ATLAS Computing in Run-2

Borut Paul Kersevan

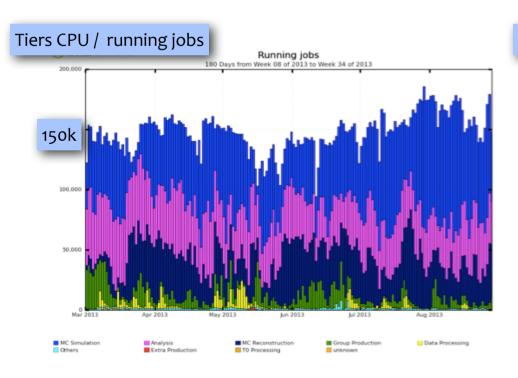
ATLAS Resource Utilization in 2013

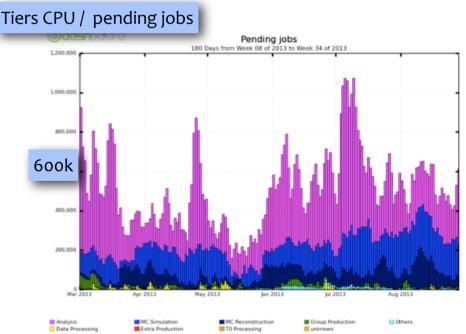


ATLAS RESOURCE USAGE IN FIRST HALF OF 2013 (RRB)

	Location	Requested	Used
CPU [kHS06]	CERN	111	111
	Tier-1	316	435
	Tier-2	360	713
Disk [PB]	CERN	9.6	8.9
	Tier-1	35 [38]	35
	Tier-2	51 [52]	48
Tape [PB]	CERN	25	29 (incl. 9 PB of ESD)
	Tier-1	42	33

- ATLAS has utilized the computing resources in its Tiers well in the last year: many thanks to sites for resources and excellent operating!
 - We manage to provide a timely throughput of analyses to meet the physics requirements.
- An ongoing effort in software development to optimise the resource utilization by reducing the CPU consumption, event sizes - for Run-2...





The Challenges of Run-2



Constraints of 'flat budget'

- Both for hardware and for operation and development
- Hardware increase from Moore's law gain, estimated at factors of 1.2/year for CPU and 1.15/year for disk

♦ Data from Run-1

 Proper data preservation and integration with Run-2 data analysis New CPU architectues: less memory / core

♦LHC operation

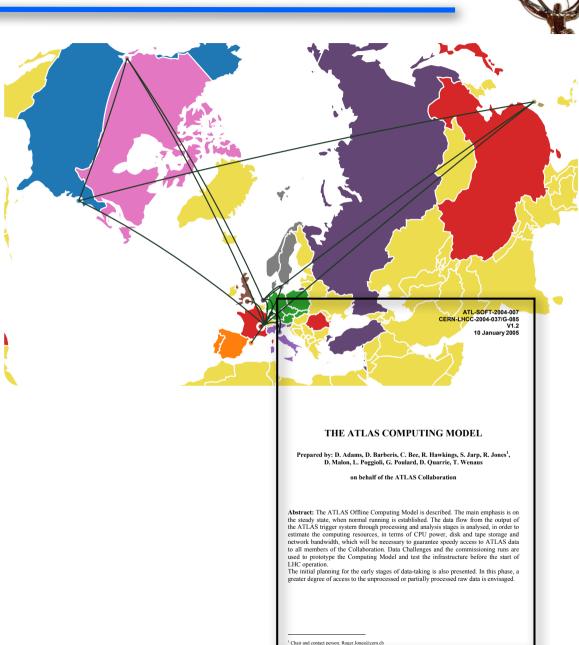
- HLT rate 1 kHz
- Pile-up up above 30
- 25 ns bunch spacing
- Centre-of-mass energy x ~2

→ ~new detector

 Needs to be incorporated in simulation and reconstruction.

