



### RE-THINKING PARTICLE TRANSPORT IN THE MANY-CORE ERA

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# A STEP FORWARD... (FUNCTIONALITY)



- The LHC experiments have expressed the requirement of a continuum spectrum from very detailed to fast simulation
- The idea is to develop a new framework integrating various levels of fast and detailed simulation
  - Keeping services, geometry, I/O and scoring the same (as far as possible)
  - The model is the one of ROOT VMC, but substantially extended
- This will be essential for the new generation of experiments



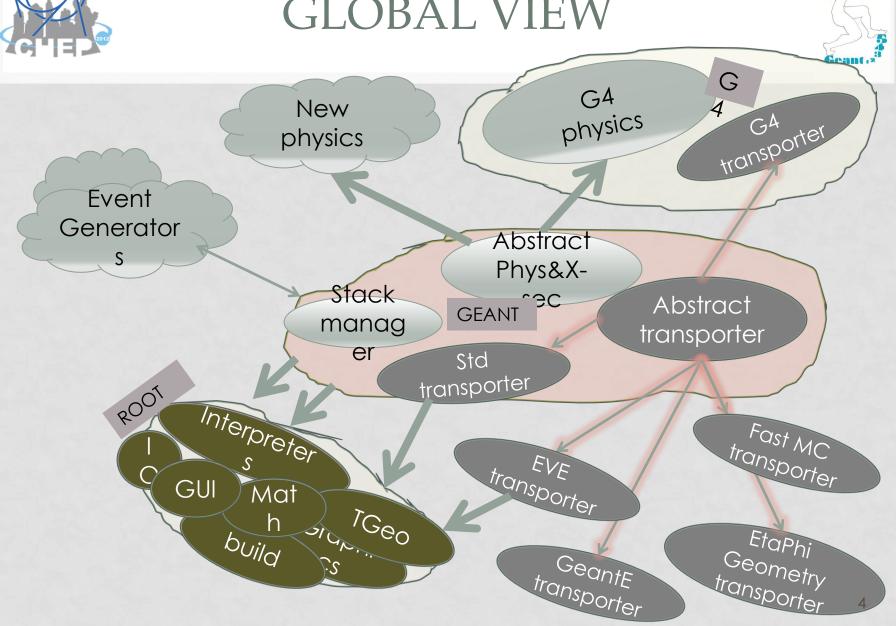
# A STEP FORWARD... (SPEED)



- We have been running Monte Carlo simulations for decades
  - A very large experience has been gained, reflected in the quality of the physics, the complexity of the setups and the precision of the transport
  - State of the art full simulation in a sequential approach...
- CPU frequency has an upper limit, technology has made a turnaround
  - Not yet followed by simulation software
- After a serious investigation, it becomes clear that a rather complete rethinking of the code becomes necessary to exploit the new hardware



