

On-Demand Column-Joining for End-User Analysis

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NCSA | National Center for Supercomputing Applications





Motivation

2





- As the LHC moves into the HL-LHC era, the volume of data to be stored and processed will grow significantly (x10)
 - CMS AOD (450kB/event) Limited Availability, high processing costs, data: C++ classes
 - MiniAOD (45kB/event) Suitable for nearly all analyses, still large, still significant processing, data: C++ classes
 - NanoAOD(4kB/event) Suitable for half of analyses, analysis-ready, data: primitives





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- Several competing needs create an impedance mismatch ٠

2

- Disk space comes at a premium •
- High throughput requires high availability and duplication across sites around the world
- Ntuples
- Fast turnaround (for analyzers) is paramount to getting the science **done**!



• Traditional analysis workflows tend to **duplicate information** from large data-tiers (Mini/AOD) via custom "Ntuples", in order to create more streamlined but self-contained input data - for the half of analyses able to use NanoAOD (a "generalized" ntuple), it's nearly optimal and can obviate the need for intermediate



Data Duplication in a Typical Analysis

3

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- An analysis may be able to use NanoAOD(-like inputs), but must store expensive ML outputs
 - Typical approach: Duplicate all necessary input data from NanoAOD + added ML information into a custom NanoAOD (May permit dropping columns or certain events, but these decrease flexibility)



Data Duplication in a Typical Analysis

- An analysis may be able to use NanoAOD(-like inputs), but must store expensive ML outputs
 - Typical approach: Duplicate all necessary input data from **NanoAOD + added ML** information into a custom NanoAOD (May permit dropping columns or certain events, but these decrease flexibility)
- An analysis may have 90% of data needs met by NanoAOD, but the additional requirements drive it to use MiniAOD or AOD
 - Custom NanoAOD variant (superset of central variation) or custom NTuple format created from larger datatier (labor and computeintensive), **duplicating** a significant amount of centrally-stored events in Nano and Mini formats (inefficient disk utilization)

4



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Source A Source B (TTree) Conversion (ServiceX) ** Source A* (Parquet)





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Trino (as described by o1-mini)

Trino is a high-performance, distributed SQL query engine designed for running interactive analytic queries against various data sources of all sizes. Originally developed by Facebook under the name **Presto**, Trino was forked in 2020 by the original creators to foster a more open and community-driven development model. Trino has since evolved into a robust, open-source project maintained by the **Trino Software Foundation**.

• Key Features of Trino:

- Distributed Architecture
- SQL Compatibility
- Federated Querying
- Performance Optimization
- Extensibility and Customization:
 - Plugin Architecture: Users can develop **custom connectors** and functions to extend Trino's capabilities.
 - Community-Driven: Being open-source, it benefits from contributions and innovations from a broad community of developers and organizations.
- Security and Access Control:
 - Authentication and Authorization: Supports various security protocols and integrates with enterprise security systems.
 - Data Encryption: Ensures data privacy through encryption in transit and at rest.

Trino, a query engine that runs at ludicrous speed

Fast distributed SQL query engine for big data analytics that helps you explore your data universe.







6

Source File 1



Source File A

A:1 A:2 ... A:M

- We've prepared several benchmark datasets using CMS OpenData
 - Source datasets ranging from ~2GB to ~500GB (converted to parquet, ZSTD:5, full NanoAOD)
 - GNN Inference (parquet, 8 scalar-float columns / event)
 - fully-aligned, intra-file-reversed, intra-fileshuffled, globally-shuffled variants
 - Testing various combinations (from a few scalar columns from source + inference, to dozens of ragged fields in source + all 8 inference columns)



Source File 1 Inference File 1



1:1 1:2 . . .

Inference File A Source File A





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Source File 1 Inference File 1 1:N 1:1 1:1 1:N-1 1:2 1:2 1:N 1:N 1:1

Inference File A Source File A



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Source File 1 Inference File 1



Source File A Inference File A

A:98 A:1 A:M A:1 A:2 A:M-1 A:2 A:3 A:* A:M A:M A:1 Shuffled Aligned Reversed 6

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Source File 1 Inference File 1



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A:1 A:98 A:M A:1 A:2 A:M-1 A:2 A:3 A:M A:* A:M A:1 Shuffled Aligned Reversed 6

3:5 A:M 8:104 . . . 1:1 9:48 . . . A:4 * *

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Globally Shuffled Nick Manganelli - CHEP 2024 - Krakow