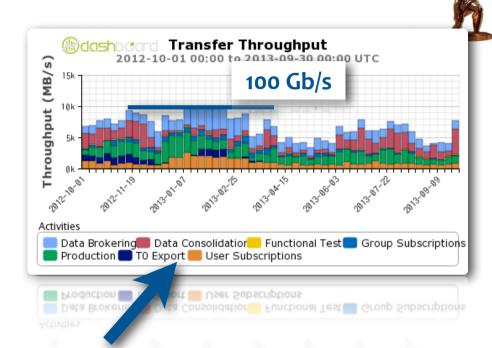
Initial Computing Model (2005)

- Derived from MONARC ('99) model
- CERN-To the center
- 10 T1s connected by dedicated 10Gb/s links (LHCOPN)
- O(100) T2s each attached to a T1
- The data flows along the hierarchy
- Insufficient networking assumed
- Hierarchy of functionality and capability



From 2010 to 2013: some of the many changes

- Hide grid complexity from users,
 simplifications, less middleware dependence
- Caching opposed to centralized DB
 - Conditions data access from any site, not only at T1s (Frontier)
 - No more need to pre-install software releases at sites (CVMFS)
- Dynamic data placement and deletion based on popularity
 - Better usage of disk space
 - Reduced job waiting times
- T2 \rightarrow N-T1s & T2 \leftrightarrow T2 exchanges



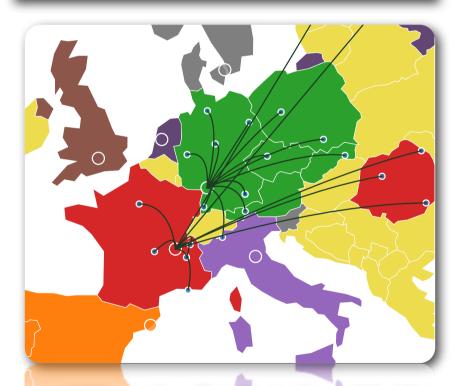
Network performing over expectations

2011 LHCONE:
Dedicated network between (some) WLCG sites



2010

Planned data distribution
Jobs go to data
Multi-hop data flows
Poor T2 networking across regions



~20 AOD copies distributed worldwide

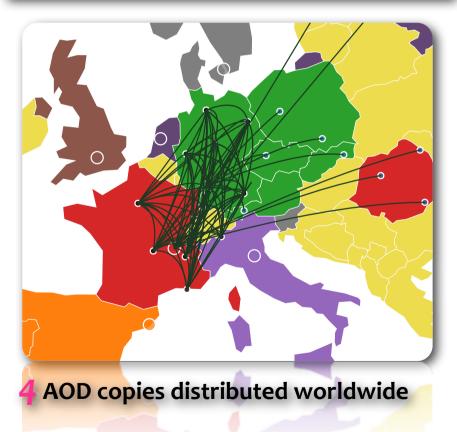
2010

Planned data distribution
Jobs go to data
Multi-hop data flows
Poor T2 networking across regions





~20 AOD copies distributed worldwide



Networking Potential

- Networking is the one item that will most probably continue its progress & evolution further..
- ... and reduce the cost fastest.
 - In terms of bandwidth increase.
 - In terms of new technologies (NaaS Network as a (virtual) Service?)



Advanced Networking for HEP, Research and Education in the LHC Era **Harvey B Newman** California Institute of Technology Remarkable Historical ESnet Traffic Trend **ESnet Traffic Increases Actual Dec** 2012 100,000 1 10X Each 4.25 Yrs, for 20+ Yrs 12 PBy/mo 15.5 PBytes/mo. in April 2013 Apr 2007 1 PBv/mo. The Trend Continues 10,000 Nov 2001 1000 100 TBv/mo. Avg. Annua Jul 1998 Terabytes / month Growth: 72% 10 TBy/mo. Oct 1993 1 TBy/mo. 53 months **Projection** 40 months 100 PBv/mo Actual 57 months Exponential fit + 12 month projection W. Johnston, G. Bell 38 months Log Plot of ESnet Monthly Accepted Traffic, January 1990 - December 2012,

https://indico.cern.ch/sessionDisplay.py?sessionId=0&confld=214784#20131017

If so, we could design servers without needing NIC cards – no difference

With HEP applications like FAX, file systems or memory can be mounted

SDN could revolutionize how computing is done, are we ready for that?

