Neural Networks learning landscapes and post-quantum cryptographic primitives : a statistical physics approach

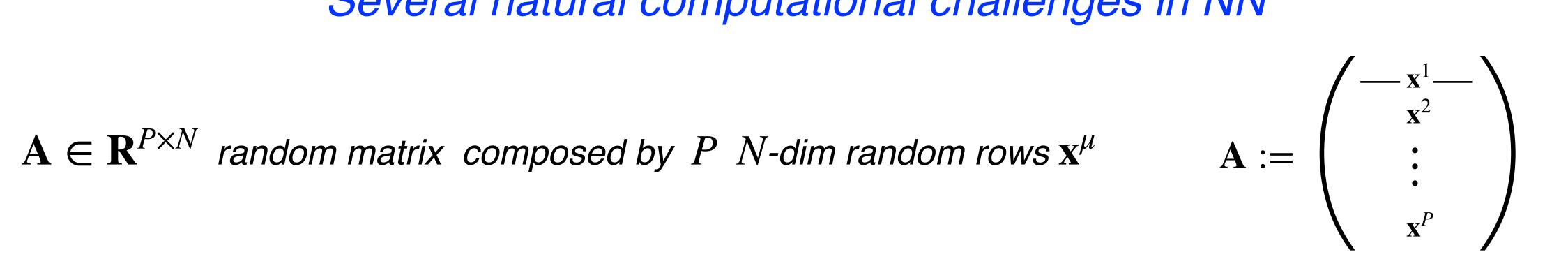
Marco Benedetti, Andrej Bogdanov^{*}, Enrico Malatesta, Marc Mezard, Gianmarco Perrupato, Alon Rosen, Nikolaj I. Schwartzbach, and Riccardo Zecchina

<u>Department of Computing Sciences, Bocconi University, Milan</u>

*University of Ottawa

Work in progress!

Several natural computational challenges in NN

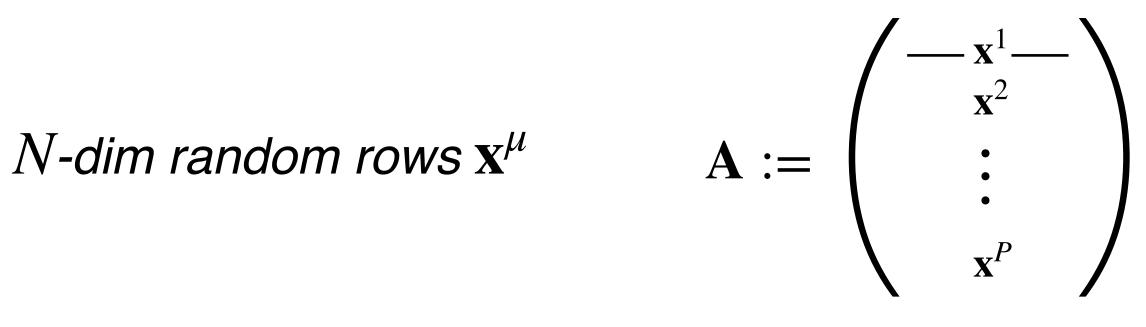


 $\mathbf{A} \in \mathbf{R}^{P \times N}$ random matrix composed by *P N*-dim random rows \mathbf{x}^{μ}

weights $W \in \{-1,1\}^N$ such that $y_A(W) = y$, assuming such W exists.

•Teacher-student: given $\mathbf{A} \in \mathbf{R}^{P \times N}$ and the labels $\hat{\mathbf{y}} = y_{\mathbf{A}}(W) \in \{-1,1\}^{P}$ for

Several natural computational challenges in NN

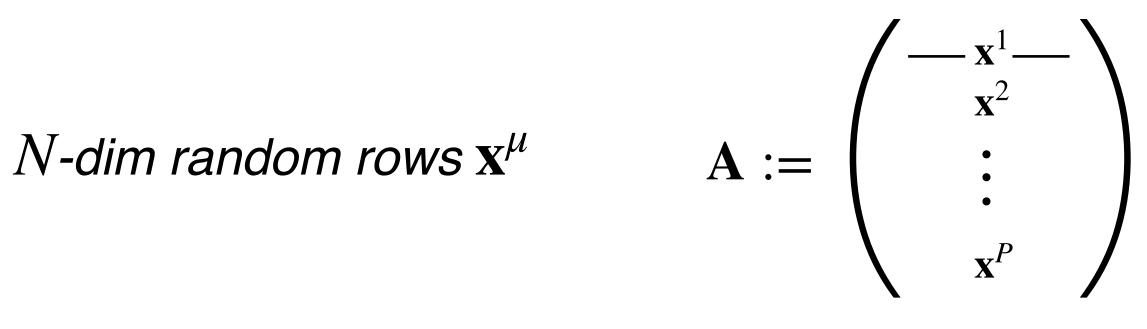


- •Inversion (learning): given $\mathbf{A} \in \mathbf{R}^{P \times N}$ and the labels $\mathbf{y} \in \{-1,1\}^P$, find any set of
- uniformly sampled $W \in \{-1,1\}^N$, find any $W' \in \{-1,1\}^N$ such that $y_A(W') = \hat{\mathbf{y}}$

 $\mathbf{A} \in \mathbf{R}^{P \times N}$ random matrix composed by *P* N-dim random rows \mathbf{x}^{μ}

- •Inversion (learning): given $\mathbf{A} \in \mathbf{R}^{P \times N}$ and the labels $\mathbf{y} \in \{-1,1\}^P$, find any set of weights $W \in \{-1,1\}^N$ such that $y_A(W) = y$, assuming such W exists.
- •Teacher-student: given $\mathbf{A} \in \mathbf{R}^{P \times N}$ and the labels $\hat{\mathbf{y}} = y_{\mathbf{A}}(W) \in \{-1,1\}^{P}$ for uniformly sampled $W \in \{-1,1\}^N$, find any $W' \in \{-1,1\}^N$ such that $y_A(W') = \hat{\mathbf{y}}$
- Collision finding: given $\mathbf{A} \in \mathbf{R}^{P \times N}$, find any two $W \neq W' \in \{-1,1\}^N$ such that $y_{\mathbf{A}}(W) = y_{\mathbf{A}}(W')$ (unexplored so far).

Several natural computational challenges in NN



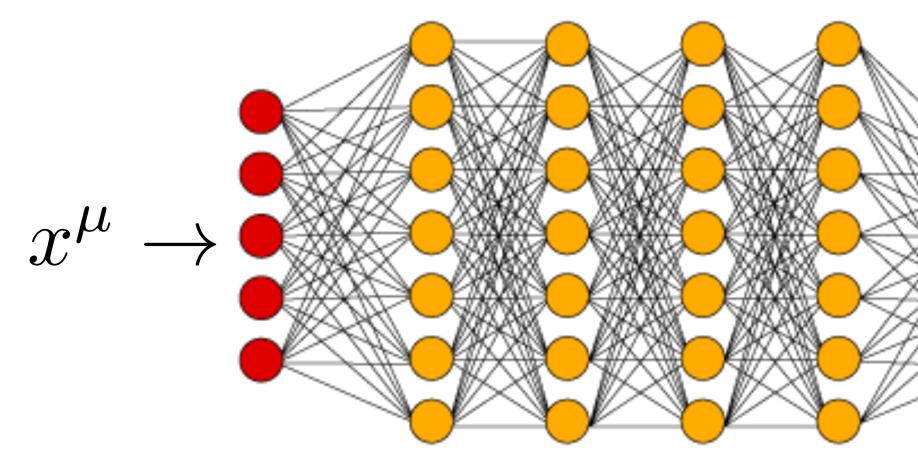
Plank of the talk

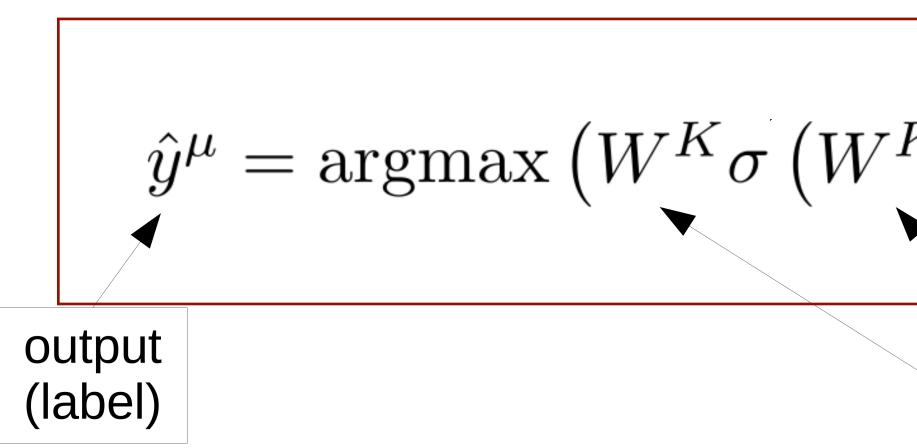
- Local entropy in non convex Neural Networks
- The Overlap Gap Property (OGP) and limiting performance of stable algorithms
- Random functions and post-quantum cryptography
- Collision Robust Hash function from random NN and their OGP transition

Training large deep neural network is in principle a non-convex hard computational problem.

- based algorithms)
- 2. Lead to solutions which have good generalisation properties
- 3. "Benign" overfitting even in presence of noise!

- Evidence about learning huge data sets with largely overparametrized networks:
 - 1. Algorithmically easy for relatively simple algorithms (e.g. gradient





Given an input vector of size N, the network computes an output by alternating layers of linear transformations with non-linear activation functions.

Training set:
$$\{(x^{\mu}, y^{\mu})\}_{\mu=1,...,n}$$

 \hat{y}^{μ} with y^{μ}
 $\Delta^{\mu} \stackrel{\circ}{=} (W^{K} \cdot \sigma_{K}(x^{\mu})) y^{\mu}$
pre-activations at layer K

$$K^{-1}\sigma \left(\dots \sigma \left(W^{2}\sigma \left(W^{1}x^{\mu}\right)\right)\right))$$
weights
(matrices) input
(vector)





Energy function and surrogate energy functions

- Energy = "0-1 loss": number of errors on the training set (not differentiable)

$$\mathcal{L}_{NE} = \sum_{\mu} (1 - \delta(\hat{y}^{\mu}, y^{\mu}))$$

- Surrogate differentiable energies

Mean Square error

 $\mathcal{L}_{MSE} = \sum ($

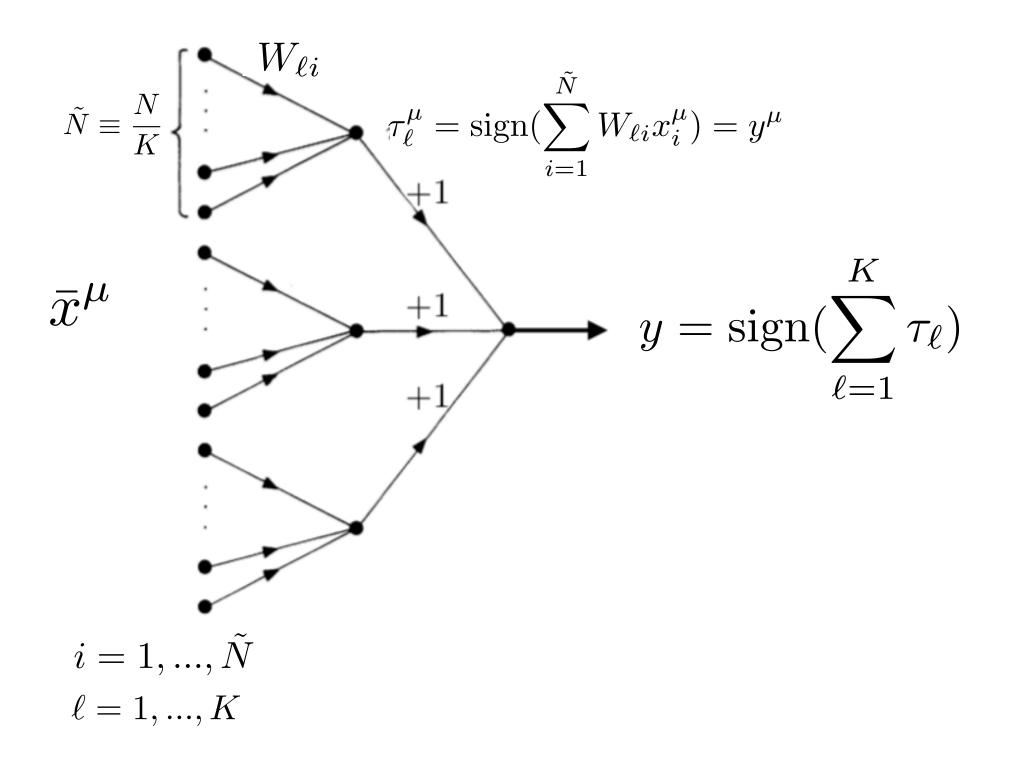
Cross-entropy: softmax

$$(\hat{\Delta}^{\mu} - \Delta^{\mu})^2$$

$$\mathcal{L}_{CE} = -\sum_{\mu} (\hat{\Delta}_{y^{\mu}}^{\mu} - \log \sum_{k} \exp \gamma \hat{\Delta}_{k}^{\mu})$$

 $\Delta^{\mu} \stackrel{\circ}{=} \left(W^K \cdot \sigma_K(x^{\mu}) \right) y^{\mu}$





Non convex also for K=1 Results generalise to networks with continuous weights

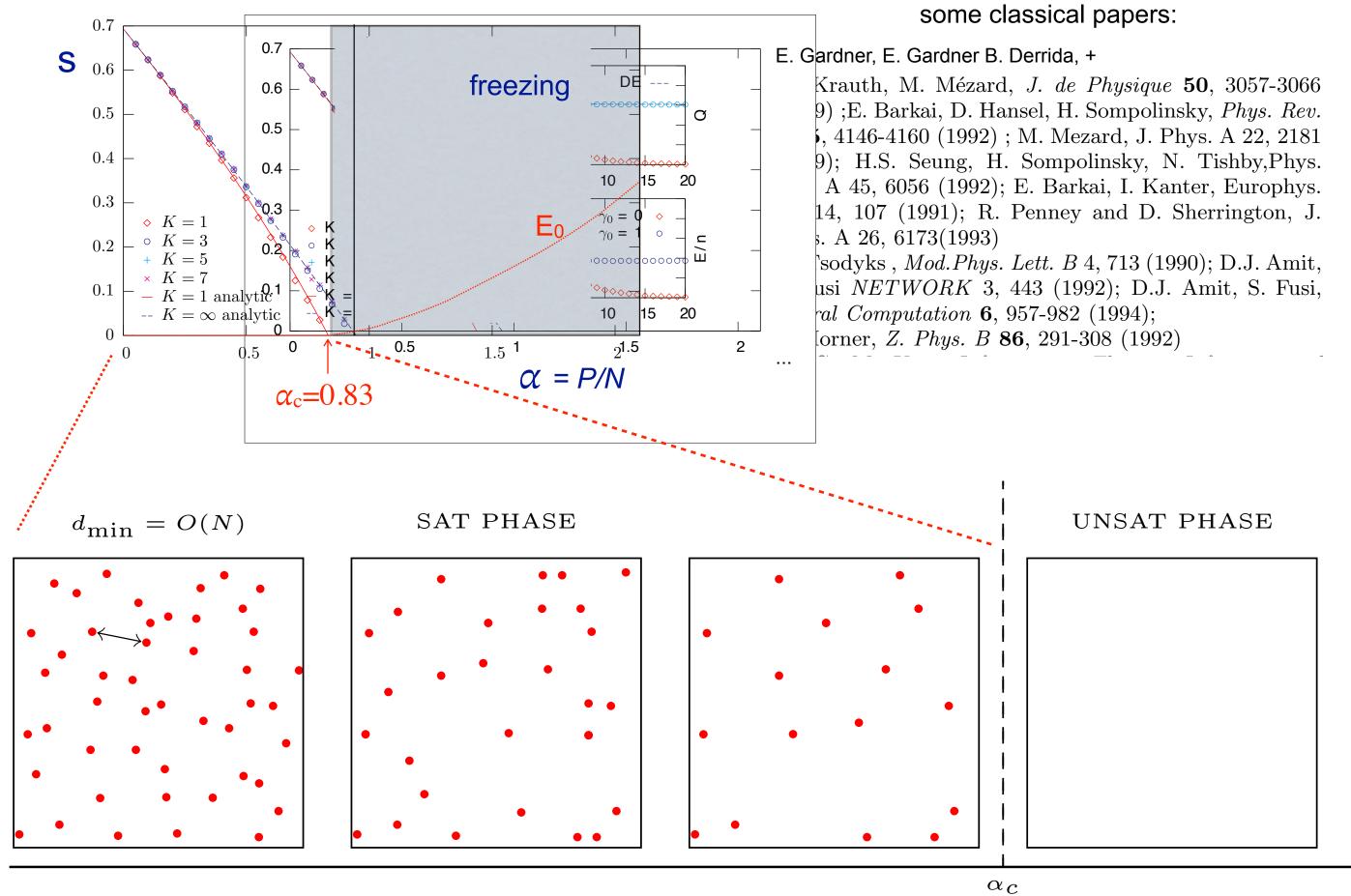
Simplest non convex neural device : 1-hidden layer, i.i.d. random associations

training set: $\{(\bar{x}^{\mu}, y^{\mu})\}\ \mu = 1, ..., P = \alpha N$ $x_{\ell i}^{\mu} = \pm 1$ (*i.i.d.* p = 1/2) $y^{\mu} = \pm 1$ (*i.i.d.* p = 1/2) control parameter: $\alpha = \frac{\text{\# patterns}}{\text{\# weights}}$



In the large N, P limit (with $\alpha = P/N$ fixed):

- $\forall \alpha > 0$ typical solutions are isolated (Huang, Kabashima (2014));
- Rigorous Proofs: Abbe, Li, Sly (2021), Perkins, Xu (2021), Nakajima Sun (2022).



Learning in the K=1 binary perceptron

• the space of solution splits into separated states of vanishing entropy (Gardner, Derrida, (1988); Krauth, Mézard (1989));

What has early Stat. Mech. brought to the field?

