

Performance of a Cluster File System Backed by an HDD-Based Data Storage System Under True Concurrent Read-Write Load



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Outline

- Starting Point
- Initial Acceptance Benchmark, Results
- Improved Benchmark, Results
- Comparison, Discussion
- Conclusion
- ■Q & A



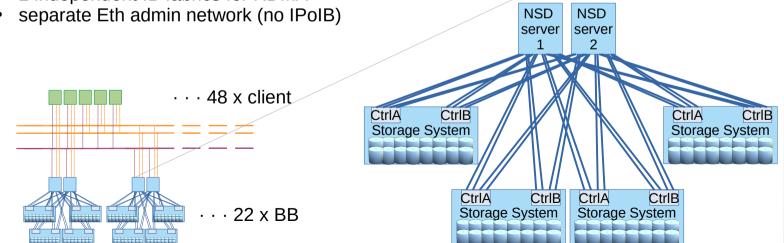




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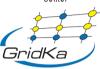
Starting Point - Environment

- IBM® Storage Scale (AKA "GPFS" AKA "Spectrum Scale" AKA "Elastic Storage" AKA "Tiger Shark")
- Traditional building blocks (BB): 22 x (2 NSD server + 4 x Seagate Exos® CORVAULT™ Storage Systems)
- Seagate Exos® CORVAULT™: 106 x 18TB HDD, ADAPT disk groups 16+2, SAS
- 48 NSD clients
- Servers: SuperMicro H12SSL-i, ConnectX6 IB (2x100Gbps)
- 2 independent IB fabrics for RDMA









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Starting Point - Initial Acceptance Test Definition

Freshly delivered system to be acceptance-tested.

Test defined in tender invitation (and consequently in contract): Use gpfsperf (tool provided with GPFS, currently IBM® Storage Scale):

- At once, on each client, start one process reading and one writing.
- each process to use a separate file of 100GB
- After completion of last process, add up the rates reported by all writing and all reading processes, resp., to compute the accumulated write and read rates.

Requirement for storage capacity-specific rates (data space in file system):

$$r = R / C > 3.4GB s^{-1}PB^{-1} (= 3.4 \times 10^{-6} s^{-1})$$

to be met both for writes and reads, resp. (hence also $r_R + r_W > 6.8 \text{GB s}^{-1} \text{PB}^{-1}$).







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Starting Point - Filesystem

352 NSDs (8 per Storage System, 4 per Disk Group), 70PB (64PiB)

mmlsconfig (some)

numaMemoryInterleave yes nsdMaxWorkerThreads 3842 nsdMinWorkerThreads 3842 scatterBufferSize 256K workerThreads 1024 ignorePrefetchLUNCount yes verbsRdmasPerNode 1024 verbsRdmasPerConnection 16 prefetchLargeBlockThreshold 4M nsdbufspace 25 pagepool 64G maxMBpS 88000 prefetchPct 3 verbsPorts mlx5_0/1/1 mlx5_1/1/2







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mmlsfs (some) value description flag 16384 Minimum fragment (subblock) size in bytes - T 32768 Indirect block size in bytes Default number of metadata replicas - m Maximum number of metadata replicas - M Default number of data replicas -r -R Maximum number of data replicas -j scatter Block allocation type Estimated number of nodes that will mount file system -n 48 8388608 **Block size** -B -V 29.00 (5.1.5.0) File system version -E Exact mtime mount option Yes -S relatime Suppress atime mount option --log-replicas Number of log replicas Number of subblocks per full block --subblocks-per-full-block 512 Disk storage pools in file system -P system; gpfstest

Starting Point - Expectations

Read and Write Rates to be Nearly Identical

Read Rates Slightly Higher

A common belief amongst storage non-experts.







