

# Lectures on Quantum Computing

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Kingdom of Morocco  
Ministry of Higher Education,  
Scientific Research and Innovation

## THE EIGHTH BIENNIAL AFRICAN SCHOOL OF FUNDAMENTAL PHYSICS AND APPLICATIONS (ASP2024)



Co-organized by Cadi Ayyad University and Mohammed V University  
at Faculty of Science Semlalia, Marrakesh, Morocco

**April 15<sup>th</sup>–19<sup>th</sup> and July 7<sup>th</sup>–21<sup>st</sup>, 2024**

# Outline of the lectures

- 1 Lecture 1 : Introduction to Quantum Computing
- 2 Lecture 2 : Basics of Quantum Computing

# UNESCO International Year of Quantum Science and Technology IYQ2025



## INTERNATIONAL YEAR OF Quantum Science and Technology

100 years of quantum is just the beginning...

An international partnership of major scientific bodies and academies is preparing a resolution for the 2024 General Conference of the United Nations Educational, Scientific and Cultural Organization (UNESCO) and the 2024 General Assembly of the United Nations to proclaim 2025 the **International Year of Quantum Science and Technology**. This year-long initiative would celebrate the profound impacts of quantum science on technology, culture, and our understanding of the natural world.



# Classical Information : Shannon Theory

- Defined the quantity of information produced by a source by a formula similar to the equation that defines thermodynamic entropy in physics :



Claude Shannon  
1916-2001

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- Analyzed the ability to send information through a communications channel, proving the existence of a maximum transmission rate that could not be exceeded (bandwidth).
- Demonstrated mathematically that even in a noisy channel with a low bandwidth, essentially perfect, error-free communication could be achieved by keeping the transmission rate within the channel's bandwidth and by using error-correcting schemes (redundancy)

# Noiseless and noisy Shannon theorems

2

- Noiseless channel case:



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  - FRCN SCHL F PHSCS HS NTRSTNG  
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  - QUANTUM INFORMATION THEORY IS FUN

# Shannon Theory

For more details about Shannon Information Theory :

arXiv:1106.1445v8 [quant-ph] 14 Jul 2019

## From Classical to Quantum Shannon Theory

Mark M. Wilde

Hearne Institute for Theoretical Physics  
Department of Physics and Astronomy  
Center for Computation and Technology  
Louisiana State University  
Baton Rouge, Louisiana 70803, USA

July 16, 2019

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- The purpose of Quantum Information Theory is precisely to take advantage of these properties in order to perform tasks which are impossible to realize with classical computers.
- The most known applications of quantum information are [Quantum Computing](#), with the focus on the physical implementation of a universal quantum computer, and [Quantum Cryptography](#) for the secure transmission of information.

# Quantum computing vs classical computing



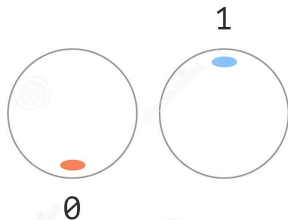
SpinQ Gemini Quantum  
computer



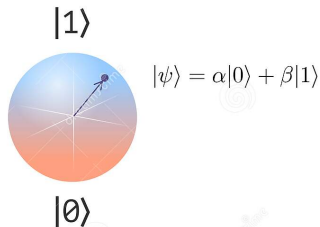
Laptop classical computer

# Classical bit vs Quantum bit

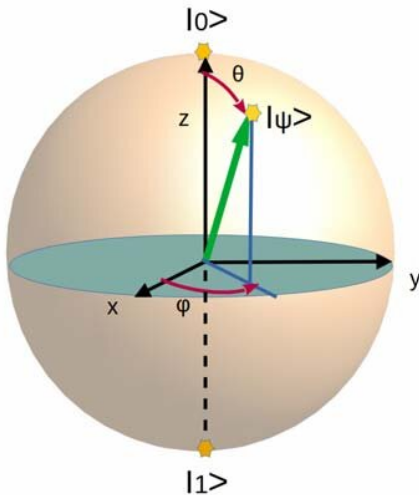
Bit



Qubit

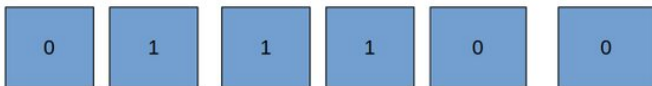


# Qubit representation: Bloch Sphere

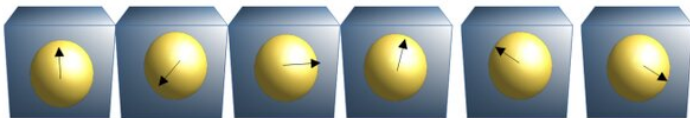


Credits : S. Simonović, Adv. Tech. and Mat. **46-2** , 24-31 (2021)

