

Alexandria Quantum Computing Group (AleQCG)

CENTER OF EXCELLENCE FOR QUANTUM COMPUTERS

Basics of Quantum Computing Day 3

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Description

Quantum Communication Protocols

Outline

- ⚫ **Light Polarization**
- ⚫ **Quantum Dense Coding**
- ⚫ **Quantum Teleportation**

Light Polarization

Optical Fiber

Photon polarization as a qubit

Quantum Dense Coding -1

Is used to send two classical bits using a single quantum channel.

Sender

Alice. Alice receives two classical bits, encoding the numbers 0 through 3. Depending on this number Alice performs one of the transformations $\{I, X, Y, Z\}$ on her qubit of the entangled pair ψ_0 . Transforming just one bit of an entangled pair means performing the identity transformation on the other bit. The resulting state is shown in the table.

Alice then sends her qubit to Bob.

Receiver

Bob. Bob applies a controlled-NOT to the two qubits of the entangled pair.

Bob now applies H to the first bit:

Quantum Teleportation - 1

In used to teleport an unknown quantum state using two classical channels.

Quantum Teleportation -2

Quantum Teleportation - 3

Alice. Alice has a qubit whose state she doesn't know. She wants to send the state of ths qubit

 $\phi = a|0\rangle + b|1\rangle$

to Bob through classical channels. As with dense coding, Alice and Bob each possess one qubit of an entangled pair

$$
\psi_0 = \frac{1}{\sqrt{2}}(|00\rangle + |11\rangle).
$$

Alice applies the decoding step of dense coding to the qubit ϕ to be transmitted and her half of the entangled pair. The starting state is quantum state

$$
\phi \otimes \psi_0 = \frac{1}{\sqrt{2}} (a|0\rangle \otimes (|00\rangle + |11\rangle) + b|1\rangle \otimes (|00\rangle + |11\rangle))
$$

=
$$
\frac{1}{\sqrt{2}} (a|000\rangle + a|011\rangle + b|100\rangle + b|111\rangle),
$$

of which Alice controls the first two bits and Bob controls the last one. Alice now applies $C_{not} \otimes I$ and $H \otimes I \otimes I$ to this state:

$$
(H \otimes I \otimes I)(C_{not} \otimes I)(\phi \otimes \psi_0)
$$

= $(H \otimes I \otimes I)(C_{not} \otimes I)\frac{1}{\sqrt{2}}(a|000\rangle + a|011\rangle + b|100\rangle + b|111\rangle)$
= $(H \otimes I \otimes I)\frac{1}{\sqrt{2}}(a|000\rangle + a|011\rangle + b|110\rangle + b|101\rangle)$
= $\frac{1}{2}(a(|000\rangle + |011\rangle + |100\rangle + |111\rangle) + b(|010\rangle + |001\rangle - |110\rangle - |101\rangle))$
= $\frac{1}{2}(|00\rangle(a|0\rangle + b|1\rangle) + |01\rangle(a|1\rangle + b|0\rangle) + |10\rangle(a|0\rangle - b|1\rangle) + |11\rangle(a|1\rangle - b|0\rangle))$

Alice measures the first two qubits to get one of $|00\rangle$, $|01\rangle$, $|10\rangle$, or $|11\rangle$ with equal probability. Depending on the result of the measurement, the quantum state of Bob's qubit is projected to $a|0\rangle + b|1\rangle$, $a|1\rangle + b|0\rangle$, $a|0\rangle - b|1\rangle$, or $a|1\rangle - b|0\rangle$ respectively. Alice sends the result of her measurement as two classical bits to Bob.

Note that when she measured it, Alice irretrievably altered the state of her original qubit ϕ , whose state she is in the process of sending to Bob. This loss of the original state is the reason teleportation does not violate the no cloning principle.

Bob. When Bob receives the two classical bits from Alice he knows how the state of his half of the entangled pair compares to the original state of Alice's qubit.

Quantum Algorithms

Outline

- ⚫ Quantum Parallelism
- ⚫ Superposition Preparation
- Parallel Evaluation of a function
- Marking Solutions
- ⚫ Grover's Quantum Search Algorithm

Quantum Algorithms

⚫What class of computations can be performed using quantum circuits?

● How does that class compare with the computations which can be performed using classical logical circuits?

■ Can we find a task which a quantum computer may perform better than a classical computer?

Quantum Parallelism

- ⚫ *Quantum parallelism* is a fundamental feature of many quantum algorithms.
- ⚫ Allows quantum computers to evaluate a function *f*(*x*) for many *different* values of *x* simultaneously.