Cluster Configuration

- Heterogeneous Cluster ٠
 - 3 Nodes: Intel Xeon Silver 4212 @ 2.20GHz (22 cores) - 88GB RAM - Ceph on NVME storage
 - 1 Node: Intel Xeon Silver 4210 @ 2.20GHz (39 cores) - 100GB RAM - Ceph on NVME storage

Exploratory Environment at FNAL

- Store

- Shared resource
 - 10 Workers for Trino on Cluster

7

Nick Manganelli - CHEP 2024 - Krakow

See more in Ben's slides: https://indico.cern.ch/event/1369601/contributions/5883602/





```
CREATE TABLE single_top_s_chan_join
WITH (
 format = 'PARQUET',
 external_location = 's3a://servicex/nanotest/parquet/AGC/single top s chan/join out/
AS
 SELECT single top s chan.run, single top s chan.event, single top s chan.luminosityBlock,
  Electron pt,
  Electron eta,
  Electron_phi,
  Electron_cutBased,
                                                                                         Dataset wit
                                                                                    Electron_ip3d,
  Electron_sip3d,
  Electron mass,
                                                                                         ragged prin
  Electron_pfRellso03_all,
  Electron_pfRellso03_chg,
                                                                                         (3.7GB sou
  Muon_pt,
  Muon eta,
  Muon phi,
  Muon mass,
  Muon tightld,
  Muon_ip3d,
                                                                                        Encouragin
                                                                                    •
  Muon sip3d,
  Muon_pfRellso04_all,
                                                                                         invariant*
  Jet_mass,
  Jet_pt,
  Jet eta,
                                                                                         joined
  Jet_phi,
  Jet_jetId,
  Jet_btagCSVV2,
  Jet_btagDeepFlavB,
  Jet btagDeepFlavCvB,
  Jet_btagDeepFlavCvL,
  Jet_btagDeepFlavQG,
  Jet chEmEF,
  Jet chHEF,
  Jet muEF,
  Jet neEmEF,
  Jet neHEF
  Jet_puldDisc,
                                                                                                2.2M
  Jet_qgl,
  Jet_rawFactor,
  Jet_bRegCorr,
  Jet bRegRes.
  Jet_electronIdx1
  Jet electronIdx2,
  Jet muonIdx1,
  Jet muonldx2,
  GNN_p1, GNN_p2, GNN_p3, GNN_p4
 FROM single_top_s_chan
 JOIN single_top_s_chan_infer ON
  single_top_s_chan.run = single_top_s_chan_infer.run AND
  single_top_s_chan.luminosityBlock = single_top_s_chan_infer.luminosityBlock AND
  single_top_s_chan.event = single_top_s_chan_infer.event;
```

Nick Manganelli - CHEP 2024 - Krakow

Small Dataset Join Benchmark

single_top_s_chan_infer

| | select count(*) from single_top_s_chan_join; _col0 2867199 | | |
|---------------------------------------|---|--|--|
| events/s | single_top_s_chan_infer_globalshuffle Query 20240927_184518_00039_au69m, FINISHED, 11 nc Splits: 859 total, 859 done (100.00%) 15.00 [34.4M rows, 1.5GB] [2.29M rows/s, 102MB/s] | | |
| | 2867199 (1 row) | | |
| | <pre>select count(*) from single_top_s_chan_join; _col0</pre> | | |
| | single_top_s_chan_infer_intrafileshuffle Query 20240927_184318_00036_au69m, FINISHED, 11 nc Splits: 862 total, 862 done (100.00%) 15.49 [34.4M rows, 1.5GB] [2.22M rows/s, 98.9MB/s] | | |
| to permutations being | 2867199 (1 row) | | |
| ig result: seemingly | trino:servicex> select count(*) from single_top_s_ col0 | | |
| | single_top_s_chan_infer_reversed Query 20240927_184139_00033_au69m, FINISHED, 11 nc Splits: 862 total, 862 done (100.00%) 14.54 [34.4M rows, 1.48GB] [2.37M rows/s, 104MB/s] | | |
| 1000000000000000000000000000000000000 | 2867199 (1 row) | | |
| th 2.8M events (rows), | trino:servicex> select count(*) from single_top_s_ _col0 | | |
| | Query 20240927_183428_00028_au69m, FINISHED, 11 nc Splits: 862 total, 862 done (100.00%) 16.64 [34.4M rows, 1.48GB] [2.07M rows/s, 91.3MB/s | | |

* stat fluctuations, warm-caching to be eliminated as sources of differences in high-stat testing













Large Dataset Join Benchmark





Large Dataset Join Benchmark

Largest dataset of 500GB (ttbar) crashed • trino deployment*





9

Incorrect resource limits set for trino in OKD *

Large Dataset Join Benchmark

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Incorrect resource limits set for trino in OKD *