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[&]quot;Writing is more work than just reading"

[&]quot;Need to calculate redundancy data (RS coding/ Erasure coding)"

[&]quot;Reads access less data than writes (with 16+2 coding just 16/18 of what writes do)"

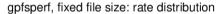
Initial Acceptance Benchmarks

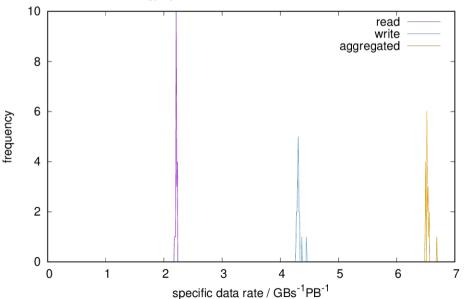
Requested: Read and Write Rate Both > 3.4 GB PB⁻¹ s⁻¹

Filesystem Size: $70.1 \text{ PB} (63.7 \text{ PiB}) = \text{rates} > 238 \text{GBs}^{-1} (222 \text{ GiBs}^{-1}, 2.27 \text{ x} 10^5 \text{ MiBs}^{-1})$

Aggr

(Read + Write): (6.493 ... 6.693) GB PB⁻¹ s⁻¹ # Avg: 6.533 GB PB⁻¹ s⁻¹





Read :

(2.188 ... 2.247) GB PB-1 s-1

Avg. 2.219 GB PB⁻¹ s⁻¹

Write:

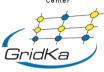
(4.249 ... 4.457) GB PB⁻¹ s⁻¹

Avg. 4.314 GB PB⁻¹ s⁻¹

Increasing the number of Read processes helps, maybe?



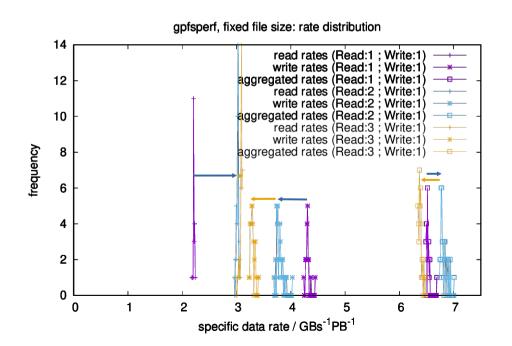




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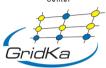
Acceptance Benchmarks - Increasing the Number of Read Processes

Requested: Read and Write Rate Both > 3.4 GB PB⁻¹ s⁻¹ Increasing the number of Read processes helps, maybe ?









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Initial Acceptance Benchmarks

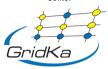
Write = 2 x Read



read processes run twice as long as writes

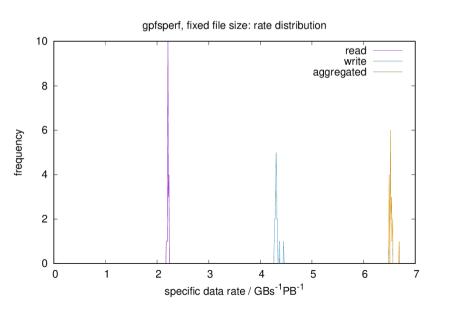






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parallel execution of reads and writes?

Not at all!

Far from it!

Improved Benchmarks

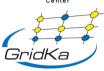
- Time constraints instead of Size constraints
- gpfsperf tool allows to set a run time limit
- Running times of > 100s reduce error due to jitter in processes' starting time (mind: we do not use MPI but simply start processes by ssh commands)
 Used: 120.0s for sole reads, 180.0s for write and read-write).
- Varied Parameters: number of clients, number of processes per client (caching effects seen when using multiple threads in one process instead), number of processes N means: N writing and N reading processes:

```
for IOSZ in IOSIZES ## 256kiB , (256kiB 2048kiB) on Jan 6,2023
  for N_CLIENTS in NUMS_CLIENTS
    for N_PROCS in NUMS_PROCS # 1 4 16
       for OP in OPS # create(=write) read rw(=read-write)
       run OP gpfsperf with N_PROCS processes on N_CLIENTS using IOSZ and timeout
```

Case of 1 process on 47 clients is closest to what the original benchmark suggested.