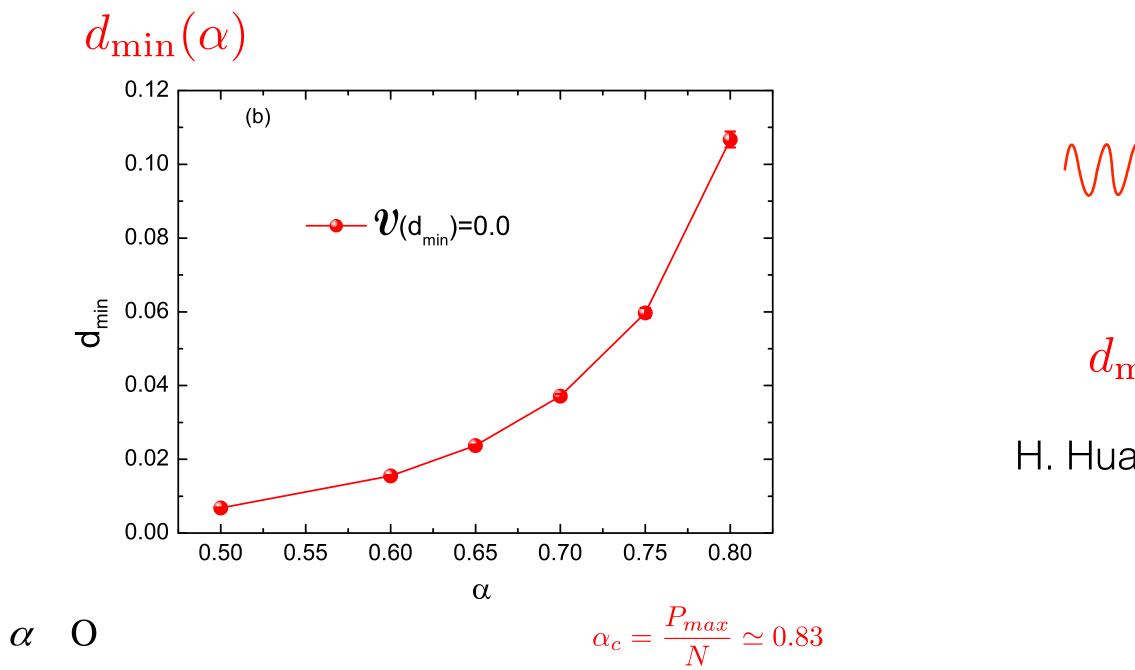
- Phase transitions
- Probabilities in large dimensions
- Dynamics
- New algorithms for convex perceptrons
- ...

lpha

Geometry of the space of solutions:

Franz-Parisi potential: entropy at distance **d**, sampling from typical solution **J**

$$F(x) = \left\langle \frac{1}{Z(T')} \sum_{\mathbf{J}} \Theta\left(\frac{1}{\sqrt{N}} \sum_{i=1}^{N} J_i \xi_i^{\mu}\right) \ln \sum_{\mathbf{w}} \Theta\left(\frac{1}{\sqrt{N}} \sum_{i=1}^{N} w_i \xi_i^{\mu}\right) e^{x \mathbf{J} \cdot \mathbf{w}} \right\rangle_{\boldsymbol{\xi}}$$



Golf course for any α ?

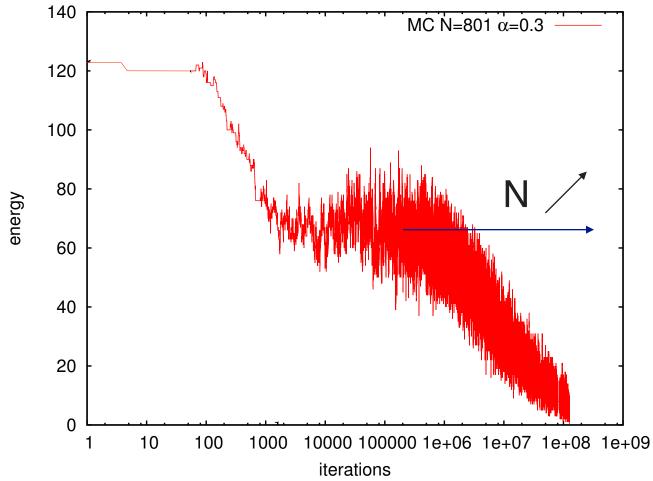
Efficient learning impossible ?

However other algorithms find solution efficiently!

$$\alpha$$
 α

 $d_{\min}(\alpha) \sim O(N)$

H. Huang, Y. Kabashima (2014)

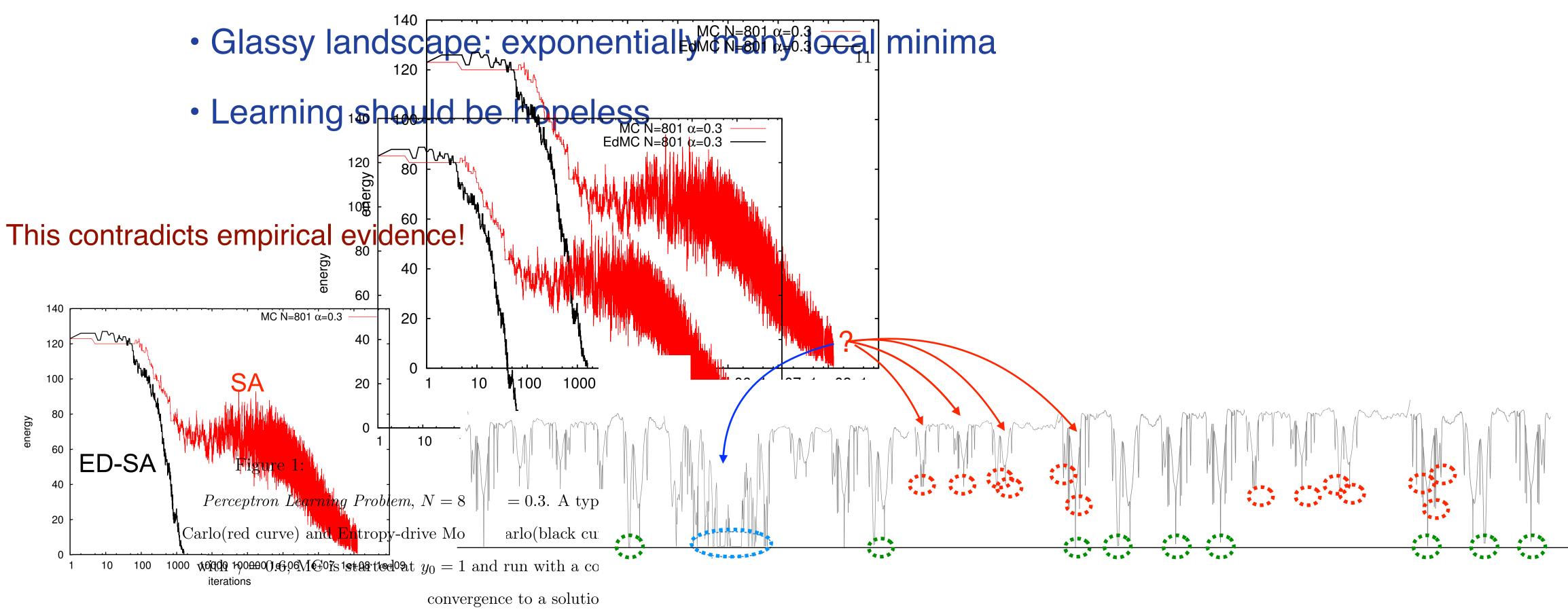


Sampling from the Gibbs distribution is not a good algorithm (as expected)



The learning problem is predicted to be typically computationally difficult

Typical global minima are isolated (mutual distance of O(N))



We performed extensive simulations and studied the scaling properties of EdMC in contrast to simulated annealing. Figure 2 is a log-log plot of the number of iterations $n_{E=0}$ to



Braunstein, Z. PRL 2006



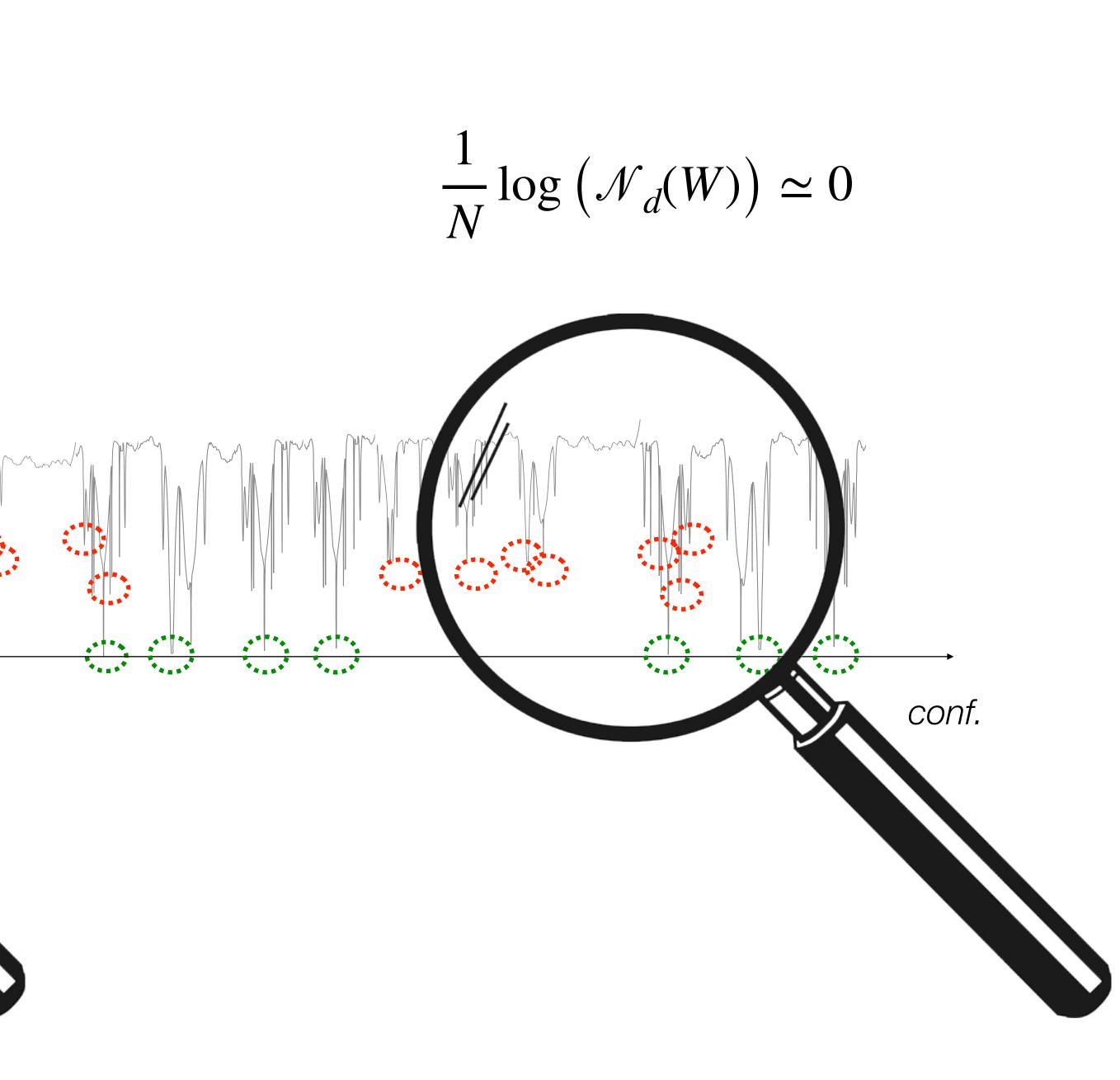
(perception) maps vectors qi, 11 suipuls The veelongits binary outputs as $\tau(W,\xi) = \text{sign}(W,\xi)$, where $W \in \{-1,1\}^N$ is the vector of synaptic weights. Given the property of the synaptic weights in the synaptic weights in the synaptic weights. 0.005 input patterns ξ^{μ} with $\chi espending N deside the put s$

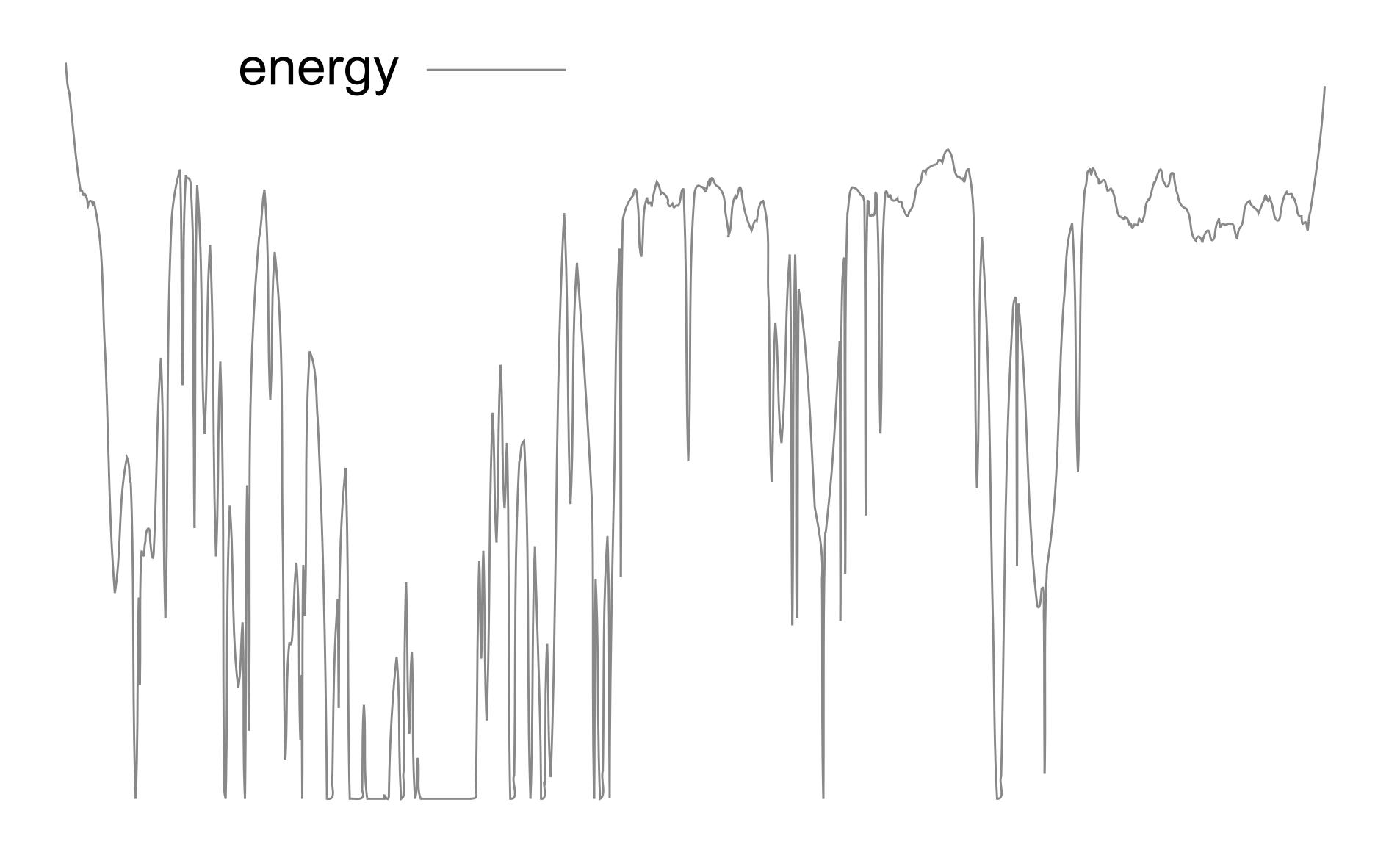
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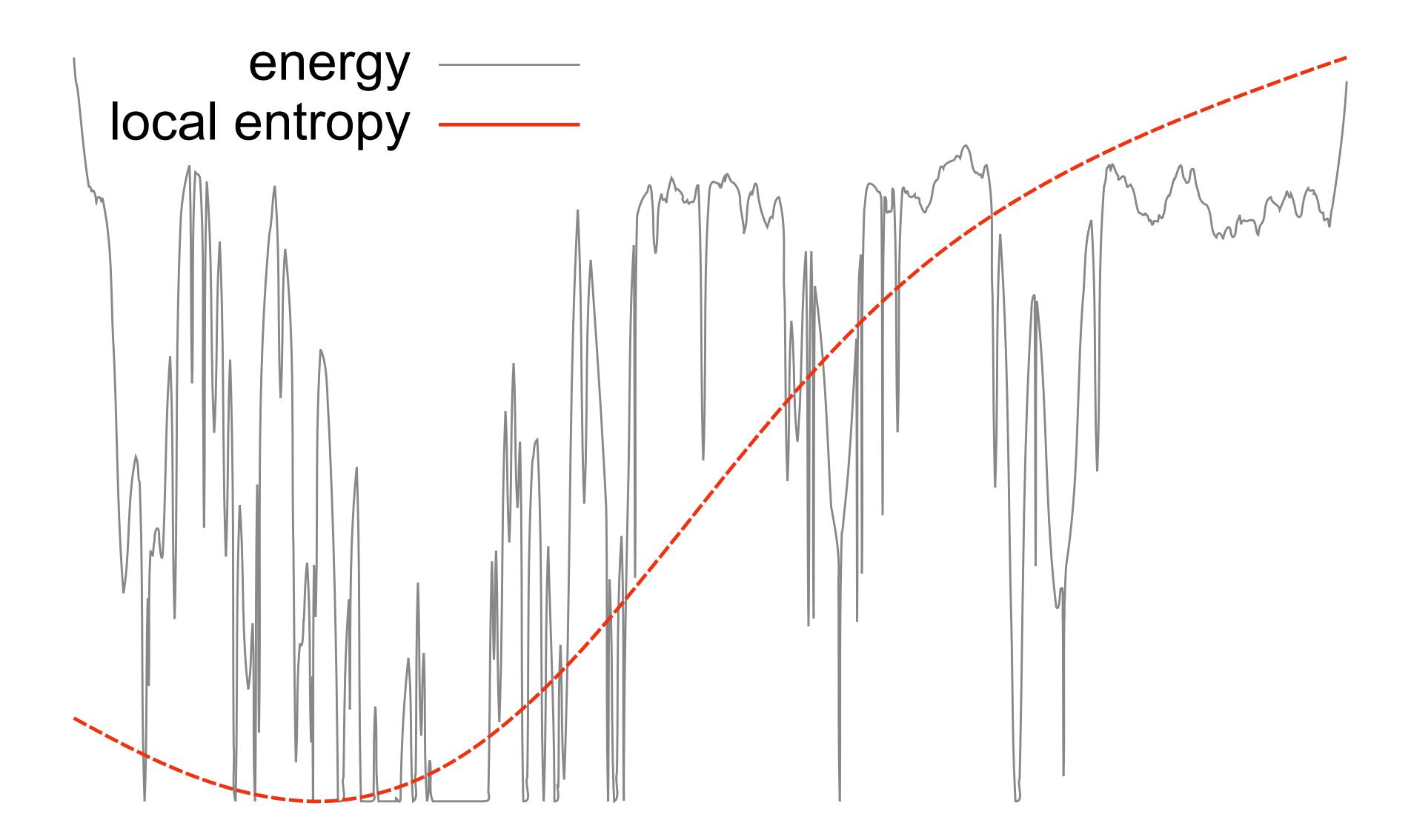


