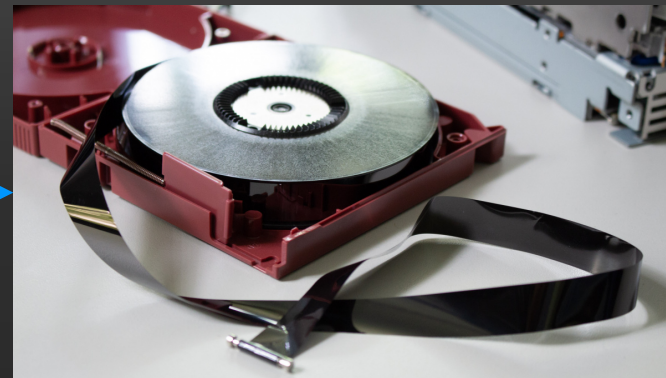
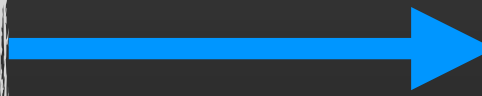
blue = scheduler DB tasks

Scheduler DB Implementation Architecture



Design Allows Multiple Scheduler DB Implementations

- use CTA generic RDBMS interface for Postgres DB implementation
- workflow oriented tables (Archive/Retrieve/Repack), views, sequences, etc.
(file transfer job = row in a table)



Archival

- ✓ queueing
- ✓ job summaries per mount decision
- ✓ job fetching for transfer
- ✓ management of failures and retries
- ✓ transfer status reporting
- ✓ improvements in CTA rdbms layer

Lower Granularity Locking

- ✓ drive deciding to mount a tape
 - lock per logical library
(prevents empty mounts)
 - ✓ fetching the jobs from the queue
 - lock per tape pool per workflow
(avoids interwoven row sets per drive)
- ✓ no global lock on scheduling anymore !**

Functional Testing of Archival Workflow



Setup

- Full System deployment in Minikube (CTA+EOS)
 - 1 disk for buffer (EOS MGM & FST)
 - 1 disk for tape drive
- External Catalogue: Oracle DB
- External Scheduler: PostgresDB