Superposition Preparation

- ⚫ This procedure can easily be generalized to functions on an arbitrary number of bits, by using a general operation known as the *Hadamard transform*, or sometimes the *Walsh– Hadamard transform*.
- For $n = 2$, $H^{\otimes 2} |00\rangle$

$$
\begin{aligned}\n\left|0\right\rangle &\xrightarrow{\hspace{1em}} \left(\frac{|0\rangle + |1\rangle}{\sqrt{2}}\right) \left(\frac{\left|0\right\rangle + |1\rangle}{\sqrt{2}}\right) \\
\left|0\right\rangle &\xrightarrow{\hspace{1em}} \left(\frac{|0\rangle + |1\rangle}{\sqrt{2}}\right) \\
&= \frac{|00\rangle + |01\rangle + |10\rangle + |11\rangle}{2}\n\end{aligned}
$$

Superposition Preparation

• For $n = 3$, $H^{\otimes 3} |000\rangle$ $\frac{-\boxed{H}}{-\boxed{H}} = \frac{\left(\ket{0}+\ket{1}}{\sqrt{2}}\right)\left(\frac{\ket{0}+\ket{1}}{\sqrt{2}}\right)\left(\frac{\ket{0}+\ket{1}}{\sqrt{2}}\right)$ $\frac{1}{|H|} = \frac{1}{2} \left| \frac{1}{2} \left(\frac{1}{2} \left| 0 \right| 0 \right) + \left| 0 \right| 0 \right| + \left| 0 \right| 0 \right| + \left| 1 \right| 0 \right| + \left| 1 \right| 0 \right| + \left| 1 \right| 1 \right|$ $2\sqrt{2}$

● For arbitrary n>0

Boolean Quantum Circuits $=\overline{x_0}x_1+x_0x_2$

 x_0

 $\overline{0}$

 x_I

 x_2

Quantum circuit

Parallel Evaluation of *f*

Tracing the Circuit

$$
\psi_0 \rangle = |000\rangle \otimes |0\rangle
$$

\n
$$
\psi_1 \rangle = \left(\frac{|0\rangle + |1\rangle}{\sqrt{2}}\right) \left(\frac{|0\rangle + |1\rangle}{\sqrt{2}}\right) \left(\frac{|0\rangle + |1\rangle}{\sqrt{2}}\right) \otimes |0\rangle
$$

\n
$$
= \left(\frac{|000\rangle + |001\rangle + |010\rangle + |011\rangle + |100\rangle + |101\rangle + |110\rangle + |111\rangle}{2\sqrt{2}}\right) \otimes |0\rangle
$$

\n
$$
= \left(\frac{|000,0\rangle + |001,0\rangle + |010,0\rangle + |011,0\rangle + |100,0\rangle + |101,0\rangle + |110,0\rangle + |111,0\rangle}{2\sqrt{2}}\right)
$$

\n
$$
\psi_2 \rangle = \left(\frac{|000,0\rangle + |001,0\rangle + |010,1\rangle + |011,1\rangle + |100,0\rangle + |101,1\rangle + |110,0\rangle + |111,1\rangle}{2\sqrt{2}}\right)
$$

\n
$$
= \frac{1}{2\sqrt{2}}(|000\rangle + |001\rangle + |100\rangle + |110\rangle) \otimes |0\rangle + \frac{1}{2\sqrt{2}}(|010\rangle + |011\rangle + |101\rangle + |111\rangle) \otimes |1\rangle
$$

Effect of Entanglement

• If we measure the extra qubit and finds |0>, then the system collapses to

$$
\frac{1}{2}(|000\rangle+|001\rangle+|100\rangle+|110\rangle)|\otimes|0\rangle
$$

• If we measure the extra qubit and finds |1>, then the system collapses to

$$
\frac{1}{2}(|010\rangle+|011\rangle+|101\rangle+|111\rangle)\otimes|1\rangle
$$

General Form (Marking by Entanglement)

Entanglement

Marking the Solutions by Phase Shift

$$
|\psi_0\rangle = |000\rangle \otimes |1\rangle
$$

\n
$$
|\psi_1\rangle = \left(\frac{|0\rangle + |1\rangle}{\sqrt{2}}\right) \left(\frac{|0\rangle + |1\rangle}{\sqrt{2}}\right) \left(\frac{|0\rangle + |1\rangle}{\sqrt{2}}\right) \otimes \left(\frac{|0\rangle - |1\rangle}{\sqrt{2}}\right)
$$

\n
$$
= \left(\frac{|000\rangle + |001\rangle + |010\rangle + |011\rangle + |100\rangle + |101\rangle + |110\rangle + |111\rangle}{2\sqrt{2}}\right) \otimes \left(\frac{|0\rangle - |1\rangle}{\sqrt{2}}\right)
$$

\n
$$
|\psi_2\rangle = \left(\frac{|000\rangle + |001\rangle - |010\rangle - |011\rangle + |100\rangle - |101\rangle + |110\rangle - |111\rangle}{2\sqrt{2}}\right) \otimes \left(\frac{|0\rangle - |1\rangle}{\sqrt{2}}\right)
$$