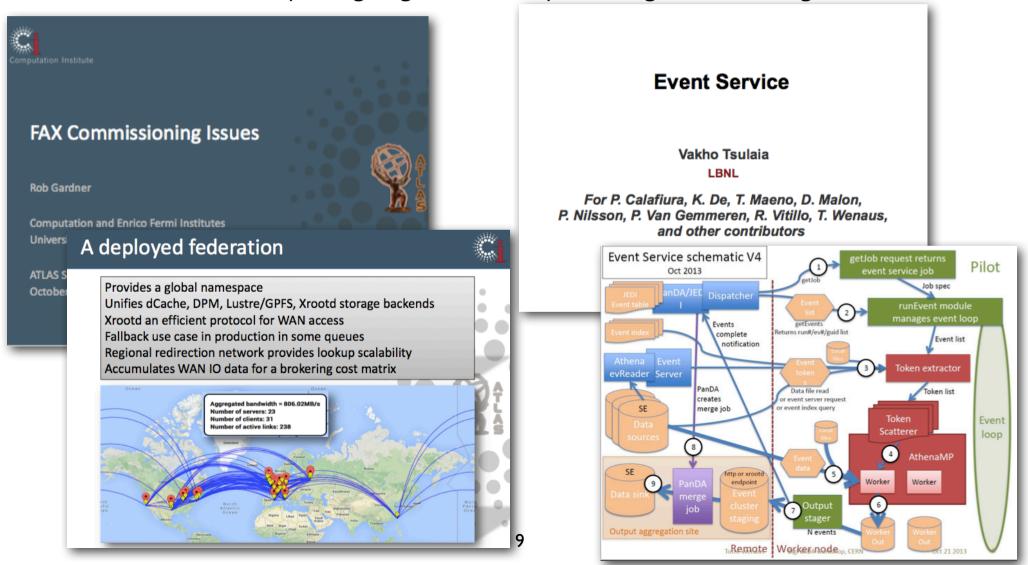
With SDN, can effectively route IP and non-IP protocols (like ROCE)

between communication within the motherboard or outside.

remotely to my chip while 'streaming data for analysis.'

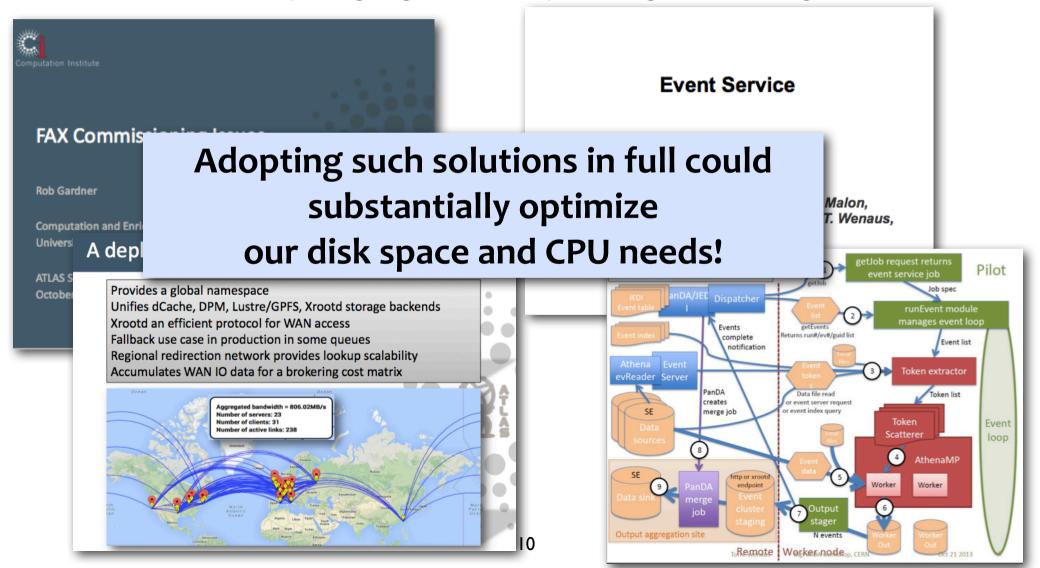
How can we profit from network evolution?

