

## SuperhistogramS — or abstract algebra for fun and profit

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#### Vague statement of the problem



Long, long ago, histograms were individual objects that were managed individually.





Now, histograms are more often used collectively, with thousands of histograms in a single fit.





# A "superhistogram" is a large collection of histograms that are meant to be interpreted together.



#### Representing a superhistogram with a directory of ordinary histograms is

- wasteful because much of the same metadata is copied in memory or on disk, and many small buffers of bin contents is less efficient than one big buffer.
- inconvenient because the object with a common meaning has to be managed as individual objects without an explicit connection. (Often in practice, they're only linked by naming conventions.)



Boost::Histogram provides a generic way to create an *n*-dimensional space with regular, variable, and categorical axes.





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A superhistogram has multiple sources, channels, and systematics.

- Not all histograms in the collection have the same number of bins or the same dimensions, but many do.
- One fill of the superhistogram would increment every histogram in the collection.

#### Something like this



```
h = SuperHist(SuperHist(
        SuperHist(
            Hist.new.Reg(100, 0, 30, name="syst-up"),
            Hist.new.Reg(100, 0, 30, name="nominal"),
            Hist.new.Reg(100, 0, 30, name="syst-down"),
            name="pt",
        ),
        SuperHist(
            Hist.new.Reg(50, -5, 5, name="syst-up"),
            Hist.new.Reg(50, -5, 5, name="nominal"),
            Hist.new.Reg(50, -5, 5, name="syst-down"),
            name="eta",
        ),
        name="data",
    ),
          # similarly for name="mc"
    . . . ,
h.fill(df) # one DataFrame row is a scalar fill operation
```

## So far, this is looking like Histogrammar







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We need a relationship that restricts the superhistogram to a useful subset of all possible trees.



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is associative:  $(x \cdot y) \cdot z = x \cdot (y \cdot z)$ has an identity: there is an  $e \in S$  such that  $e \cdot x = x$  and  $x \cdot e = x$  for all  $x \in S$ .





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(You may be familiar with groups, which are monoids without inverses.)



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- Boost::Histogram axes under Cartesian product: n axes, an n-dimensional space, combined with m axes, an m-dimensional space, forms an (n + m)-dimensional space. An axis with 1 bin could be called an identity.

A semiring is any set S with two operations, "+" and " $\times$ ", such that

- ► S under + is a monoid; let's call its identity "0".
- S under  $\times$  is a monoid; let's call its identity "1".
- ▶ + is commutative: a + b = b + a.

▶ 0 absorbs everything under  $\times$ :  $a \times 0 = 0 = 0 \times a$ .

× is distributive over +:
$$a \times (b + c) = (a \times b) + (a \times c)$$
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**Example:** natural numbers under ordinary addition and multiplication.





To build a superhistogram, we put axes together in two ways:

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StrCategory(["data", "mc"], name="src") \* Reg(100, 0, 30, name="pt") +
StrCategory(["data", "mc"], name="src") \* Reg(50, -5, 5, name="eta")

Two histograms have the same categorical axis, different regular axes.



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- The identity of  $\times$  is a one-bin axis.
- $\blacktriangleright$  + and × obey a distributive property: if *a*, *b*, and *c* are axes,

$$a \times (b + c) = (a \times b) + (a \times c)$$

represents two 2-dimensional histograms, both with the same first axis a, differing in their second axes b or c.

## A superhistogram is a jagged array



In Boost::Histogram terms, Axis a, b, and c in expressions like

$$a \times (b + c) = (a \times b) + (a \times c)$$

would share one Storage, such as Double (an array of floating-point values).



## A superhistogram is filled by DataFrames



Each fill operation increments 1 bin in each multiplicative space, and every such space in an additive collection.

		src	pt	eta	
h.fill(	0	data	22.1	-1.4	
	1	data	15.8	2.8	
	2	data	25.0	0.5	
	3	mc	19.4	-3.1	
	4	mc	28.5	-0.7	

fills the histogram with axes "src" and "pt" 5 times, and fills the histogram with axes "src" and "eta" 5 times.



# Why does it matter that superhistograms form a semiring?



- fully expanded:  $(a \times b) + (a \times c)$
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Maybe this could also simplify the interface. (Slicing across all the histograms?)



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By adding this algebraic structure, we restrict the possibilities, but in ways that have benefits to how we want to fill, store, and manipulate histograms.



# To Peter's talk!