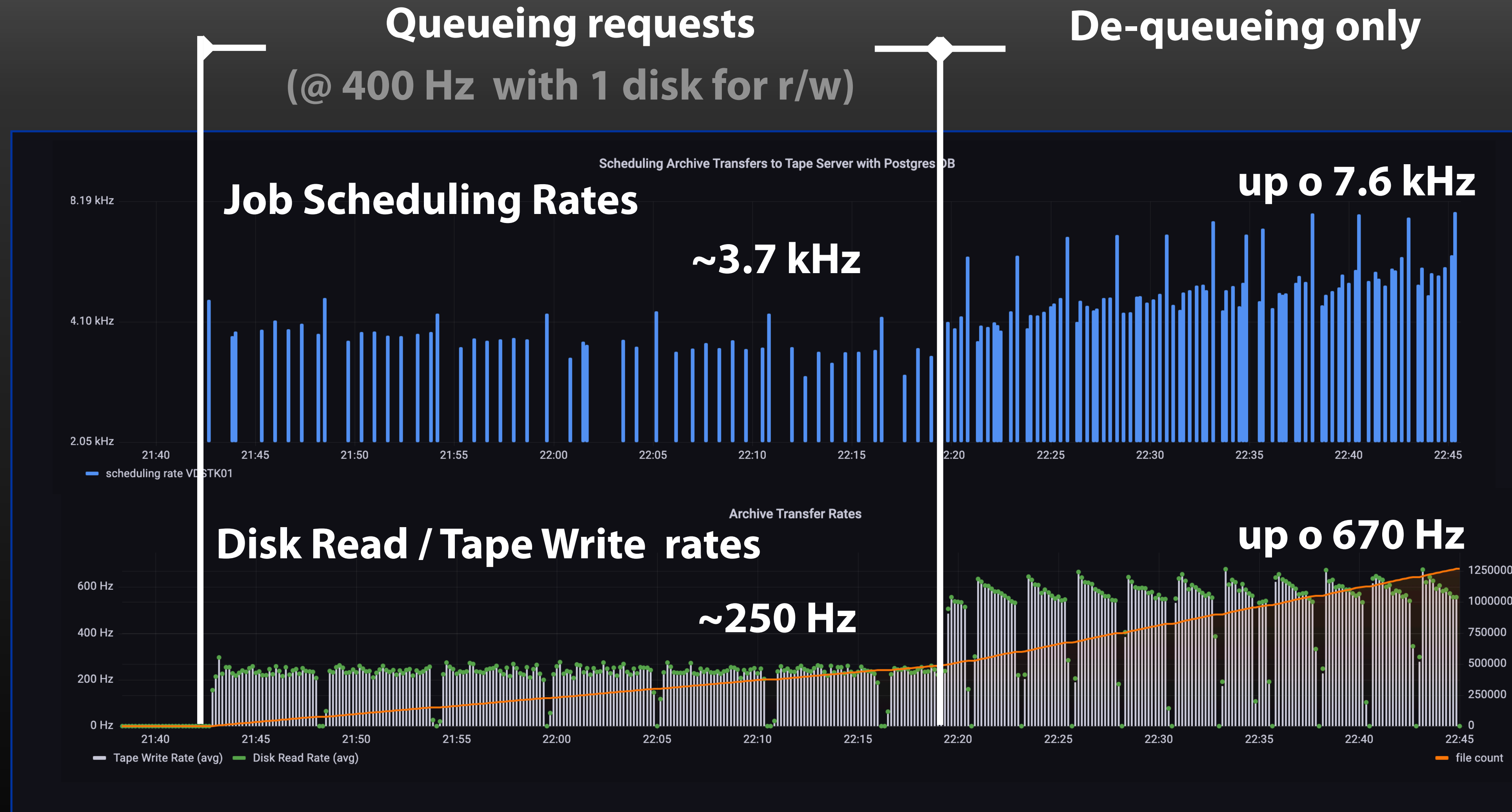
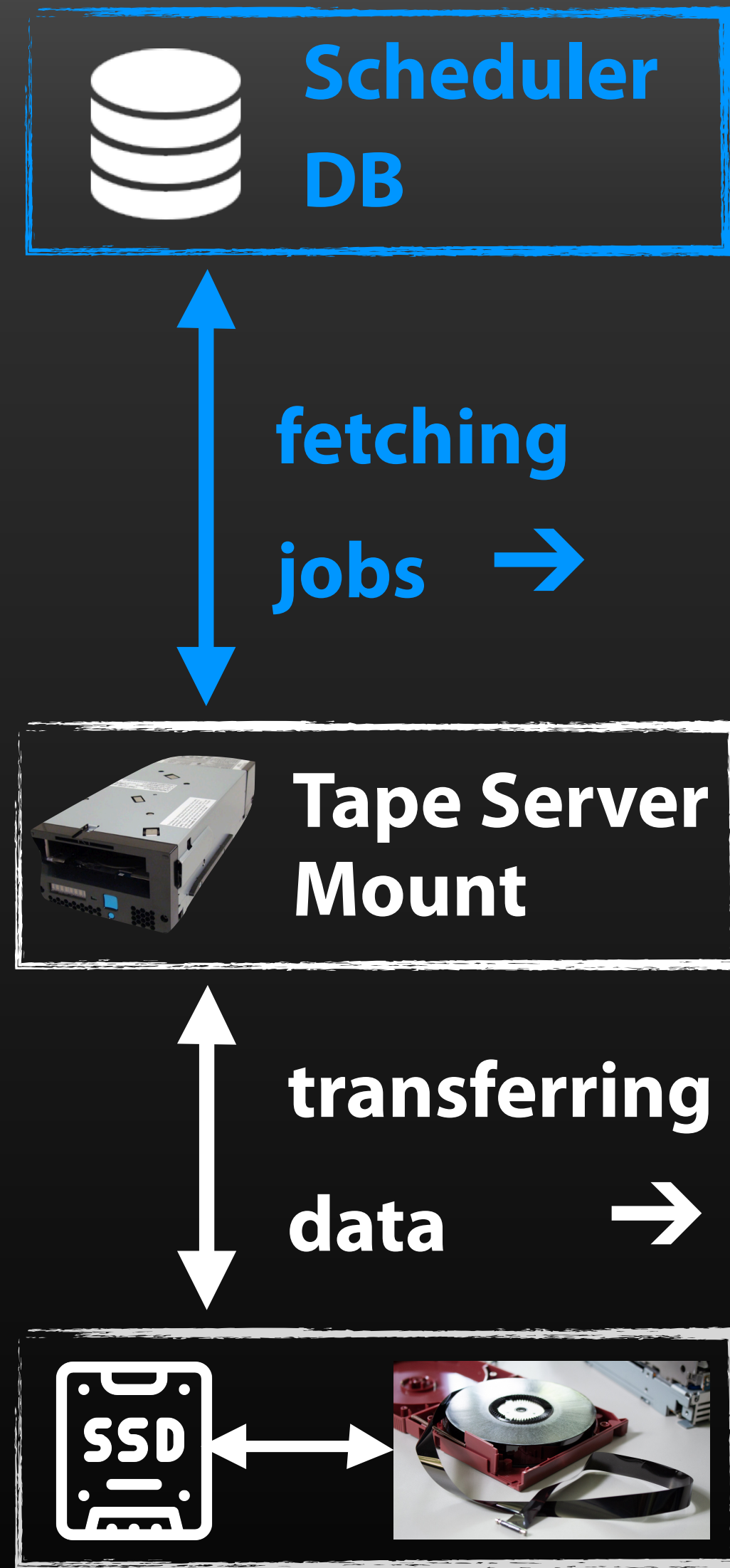
Load