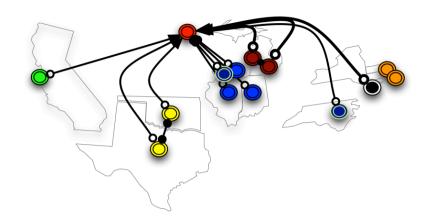
- Two interesting ATLAS initiatives ongoing:
 - data federations (FAX xrootd federation, http federation) remote file access over WAN.
 - Event Service passing single events for processing from/to storage.



How can we profit from network evolution?

- Two interesting ATLAS initiatives ongoing:
 - data federations (FAX xrootd federation, http federation) remote file access over WAN.
 - Event Service passing single events for processing from/to storage.





DISTRIBUTED STORAGE / REMOTE ACCESS

- ★ Jobs access data on shared storage resources via WAN
- ★ For a better usage of storage resources (disk prices!)
- ★ Bandwidth and stability needed
- ★ FAX (Federating ATLAS data stores using Xrootd) demonstrator, job fail-over in case of access failure for first implementation
- ★ http protocol also considered

Event Service schematic V4 Oct 2013 JEDI JEDI Event table Events complete notification Athena Event evReader Server PanDA creates merge job Data file read or event server request or event server request or event index query Token list Token extractor Scatterer Set PanDA creates merge job Data file read or event server request or event index query Token Scatterer Levent token Token list or event index query Token Scatterer Loop Set PanDA creates merge job Data file read or event server request or event index query Token Scatterer Loop Output aggregation site Token list or event server request or event se

Event service

- In development: software and distributed computing effort
- ◆ Feed Virtual Machines with short jobs (simulate one single event)
- ♦ Usages:
 - Backfilling of HPC centers
 - Opportunistic use of commercial clouds
 - Volunteer computing (ATLAS@home)

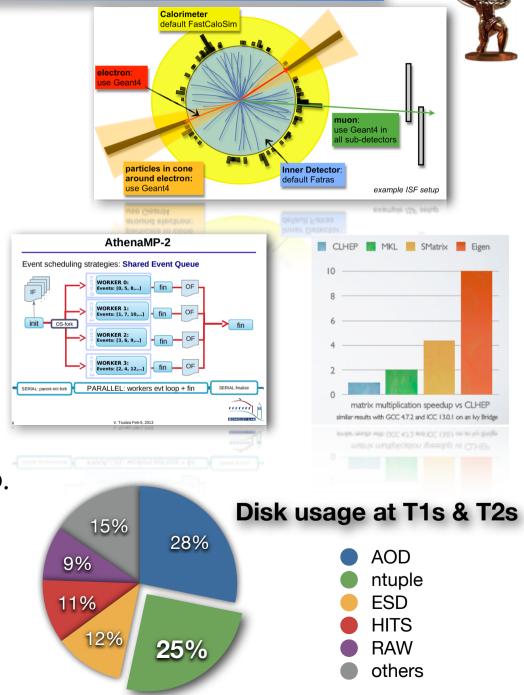
Some Limitations of Current Model & Tools



- Partitioning of resources:
 - Analysis vs Central Production,
 - T1s vs T2s.
- Difficulties of current Data Distribution Management & production systems to accommodate new use cases and technologies.
- Memory increase of MC pile-up digitization & reconstruction.
- Multitude of data format for analysis.
- Full reprocessing once a year.