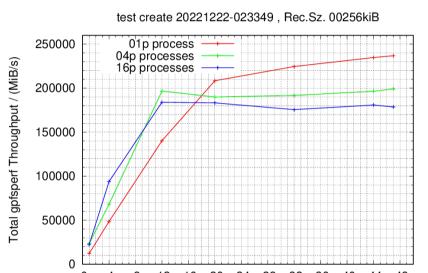




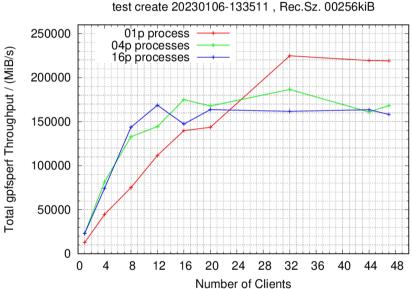


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Improved Benchmarks: sole write (create)



Number of Clients



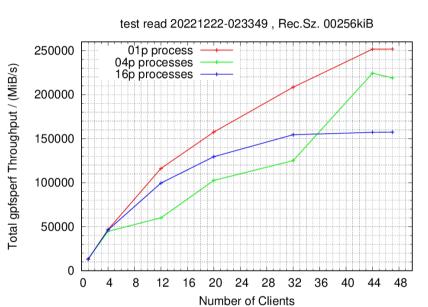


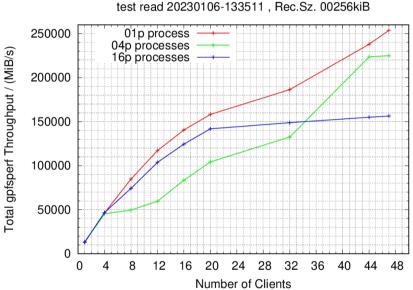




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Improved Benchmarks: sole read







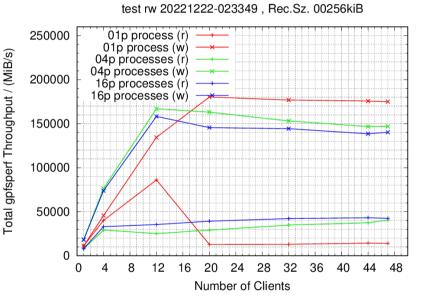


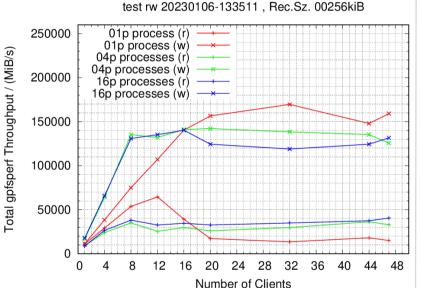


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Improved Benchmarks : Concurrent Read, Write











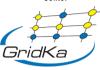


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Improved Benchmarks : Concurrent Read + Write (Sum)



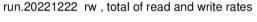


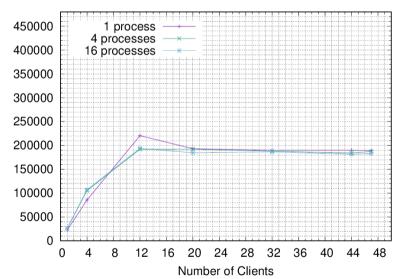


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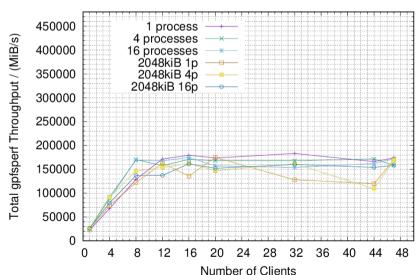
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Load





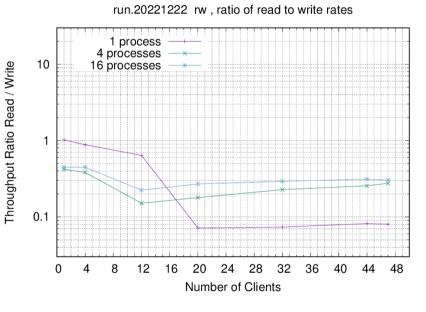
Total gpfsperf Throughput / (MiB/s)

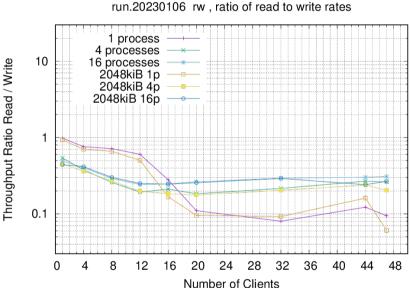


run.20230106 rw, total of read and write rates

Improved Benchmarks: Concurrent Read-to-Write (Ratio)