- 1M files, 128 B each
- 1 tape pool (30 MB tapes)

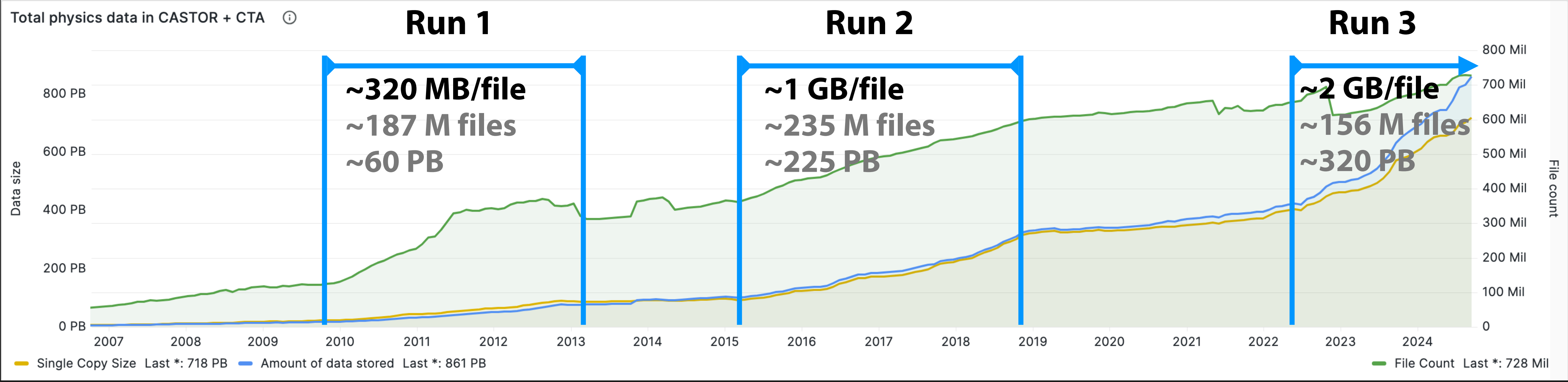


Kubernetes deployment for performance and scalability tests is in works !

Functional Test of Archive Scheduling



Minimal File Rate Requirements



Conservative estimates

- file sizes grow
- keeps minimal rate down

| | RUN 3 | RUN 4 |
|-----------------|-----------|--------------|
| CTA SLA | 50 GB / s | 125 GB / s |
| Avg. File Size | 2 GB | 2 GB |
| Avg. Throughput | 25 Hz | 63 Hz |

Postgres DB Scheduling
several times faster than ~ 100 Hz = peak rate of Run 3

Implementation and Deployment Roadmap



Implementation:

Code Optimisation:

Retrieve and Repack

Performance & Scaling Tests

Deployment:

Kubernetes Test Cluster

Repack Instance

Production

Q4 2024

Q1 2025

Q2 2025

Q3 2025

Q4 2025

Beyond Q1 2026

Not planning to replace ObjectStore on a timescale of Run 3



Next Scheduler DB Evolution

- ✓ meets with Run 4 challenges
- ✓ satisfies the needs of larger CTA Community
- ✓ decreases future development and operational costs

Implementation Status

- ✓ PostgreSQL as CTA Scheduler DB for Archival
- Workflow is functional
- ✓ Retrieve and Repack are coming next
- ✓ improved locking granularity

Long Term Prospects

- ✓ tape supply logic part of scheduling
- ✓ refactoring and cleaner handling of several workflows:
(leveraging DB features)
 - request/file deletion
 - multiple retrieves per file
 - priority queues

Thank you !



Thank you for your attention

**... special Thanks to few of my CERN colleagues for their help
(David S., Elvin S., Joao A., Julien L., Michael D., Steven M., Pablo O. C., Vlado B.)**

I welcome any questions or comments !