# Classical and quantum register



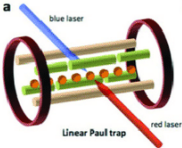
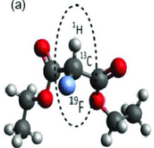
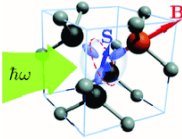
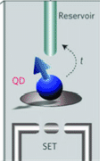
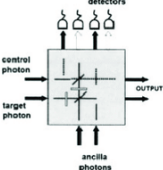
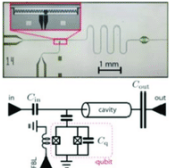
**CLASSICAL REGISTER – CAN CONTAIN ONLY ONE VARIATION OF 0 AND 1**



**QUANTUM REGISTER – CAN CONCURRENTLY CONTAIN ALL VARIATIONS OF 0 AND 1**

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# Qubit technologies

<p><b>Ion trap</b></p>  <p>Scientific Reports 4, 3589 (2014)</p>	<p><b>NMR</b></p>  <p>Sci. China Phys. Mech. Astron. 59:630302 (2016)</p>	<p><b>NV center</b></p>  <p>Phys. Rev. B 86, 125204 (2012)</p>
<p><b>Quantum dot</b></p>  <p>4 Nature Nanotechnology 9, 981–985 (2014)</p>	<p><b>Linear optical</b></p>  <p>J. Opt. Soc. Am. B, 24, 2, 209-213 (2007)</p>	<p><b>Superconducting</b></p>  <p>Ann. Phys. (Berlin) 525, 6, 395–412 (2013)</p>

...  
 many  
 more

Credit : Amundson and Sexton-Kennedy, EPJ Web of Conferences **214** 09010 (2019)

# Classical gates

## Logic Gate Symbols



OR



NOR



AND



NAND



XOR



XNOR



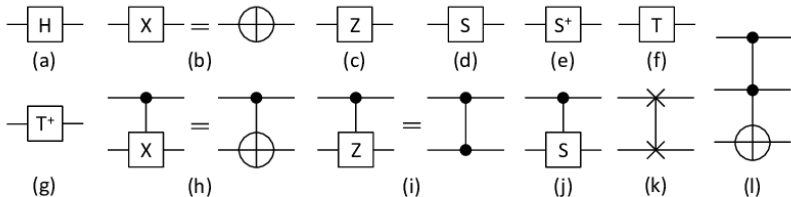
Buffer



NOT



# Quantum gates



# Classical computing

- Given a register of  $n$  bits (for example one byte), what operations can we do with it?

0	1	1	0	0	0	0	1
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|---|---|---|---|---|---|---|---|

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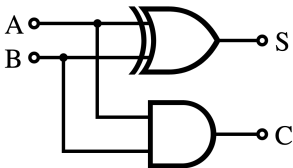
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 \end{array}$$

- Decimal:  $193+139=332$

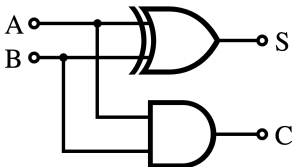
## Half-adder Circuit

- The half-adder circuit is built from an *XOR* gate and a *AND* gate (we will come back to this)



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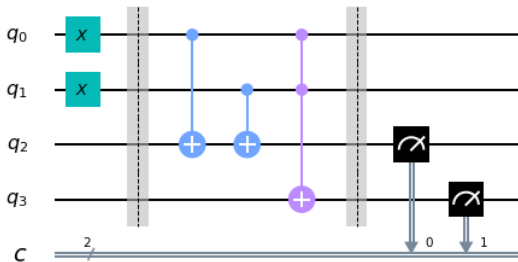


