

Running ATLAS workloads within massively parallel distributed applications using Athena Multi-Process framework (AthenaMP)

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History of AthenaMP

- Motivation
 - ATLAS reconstruction is memory-hungry
 - We needed to have a mechanism for optimizing memory footprint without touching the algorithmic code-base
- AthenaMP leverages Linux Copy-On-Write for sharing memory pages between processes, which were forked from the same master process
 - Thus the memory sharing comes "for free"
- Originally implemented as Python layer
 - Presented at CHEP 2009:
 "Harnessing multicores: strategies and implementations in ATLAS", S.Binet et al.
- Later on completely **rewritten in C++** to follow Gaudi Component Model
- Currently being actively used for running ATLAS production jobs on multi-core resources on the Grid

Workflow

Schematic View of ATLAS AthenaMP



Workflow (contd.)

- The master process goes through the initialization phase and then forks N subprocesses (workers)
- By **delaying fork as much as possible**, we increase the amount of memory shared between the workers
- AthenaMP implements **different strategies for scheduling workload to the worker processes.** Each of these strategies is implemented by a specialized **Gaudi AlgTool**
- The workers **retrieve event data independently from each other**. They process events and write their own output files to the disc
- After all events assigned to the given job have been processed, the **master usually proceeds with merging workers' outputs**

Assigning workloads to the worker processes

- The complexity of ATLAS Event Data Model does not allow us to directly exchange event data objects between processes
 - The only exception is RAW data reconstruction, where we can pass around event data in the form of void* memory buffer
 - We intend to work on the ATLAS persistency infrastructure in order to address this issue (see "Future Developments" section later in this talk ...)
- For the time being, AthenaMP delivers workload to the worker processes in the form of **event identifiers** either **integers** (event position in the input file) or **strings** (unique event token)
- For each of the above scenarios AthenaMP uses a **special auxiliary process**, which distributes event identifiers/data between worker processes
- The sub-processes of AthenaMP communicate to each other using IPC mechanisms (**shared memory, shared queue**)

Saving memory

Memory (GB)

ATLAS Preliminary. Memory Profile of MC Reconstruction



- RSS of 8 serial jobs compared to RSS of one AthenaMP job with 8 workers
- Actual memory savings depend on the job type and configuration
 - In this example: AthenaMP reduces overall memory footprint by 45% at the reconstruction step

• This plot has been obtained by running test jobs on otherwise empty machine and profiling system memory with free

Distributing event numbers. Shared Queue

- AthenaMP uses **shared event queue** for distributing event numbers either individually or in chunks between worker processes
- The events are assigned to workers on the "first come first served" basis
 - A worker pulls new event from the queue after it finished processing the previous one
- Such dynamic distribution of the workload guarantees load balancing on the workers
- Shared event queue is the default strategy for running AthenaMP jobs on the Grid

AthenaMP on the Grid



- Today ATLAS runs a substantial fraction of its production workloads on the Grid using AthenaMP
- The workloads include (but are not limited to) Geant4 simulation and simulated data reconstruction

 The plot shows the number of CPU-cores used by ATLAS production jobs – serial and MP – on the Grid in March 2015

Fine grained Event Service

- **Event Service** a new approach to event processing in ATLAS
 - Job granularity changes from files to individual events
 - Deliver only those events to a compute node, which will be processed there by the payload application, don't stage in entire input files
- Event Service is agile and efficient in exploring diverse, distributed, potentially short-lived (opportunistic) resources: "conventional resources" (Grid), supercomputers, spot market clouds, volunteer computing
- For more details about Event Service see the presentation by Torre Wenaus at CHEP2015: "The ATLAS Event Service: A new approach to event processing" (Contribution #183)
- Event Service uses AthenaMP as **payload application for event processing**

AthenaMP and Event Service



- For the Event Service
 AthenaMP uses the strategy
 of distributing event tokens
 to the worker processes
- The tokens are retrieved from external source by a specialized AthenaMP subprocess Token Scatterer
- Workers retrieve event data using the token. Data may be local or remote
- AthenaMP is configured to write new output file for each processed event range
- The latter functionality is implemented by the Output File Sequencer a new mechanism developed in ATLAS, which is currently being used only by AthenaMP within Event Service



AthenaMP and Yoda



- Yoda MPI-based implementation of the Event Service, designed to run on supercomputer systems with no internet connection from the compute nodes
- For more information about Yoda see the presentation by V. Tsulaia at CHEP2015: "Fine grained event processing on HPCs with the ATLAS Yoda system" (Contribution #140)

• Payload components (AthenaMP) of the conventional event service and Yoda are absolutely identical V.Tsulaia, ATLAS, CHEP 2015

Running at large scale



Future developments: shared I/O workers



Summary

- AthenaMP has successfully passed the physics validation procedure and now is being widely used for running ATLAS workloads on the Grid
- Originally developed for optimizing memory footprint of reconstruction jobs, AthenaMP later became an efficient mechanism for processing fine grained workloads
- Recent tests on supercomputers demonstrated, that AthenaMP can run efficiently within distributed applications at large scale