- **Archive of the physics data**
- **Provisioned capacity: ~1.18 EB**

- **Libraries:**
 - **4 x IBM TS4500**
 - **2 x Spectra Logic TFinity**
- **Drives:**
 - **40 x TS1170, 46 x IBM1160**
 - **88 x LTO-9, 10 x LTO-8**
- **Media:**
 - **150 PB on 3592JF, 150 PB on 3592JE, 227 PB on 3592JD**
 - **551 PB on LTO-9, 17 PB on LTO-8, 59 PB on LTO-7M**



- **Backup of the business data**
- **Licensed capacity: ~15 PB**

- **Libraries:**
 - **1 x IBM TS4500 (partitioned)**
 - **1 x Spectra Logic TFinity (partitioned)**
- **Drives:**
 - **10 x LTO-9**
 - **10 x LTO-8**
- **Media:**
 - **12 PB on LTO-8**
 - **11 PB on LTO-7M**



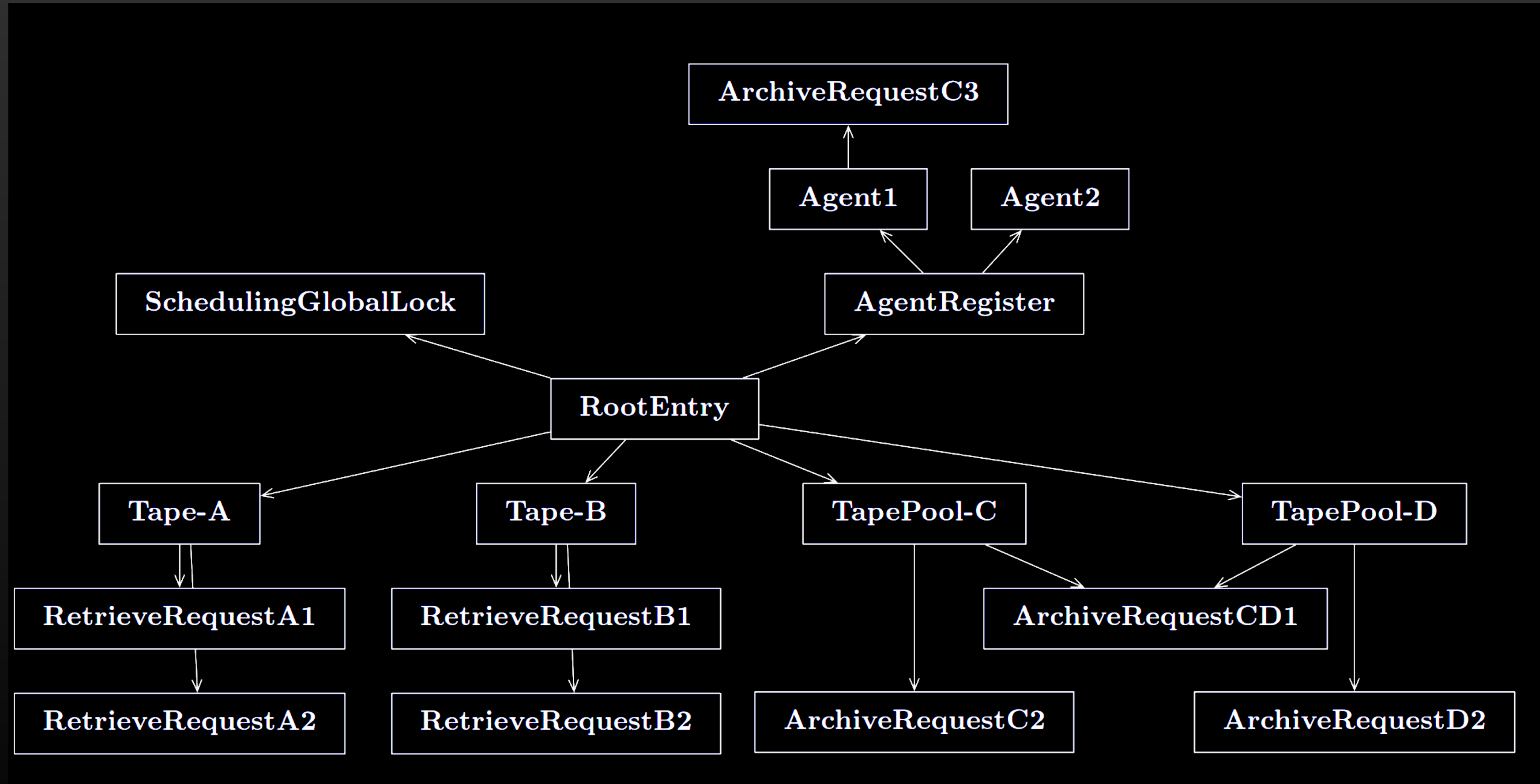
Challenges

- **independent solutions for ObjectStore-coupled scheduler logic**
- **high performance, reliability and scalability IF**
 - ▶ **DB features exploited smartly !**
(e.g. not counting all rows for every query)
 - ▶ **requires optimisation efforts per use-case**
 - ▶ **relies on diligence of developer with DB queries and DB configuration**
- **LHC Run 4 !**

Scheduler DB Implemented as ObjectStore



Architecture



IO Limited Sanity Check



1 thread queueing
2 threads writing to tape