- It takes two bits  $A$  and  $B$  as input and delivers 2 outputs, the sum  $S$  and the carry  $C$ ,

bit 1	bit 2	sum	carry
0	0	0	0
0	1	1	0
1	0	1	0
1	1	0	1

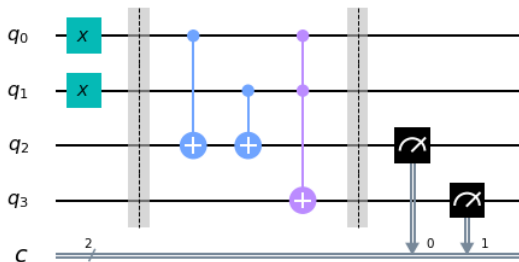
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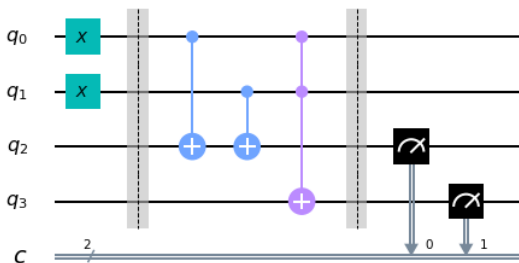
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- For now, it seems complicated, but at the end of this talk, this kind of circuit will become clear.
- But already know at this level that this circuit is formed by quantum gates  $X$ ,  $CNOT$  and  $Toffoli$ !
- Question** : What would be the result of  $(193 + 139)$  if we use a quantum full-adder ?

# Supercomputers

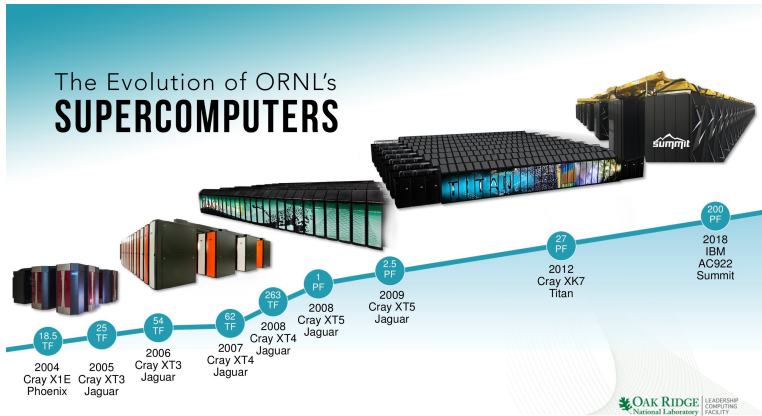


ENIAC 1945  
50 KFLOPS  
167 m<sup>2</sup>, 150 kW



IBM Summit 2018  
200 PFLOPS  
873 m<sup>2</sup>, 13 MW  
219 kms of cabling

# Classical Supercomputers





# Evolution of storage capacities



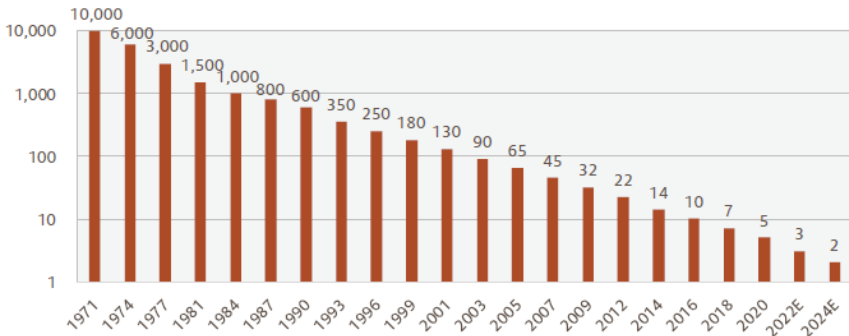
**5 MB hard disk drive - 1956**



**1 TB micro SD card - 2020**

# Limits of miniaturization: Moore's Law

**Figure 2: The incredible shrinking universe (device size in nm, log scale)**



Source: PC Magazine, Epoch Investment Partners

Note: For reference, most atoms are 0.1 to 0.5 nm in diameter

Epoch perspectives, 11 February 2021

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- Yes, to a certain extent.
- However, for some problems we reach a limit. For example, **factoring large numbers**, used to encrypt messages (RSA protocol) and ensure the security of communications and transactions related to e-commerce, electronic signatures, etc.