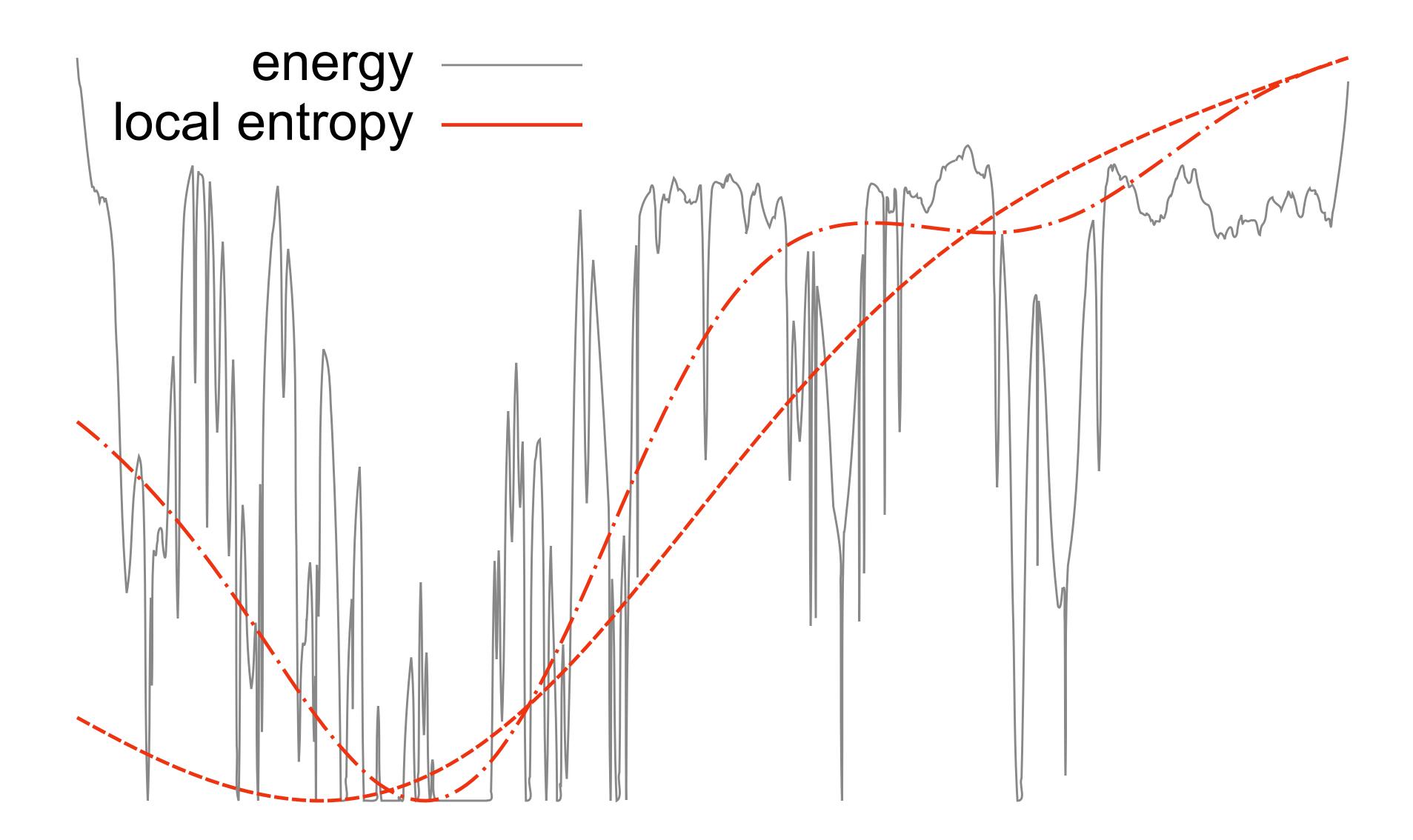
distance from

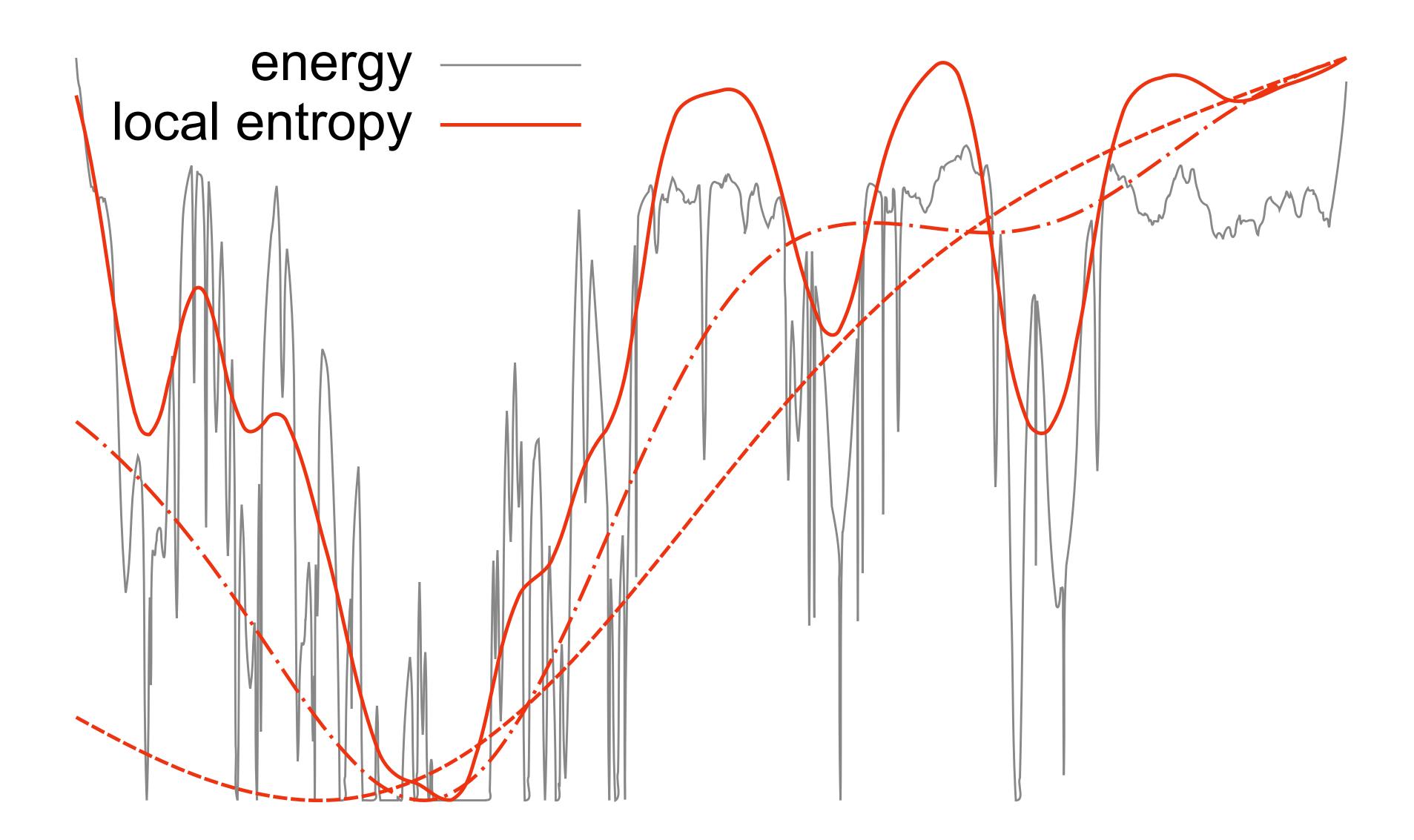
 $\frac{1}{N}\log\left(\mathcal{N}_d(W)\right) > 0$ d***********









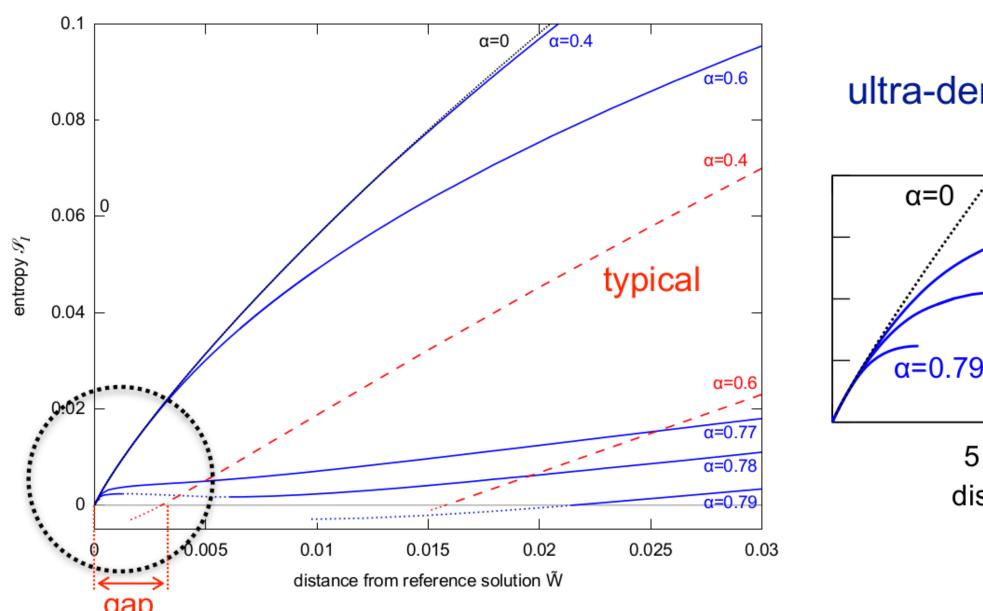


Finite temperature version (not only zero error states) : free local entropy

$$P(\tilde{W}) = \frac{e^{y\Phi(\tilde{W},\beta,\gamma)}}{Z} \qquad P(\tilde{W}) = \frac{e^{-y\mathcal{E}_d(\tilde{W},d)}}{Z}$$

Find \tilde{W} that maximises the local free entropy the: $\operatorname{argmin}_{\tilde{W}}\langle \Phi(\tilde{W},\beta,\gamma)\rangle$

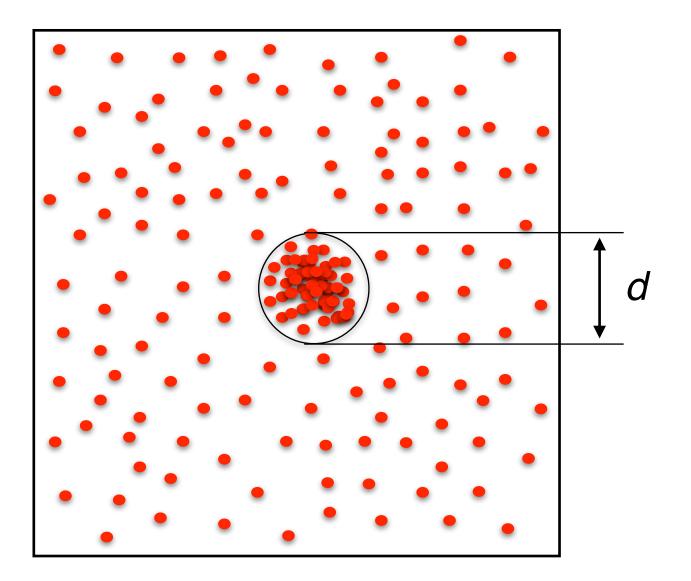
P



is of solutions in the Binary Perceptron

$$(\beta \to \infty)$$

$$(\tilde{W}) = \lim_{y,\beta \to \infty} \frac{e^{y\Phi(\tilde{W},\beta,\gamma)}}{Z}$$



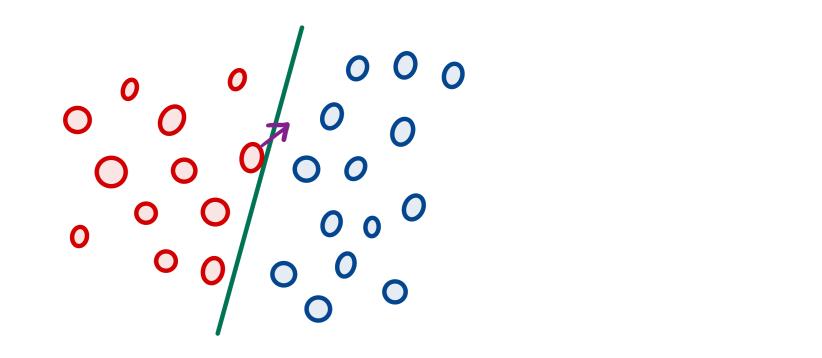
ultra-dense cluster

Geometrical phase discontinuous transition

$$\alpha_u \simeq 0.77$$



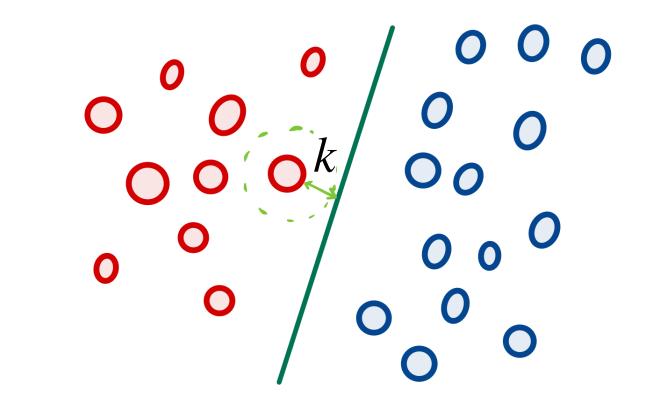
 $\mathbb{X}_{\xi,F}(W)$ Solutions with margin k:



$$E = \sum_{\mu=1}^{P} \Theta \left(-y^{\mu} \sigma_{\text{out}}^{\mu} \right)$$

How to connected the results with the "traditional" maximum margin studies ?

$$;\kappa) = \prod_{\mu=1}^{P} \Theta \left(y^{\mu} \sigma_{\text{out}}^{\mu} - \kappa \right)$$

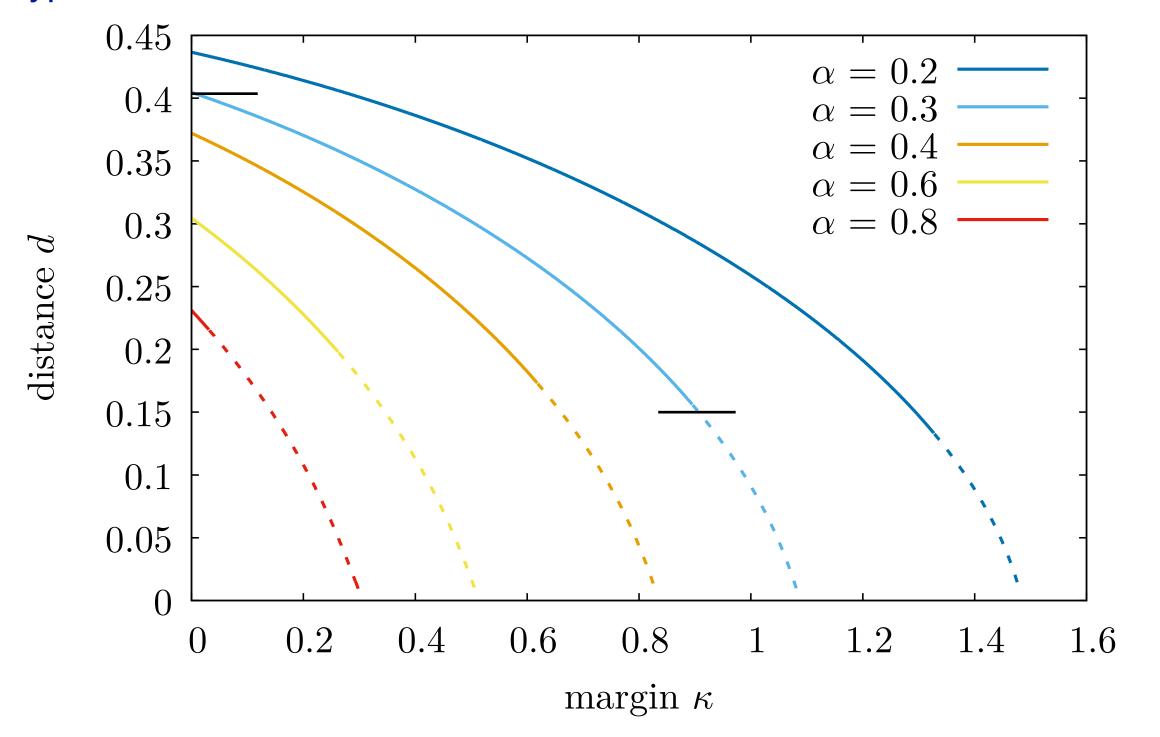


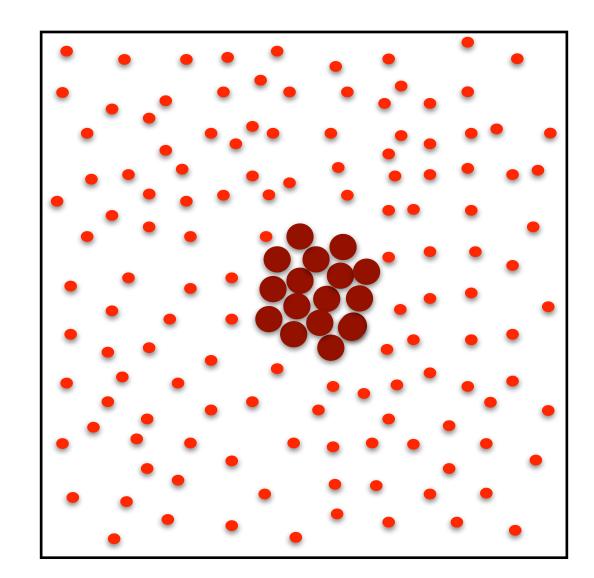
$$E = \sum_{\mu=1}^{P} \Theta \left(-y^{\mu} \sigma_{\text{out}}^{\mu} + \kappa \right)$$

High margin solutions are less but tend to be much closer to each other!

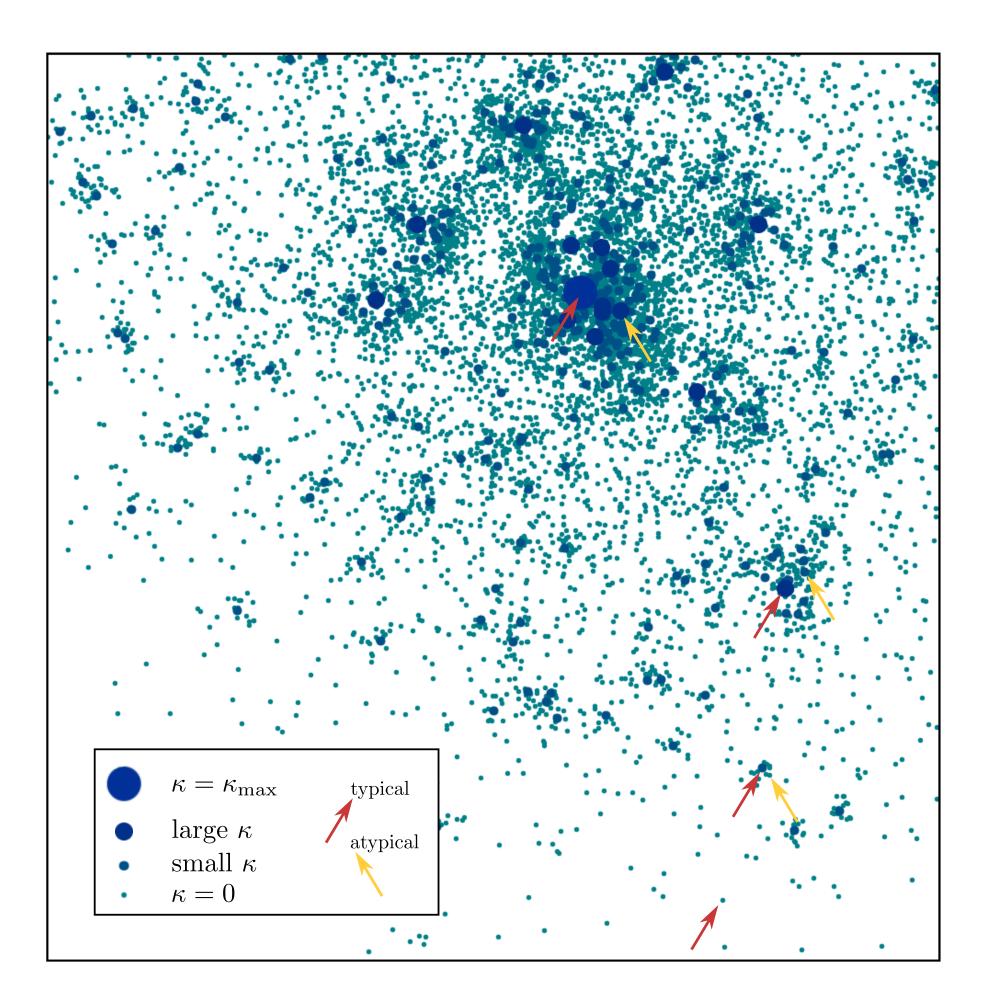
The lines change from solid to dashed when the entropy of solutions becomes negative, i.e. when $\kappa = \kappa_{max}$

typical distance





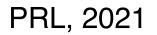
A wide flat minima arises by the coalescence of (atypical) high margin minima!

