Dynamic Federations

Global Access to HTTP/WebDAV Storage Elements

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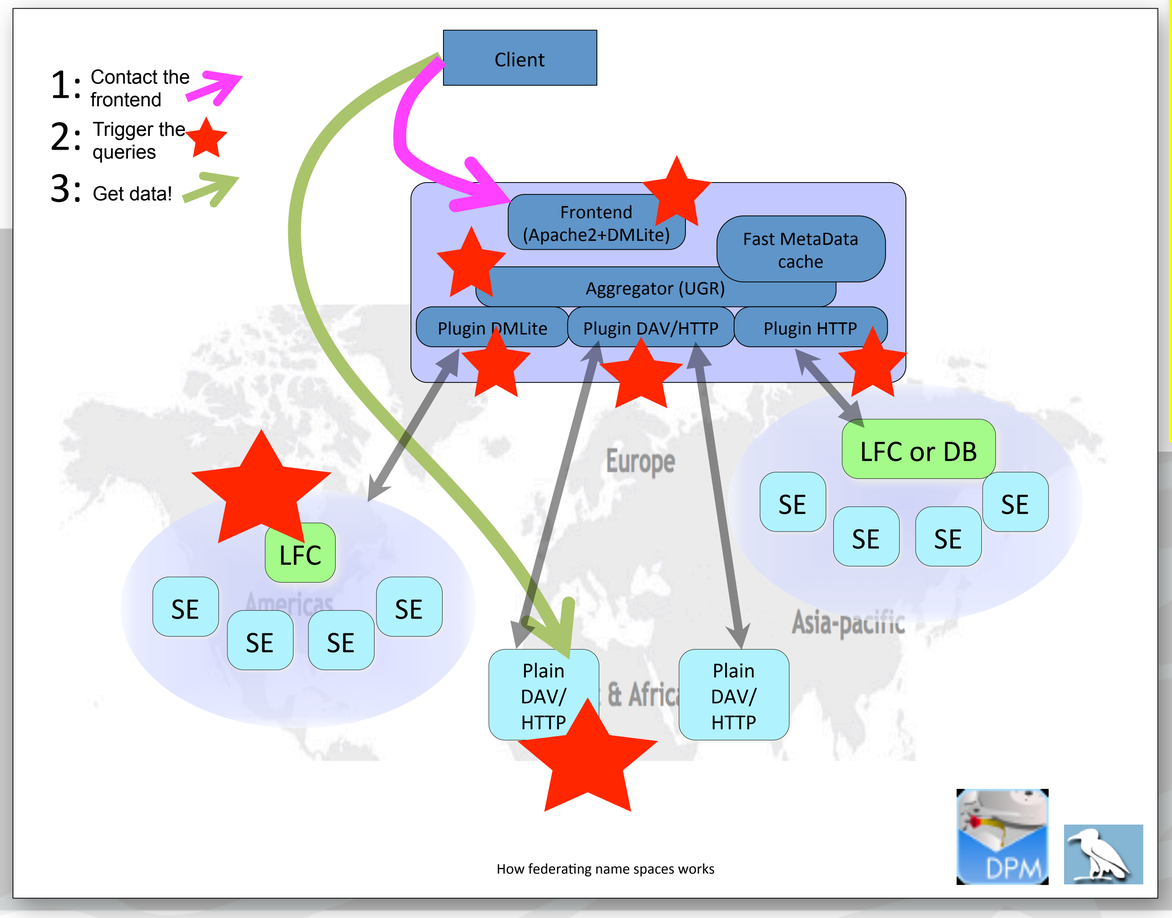
Currently data lives on islands of storage
Catalogues are the maps, FTS/gridFTP are the delivery companies, experiment frameworks populate the island
Jobs are directed to places where the needed data is (or should be) .....
Almost all data lives on more than one island
Common assumptions:
perfect storage (unlikely to impossible) perfect experiment workflow and catalogues (unlikely)
Strict data locality has limitations
e.g. a single missing file can derail the whole job or series of jobs

Make different storage clusters be seen as one

Make global file-based data access seamless

Fully support metadata browsing

No strange APIs, everything looks natural
Use dynamic systems that are easy to setup/maintain:
No complex metadata persistency
No DB babysitting (and lives well with the experiment's metadata repositories)



How federating name spaces works

Federations can help
- **Failover between islands**
- **Simplicity for personal interactive access, anything is anywhere**
- **'diskless' sites where data is nearby**
- **Storage sharing, eg 3 collaborating sites, each with 1/3 of required files**
- **Seamless integration of cloud storage**
- **Reduced data management for analysis, files fetched when needed**
- **Site 'overflow' (when jobs are waiting too long in a site queue, balance IO and CPUs)**

We federate (meta)data repositories that are 'compatible'
Name space (modulo simple prefixes)
Permissions (they don't contradict across sites)
Content (same key or filename means same file)

Dynamic, transparent metadata discover
looks like a unique, very fast file metadata system
properly presents the aggregated metadata views
redirects clients to the geographically closest endpoint

