traccc Integrating the Alpaka framework

Ryan Cross GridPP51 & SWIFT-HEP07 2024/03/27



Overview

This talk will cover:

- 1. traccc.
- 2. Cross-Platform Abstraction Libraries.
- 3. Where We Are
- 4. Current Work
- 5. What comes next?

A Common Tracking Software

ACTS is a generic, experiment independent framework/software toolkit, written in C++. Through it, you can get algorithms for track reconstruction that can be used in any experiment, agnostic of any technical details (detector tech, design and event processing framework).

It has been designed in a thread-safe manner, with support for parallel code execution and optimised data structures for speeding up the many linear algebra operations used throughout the code base.



ACTS R&D Projects

Many of the core algorithms in ACTS have been ported to CUDA and SYCL, but there is a limit as to how far this can go. Full offloading is difficult, with some of the event data model and geometry not being the most GPU-friendly.

To tackle this, ACTS has launched several R&D projects:

- traccc Tracking Algorithms on the GPU.
- detray A GPU based Geometry Builder.
- algebra-plugin Provides varying algebra plugins for the other projects.
- vecmem A GPU Memory Management Tool for the other projects.



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traccc specifically, is aiming to establish a sensible event data model and algorithms that are able to exploit parallelisation architecture, whilst relying heavily on the other projects.



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• **SYCL** is a higher level programming model, developed by the Khronos group (OpenCL/OpenGL/Vulkan and more). It defines an abstraction layer that enables code for heterogeneous processors via a 'single-source' style in standard C++. Supports many backends: CUDA, AMD GPUs, Intel GPUs, OpenMP, MPI, Vulkan, std::thread, OpenCL and more.



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- Kokkos is C++ based programming model, which provides methods that abstract away details of parallel execution and memory management, such that code can be written for many shared-memory programming models in a unifed way. Supports CUDA, HIP, SYCL, HPX, OpenMP and std::thread.



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- **alpaka** is a header-only C++ 17 abstraction library for accelerator development. It aims to provide performance portability across a range of accelerators through the abstraction of the underlying levels of parallelism. Support CUDA, OpenMP, std::thread, TBB, HIP and OpenAcc.

Kokkos al Baka

Cross-Platform Abstraction - How?

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Get an accelerator device:

```
accelerator = getAcceleratorDevice();
queue = getDeviceQueue(accelerator);
```

Define an operation for the device to perform:

```
job = [](auto accelerator, auto config, auto items) {
    auto item = items[getThreadIndex()];
    ...
};
```

Run the jobs in parallel:

```
queue.submit(job, configuration, items);
queue.wait();
```

Why alpaka?

I've just outlined three projects that support the "write once, support many" paradigm, and both SYCL and Kokkos are already implemented in traccc, with differing levels of functionality. So why a third?

alpaka was chosen as a possible candidate for a few reasons:

- **Simplicity**: alpaka is a lightweight, header-only library, which makes integration into traccc very easy, as well as it being written in the same modern C++17 as traccc/acts.
- **Familiarity**: The alpaka abstraction model is very similar to the CUDA grid-blocks-thread model, making writing code for alpaka simple, and familiar for those with CUDA experience, whilst also providing a CPU and non-CUDA based implementation.
- **Community Support**: alpaka has been used extensively at CMS, including in cms-sw and their HLT achieving performance close to that of the native CUDA codebase, from a single source code that can be utilised on many devices.

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The spacepoint binning gave me a first look at development with Alpaka, as well as developing inside of traccc/ACTS. My previous slides, given at a UK SWIFT-HEP / GRIDPP meeting, give a bit of a better overview of that work, as well as some more basic comparisons of Alpaka vs CUDA.

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Most of that work was merged recently as part of PR (#451), with a secondary PR incoming to include the latest throughput examples and general tidying up.

from the spacepoint hinning Processing Time Per Event cpu Thi cuda alpaka but nati 10² Mos lates 10¹ Time (ms) 100 10^{-1} 10-2 50 200 250 100 150 300 0 Dataset (ttbar_muX)

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- Continue the porting process of further algorithms, bringing the Alpaka implementation closer to completely matching the CUDA one.

Each of these pieces of work are at different stages of completion, and I'll go into a touch more detail on them each now.



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The end goal here is that we can have the full examples running on both Nvidia and AMD, with only a single flag change at compile time to target the relevant architectures. All the changes made for this work also help improve the code for potential later test with other accelerators.

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It isn't seen at all in the examples that only run a single event, but only in those that run many events in sequence, at least when using the CUDA back-end. I'm also doing some testing with a CPU, single-threaded back-end to see if I can reproduce the error there, which would make debugging a lot easier (at least compared to hundreds / thousands of CUDA kernels running).

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Work is slowly moving here, first testing within thrust with a CUDA target only, to slightly delay working out the best / most appropriate approach to replace the Thrust code.

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- Extend the testing and verifying work once HIP is working, to ensure that CUDA and HIP continue to work.
- Finally, more in-depth benchmarking of the Alpaka implementation, to help understand if / where bottlenecks are, and if there is anything in our Alpaka code that needs improving.

Conclusion

In Conclusion:

- traccc is a R&D effort as part of the ACTS project, working on exploiting GPUs and other accelerators to speed up tracking across a range of experiments.
- As part of that, many different acceleration abstraction libraries have been implemented, with alpaka being the newest.
- alpaka has good support already in HEP, and its parallelisation model make it a strong candidate for being the general purpose abstraction library.
- This talk gives a brief overview of the already completed work porting algorithms to utilise Alpaka in traccc.
- More work in ongoing to verify alpaka with non-CUDA targets, improve the robustness of the alpaka implementation, and further complete porting of algorithms to alpaka.

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Backup Slides



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