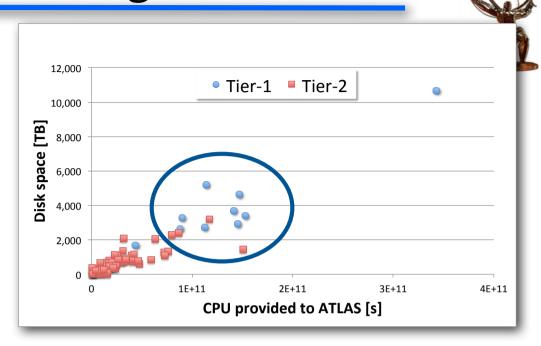
Working towards Solutions

- Simulation: CPU
 - Integrated Software Framework.
- Reconstruction : Memory & CPU
 - Parallelism, code speedup.
 - MP solution to reduce memory footprint.
- Analysis Model: multiplication of data formats
 - Common analysis data format, xAOD.
 - Streamlining the analysis flow.



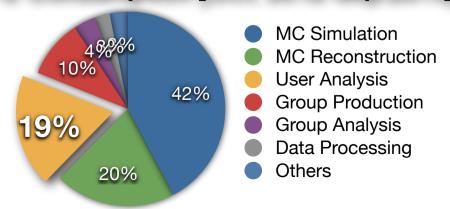
Data Processing

- Optional extension of first pass processing from T0 to T1s in case of resource shortage at To
- T1s and some T2s used for the most demanding workflows: high memory and I/O intensive tasks
- Data reprocessing & MC reconstruction also performed at SOME T2S
- Still one full reprocessing from RAW / year, but multiple AOD2AOD reprocessings / year
- Derivation Framework (train model) to centrally produce TB size data samples for analyses



Some T2s are equivalent to T1s in term of disk storage & CPU power

CPU Consumption [Oct. 2012-Sep.2013]



Data Placement in Run-2



- Initially **2 copies** of analysis data formats (xAOD: one at T1s, one at T2s):
 - Already being implemented to gain disk space.
- Non popular data will be archived to tape at T1s:
 - Further refinement of the popularity monitoring and data placement:
 - In October we managed to recover 9% of disk space occupied by data never accessed over last 9 months.
 - Minimal number of copies on disk not guaranteed.
 - User access to data on tape granted through centralized tools.
- Investigate data access patterns to provide information to sites to optimize the site hardware configurations (and cost):
 - Low access/high access on disk (caching of popular data)
 - Low access/high access on tape (different tape technologies & cost)

New production system: ProdSys-2

PanDA+JEDI+DEfT



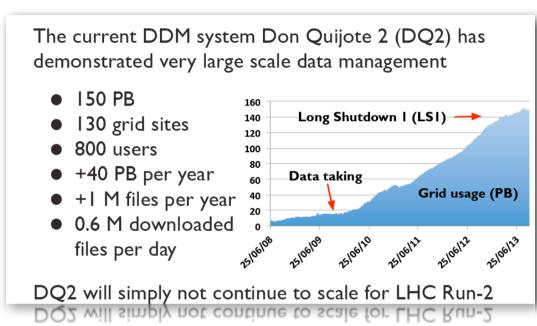
- Same engine for analysis and production
 - Currently analysis vs production shares managed by Sites not by ATLAS.
 - Better reactivity to analysis load
- Data traffic minimized.
- Optimized job to resource matching.



Data Distribution Management & databases



- New Data Distribution Management system: Rucio will replace DQ2
 - New scalable architecture.
 - File level functionality instead of dataset.
 - Built-in data replication policy for space and network optimization.
 - Multi-protocol (http,...).
- Database infrastructure: simplification and streamlining.





http://rucio.cern.ch

Opportunistic resources

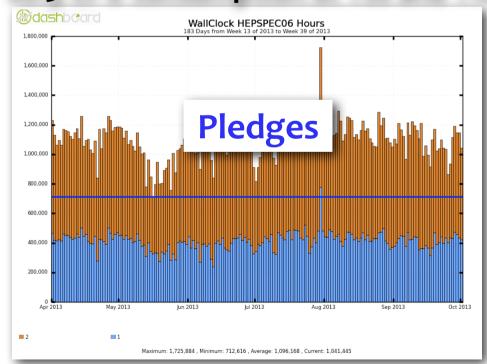


CPU consumption above pledges both at T1s and T2s

- Additional solutions :
 - HLT farm at P1.
 - Cloud computing.
 - Large HPC (High Performance Computing) centers.
 - Volunteer computing: ATLAS@home.