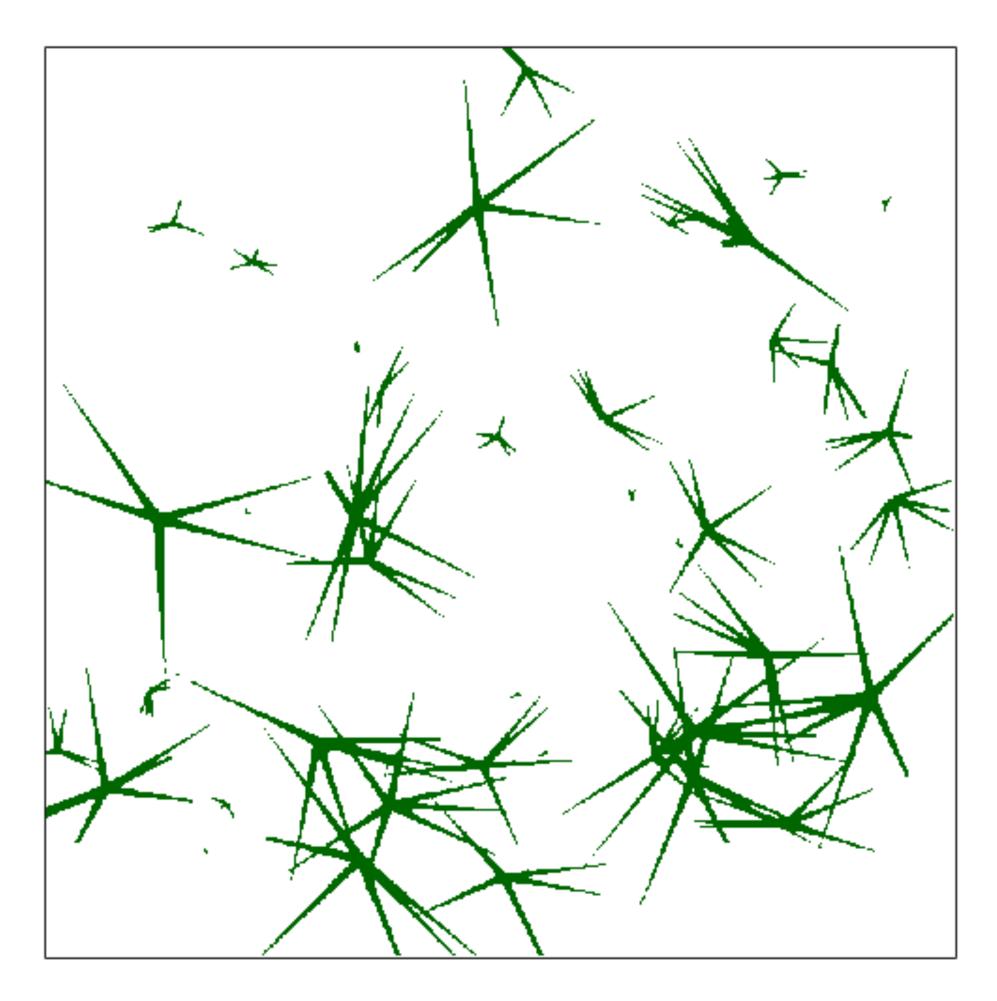


Unveiling the structure of wide flat minima in neural networks

Carlo Baldassi,¹ Clarissa Lauditi,² Enrico M. Malatesta,¹ Gabriele Perugini,¹ and Riccardo Zecchina¹ ¹Artificial Intelligence Lab, Bocconi University, 20136 Milano, Italy ²Department of Applied Science and Technology, Politecnico di Torino, 10129 Torino, Italy

arXiv:2107.01163





Binary perceptron: efficient algorithms can find solutions in a rare well-connected cluster

Emmanuel Abbe * Shuangping Li[†] Allan Sly[‡]

arXiv:2111.03084

Algorithmic follow-up

Local free entropy:

Z(y,Large-deviation partition function:

Assume y integer:

 $Z(y,\gamma,\beta',\beta) = \sum_{i=1}^{n}$ $W, \{$

 $\phi(W,\gamma,\beta) = \log \sum e^{-\beta \mathcal{L}_{NE}(W') - \frac{\gamma}{2}d(W,W')}$ W'

$$\gamma, \beta', \beta) = \sum_{W} e^{-\beta' \mathcal{L}_{NE}(W) + y \phi(W, \gamma, \beta)}$$

$$\sum_{W_a} e^{-\beta' \mathcal{L}_{NE}(W) - \beta \sum_{a=1}^{y} \mathcal{L}_{NE}(W_a) - \frac{\gamma}{2} \sum_{a=1}^{y} d(W, W_a)}$$

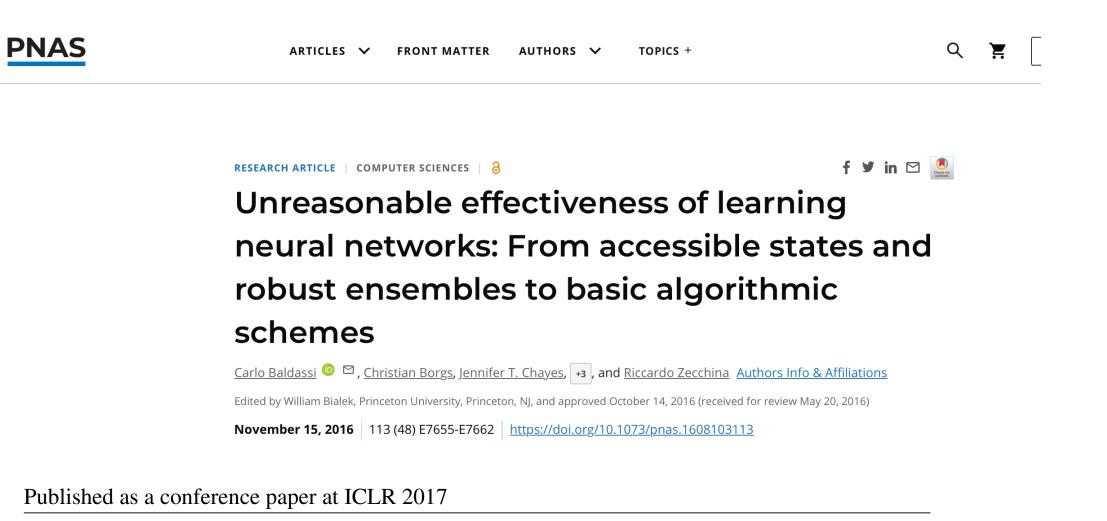
center y (real) replicas

on

Local entropy algorithms

- Local Entropy driven or replicated Simulated Annealing
- Replicated Message-Passing (Belief Propagation)
- Replicated Stochastic Gradient Descent (SGD)
- Entropy-SGD: Langevin dynamics to estimate local entropy
- Replicated Greedy Algorithms
- Sharpness Aware Minimization
- Stochastic weights +gradient on the probabilities
- Quantum Annealing delocalization mechanism for finding NN ground states

some known algorithms for DNNs



ENTROPY-SGD: BIASING GRADIENT DESCENT INTO WIDE VALLEYS

Pratik Chaudhari¹, Anna Choromanska², Stefano Soatto¹, Yann LeCun^{3,4}, Carlo Baldassi⁵, Christian Borgs⁶, Jennifer Chayes⁶, Levent Sagun³, Riccardo Zecchina⁵

Published as a conference paper at ICLR 2021



SHARPNESS-AWARE MINIMIZATION FOR EFFICIENTLY IMPROVING GENERALIZATION

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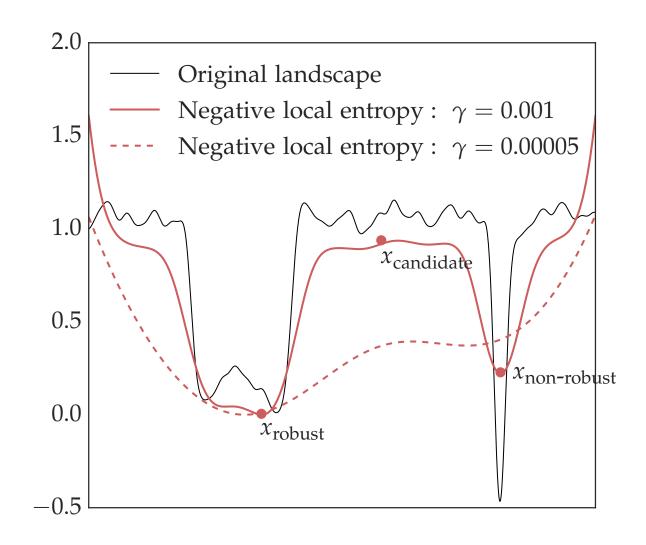
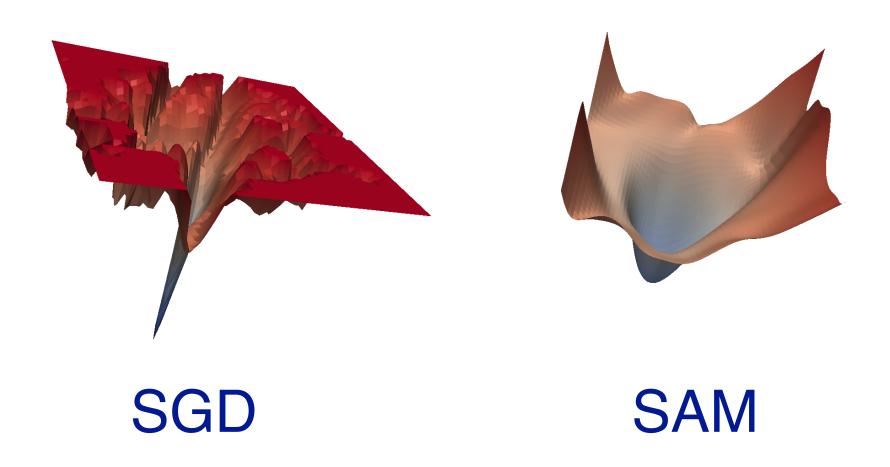


Figure 2: Local entropy concentrates on wide valleys in the energy landscape.



Similar analytical results hold for 1-hidden layer NN with continuous weights and for overparametrized NN

High Local Entropy regions - Wide Flat Minima (WFM)

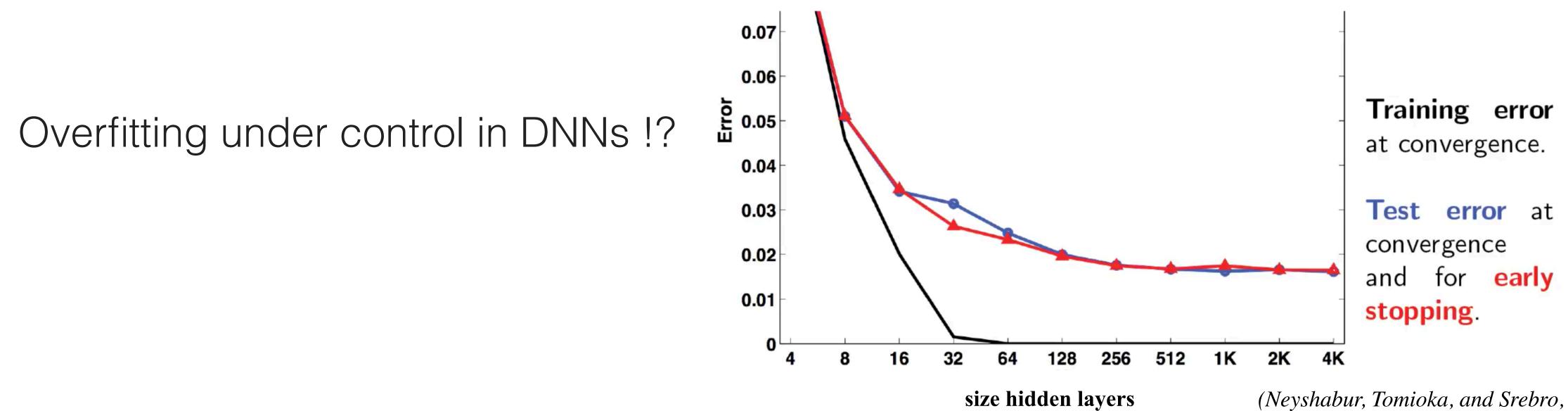
Analytical/numerical results: they have good generalisation properties

Analytical Results: Rare Wide Flat Minima (WFM) exist in non-convex networks with continuous weights storing random patterns.

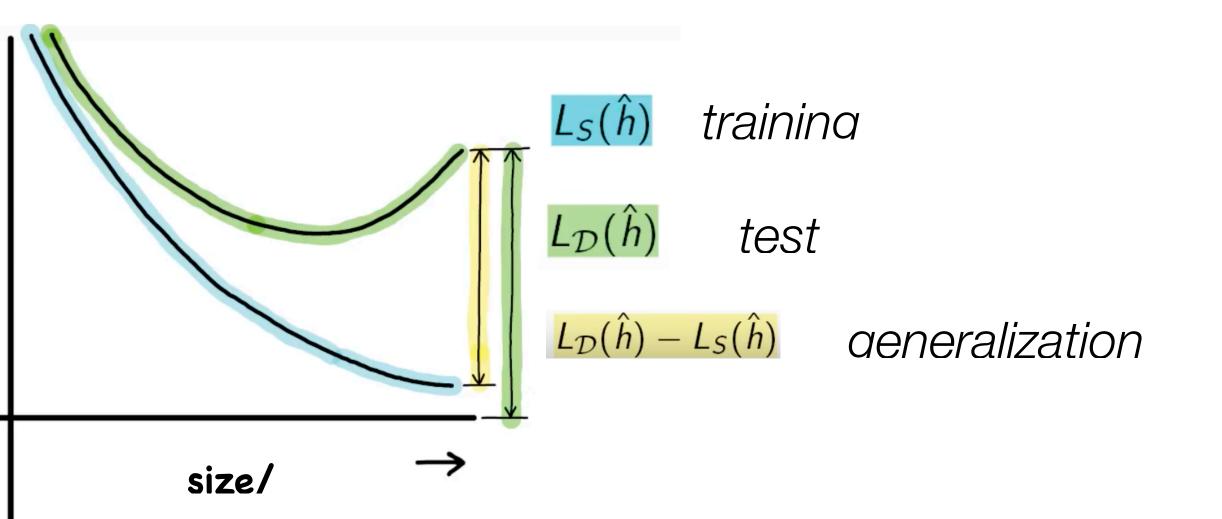
(Baldassi, Pittorino, Zecchina, PNAS 2019)

Classical overfitting problem

Using stochastic gradient descent (SGD), trained networks of increasing size on MNIST until convergence.

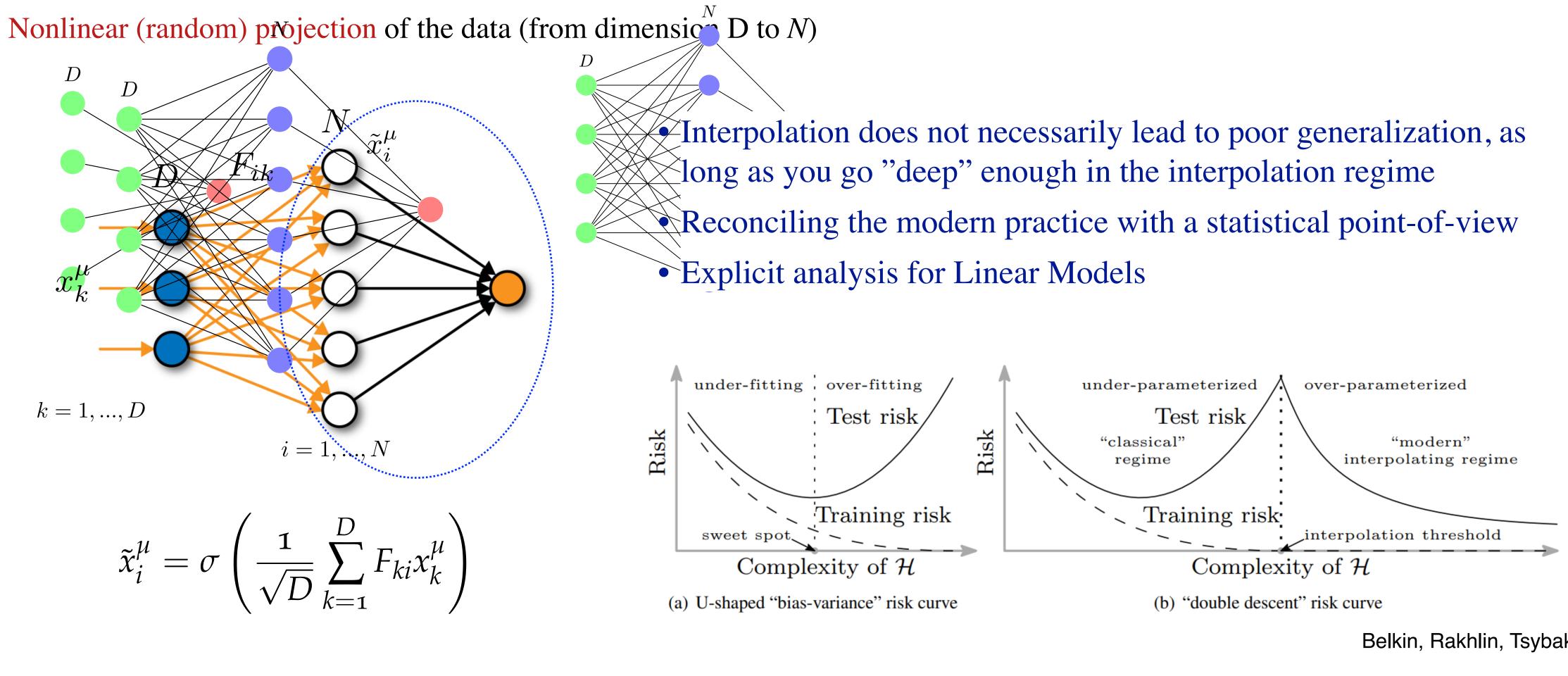


error



(Neyshabur, Tomioka, and Srebro, 2014)