TODAY we can federate multiple, worldwide instances of:
dCache DAV/HTTP
DPM DAV/HTTP, LFC DAV/HTTP
Cloud DAV/HTTP
Native LFC and DPM databases (through DMLite used as a client)
Can be extended to other metadata sources

The system also can load a Geo plugin
Gives a geographical location to replicas and clients
Allows the core to choose the replica that is closer to the client
The one that's available uses GeoIP (free)* other implementations are possible

* "This product includes GeoLite data created by MaxMind, available from http://maxmind.com/"

No central catalogue
inconsistencies, by design
(central catalogue is not needed in a dynamic federation)

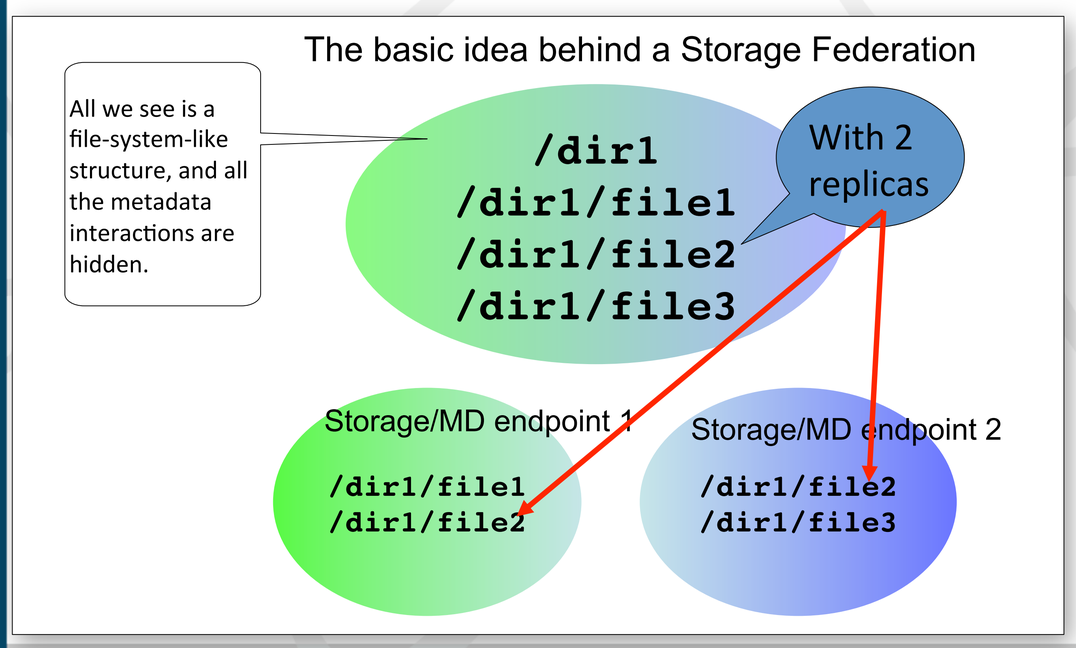
High Performance

Can be used by client tools that everybody knows
focus on HTTP/DAV, we can use it from a smartphone.
Base everything on open 'just works' technologies.
Use your browser to browse your data, anywhere, make the GRID jobs use the same
See poster "Web enabled data management with DPM&LFC"

Technology

The core component is a plugin-based component called "Uniform Generic Redirector" (*Ugr*). It can plug into an Apache server thanks to the *DMLite* and *DAV-DMLite* modules (by IT-GT). Architecturally, Ugr acts as an information feeder for *DMLite*.

Ugr internally splits the queries into parallel tasks of information location
Composes on the fly the aggregated metadata views by managing these tasks and delaying clients the minimum amount of time that is necessary
Never stacks up latencies! No fixed delays.
Able to handle file listings and metadata
Able to **redirect** clients to replicas
By construction, the internal workspace is a data structure that models a partial, volatile namespace
An LRU purging policy makes a fast, in-memory, 1st level cached namespace.



Full parallelism, high performance
No limit to the number of outstanding clients/tasks
No global locks/serializations!
The endpoints are treated in a completely independent way
Thread pools, prod/consumer queues used extensively (e.g. to stat N items in M endpoints while X clients wait for some items)

Aggressive metadata caching
A relaxed, hash-based, in-memory partial name space
Juggles info in order to always contain what's needed
Stalls clients the minimum time that is necessary to juggle their information bits
Peak performance per CPU core: 0.5~1M stats/sec

High performance DAV client implementation (DAVIX)
Loaded by the core as a "location" plugin
Uses libneon w/ sessions caching
Compound list/stat operations

Seamless storage federations of:
Official Storage Elements, LFCs, catalogues…
Cached data (i.e. SQUID-like things, not registered in any catalogue)
HTTP/DAV-based servers
Cloud storage services
HTTP-enabled XROOTD/EOS clusters, sharing the data.
Local SE as a preference, give the freedom to point to an **efficient** and **reliable** global federation

Optimize redirections based on on-the-fly client-data proximity
Limit complexity: read only
Usually writes happen to well-known, close islands

Detailed performance measurements are on the way.

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