Large Dataset Join Benchmark

- Largest dataset of 500GB (ttbar) crashed • trino deployment*
- May indicate that partitioning the joins into dask-task sized elements will be a necessity to use trino as our distributed SQL engine (already a desirable element, as described later)



Benchmark Output Analysis Formats

Test of a simple pseudo-analysis running on various root TTree and parquet files, in various compression schemes





Benchmark of "analysis" read-speed Stats computed over 7 runs, 5 loops (timeit) 1,334,428 events processed

| Format | Source | How Made | Compression | Time |
|---------|-----------|------------|-------------|----------------|
| ROOT | AGC | N/A | ZLIB:1 | 11.7 s ± 389 n |
| ROOT | converted | root hadd | LZMA:9 | 11.9 s ± 551 m |
| ROOT | converted | root hadd | ZSTD:5 | 10.9 s ± 477 r |
| Parquet | converted | hepconvert | ZSTD:5 | 4.73 s ± 61 m |
| Parquet | joined | trino | GZIP | 3.99 s ± 129 r |



Coffea analysis

- coffea brings together individual scikit-hep elements needed for a full analysis, and provides schema-application, corrections, scaleout(-patterns)
 - and often the first place something is prototyped and tested before being spun out • into it's own package
- Tightly integrated with dask: •
 - User's analysis code is broadcast over datasets to create **task graphs** •
 - Typetracer setup per dataset records operations (lazy, no data executed on)
 - Task graphs are distributed to compute resources to execute •
 - Results returned to user's client (histograms, small arrays, locations of output root/ • parquet files, ...)
- Task graphs are key: allow programmatic optimization of analysis, understanding necessary inputs as mapped to requested outputs















Source A (Format A)

Source B (Format B)

Source C? (Format C)





Source A (Format A)

Source B (Format B)

Source C? (Format C)





Source A (Format A)

Source B (Format B)

Source C? (Format C)





Source A (Format A)

Source B (Format B)

Source C? (Format C)

coffea (Build joined typetracer)



Source A (Format A)

Source B (Format B)

Source C? (Format C)

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12

Dask Tasks



Source A (Format A)

Source B (Format B)

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coffea (Build joined typetracer)

12

Dask Tasks

Dataset 1, Chunk 1 (N events)







Dask Tasks

Dataset 1, Chunk 1 (N events)







Dask Tasks

Dataset 1, Chunk 1 (N events)



































Dataset Z, Chunk Y



Transform A

12







- Baseline and Optional Targets of US CMS HL-LHC R&D Program (2024) in collaboration with IRIS-HEP
 - Building Typetracer for pseudo-joined data •
 - Generation of ServiceX conversion tasks •
 - Building Trino join queries, embedding of joined-output into daskified analysis •



• Updated and expand ServiceX MiniAOD conversion/selection of auxiliary information to be joined with NanoAOD (Stretch Goal)

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- Funded by LPC Distinguished Researcher Program (2025) •
 - Native Ceph Object Store usage (Currently object -> file -> object)

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 - Native Ceph Object Store usage (Currently object -> file -> object)
- Future developments within US CMS and IRIS-HEP •
 - Kafka integration for streaming output •
 - RNTuple support in trino •

Updated and expand ServiceX MiniAOD conversion/selection of auxiliary information to be joined with NanoAOD (Stretch Goal)

14

- Augmenting with objects from MiniAOD datatier
 - Adding ParticleFlow candidates for reclustering
 - Full set of ML taggers





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 - Full set of ML taggers
- FlashSim?)
 - Versioning potential



Caching Experiment-wide objects (systematic variations, CMS ParticleTransformer,

- Augmenting with objects from MiniAOD datatier •
 - Adding ParticleFlow candidates for reclustering
 - Full set of ML taggers
- FlashSim?)
 - Versioning potential
- results (event, object classifiers)

Caching Experiment-wide objects (systematic variations, CMS ParticleTransformer,

Caching Analysis-specific non-ML (derived quantities, systematic variations) and ML

Funding Acknowledgements

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BACKUP

US CMS Software and Computing - HL-LHC R&D (WBS4)

2 Grand Challenges for HL-LHC Computing and Software

In order to effectively focus and structure the U.S. CMS R&D efforts, we organize the innovation, research, and development needs for HL-LHC computing into the four Grand Chal*lenges* that encompass the advances needed for HL-LHC computing to succeed:

(1) Modernizing Physics Software and Improving Algorithms

Exploit novel algorithms, including ML/AI, reduce algorithmic complexity, increase computational intensity, and provide core software infrastructure to enable effective use of modern hardware and accelerators. The work is organized in the following work packages:

- Core Software Framework and Software Portability
- Establish Performance Metric and Performance Baseline for Physics Software
- U.S. Contributions to the Charged Particle Tracking Software
- U.S. Contributions to Software for High Granularity Calorimeter
- U.S. Contributions to CMS Advanced Algorithms Work

17

(2) **Building Infrastructure for Exabyte-Scale Datasets**

Build infrastructure to archive, store, transfer, and provide access to exabyte-scale datasets. Explore data lakes and custodial storage: establish a technology and cost model for custodial/archival storage facilities which manages operations costs, and optimizes hardware costs. Orchestrate computational services and data access, provide intelligent

FriendTrees

- duplication issues, via the FriendTree mechanism

 - with BuildIndex (major + minor keys, i.e. run + event)
 - <u>https://root.cern/doc/master/</u> classTTree.html#a3f6b5bb591ff7a5bd0b06eea6c12b998

Traditional ROOT workflows have some functionality that can alleviate the data

• Supports joining entry-by-entry events in two separate (groups of) files

Can accommodate some situations with non-aligned (out-of-order) joins

O1-mini's description of Trees, Parquet, RNTuple

- 2. Technology Descriptions
- a. ROOT TTrees

ROOT TTrees are a fundamental data structure within the ROOT framework, widely used in HEP for storing and analyzing large datasets.

- Structure: TTrees store data in a hierarchical, tree-like structure with branches representing different variables (columns).
- Features:
 - Custom Compression: Supports various compression algorithms tailored for HEP data.
 - Provenance Tracking: Maintains metadata and provenance information essential for reproducible research.
 - Integration with ROOT: Seamless integration with ROOT's data analysis tools and C++/Python interfaces.
- b. Apache Parquet

Apache Parquet is an open-source, columnar storage format optimized for performance and interoperability across different data processing frameworks. • Structure: Stores data in column chunks, enabling efficient compression and encoding schemes.

- Features:
 - Language Agnostic: Supports multiple programming languages (e.g., Java, C++, Python).
 - Wide Ecosystem Support: Integrates with big data tools like Apache Spark, Hive, and Pandas.

• Efficient Compression: Utilizes advanced compression techniques to reduce storage footprint.

c. ROOT RNTuple

ROOT RNTuple is a newer addition to the ROOT framework, designed to address some limitations of TTrees and to provide a modern, efficient, and scalable storage solution. • Structure: Similar to TTrees but optimized for better performance and scalability.

- Features:
 - Asynchronous I/O: Enhances read/write performance by supporting non-blocking operations.
 - Better Parallelism: Designed to leverage multi-threading and distributed computing environments.

19

• Improved Data Model: Offers a more flexible and robust data schema compared to TTrees.

Trino Use Cases (as described by o1-mini)

- Use Cases of Trino
 - Interactive Analytics:
 - Ideal for business intelligence (BI) tools and dashboards that require real-time or near-real-time query responses.
 - Data Lake Querying:

 - Federated Data Access:
 - Enables organizations to query and analyze data spread across multiple heterogeneous systems from a single interface.
 - ETL and Data Transformation:
 - Can be used in Extract, Transform, Load (ETL) processes to transform and move data between systems efficiently.



• Facilitates querying large-scale data lakes stored in platforms like Amazon S3, Google Cloud Storage, or HDFS without the need for data movement.

Model o1-mini says...

Trino in the Context of Data Storage Technologies

Given your previous interest in **ROOT TTrees**, **Apache Parquet**, and **ROOT RNTuple**, it's valuable to understand how Trino interacts with these and other columnar storage formats:

- Apache Parquet:
 - formatted data stored in data lakes or distributed file systems.
 - and speeding up query execution.
- ROOT TTrees and RNTuple:
 - data transformation steps.
 - data sources.

Integration: Trino has robust support for querying Parquet files, allowing users to perform SQL queries directly on Parquet-

• Performance: Leveraging Parquet's columnar storage, Trino can efficiently read only the necessary columns, reducing I/O

• Integration: While Trino is primarily designed to work with widely-adopted data formats like Parquet, CSV, JSON, and various databases, integrating with specialized formats like ROOT TTrees or RNTuple may require custom connectors or

• Use in HEP: In High Energy Physics (HEP), where ROOT formats are prevalent, Trino could be used alongside data processing pipelines that convert ROOT data to more Trino-friendly formats for broader analysis and integration with other