Random Feature Model



 $F_{ki} \sim \mathcal{N}(0, 1)$

Neal, 1996; Balcan, Blum, Vempala 2006; Rahimi, Recht; 2008; Bach, 2016

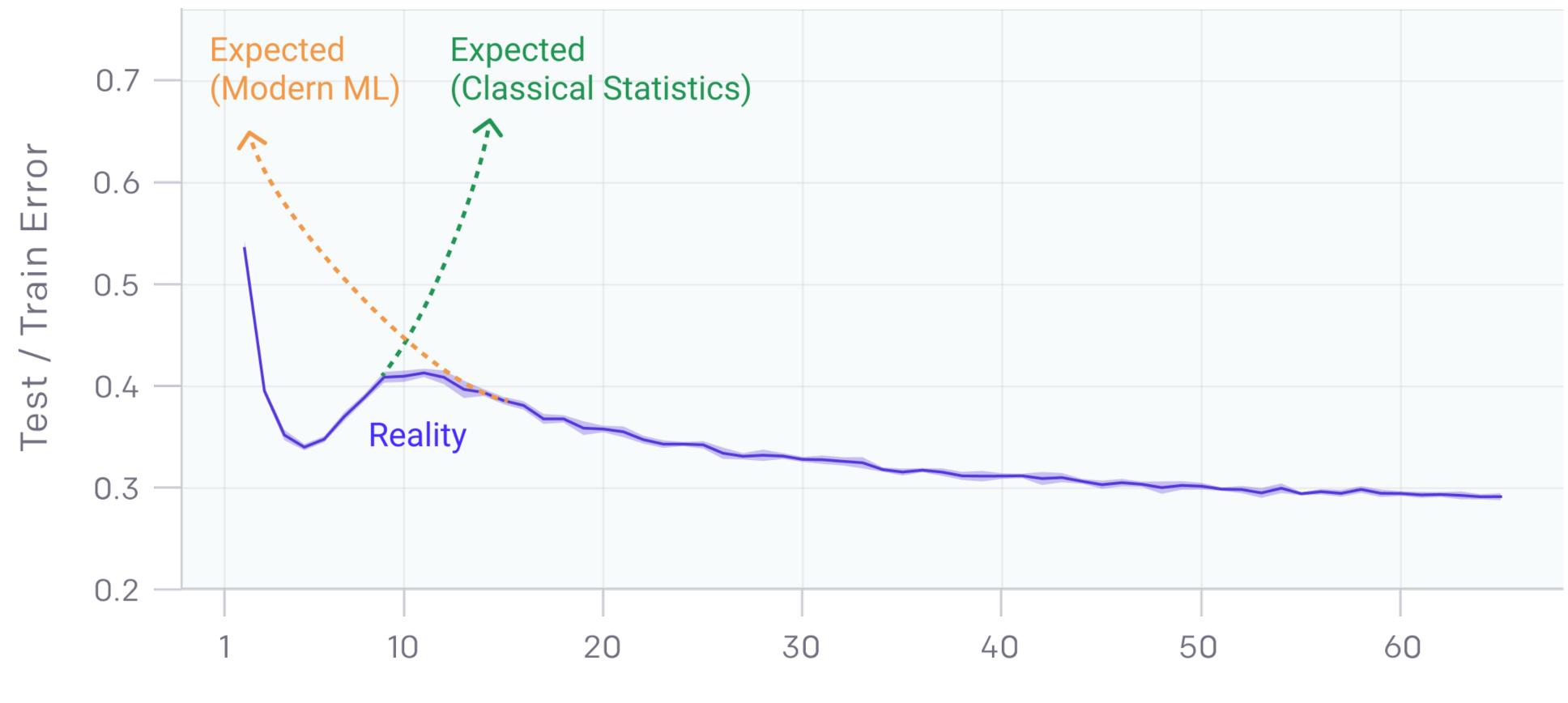
Belkin, Rakhlin, Tsybakov, 2018

• Connection with the Hidden Manifold Model

Goldt, Mezard, Krzakala, Zdeborova, 2020 Montanari, Mei, 2019



Deep Double Descent



Model Size (ResNet18 Width)

From: Deep Double Descent: Where Bigger Models and More Data Hurt P Nakkiran, G Kaplun, Y Bansal, T Yang, B Barak, I Sutskever (2019)



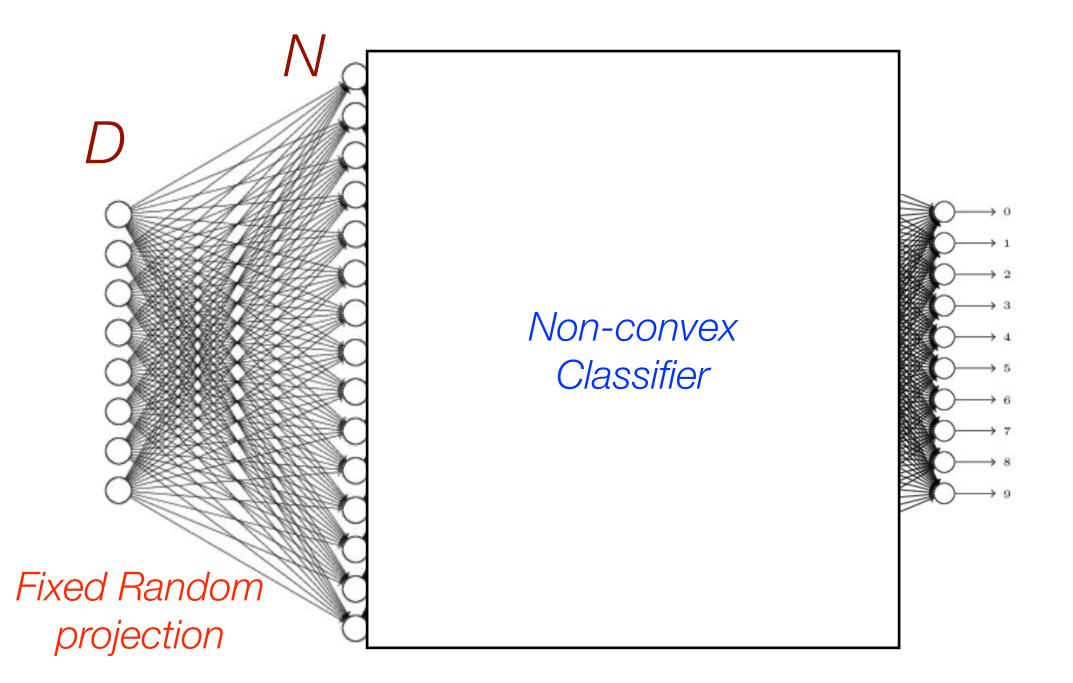
Geometric and algorithmic phase transitions on non-convex overparametrized NN

Overparametrized Regime

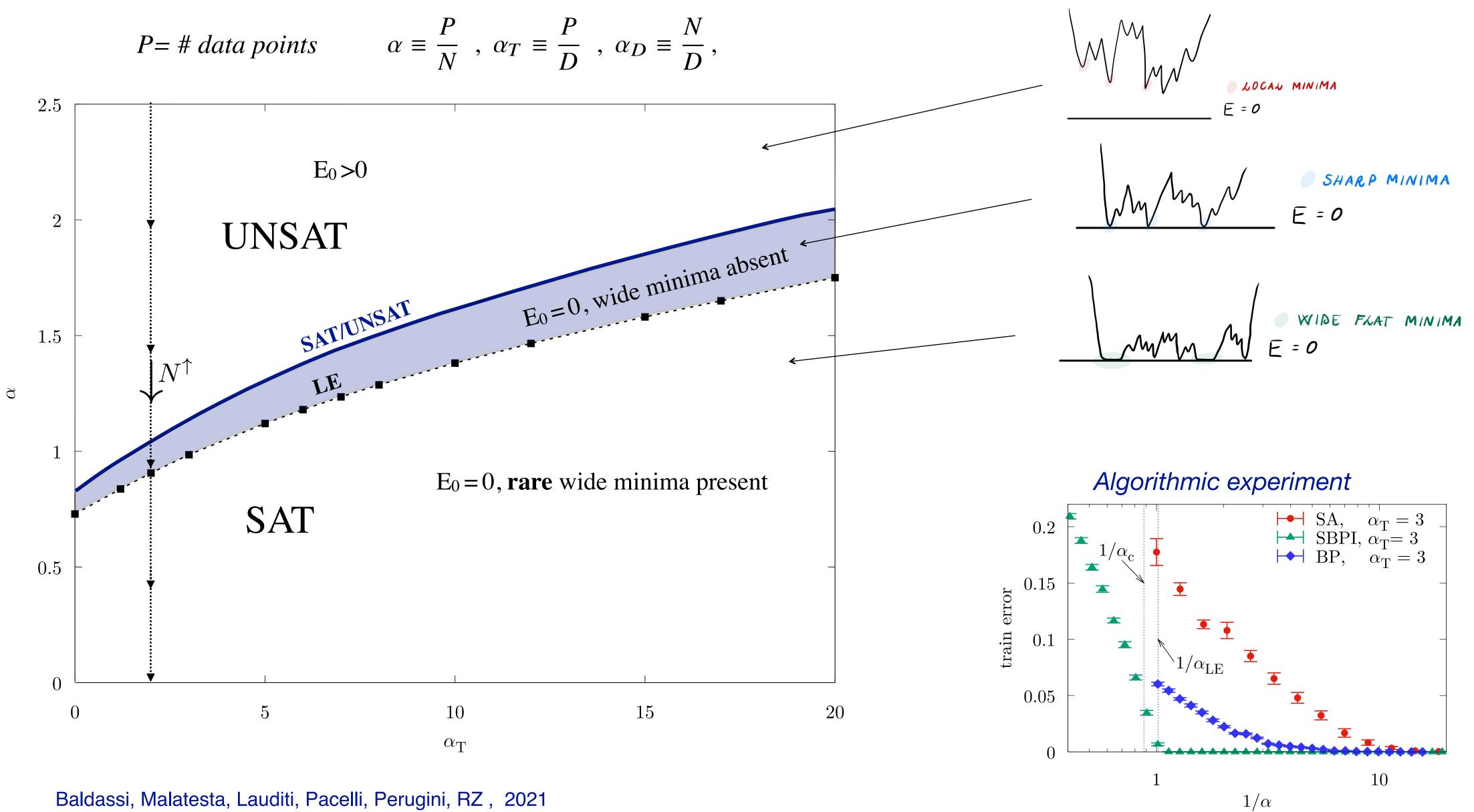
Input:
$$D \to \infty$$

 $N \to \infty$
 $\frac{D}{N} \to \psi \ll 1$

(D intrinsic dimension of the data)



Learning through atypical "phase transitions" in overparametrized neural networks



Baldassi, Malatesta, Lauditi, Pacelli, Perugini, RZ, 2021

Analytics vs Numerics in large scale DNNs

PNAS | January 7, 2020 | vol. 117 | no. 1 | 161–170

Shaping the learning landscape in neural networks around wide flat minima

Carlo Baldassi^{a,b,1,2}, Fabrizio Pittorino^{a,c}, and Riccardo Zecchina^{a,d,1,2}

Published as a conference paper at ICLR 2021

ENTROPIC GRADIENT DESCENT ALGORITHMS AND WIDE FLAT MINIMA

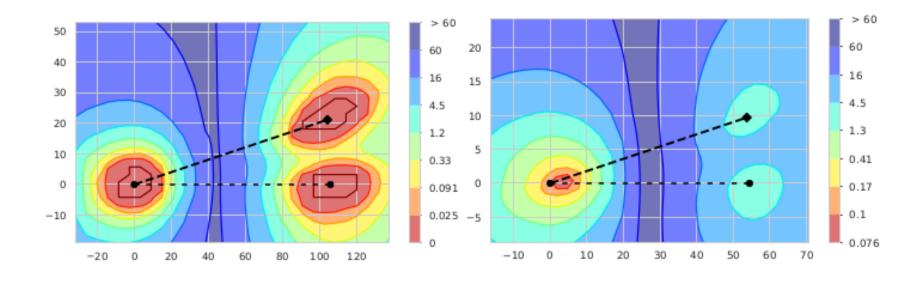
Fabrizio Pittorino^{1,2}, Carlo Lucibello¹, Christoph Feinauer¹, Gabriele Perugini¹, Carlo Baldassi¹, Elizaveta Demyanenko¹, **Riccardo Zecchina**¹

ICML 2022

Deep Networks on Toroids: Removing Symmetries Reveals the Structure of Flat Regions in the Landscape Geometry

Fabrizio Pittorino¹ Antonio Ferraro¹ Gabriele Perugini¹² Christoph Feinauer¹ Carlo Baldassi¹ **Riccardo Zecchina**¹

VGG16 on Cifar10



- Left Panel: Unnormalized
- Right Panel: Normalized
- Left Points: RSGD (finds flatter minima)
- Right Points: unaligned/aligned SGD with adversarial initialization

Difference is only visible *after* symmetry removal



The quite expensive Google experiments

(Y. Jiang, B.Neyshabur, H. M. D.Krishnan, S.Bengio, 2019)

- Training under all combination of hyperparameters and optimization resulted in a large pool of models.
- For any such model, we considered 40 complexity measures*.

Findings:

"... the relative success of sharpness-based and optimization-based complexity measures for predicting the generalization gap can provoke further study of these measures."

* complexity measure in machine learning: a quantity that monotonically relates to some aspect of generalization. Typically it depends on the trained model and the training data, but should not have access to a validation set. Lower complexity should often imply smaller generalization gap.

• Trained more than 10,000 models over two image classification datasets (CIFAR-10, Street View House Numbers).

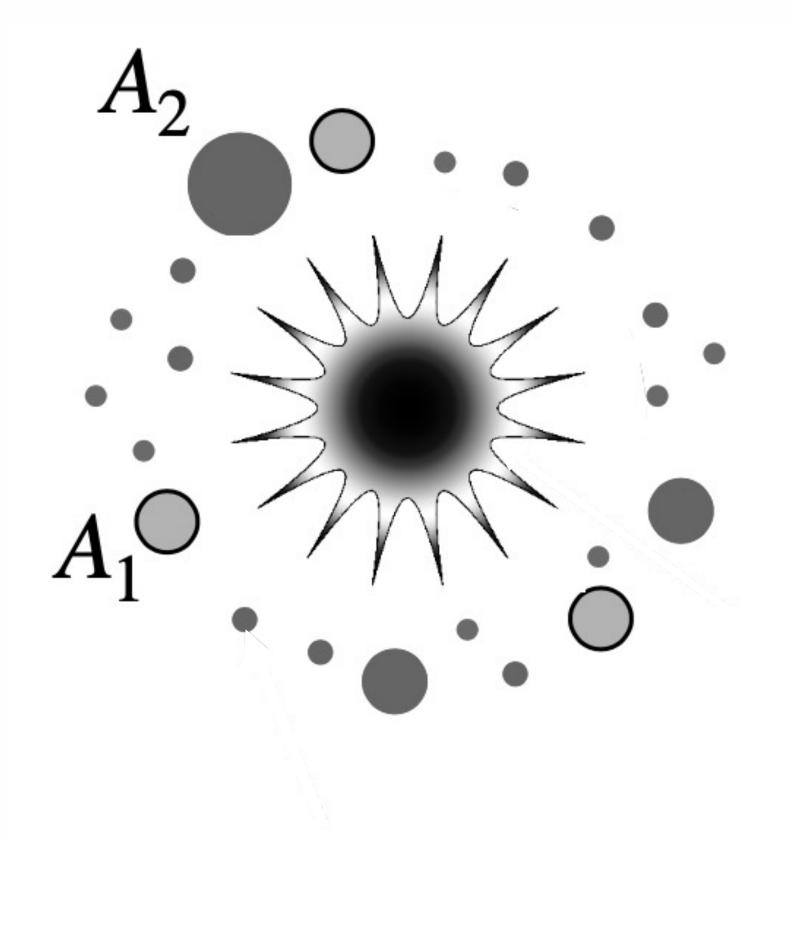




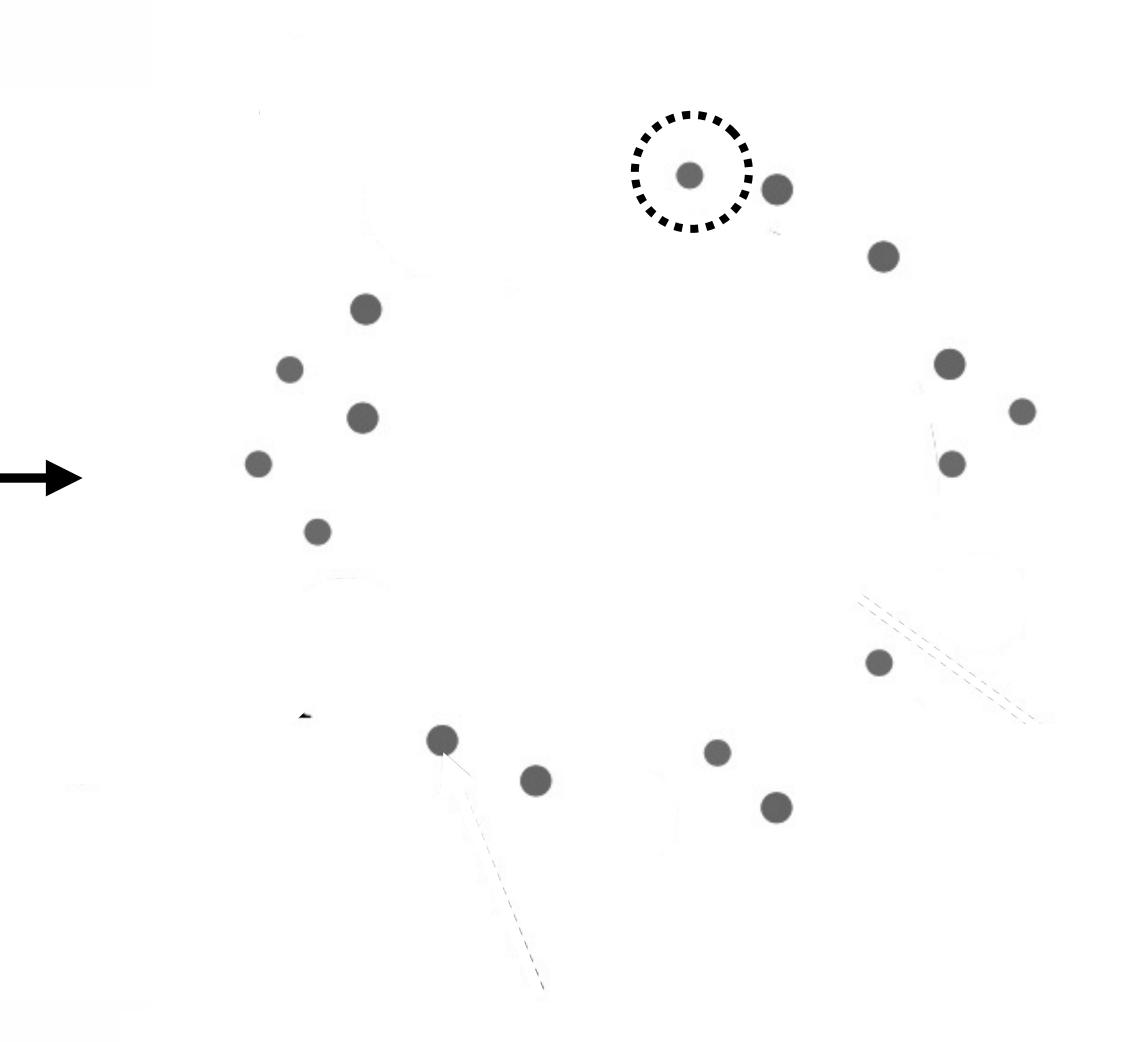
Our aim is to use random Neural Network models to build cryptographic system, based on what we know about the geometry of solutions.

The Overlap Gap Property (OGP)

OGP Definition: In high-dimensional op a "gap" structure.