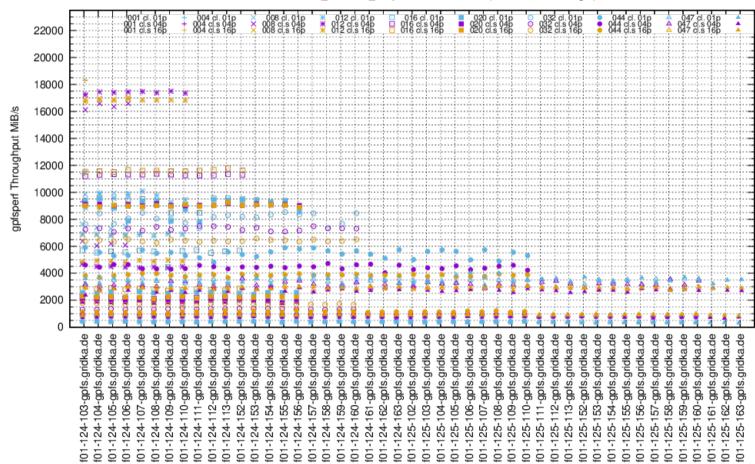




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Improved Benchmarks : Concurrent Read, Write Per-Node Rates

run.20230106 res.R 00256.M rw, per-node data rates (reads left, writes right)









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Comparison, Discussion

Requested: Read and Write Rate Both > 3.4 GB PB⁻¹ s⁻¹

Filesystem Size: 70.1 PB (63.7 PiB) => rates > 238GBs⁻¹ (222 GiBs⁻¹, 2.27 x10⁵ MiBs⁻¹)

Acceptance Test:

Read 2.219 GB PB⁻¹ s⁻¹ => **155** GBs⁻¹

GBs⁻¹ Write 4.314 GB PB^{-1} s⁻¹ => **302**

 $6.533 \text{ GB } PB^{-1} \text{ s}^{-1} \Rightarrow 458$ GBs⁻¹ Sum

Improved Test, 47 nodes, 1 process:

Read 14 GBs-1 Write 175 GBs-1 Sum 189 GBs-1

$$V_{W}=V_{R}=47 \times 100 \text{GB}$$
,
 $R_{W}=175 \text{GBs}^{-1}$, $R_{R1}=0.08*R_{W}$, $R_{R2}=250 \text{GBs}^{-1}$
 $T=T_{1}+T_{2}$ write

$T_1 = \frac{V_{\rm W}}{R_{\rm W}} = 26.86 \text{ s}$

Improved Test, 47 nodes sole read / sole wite (approx):

Read 250 GBs⁻¹ Write 230 GBs⁻¹ Sum 480 GBs⁻¹

closer to the acceptance test result but still not matching.







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Comparison, Discussion

previous result:

$$\frac{V_{\rm R}}{T} = 106.4 \text{ GBs}^{-1}$$

closer to the acceptance test result but still not matching.

 $V_W = V_R = 47 \text{ x } 100\text{GB}$, $R_W = 302\text{GBs}^{-1}$, $R_{R1} = 0.08*R_W$, $R_{R2} = 250\text{GBs}^{-1}$

$$T = T_1 + T_2$$

$$T_1 = \frac{V_{\rm W}}{R_{\rm W}} = 15.56 \text{ s}$$

$$T_2 = \frac{V_R - T_1 R_{R1}}{R_{R2}} = 17.30 \text{ s}$$

$$\frac{V_{\rm R}}{T}$$
 = 143.0 GBs⁻¹

Improved Test, 47 nodes sole read / sole wite (approx):

Read 250 GBs⁻¹

Acceptance Test: Write 302 GBs⁻¹

agrees better with the acceptance test result of 155.4GBs⁻¹ but still not matching.

Match requires R_{R2} =293GBs⁻¹ which we did never observe. Varying running times cause positive errors





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Comparison, Discussion : Acceptance Test with Adapted File Sizes

With Read / Write ratio of roughly 0.1, another test was set up:

Use fixed file size, but make files to write 10-fold the size of files to read.