# Complexity classes

- To better understand the difficulty of the factoring problem, here are some examples of mathematical problems as well as the scale laws of the number of operations  $n$  with the number of bits (or digits) as well as the complexity classes:

Problem	Operations	Class
Addition of 2 numbers of $n$ bits	$n$	P
Multiplication of 2 numbers of $n$ bits	$n^2$	P
FFT de $n$ bits	$n \log(n)$	P
Factoring a number of $n$ bits	$2^{n/2}$	NP
Travelling salesman problem ( $n$ towns)	$e^{n \log(n)}$	NPC

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- Current computer architectures are unable to deal with complex problems due to a lack of efficient algorithms.

# Cryptography

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- The key used to encrypt is accompanied by a large integer, the product of two large primes kept secret (of the order of 200 digits, see RSA numbers). To calculate the decryption key, the only known method requires knowing the two prime factors.

# RSA Protocol

- The security of the RSA system is based on the fact that it is easy to find two large prime numbers  $p$  and  $q$  (using primality tests) and multiply them to have  $N = p \times q$ , but that it would be difficult for an attacker to find these two numbers  $(p, q)$  knowing  $N$ . This system also allows the creation of digital signatures, and has revolutionized the world of cryptography.

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- Example : calculate  $11 \times 17$  and  $137 \times 211$
- find  $(p, q)$ , for  $N = 667$  and  $N = 82919$
- Answers :  $(p, q) = (23, 29)$  and  $(p, q) = (283, 293)$

# RSA numbers

- RSA-200 : made up of 200 digits in decimal

27997833911221327870829467638722601621070446786955  
42853756000992932612840010760934567105295536085606  
18223519109513657886371059544820065767750985805576  
13579098734950144178863178946295187237869221823983



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- The prime factors of RSA200 are :

$$p = 35324619344027701212726049781984643686711974001976$$
$$25023649303468776121253679423200058547956528088349$$
$$q = 79258699544783330333470858414800596877379758573642$$
$$19960734330341455767872818152135381409304740185467$$

## RSA numbers

- The current record, dating from 2009, for the largest factored number (RSA-768): formed by 232 digits in decimal

12301866845301177551304949583849627207728535695953  
34792197322452151726400507263657518745202199786469  
38995647494277406384592519255732630345373154826850  
79170261221429134616704292143116022212404792747377  
94080665351419597459856902143413

- The calculation carried out on a network of computers required approximately two years of calculation, that is to say a CPU calculation time equivalent to **2000 years** on an Opetron processor running at 2.2GHz.

# RSA-2048: the beast

25195908475657893494027183240048398571429282126204  
03202777713783604366202070759555626401852588078440  
69182906412495150821892985591491761845028084891200  
72844992687392807287776735971418347270261896375014  
97182469116507761337985909570009733045974880842840  
17974291006424586918171951187461215151726546322822  
16869987549182422433637259085141865462043576798423  
38718477444792073993423658482382428119816381501067  
48104516603773060562016196762561338441436038339044  
14952634432190114657544454178424020924616515723350  
77870774981712577246796292638635637328991215483143  
81678998850404453640235273819513786365643912120103  
97122822120720357

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- This time can be reduced by resorting to parallelization. If we accept a **calculation period of 10 years**, we would have to use a cluster of computers that would cover the surface of Tunisia several times, which would cost  $10^{18}$  USD and would require an **electric power of  $10^{12}$  megawatt**, which would exhaust all the world's fossil fuel resources in one day !!! (J. Preskill 2012)

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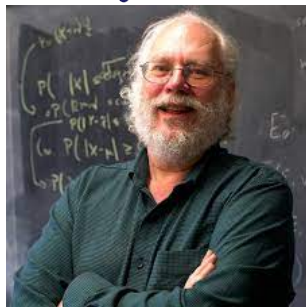
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- **Quantum information** gives hope, and it is reasonable to think that we will be able to solve the problem of factorization of large numbers and many others within a reasonable period of time.
- A major breakthrough was made in **1994** by **Peter Shor**, who developed a quantum factorization algorithm.



Peter Shor  
(ICTP Dirac medal 2017)

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- Physicists made a first rough estimate for RSA-2048 and found that with a quantum computer formed of 10,000 logical qubits and 10 million physical (superconducting) qubits, spaced 1 cm apart for the wiring, which would cost "only" 100 billion USD, and using a modest electric power of 10 MW, would get the job done in **16 hours !!!** (J. Preskill 2012)
- Shor's algorithm works thanks to quantum properties : **Superposition** and **Entanglement**

# Outline of the lectures

- 1 Lecture 1 : Introduction to Quantum Computing
- 2 Lecture 2 : Basics of Quantum Computing