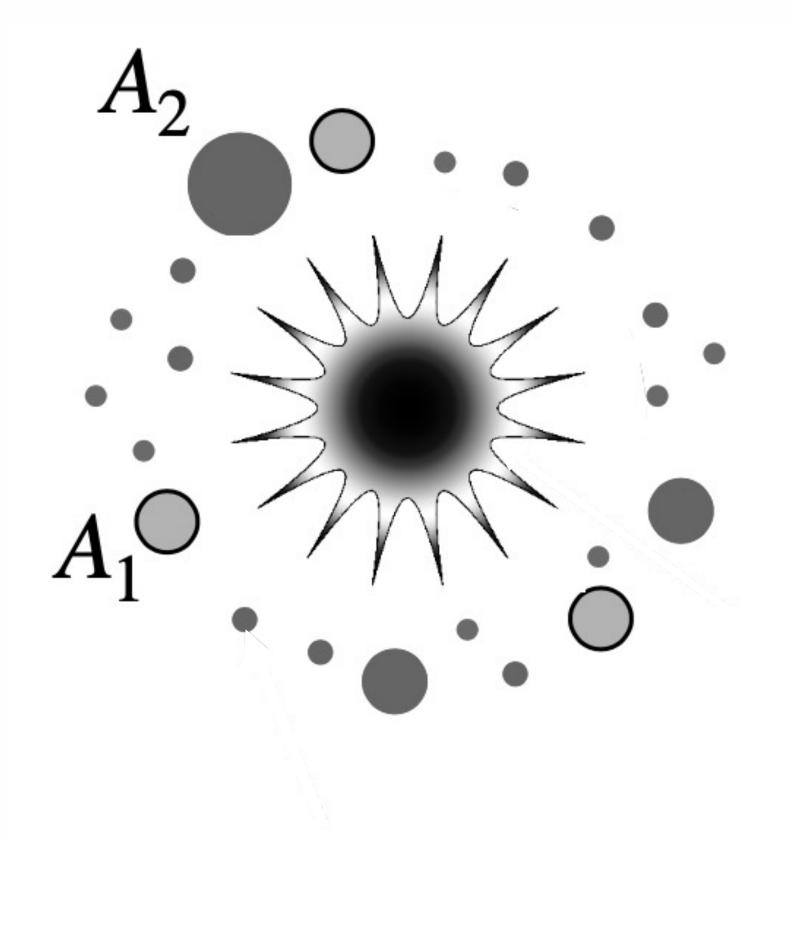


OGP Definition: In high-dimensional optimization problems, the solution space exhibits

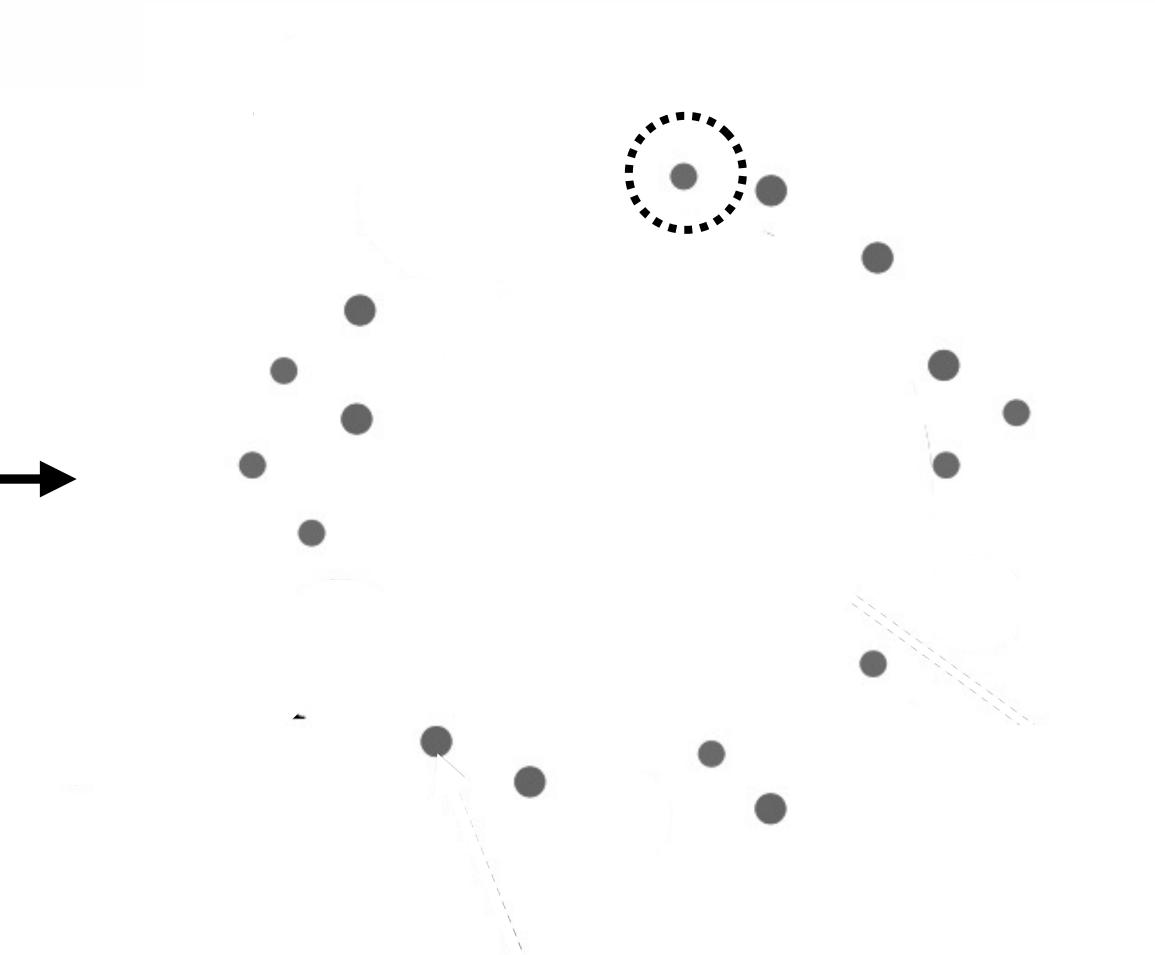


The Overlap Gap Property (OGP)

a "gap" structure.



OGP Definition: In high-dimensional optimization problems, the solution space exhibits



This scenario has to be studied for *collisions*



The Overlap Gap Property (OGP)

Theorem: (OGP and Stable algorithms), informal statement:

Let **P** be a combinatorial optimization problem exhibiting OGP.

fail to find an optimal solution with high probability.

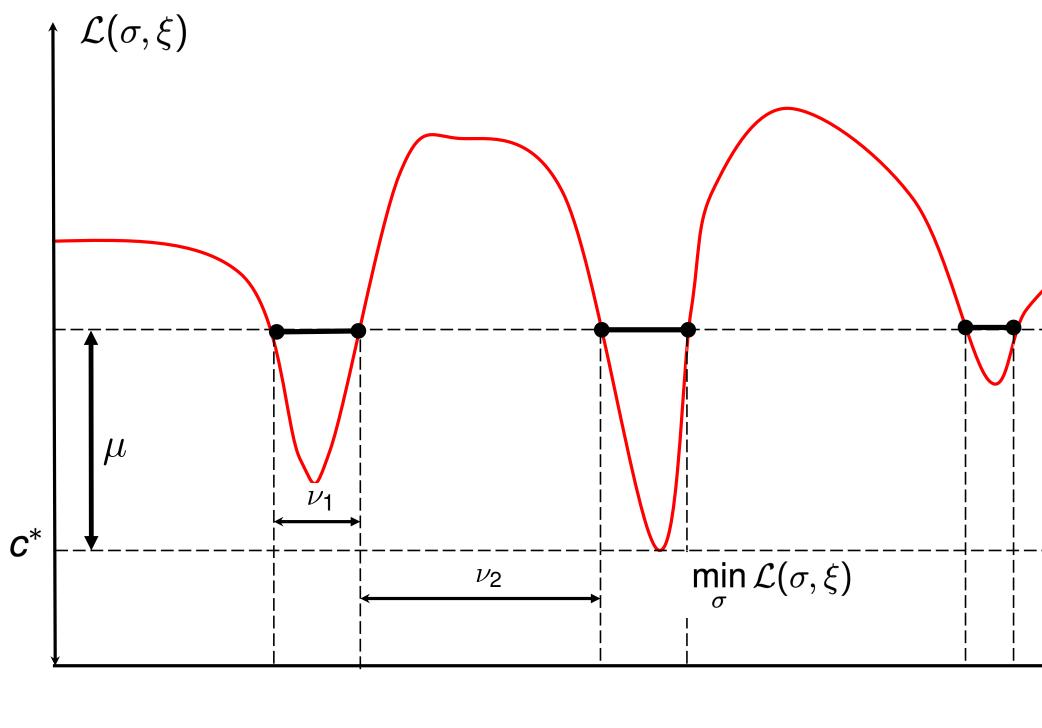
Generalised to multi-OGP: $2-OGP \Rightarrow y-OGP$ (y-OGP is sufficient)

Intuition: if there is OGP a small perturbation can result in a large change in the output, then the algorithm cannot be stable.



- If a stable algorithm is applied to solve P, then under typical conditions the algorithm will

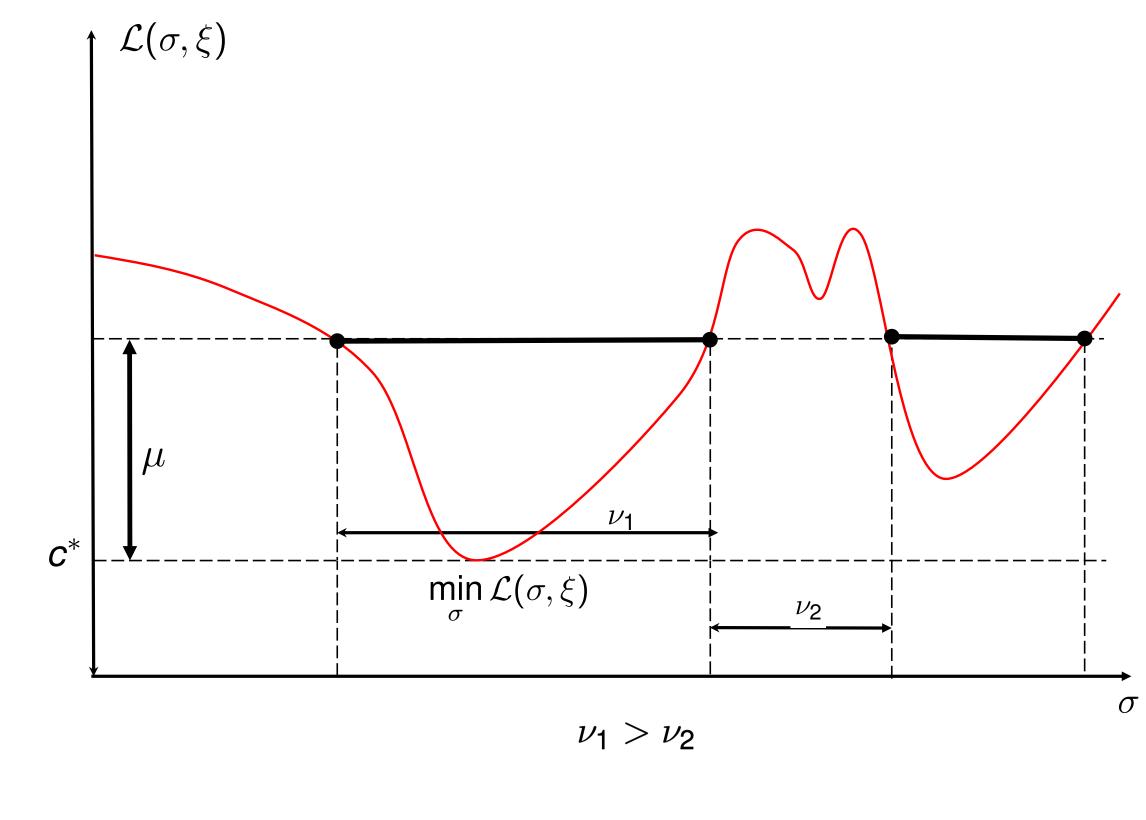




 $\nu_{1} < \nu_{2}$

 σ

Landscape exhibiting the OGP: Solutions are split into clusters, with diameter of each cluster smaller than the distance between any pair clusters.

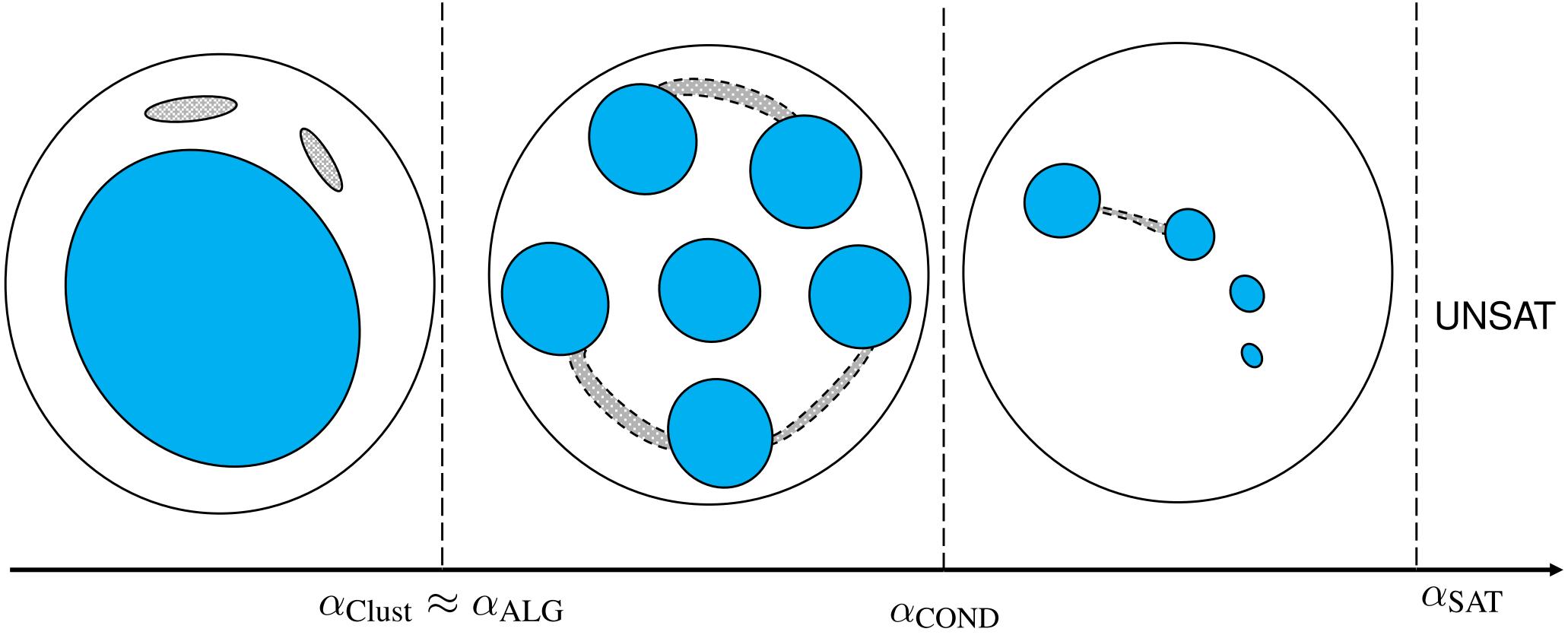


Landscape not exhibiting the OGP: diameter of one cluster is larger than distance between one pair of clusters.

from D. Gamarnik, PNAS 2018







from D. Gamarnik, PNAS 2018



What is a Stable Algorithm?

- **Definition:** it refers to an algorithm whose output does not change significantly when small, random perturbations are made to its input. The algorithm is *robust* to minor changes or noise in the input, meaning its performance or result is not highly sensitive to such perturbations.

Examples:

- Gradient Descent
- SVMs
- PCS
- Low degree polynomials (hence message-passing)
- QAOA

. . .

• Convex Optimization Algorithms (Interior-Point Methods or Simplex Method)



Post-Quantum Cryptography (PQC)

PQC: Cryptographic algorithms designed to be secure against attacks by quantum computers. Quantum computers can break widely-used cryptographic systems (e.g., RSA, ECC) by leveraging

algorithms like Shor's algorithm.

- cryptosystem.
- computers.

• Lattice-based Cryptography: Utilizes hard problems in lattice structures, such as Learning With Errors (LWE).

• Code-based Cryptography: Based on the difficulty of decoding random linear codes, such as McEliece

Hash-based Cryptography: Relies on the security of hash functions, used for digital signatures.

• Multivariate Cryptography: Solves systems of multivariate polynomial equations, considered hard for quantum





Ajtai's Function and high-dimensional Lattice problems

Lattice Problem: Consider a lattice \mathscr{L} defined as $\mathscr{L} = \{B \cdot z : z \in \mathbb{Z}^n\}$, where B is a (random) basis matrix. Ajtai's function is related to the shortest vector problem (SVP) in this lattice.

Ajtai's Theorem:

Worst-Case to Average-Case Reduction: Ajtai showed that there exists an algorithmic function f(x) that maps a random instance of a lattice problem (like SVP) to a generic solution in polynomial time if and only if the corresponding worst-case problem is solvable in polynomial time.

Hardness Guarantee: The function f(x) in Ajtai's theorem ensures that if an efficient algorithm solves the average case of this function, then it can also solve the hardest instances of the problem.

Generating Hard Instances of Lattice Problems Extended abstract M. Ajtai IBM Almaden Research Center

(1996)





Explicit case: Short Integer Solution (SIS) Problem

Find a short vector $x \in \mathbb{Z}^n$ such that $A \cdot x = 0 \mod q$ for a given random matrix A. This problem is hard under Ajtai's worst-case to average-case reduction.

Function Definition:

The SIS function could be formally written as

Observations:

- _ computers.
- not very efficient ...

 $f(A) = \{x : A \cdot x = 0 \mod q, \|x\| \text{ is small} \}.$

the hardness of these problems (and the functions derived from them) is believed to hold even against quantum



Remark:

Post-Quantum Encryption Standards

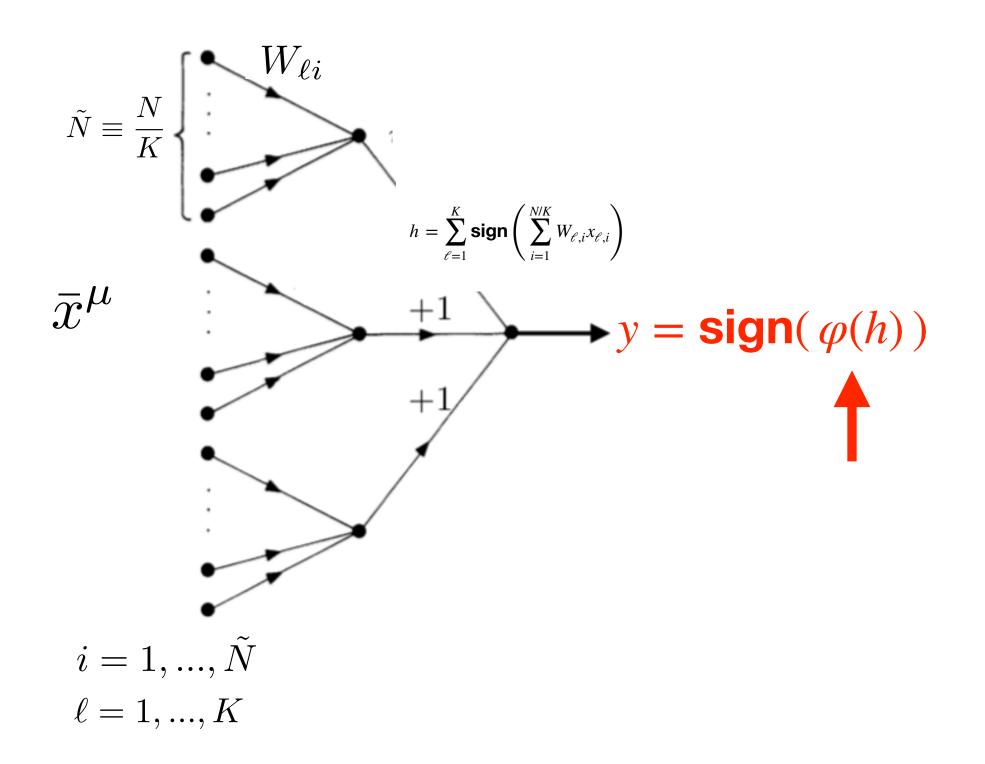
NIST PQC Standardization: The National Institute of Standards and Technology (NIST) has just approved a standard for PQC algorithms, with the first crypto system officially approved.