#### **GLOBAL VIEW**





#### MOTIVATION FOR LOOKING INTO PARALLELISM

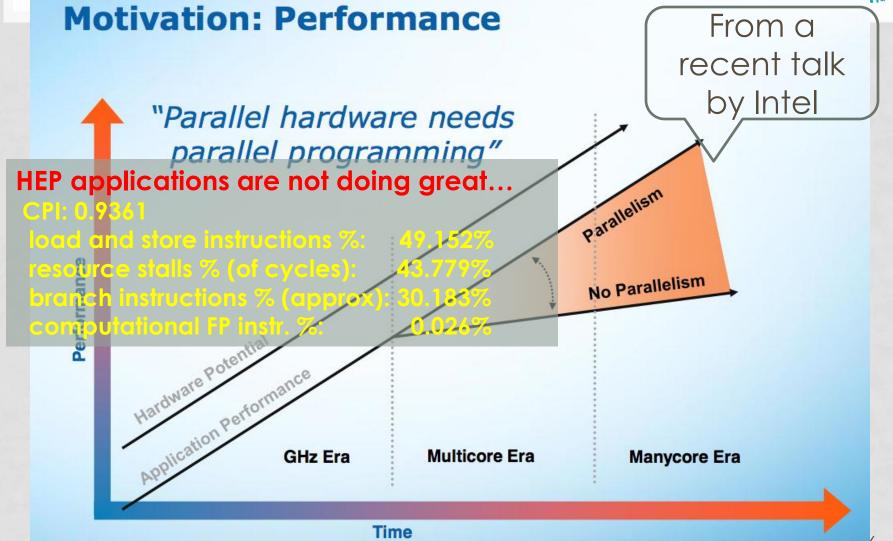


- Parallel architectures are evolving fast
  - Task parallelism in ever growing hybrid configurations
  - Instruction level parallelism (ILP) exploited more and more
    - 4 FLOP/cycle on modern processors
- HEP applications are using very inefficiently these resources
  - Running in average 0.6 instructions/cycle
  - Bad or inefficient usage of C++ leads to scattered data structures and cache misses (both instruction and data)
- We have evaluated the impact of changing the transport strategy on parallel architectures, trying to spot the weak points and a possible winning strategy



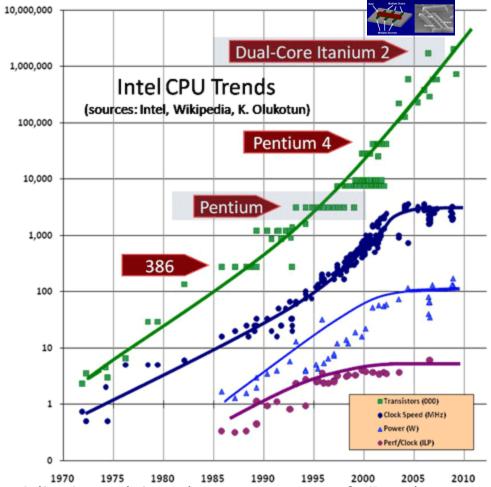
### DOING NOTHING WILL ENLARGE THE GAP







#### TRENDS...



While transistors increase was following Moore's law, frequency and power consumption was not...



Tesla k10 GPU (NVIDIA) – state of the art in GPU technology

Intel Many Integrated Core Architecture – towards highly parallel applications



### EMBARRASSING PARALLELISM?



- Data parallelism on GRID was a savior for HEP, but...
  - Resources get short when one needs to simulate x10 the size of LHC data and uses just a tiny fraction of the CPU power
  - Fast Monte Carlo is a getaway, but cannot help in many performance studies
- Event and track level parallelism: share the code and most of the RO data structures
  - Already a step forward, but does not make jobs more efficient...
- There is an additional need to merge the outputs
  - The process may take longer across different machines than the simulation itself



### RE-THINKING THE SIMULATION APPLICATION



- Vectorizing CPU-critical parts
  - Pushing common computation outside loops and optimizing CPU-intensive loops
  - Using virtualization at high levels but keeping simple structures for low level loops
  - Using efficiently vectors to minimize the scatter/gather overhead
- Parallell/O
  - Multiple buffers to output stream -> new ROOT developments
  - Asynchronous I/O using dynamic scheduling that balances the use of resources
- Using better the memory and minimizing cache misses
  - Data and code locality: transporting a single particle uses now most of the code and data
  - Needs reviewing the code and algorithms: towards a local particle transport
- Using transparently extra resources when available (GPU kernels)
  - Evolution of C++ and code instrumentation languages
- Having in mind all stages of the simulation chain from a parallel perspective
  - Including digitization and I/O
  - Avoiding synchronization issues and aiming for automatic load balancing

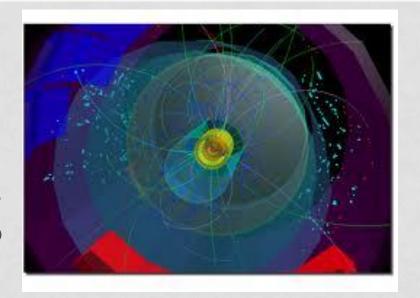
A deep re-organization of particle transport system and a complete re-design of the steering mechanisms and data structures are needed