\n
$$
= \left(\frac{1}{2\sqrt{2}}\left(|000\rangle + |001\rangle + |100\rangle + |110\rangle\right) - \frac{1}{2\sqrt{2}}\left(|010\rangle + |011\rangle + |101\rangle + |111\rangle\right)\right) \otimes \left(\frac{|0\rangle - |1\rangle}{\sqrt{2}}\right)
$$

General Form (Marking by Phase Shift)

$$
|\psi_0\rangle = |0\rangle^{\otimes n}
$$

\n
$$
|\psi_1\rangle = H^{\otimes n} |\psi_0\rangle
$$

\n
$$
= \frac{1}{\sqrt{2^n}} \sum_{x=0}^{2^n - 1} |x\rangle
$$

\n
$$
|\psi_2\rangle = U_f |\psi_1\rangle
$$

\n
$$
= \frac{1}{\sqrt{2^n}} \sum_{x=0}^{2^n - 1} (-1)^{f(x)} |x\rangle
$$

Grover's Quantum Search Algorithm

- Given a List L of N=2ⁿ items
	- O Step 1 Prepare a superposition on N items on O(log N)
	- **O** Step 2 Iterate the Amplitude Amplification for $O(\sqrt{N})$
	- ◯ Step 3- Measure the quantum register
- Classical Computers require O(N) iteration.

Grover diffusion operator

Repeat $O(\sqrt{N})$ times

Average of all Amplitudes Flip all Amplitudes around Avg

Algorithm: Quantum search

1- Register Preparation. $|W_0^{(G,1)}\rangle = |0\rangle^{\otimes n} \otimes |1\rangle$.

$$
2\text{-} Register Initialization. } \left| W_1^{(G,1)} \right\rangle = H^{\otimes n+1} \left| W_0^{(G,1)} \right\rangle = \frac{1}{\sqrt{N}} \sum_{i=0}^{N-1} |i\rangle \otimes \left(\frac{|0\rangle - |1\rangle}{\sqrt{2}} \right).
$$

3- Applying the Oracle and Changing Sign. $U_f|i\rangle \rightarrow (-1)^{f(i)}|i\rangle$, so that,

$$
\left|W_2^{(G,1)}\right\rangle = U_f \left|W_1^{(G,1)}\right\rangle = \frac{1}{\sqrt{N}} \sum_{i=0}^{N-1} |i\rangle \otimes \left(\frac{\left|0 \oplus f(i)\right\rangle - \left|1 \oplus f(i)\right\rangle}{\sqrt{2}}\right).
$$

$$
\left|W_2^{(G,1)}\right\rangle = \frac{1}{\sqrt{N}} \sum_{i=0}^{N-1} \binom{n}{i} \otimes \left(\frac{\left|0\right\rangle - \left|1\right\rangle}{\sqrt{2}}\right) - \frac{1}{\sqrt{N}} \sum_{i=0}^{N-1} \binom{n}{i} \otimes \left(\frac{\left|0\right\rangle - \left|1\right\rangle}{\sqrt{2}}\right).
$$

4- Inversion about the Mean. $G = H^{\otimes n} (2 |0\rangle \langle 0| - I_n) H^{\otimes n}$,

$$
\begin{aligned} \left| \psi \right\rangle &= \sum_{j=0}^{N-1} \alpha_j \left| j \right\rangle. \\ G \left| \psi \right\rangle &= \sum_{j=0}^{N-1} \left[-\alpha_j + 2 \left\langle \alpha \right\rangle \right] \left| j \right\rangle, \end{aligned}
$$

Algorithm: Quantum search

Iterating the Algorithm

$$
a_1^G = \frac{N-2M}{N}a_0^G + \frac{2(N-M)}{N}b_0^G, \ \ b_1^G = \frac{N-2M}{N}b_0^G - \frac{2M}{N}a_0^G.
$$

The system after $q_G \geq 1$ iterations can be written as follows,

$$
\left|W^{(G,q_G)}\right\rangle = b_q^G \sum_{i=0}^{N-1} \binom{n}{i} + a_q^G \sum_{i=0}^{N-1} \binom{n}{i},
$$

such that,

$$
M(a_q^G)^2 + (N-M)(b_q^G)^2 = 1,
$$

$$
\begin{array}{l} a_0^G=b_0^G=\frac{1}{\sqrt{N}},\\ a_q^G=\frac{N-2M}{N}a_{q-1}^G+\frac{2(N-M)}{N}b_{q-1}^G,~~b_q^G=\frac{N-2M}{N}b_{q-1}^G-\frac{2M}{N}a_{q-1}^G. \end{array}
$$

Solving these recurrence relations, the closed forms can be written as follows :

$$
a_q^G = \frac{1}{\sqrt{M}} \sin ((2q_G + 1) \theta_G), \ \ b_q^G = \frac{1}{\sqrt{N - M}} \cos ((2q_G + 1) \theta_G),
$$

where $\sin^2(\theta_G) = M/N$ and $0 < \theta_G \le \pi/2$.

Thank you