Open problem:

PQC algorithms may have larger key sizes and slower performance compared to traditional algorithms. There is a need for efficient crypto systems.

these rare events should not exist!

Collision problem



Non convex also for K=1 (perceptron) Find W such that $y_A(W) = y$ with

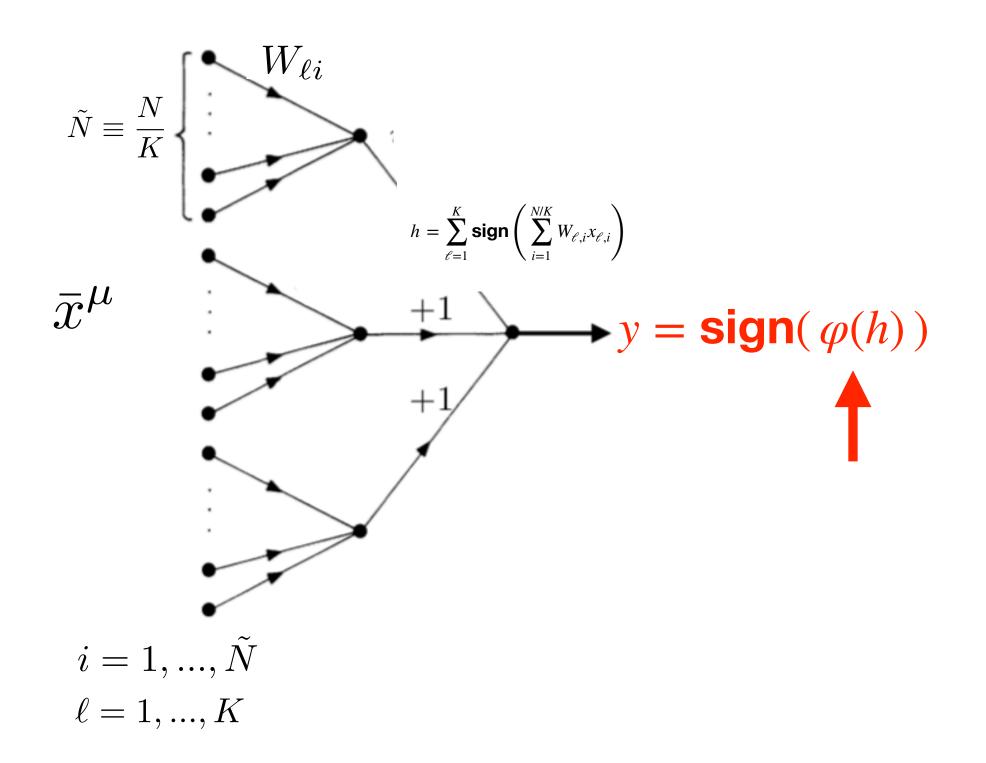
Simplest non convex neural device : 1-hidden layer, i.i.d. random associations

training set:
$$\{(\bar{x}^{\mu}, y^{\mu})\}\ \mu = 1, ..., P = \alpha N$$

 $x_{\ell i}^{\mu} = \pm 1 \ (i.i.d. \ p = 1/2)$
 $y^{\mu} = \pm 1 \ (i.i.d. \ p = 1/2)$
 $\mathbf{A} := [\mathbf{x}^1, \mathbf{x}^2, ..., \mathbf{x}^P]$
control parameter: $\alpha = \frac{\# \text{ patterns}}{\# \text{ weights}}$

h {
$$W_i = \pm 1$$
}, i.e. sign($\varphi(\sum_i W_i x_i^{\mu})) = y^{\mu}$





Non convex also for K=1 (perceptron) Find W such that $y_A(W) = y$ with

> We will need to generalize computational bounds we need for the collision problem

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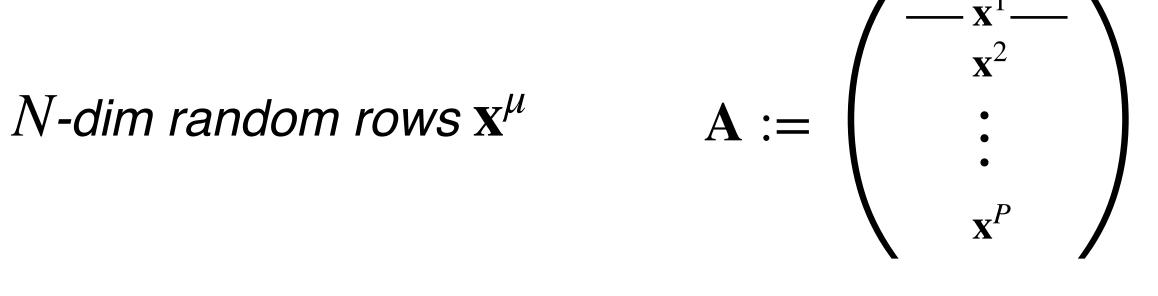
h {
$$W_i = \pm 1$$
}, i.e. sign($\varphi(\sum_i W_i x_i^{\mu})) = y^{\mu}$
this model through φ to obtain the



Computational challenges in non-convex NN:

$\mathbf{A} \in \mathbf{R}^{P \times N}$ random matrix composed by P N-dim random rows \mathbf{x}^{μ}

- •Inversion (learning): given disorder $\mathbf{A} \in \mathbf{R}^{P \times N}$ and labels $\mathbf{y} \in \{-1,1\}^P$, find any set of weights $W \in \{-1,1\}^N$ such that $y_{\mathbf{A}}(W) = \mathbf{y}$, assuming such W exists.
- •Teacher-student: given disorder $\mathbf{A} \in \mathbf{R}^{P \times N}$ and labels $\hat{\mathbf{y}} = y_{\mathbf{A}}(W) \in \{-1,1\}^{P}$ for uniformly sampled $W \in \{-1,1\}^{N}$, find any $W' \in \{-1,1\}^{N}$ such that $y_{\mathbf{A}}(W') = \hat{\mathbf{y}}$
- Collision finding: given disorder $\mathbf{A} \in \mathbf{R}^{P \times N}$, find any two $W \neq W' \in \{-1,1\}^N$ such that $y_{\mathbf{A}}(W) = y_{\mathbf{A}}(W')$ (unexplored so far).



Collision finding:

distinct $W \neq W'$ such that $y_A(W) = y_A(W')$.

Collision Resistant Hash Functions

time algorithm $A(\cdot)$ and any constant c > 0, it holds that,

$$\Pr_{A,h\in_R\mathscr{H}} \left[h(x) = h(y) \land x \neq y \right]$$

where the randomness is taken over a uniform random choice of h, and the random coins used by A.

The input is simply the function y_{A} itself, and the problem is to find a collision, defined as a pair of

Def.: A hash function family $\mathcal{H} = \{h : X \to Y\}$ is said to be *collision resistant*, if for any polynomial-

$$(x, y) \leftarrow A = o(n^{-c}),$$

The Generalised Binary Perceptron model(s)

where
$$h \equiv \frac{1}{\sqrt{N}} \sum_{i} w_i x_i$$
, w_i are binary va

In order to fit our model to a set of random inputs x^{μ} and labels y^{μ} , we need to impose that the stability Δ^{μ} is larger than zero

$$\Delta^{\mu}(\boldsymbol{w}) \equiv y^{\mu}\varphi\left(-\frac{1}{2}\right)$$

$\hat{y} = \operatorname{sign}(\varphi(h))$

riables, and \boldsymbol{x} is a N-dimensional pattern

 $\left(\frac{1}{\sqrt{N}}\sum_{i}^{}w_{i}x_{i}^{\mu}\right)\geq0\,,$

for any $\mu \in [P]$.

Construction of a Hash Function Using a random NN

 $y: \mathbf{B}^N \to \mathbf{B}^P$

Random Function Generation: Generate a set of $P = \alpha N$ random patterns

Input: Consider an input vector W from the space of possible inputs (e.g., a message or file).

Hash Function: The hash function based on the GRW model can be defined as follows.

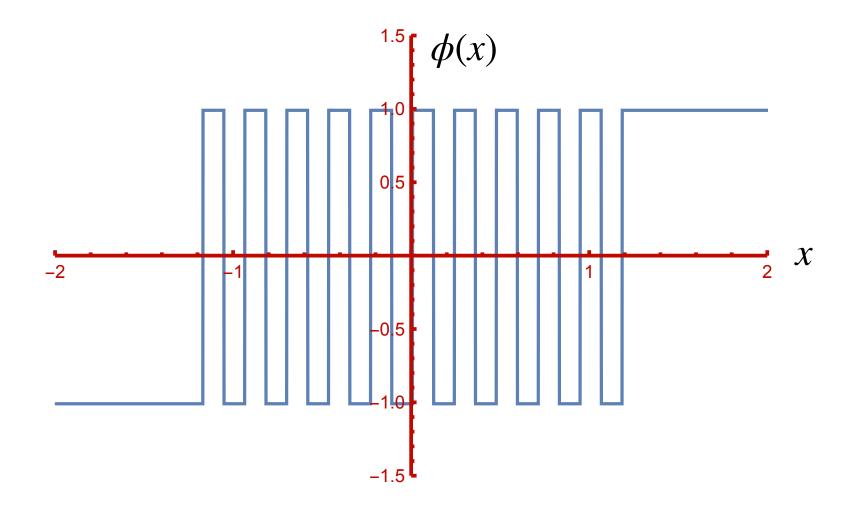
$$y_{1}(x) = \operatorname{sign}\left(\varphi\left(\frac{1}{\sqrt{N}}\sum_{i}w_{i}x_{i}^{(1)}\right)\right)$$
$$y_{2}(x) = \operatorname{sign}\left(\varphi\left(\frac{1}{\sqrt{N}}\sum_{i}w_{i}x_{i}^{(2)}\right)\right)$$
$$\vdots$$
$$\vdots$$
$$y_{P}(x) = \operatorname{sign}\left(\varphi\left(\frac{1}{\sqrt{N}}\sum_{i}w_{i}x_{i}^{(P)}\right)\right)$$

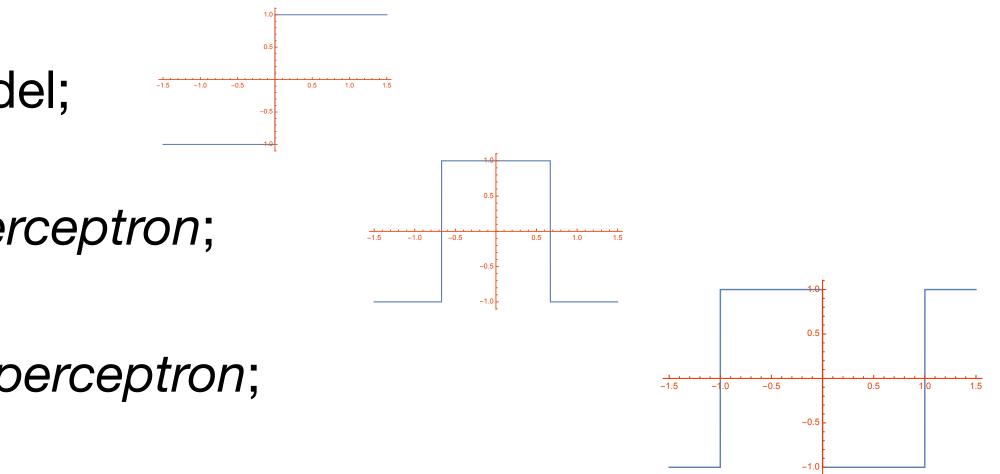
Some relevant examples of non-linearities

- $\varphi(h) = h$ standard binary perceptron model;
- $\varphi(h) = \kappa |h|$ and $y^{\mu} = 1$, symmetric perceptron;
- $\varphi(h) = (h \gamma)h(h + \gamma)$ reversed wedge perceptron;

•
$$\varphi(h) = \prod_{l=-K}^{K} \left(h + \frac{l\gamma}{K} \right)$$

generalizatized reverse wedge perceptron, with K oscillations in $[-\gamma, \gamma]$.





Indicator function $\mathbb{X}_{x}(w;\kappa) \equiv \prod^{r} \Theta(\Delta^{\mu}(w)),$ $\mu = 1$ $x_i^{\mu} \sim \mathcal{N}(0,1)$ and $\alpha \equiv \frac{P}{N}$.



The problem of Collisions

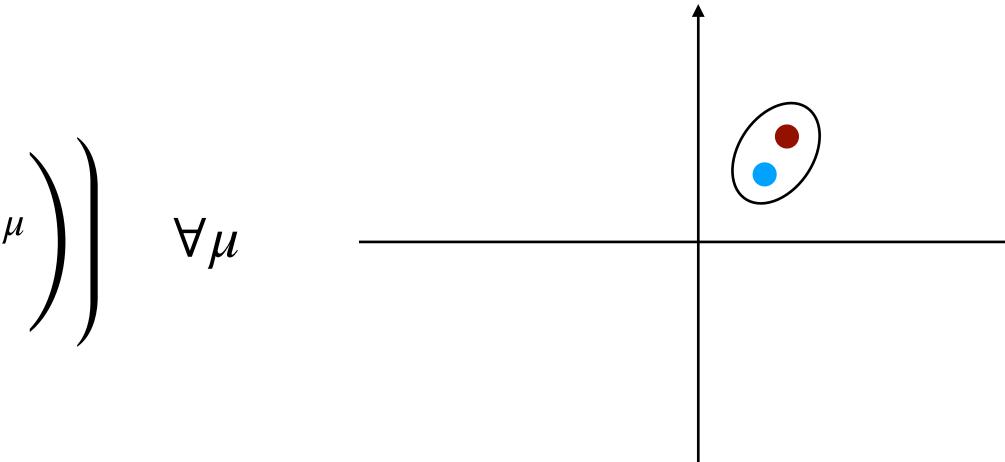
Given
$$x^{\mu}$$
, $\mu = 1, ..., P$, and $x_i^{\mu} \sim \mathcal{N}(0,1)$, find $w^{(1)}$, $w^{(2)}$ s.t.

$$\operatorname{sgn}\left(\varphi\left(\frac{1}{\sqrt{N}}w^{(1)}\cdot x^{\mu}\right)\right) = \operatorname{sgn}\left(\varphi\left(\frac{1}{\sqrt{N}}w^{(2)}\cdot x^{\mu}\right)\right) \quad \forall \mu$$

Indicator function $X_x(c)$:

$$\mathbb{X}_{\boldsymbol{x}}(\boldsymbol{c}) \equiv \prod_{\mu=1}^{P} \sum_{y^{\mu}} \Theta\left(y^{\mu}\varphi\left(\frac{1}{\sqrt{N}}\boldsymbol{w}^{(1)} \cdot \boldsymbol{x}^{\mu}\right)\right) \Theta\left(y^{\mu}\varphi\left(\frac{1}{\sqrt{N}}\boldsymbol{w}^{(2)} \cdot \boldsymbol{x}^{\mu}\right)\right) \qquad \text{with } \boldsymbol{c} \equiv (\boldsymbol{w}^{(1)}, \boldsymbol{w}^{(2)})$$

Partition function of collisions is $Z_x \equiv \int d\mathbf{c} \, \mathbb{X}_x(\mathbf{c}; \kappa_c)$ where $d\mathbf{c} \equiv dw^{(1)} dw^{(2)}$.



Geometric landscape collisions

Local entropy of collisions $\widetilde{w}_1, \widetilde{w}_2$:

$$\ln \mathcal{N}_{\boldsymbol{\xi}}(\widetilde{\boldsymbol{w}}_1, \widetilde{\boldsymbol{w}}_2; d) \equiv \ln \left[d\boldsymbol{w}_1 d\boldsymbol{w}_2 \, \mathbb{X}_{\boldsymbol{\xi}}(\boldsymbol{w}_1, \boldsymbol{w}_2) \, \delta\left(d\left[(\widetilde{\boldsymbol{w}}^{(1)}, \widetilde{\boldsymbol{w}}^{(2)}), (\boldsymbol{w}^{(1)}, \boldsymbol{w}^{(2)}) \right] - d \right) \right]$$

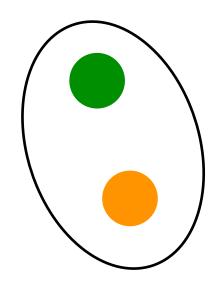
where $d\left[(\widetilde{w}^{(1)}, \widetilde{w}^{(2)}), (w^{(1)}, w^{(2)}]\right]$ is a permutation invariant distance between two collisions.

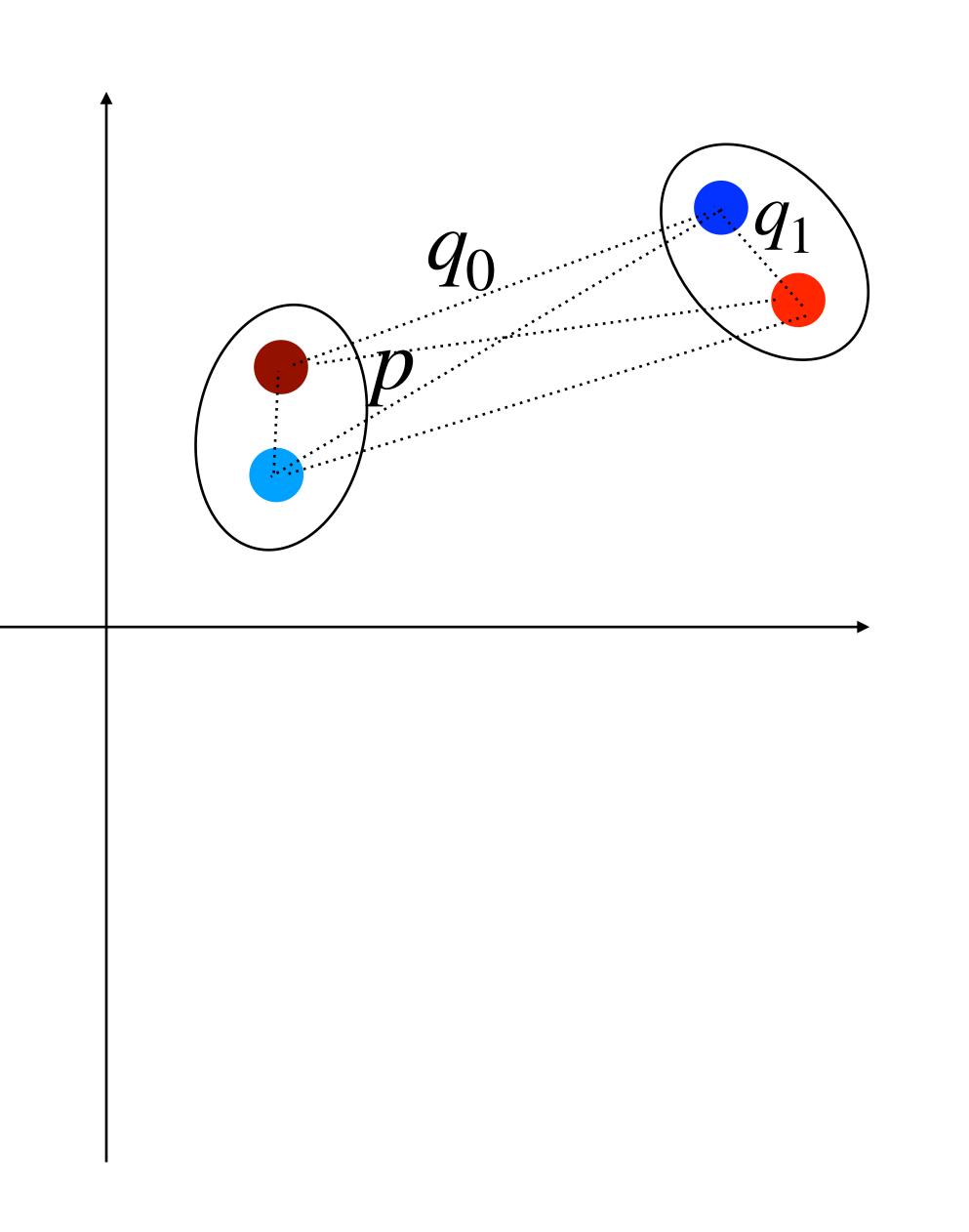
Consider
$$c_a = (w_a^{(1)}, w_a^{(2)})$$
 and $c_b = (w_b^{(1)}, w_b^{(2)})$ with $a \neq b$.