### THE CLASSIC TRANSPORT APPROACH



- There is a main event loop pushing primary tracks to the transport stack
- Track loop where particles are being extracted from the stack followed one by one to the exit of the detector, by simulating the physics processes along the way

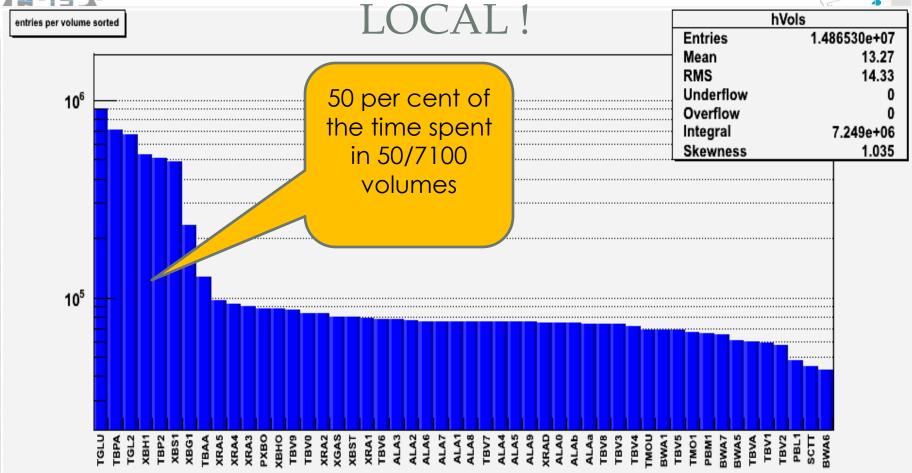


- During each transport step the user code for hit generation is called
- At the end of each event, a digitization procedure is using the specific detector response to convert the hit structure into digits, which are typically dumped to file



### GOOD NEWS: HEP TRANSPORT IS MOSTLY



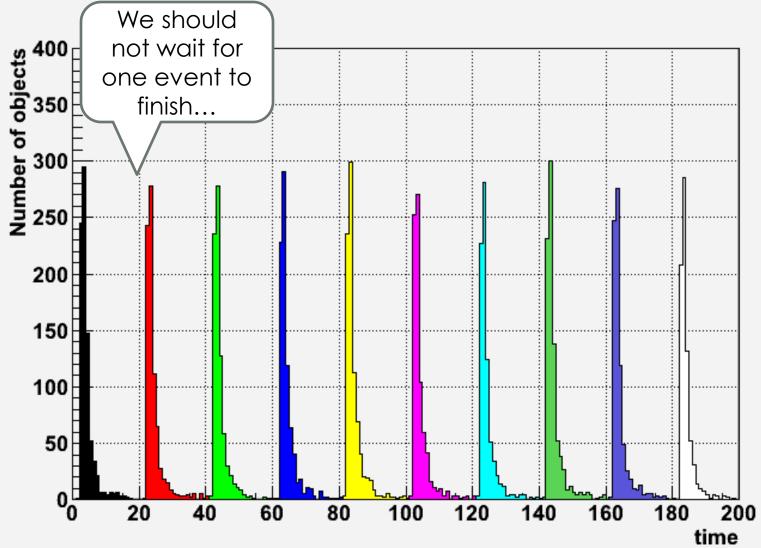


ATLAS volumes sorted by transport time. The same behavior is observed for most HEP geometries.



#### BAD NEWS: TAILS... TAILS...







### A PLAYGROUND FOR NEW IDEAS



- We started with a simple prototype to help exploring some of these issues
  - Simple idea: minimal physics to start with, realistic HEP geometry: can we implement a
    parallel transport model exploiting locality?
  - Base work unit: a vector of particles
- Events and tracks are independent
  - Mixing tracks from different events to avoid tails and have reasonably-sized vectors
  - Study how does scattering/gathering impact the simulation data flow?
  - Can we achieve a good load balancing?
- Toy physics at first, more realistic EM processes to be integrated soon
  - The application should be tuned based on realistic numbers
- New transport model more "detector element"-oriented, profiting from the cached data structures
  - geometry and x-section wise
- Re-design the particle stack and the I/O
- Re-design transport models from a "plug-in" perspective
  - E.g. ability to use fast simulation on per track basis