Sites and Funding Agencies provide **MOPE** than pledged resources (thank you!)

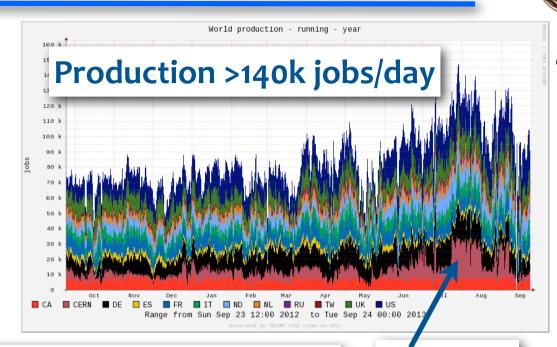




HLT farm at P1

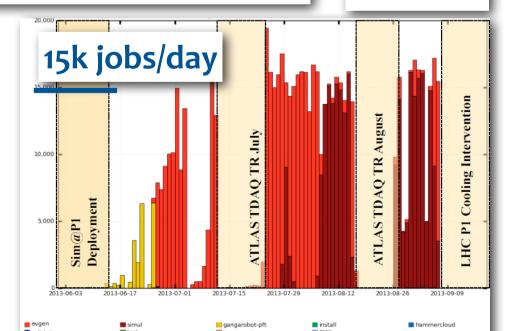
- HLT farm 'cloudified' mid-2013
 - Reached >15k concurrent simulation jobs
- Switch between trigger and simulation mode operational.
- Availability in Run-2:
 - for MC production during shutdowns or LHC technical stops if/when no other TDAQ activities.
 - ~30% over a Run-2 year?

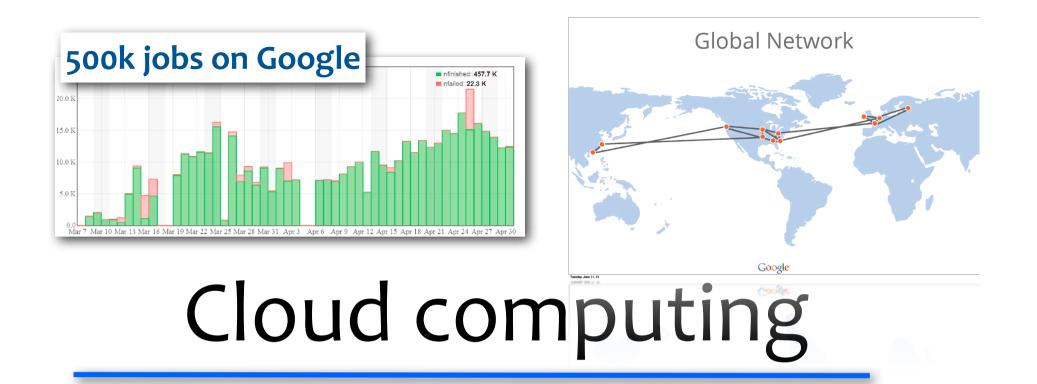
BNL/T1+IT/SDC+ATLAS/TDAQ experts: many thanks!



Equivalent to a T₁ or a large T₂







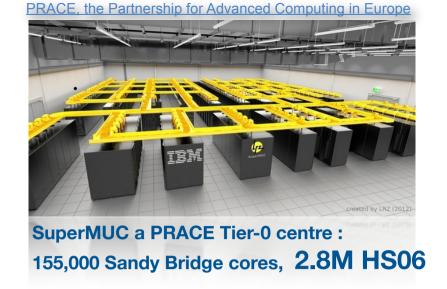
On going R&D on academic clouds and Amazon or Google (AUS,CA, US,...)