$$d\left(\boldsymbol{c}_{a},\boldsymbol{c}_{b}\right) = \min_{\boldsymbol{\pi}\in\mathcal{S}_{2}}\frac{1}{2}\sum_{s=1}^{2}d\left(\boldsymbol{w}_{a}^{(s)} - \boldsymbol{w}_{b}^{\boldsymbol{\pi}(s)}\right) = \min_{\boldsymbol{\pi}\in\mathcal{S}_{2}}\frac{1}{2}\sum_{s=1}^{2}\frac{1}{4N}\sum_{i=1}^{N}\left(\boldsymbol{w}_{ai}^{(s)} - \boldsymbol{w}_{bi}^{\boldsymbol{\pi}(s)}\right)^{2} \quad \text{with} \\ = \min_{\boldsymbol{\pi}\in\mathcal{S}_{2}}\frac{1}{2}\sum_{s=1}^{2}\frac{1}{2}\left(1 - \frac{1}{N}\boldsymbol{w}_{a}^{(s)} \cdot \boldsymbol{w}_{b}^{\boldsymbol{\pi}(s)}\right) = \frac{1}{4}\max_{\boldsymbol{\pi}\in\mathcal{S}_{2}}\sum_{s=1}^{2}\left(1 - q_{s\boldsymbol{\pi}(s)}^{ab}\right) = \frac{1}{2}(1 - p)$$

(thanks to the symmetry, we have that the overlap $q_{1\pi(1)}^{ab} = q_{2\pi(2)}^{ab}$, and we can choose π to be the identity)

n p the overlap on the diagonal he overlap matrix q_{st}^{ab}





Free entropy $\phi_{v}(d)$ in the annealed approximation, i.e.

$$\phi_y(d) \leq \phi_y^A(d)$$

Since $\mathcal{N}_{v}(d; x)$ is a non-negative and integer valued random variable, by Markov inequality we get

 $P(\mathcal{N}_{v}(d; \mathbf{x}) > 0)$

 $\alpha \geq \alpha_c^{UB}(d) \Rightarrow$ no collisions at a fixed distance d to one another.

Note: that $\alpha_c^{UB}(d)$ is only an upper bound to the true value (i.e. it might be that the true α_c is lower than that).

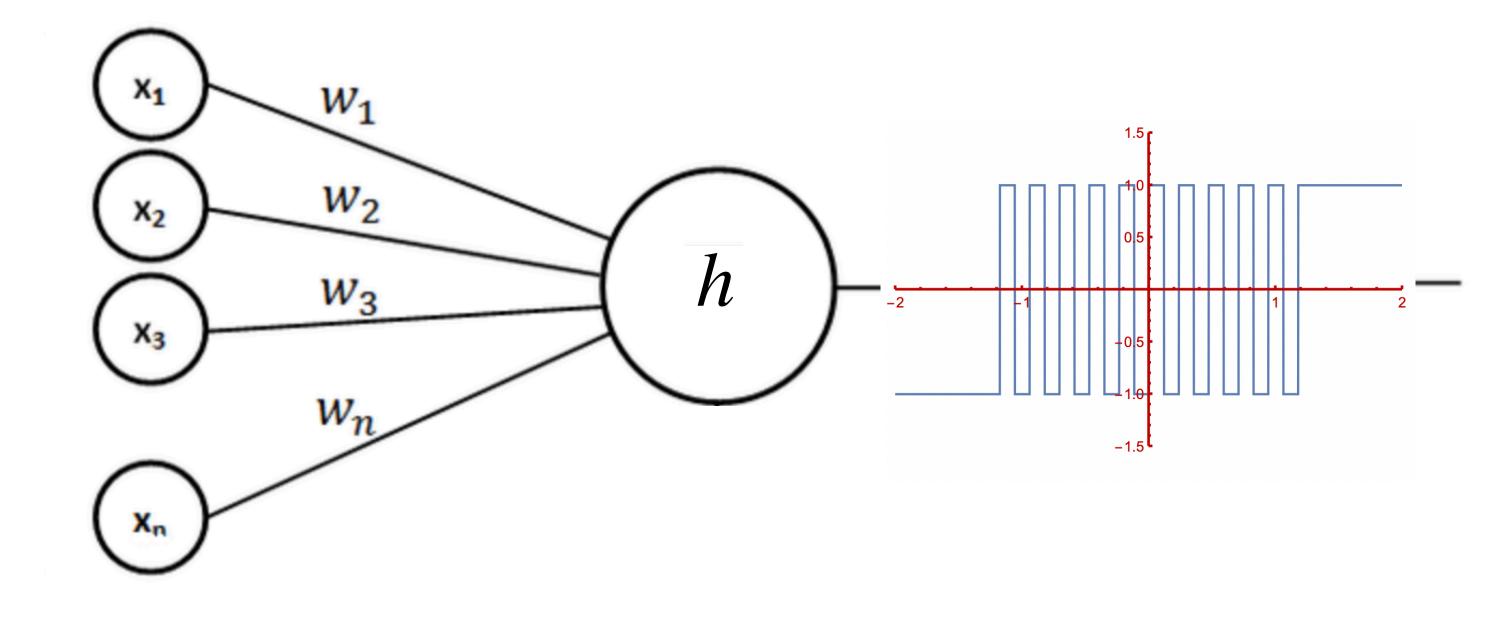
$$= \lim_{yN \to \infty} \frac{1}{yN} \ln \mathbb{E}_{x} \mathcal{N}_{y}(d;x)$$

$$\leq \mathbb{E}_{x} \mathcal{N}_{y}(d; x) = e^{y N \phi_{y}^{A}(d)}$$

If $\phi_v^A(d) \leq 0$ for $\alpha \geq \alpha_c^{UB}(d)$ then $P(\mathcal{N}_v(d; \mathbf{x}) > 0) \to 0$ for large N.

Generalised Reverse Wedge model

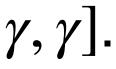
$$\varphi(h) = \prod_{l=-K}^{K} \left(h + \frac{l\gamma}{K} \right)$$



In the large K limit, the computation simplifie

generalizalized reverse wedge perceptron, with K oscillations in $[-\gamma, \gamma]$.

es:
$$\lim_{N \to \infty} \frac{K}{N} \to 0$$
, and next $K \gg 1$



Statistical physics of highly non-convex random systems and Crypto are very close

Spin glass theory used for crypto-systems design

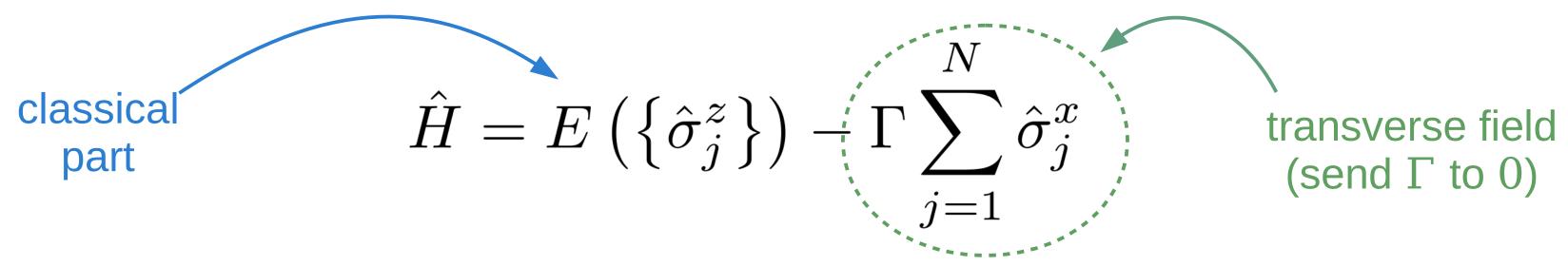
These are just first steps, we conjecture we can prove CRH w.r.t. stable algorithms

Marco Benedetti, Andrej Bogdanov, Enrico Malatesta, Marc Mezard, Gianmarco Perrupato, Alon Rosen, Nikolaj I. Schwartzbach , and Riccardo Zecchina

Conclusions

Quantum Annealing for non convex learning devices

- Quantum annealing strategy: use quantum fluctuations (rather than thermal fluctuations) to overcome energetic barriers
 - Classical energy function + quantum perturbation, slowly send the ____ perturbation to zero



• Thus far: unclear if "true" QA really helps, compared to standard annealing, in any relevant concrete scenario

QA detour



QA: Suzuki-Trotter transformation

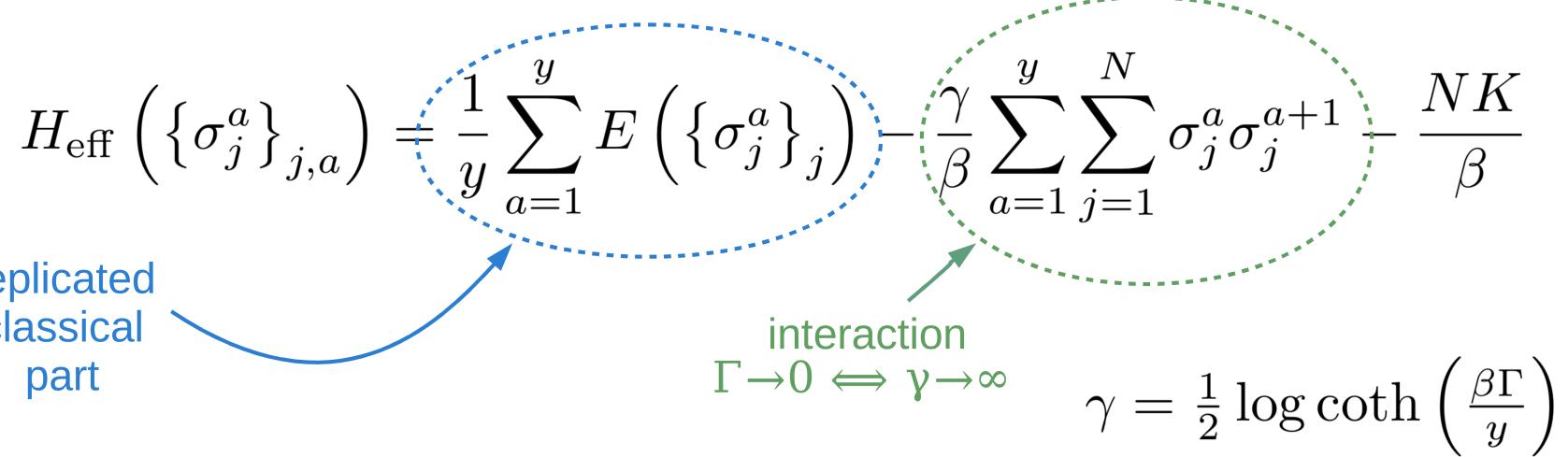
Hamiltonian (with infinite replicas, $y \rightarrow \infty$)

replicated classical part

• Can be simulated with MCMC (finite $y) \rightarrow$ Quantum Simulated Annealing (QSA)

QA detour

• Partition function transformation \rightarrow "effective" replicated classical





Quantum annealing vs Robust ensemble

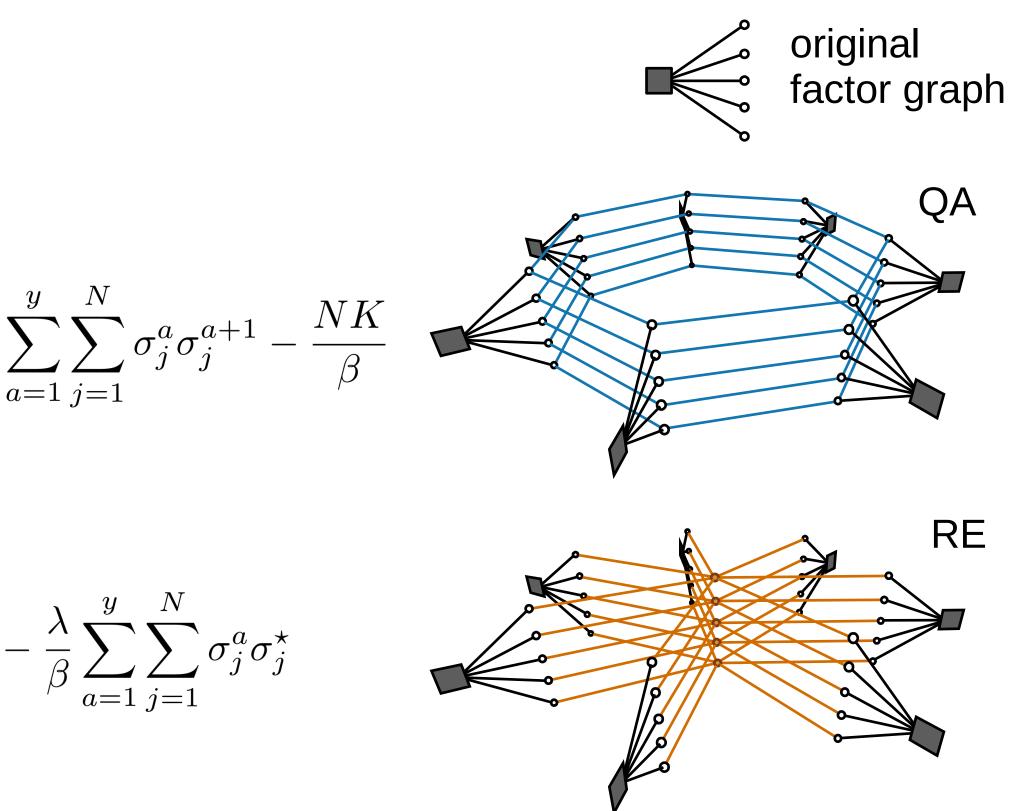
 \bullet to the robust ensemble description...

$$H_{\text{eff}}\left(\left\{\sigma_{j}^{a}\right\}_{j,a}\right) = \frac{1}{y}\sum_{a=1}^{y} E\left(\left\{\sigma_{j}^{a}\right\}_{j}\right) - \frac{\gamma}{\beta}\sum_{a=1}^{y}\sum_{j=1}^{N} \sum_{j=1}^{N} \left(\left\{\sigma_{j}^{a}\right\}_{j}\right) - \frac{\gamma}{\beta}\sum_{a=1}^{y}\sum_{j=1}^{N} \sum_{j=1}^{N} \sum_{j=1}^{N} \left(\left\{\sigma_{j}^{a}\right\}_{j}\right) - \frac{\gamma}{\beta}\sum_{a=1}^{y}\sum_{j=1}^{N} \sum_{j=1}^{N} \sum_$$

$$H_{\text{eff}}^{\text{RE}}\left(\sigma^{\star}, \left\{\sigma_{j}^{a}\right\}_{j,a}\right) = \sum_{a=1}^{y} E\left(\left\{\sigma_{j}^{a}\right\}_{j}\right) - \frac{\lambda}{\beta} \sum_{a=1}^{y} E$$

QA detour

Effective Hamiltonian after Suzuki-Trotter transformation: very similar

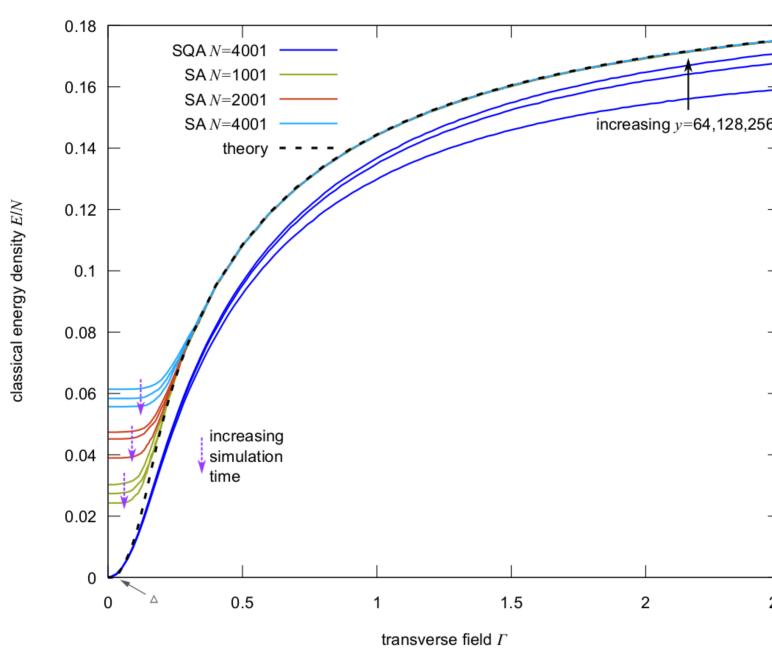


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QSA on binary neural networks study

- true QA in small instances
- annealing a physical device would work in $\sim O(1)...)$ (DWave-like)
- regions

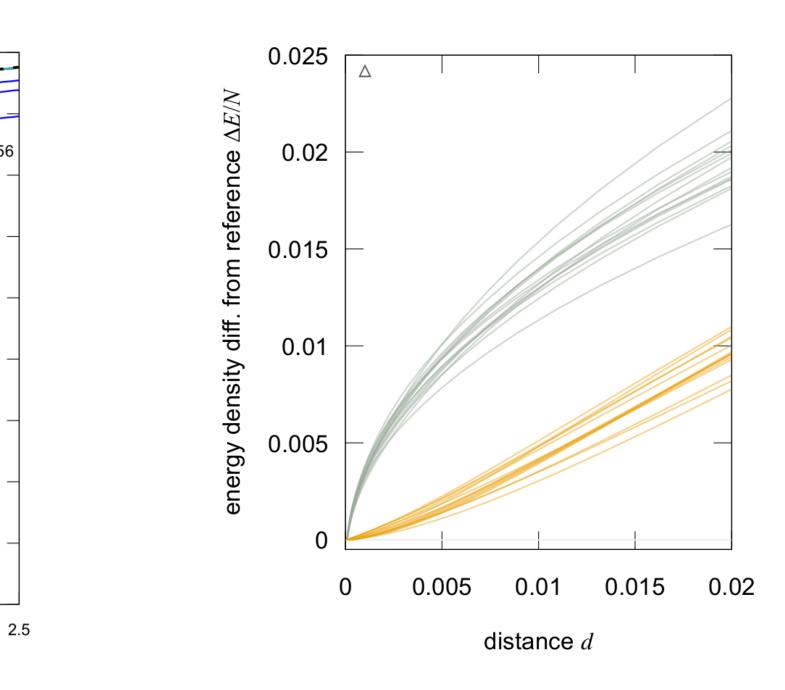


QA detour

Analytical calculations + numerical experiments + comparison with

Ends up in the dense states (exponential speed-up w.r.t. thermal

QA lowers kinetic energy by delocalizing \rightarrow favors dense



C. Baldassi, R. Zecchina, PNAS 2018