# VOLUME-ORIENTED TRANSPORT MODEL

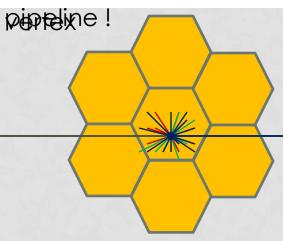
- We implemented a model where all particles traversing a given geometry volume are transported together as a vector until the volume gets empty
  - Same volume -> local (vs. global) geometry navigation, same material and same cross sections
  - Load balancing: distribute all particles from a volume type into smaller work units called baskets, give a basket to a transport thread at a time
- Particles exiting a volume are distributed to baskets of the neighbor volumes until exiting the setup or disappearing
  - Like a champaign cascade, but lower glasses can also fill top ones...
  - No direct communication between threads to avoid synchronization issues



### THE TRANSPORT PROTOTYPE



More events better to cut event tails and fill better the



Associate a set of "baskets" to each geometrical logical volume

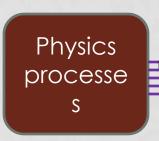


#### A FIRST APPROACH

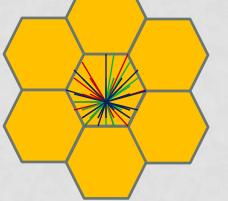


Each thread transports its basket of tracks to the boundaries of the current volume. Move crossing tracks to a buffer, then queue picks-up the next basket from the queue

Transport threads pick-up baskets from the work



Particles(i<sub>0</sub>,...,i



Geometr y transport Particles(i<sub>0</sub>,...,i

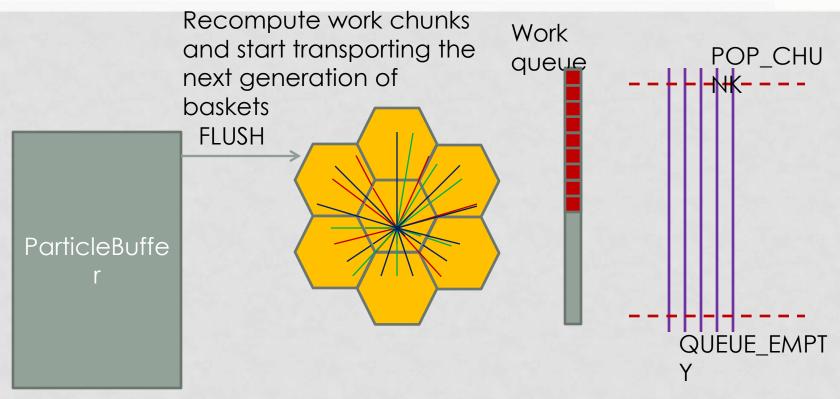
Scatter all injected tracks to baskets. Only baskets above some threshold are transported.

Physics processes and geometry transport called with vectors of particles



### FIRST VERSION REQUIRED SYNCHRONIZATION...





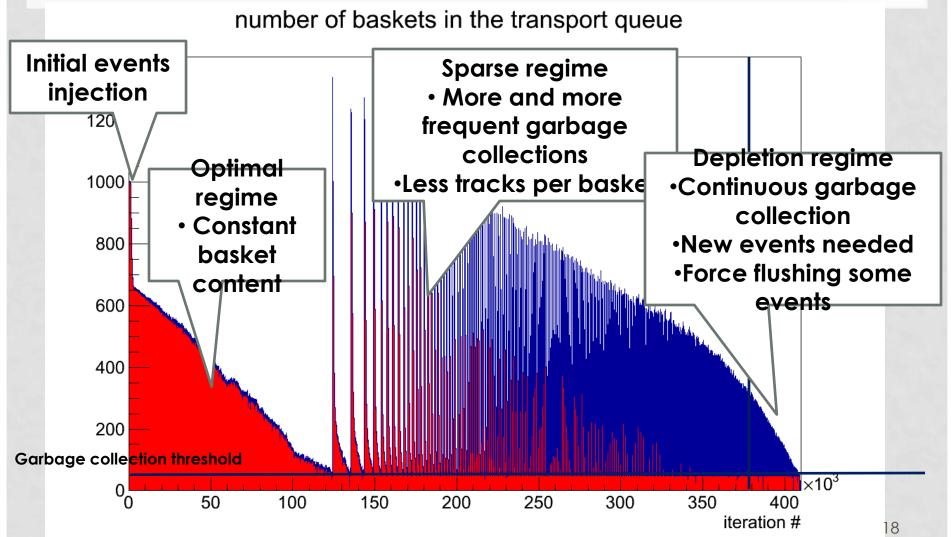
Synchronization point: flush transported particle buffer and sort baskets according content