Issues with long jobs and I/O

Plan for use 'academic' clouds and opportunistic use of 'cheap' commercial is possible

Some cloud computing providers start to propose cost-competitive offers (with some limitations)

HPC (High-Performance Computing) resources



WLCG 2013 T0/1/2 pledges ~2.0M HS06

Large investments in many countries: from Peta to Exa scales initiatives[1]

Latest competitive supercomputers are familiar Linux clusters

 Large number of spare CPU cycles are available at HPCs which are not used by 'standard' HPC applications

Projects to use idle CPU cycles at HPC centers in US, China & DE

Demonstrators working for simulation & event generation

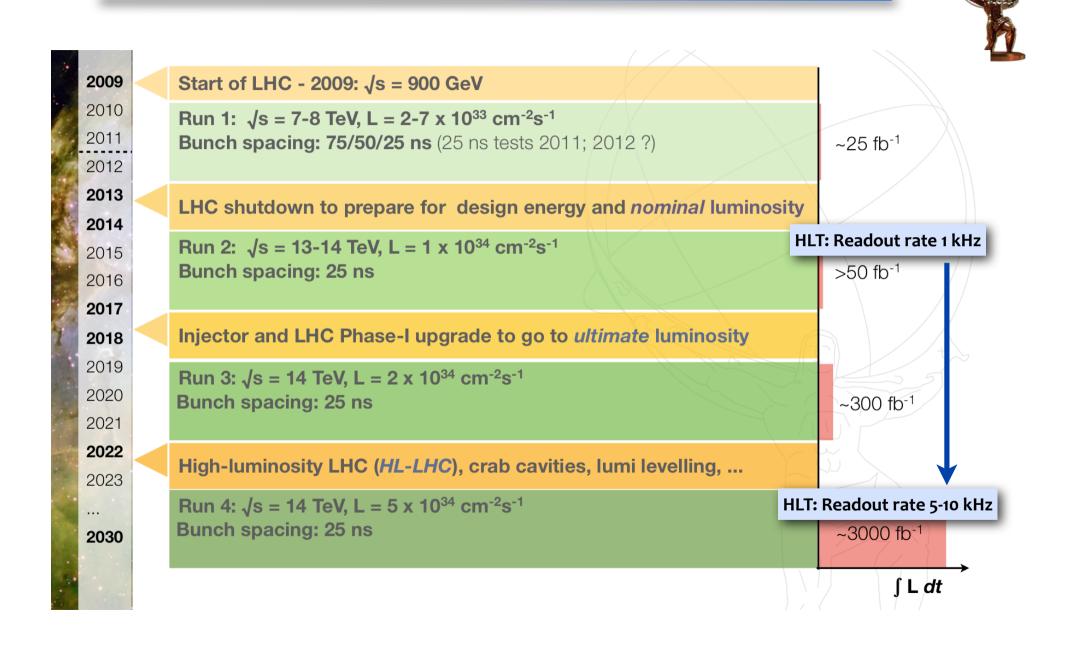
Difficult to use HPC centers for I/O intensive applications

Outbound connectivity of HPC centers may also be an issue

Some T2s plan to provide pledges resources on shared HPC facilities

Might endanger traditional HEP computing budget

LHC Upgrade Timeline - the Challenge to Computing Repeats periodically!



Looking further in the future: in 10 years?



• The answer to this has strong political/financial components, which are hard to predict. Still..

R. Walker HPC Resource Examples

- SuperMUC, Munich
 - 155,000 Sandy Bridge cores, 2.8M HS06
 - ATLAS 2013 T1/2 pledges ~ 730K HS06
 - Suse Enterprise Linux 11, 2GB/core
 - warm water cooling
 - 40°C inlet. 70°C outlet used to heat building
- Hydra, MPI, Munich
 - 'similar' cluster in spec and scale
 - due Summer 2013. 10k core integration system in place now
- · MOGON, Mainz
 - 34k cores SL6

The main item that does not have many solutions and gives a severe constraint is our data storage:

We need reliable and permanent storage under ATLAS control.