Acceptance Test

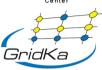
Read 2.219 GB PB-1 s-1 => **155.4 GBs-1** Write 4.314 GB PB-1 s-1 => **302.2 GBs-1** Sum 6.533 GB PB-1 s-1 => **457.6 GBs-1**

Acceptance Test, Adapted File Sizes:

Read 0.28 GB PB-1 $s^{-1} \Rightarrow 19.6 GBs^{-1}$ Write 2.92 GB PB-1 $s^{-1} \Rightarrow 204.6 GBs^{-1}$ Sum 3.20 GB PB-1 $s^{-1} \Rightarrow 224.2 GBs^{-1}$







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Comparison, Discussion: Acceptance Tests with Adapted File Sizes -Write Rate and Caching Effects



 $V_{\rm W}$ = 47 x 100GB = 4.7TB

Storage Cache:

 $44 \times 2 \times 16GB = 1.4TB$

Storage cache absorbs about 30% of the data to write. Resulting systematic rate error: 1/(1-0.3) = 1.43

Adapted File Size Test:

 $V_W = 47 \times 1000 \text{GB} = 47.0 \text{TB}$

Storage cache absorbs about 3% of the data to write. Resulting systematic rate error: 1/(1-0.03) = 1.03Reduction of systematic rate error by (1-0.3)/(1-0.03) = 0.72

In reality: 302.2GBs⁻¹ vs 204.6GBs⁻¹ means relative reduction by 0.68 Mind further caching layers (HDDs 16/18x106x44x256MB = 1.06TB)







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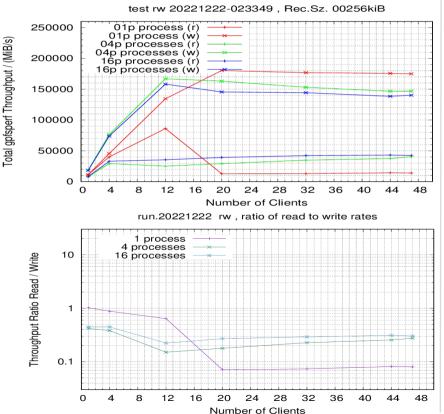
Comparison, Discussion: Read-Write-Imbalance

With increasing load, writes starve reads.

At low load, read and write rates are

balanced.

An uncontrolled stream of IO requests will get a rate reciprocal to the I/O request service times.









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Comparison, Discussion : I/O service times

Karlsruhe Institute of Technology





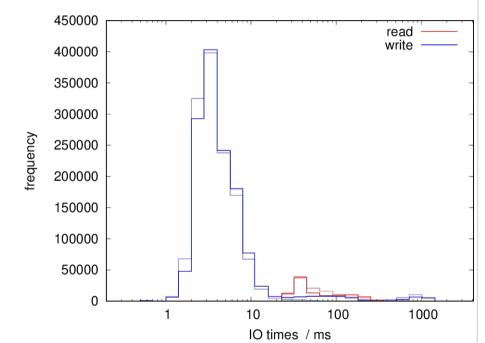
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IO statistics taken at two different times during Acceptance Benchmark with Adapted File Sizes (run time 248s, data of lighter curves taken at sec 98, data of darker curves taken at sec 175)

Write service times are about 1/10 of read service times.

=> the imbalance between read and write rate is caused by the storage system layer (controllers)!



Comparison, Discussion : Read-Write Imbalance

WIORs occupy controller cache while RIORs wait for data from HDDs

Cached WIORs are acknowledged to server immediately, cached data become precious and need to be preferrably flushed to disk.

Control of IO

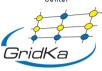
- either at filesystem level or
- above (in the application)

Real Life: Mostly no permanent pressure for read and write IOs at the same time.

BUT IF: BEWARE!







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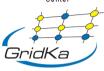
Conclusion

- Storage deserves second thoughts sometimes :-)
- Current HDD-based storage (with write caches) behaves very assymetric for reads and writes (supposed to be general, not vendor-specific!) .
- Assymetry becomes important in case of true competion writes starve reads. Rate control should be done on application level, if required.
- For flash-based storage that assymetry should be much smaller (if existing at all)
- Do not forget about caches (at various levels) and their effects, use sufficiently large data volumes
- IO benchmarks can be done easily with simple tools like gpfsperf subject to correct usage.









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