Generation =
Pop work chunks
until the queue is
empty



### BASKET POPULATION EVOLUTION







#### LESSONS LEARNED



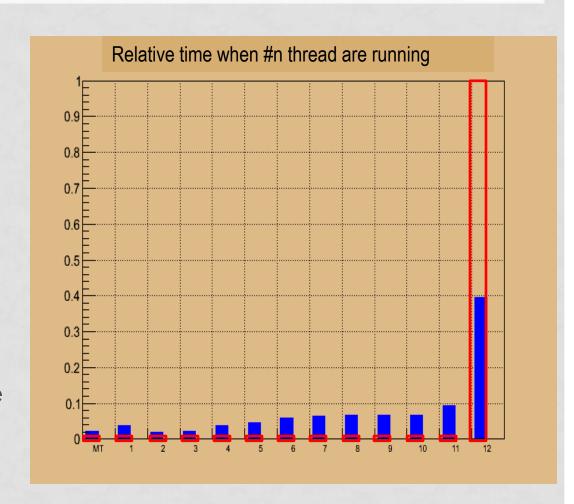
- A model requiring synchronization stages is not appropriate
  - Very large Amdahl effects, increasing in the track depletion stage
- Balancing basket populations in time is not trivial
  - Events need to be injected in the system to compensate the basket inefficiencies
    - Will cumulate hits and make memory grow
  - Hits from early introduced events need to be evacuated to the digitization and I/O threads
- We need a good estimate of the percentage of work that can be done with "efficient" vectors
  - A model including realistic physics, digitization and I/O will be needed



### CONCURRENCY IN THE FIRST APPROACH



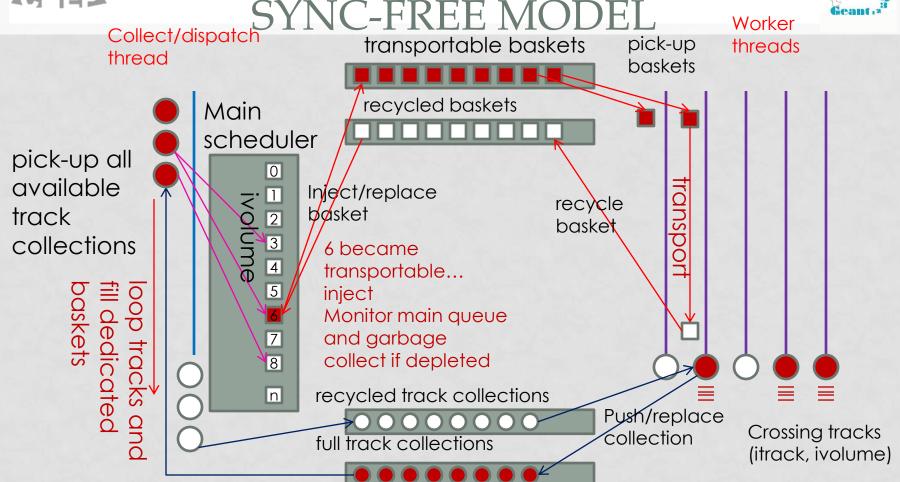
- Ideally all workers should be in running state during the processing phase, so the distribution should peak for the number of workers
- Synchronization becomes critical during the particle depletion regime, when particle baskets are nonoptimally filled and garbage collections more often