From another perspective, with the network evolution (and federated storage, event service..) 'local' becomes re-defined (again).

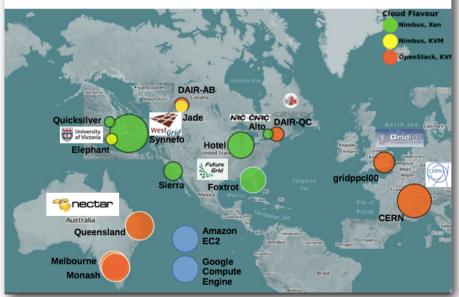
No need for local storage?

Consolidate to a fewer (20|10|..1?) main storage points? Cheaper?

Even today, our CPU capacities fit into one super-computing center.
Will we get fractions of Super-Computer CPUs?

What will be the impact of IaaS (Cloud) techologies?
Will we get cheques for commercial cloud use (again, supercomputers..)?

The "Grid of Clouds"





Ian Gable, Ryan Taylor - Sep. 2013

Summary & Outlook



- ◆A lot of experience acquired in 3 years of LHC data taking.
- Run2 will put high pressure on hardware and human resources.
- Solutions under development and manpower is needed.
- New computing model and its components will be tested during 2014 data challenge.
- LHC & ATLAS upgrades also mean resources for software& computing.

Backup



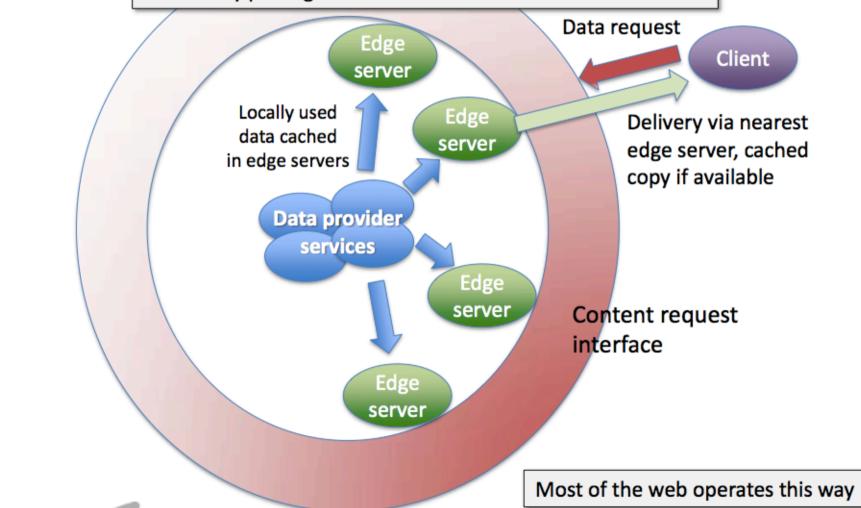
More on Future Data Acess...



The Content Delivery Network Model

T. Wenaus

Content delivery network: deliver data quickly and efficiently by placing data of interest close to its clients



More on Future Data Acess...



The Content Delivery Network Model

A growing number of HEP services are designed to operate broadly on the CDN model

T. Wenaus

Service	Implementation	In production
Frontier conditions DB	Central DB + web service cached by http proxies	~10 years (CDF, CMS, ATLAS,)
CERNVM File System (CVMFS)	Central file repo + web service cached by http proxies and accessible as local file system	Few years (LHC expts, OSG,)
Xrootd based federated distributed storage	Global namespace with local xrootd acting much like an edge service for the federated store	Xrootd 10+ years Federations ~now (CMS AAA, ATLAS FAX,) See Brian's talk
Event service	Requested events delivered to a client agnostic as to event origin (cache, remote file, on-demand generation)	ATLAS implementation coming in 2014
Virtual data service	The ultimate event service backed by data provenance, regeneration infrastructure	Few years?