# EVOLUTION OF THE PROTOTYPE TO A







= real\_time[1]/real\_time[n]

dn-peeds

5

10

#### PRELIMINARY BENCHMARKS



RealTime

24 15.24

5.945

**Entries** 

Mean

RMS

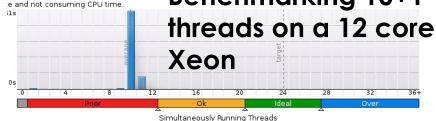
HT mode

Real speed-up 12 core x 2 HT, 1 collector

15

20

This histogram represents a breakdown of the Elapsed Time. It visualizes the percentage of the wall time the specific nber of threads were running simultaneously. Threads are considered running if they are either actually running on a Jor are in the runnable state in the OS scheduler. Essentially, Thread Concurrency is a measurement of the numb nreads that were not waiting. Thread Concurrence in the concurrence in t



#### 'U Usage Histogram 🕒

Thread Concurrency Histogram

histogram represents a breakdown of the Elapsed Time. It visualizes what percentage of the wall time the specific nber of CPUs were running simultaneously. CPU Usage may be higher than the thread concurrency if a thread is cuting code on a CPU while it is logically waiting.



Simultaneously Utilized Logical CPUs

unction	Wait Time by Utilization	Wait Count	Spin Time	Module	Object Type	Object Creation Module and Function	*	⊇ 1 stack(s) selected. Viewing ⇒ Current stack is 100.0% of sel
1	73.005s	11 11	0ms 0ms	[Unkno	Constant Constant	[Unknown] [Unknown]	1	lbThread.so/TPosixCondition::Wa
inown] riable 0xf ndition::W		607,511 607,511	0ms	[Unkno	Constant Condition Variable Condition Variable		۱	IbThread.so/TCondition::Wait(voi IbGeant.so/concurrent_queue:x IbGeant.so/WorkloadManager::C
ndition:/a known]		579,522 27,989	0ms	libThrea	Condition Variable Condition Variable	libThread.so/TPosixCondition::Wait	۱	libThread.soiTThread::Function(v libpthread-2.11.1.soistart_threa
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10cf4 riable 0xc 14dac	2.6275 0 0.510s 0.042s	150,120 2,164 73	290.145ms 0ms		Mutex Condition Variable Mutex	libThread.so!TPosixMutex:Lock libThread.so!TPosixCondition::Wait libThread.so!TPosixThread::Exit	i	1.

Locks and waits: some overhead due to

transitions coming from exchanging nthreads

baskets via concurrent queues



### A CONSIDERABLE EFFORT TO MOVE FORWARD...



- HEP simulation cannot afford ignoring the technology trends
  - LHC computing needs largely demonstrate it
  - The software inefficiency becomes more and more expensive
- The move towards efficiency is long and tedious
  - Even if all the components (physics, geometry, ...) are written
  - We started to look under the hood, and tried out few ideas, but most of the work is still to be done
- We need to constantly monitor the progress
  - GEANT 4 MT will be a good reference system once the physics will be comparable
  - Even between different versions of the prototype



#### **FUTURE PLANS**



- Continue the investigation of parallel event transport
- Develop more realistic physics models
- Integrate fast simulation
- Aiming at a working prototype in 2015
- The activity will start in earnest in September 2012
- As with other large successful projects (ROOT, G4), this will be an international collaboration
  - We count on the HEP community (LHC, ILC, CLIC, FAIR, ...) for help and feedback



#### **CONCLUSIONS**



- The new generation (Geant5) of detector simulation programs will have to
  - Integrated seamlessly fast and detailed simulation at different levels
  - Make efficient use of parallelism at different levels
  - Capitalizing on the large Geant 1-4 experience
- A prototype is being built at CERN, which will require collaboration with the HEP Community at large
- The first results are interesting and our learning curve very steep (!)
- Stay tuned...