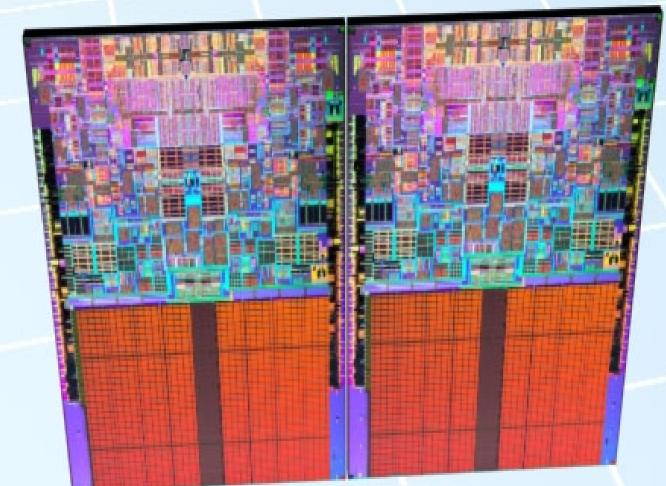


Core Performance



Paolo Calafiura Lawrence Berkeley National Lab CHEP 09 – March 28 2009



In this talk...

Efficient use of CPU and memory – Performance Monitoring and Optimization Unix programming environment – i686, Linux, C/C++

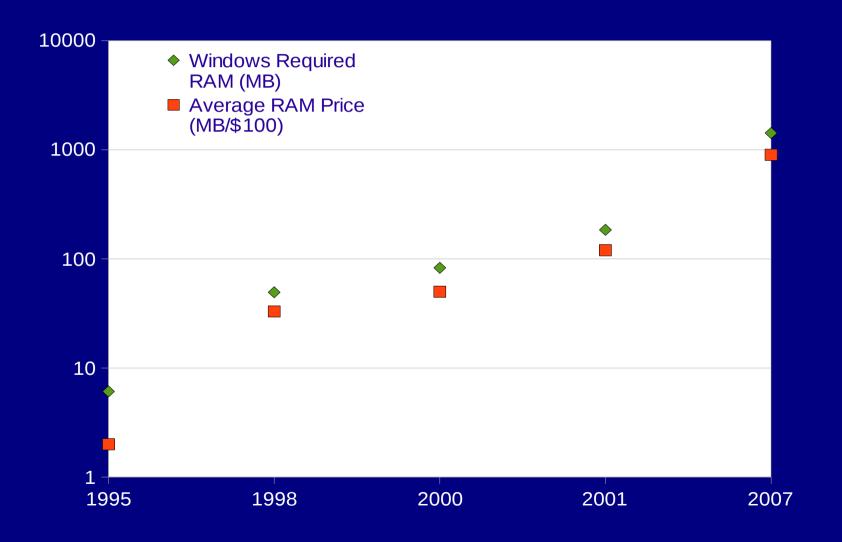
"Rules of Optimization:

- 1. Don't do it
- 2. (for experts only) Don't do it yet "
- (M. A. Jackson)
- "More computing sins are committed in the name of efficiency (without necessarily achieving it) than for any other single reason – including blind stupidity"
 - (W. A. Wulf)
- "Preliminary optimization is the root of all evil"

(Donald Knuth)

Is Efficiency Effective?

Software Horror Vacui

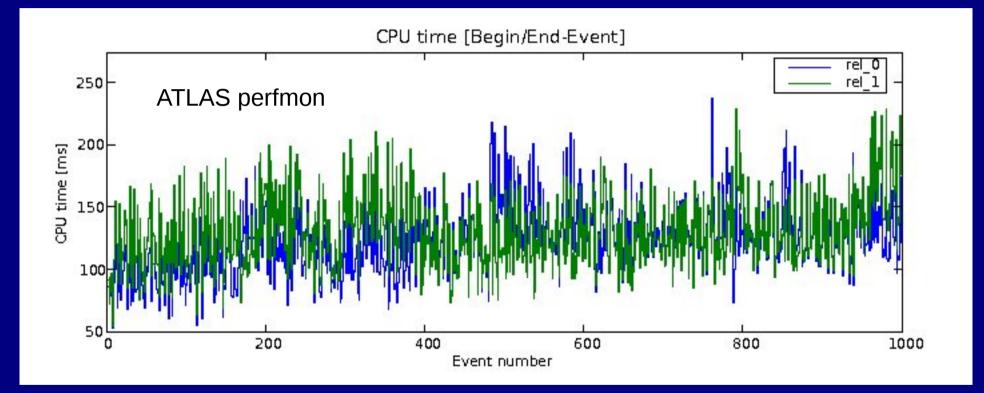


Hungrier than Vista

LHC reconstruction pushing against 2GB/core typical GRID node

Cai Guo-Qiang

Performance Monitoring

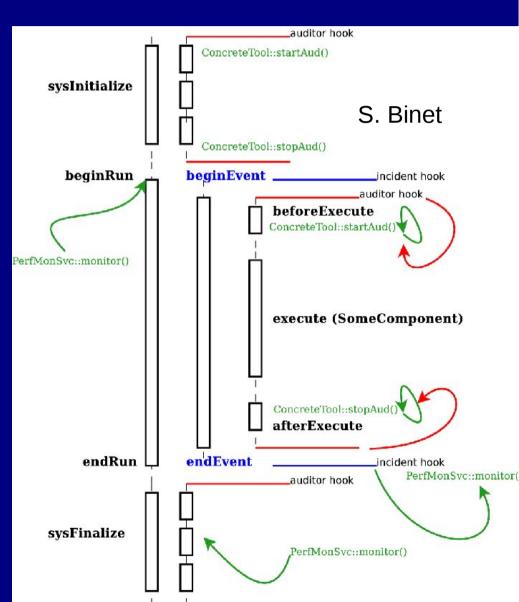


Information is the most valuable commodity (Gordon Gekko)

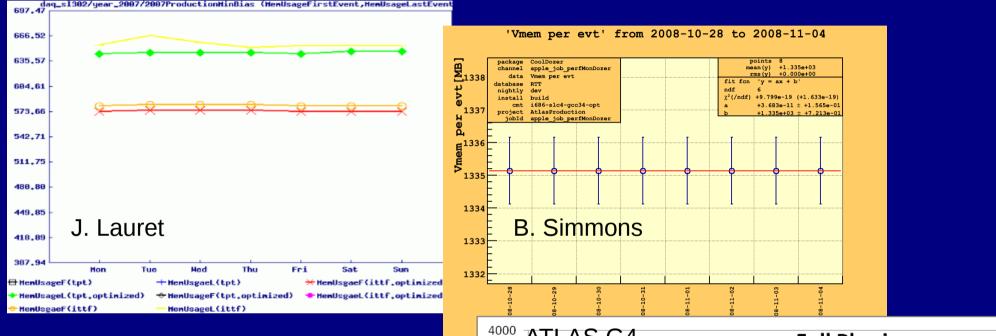
Performance Monitoring in Practice

- Integrate into framework Collect data per component
 - CPU
 - user/sys/wall
 - Memory
 - VMEM, RSS, leaks, mallocs

– Developer-defined

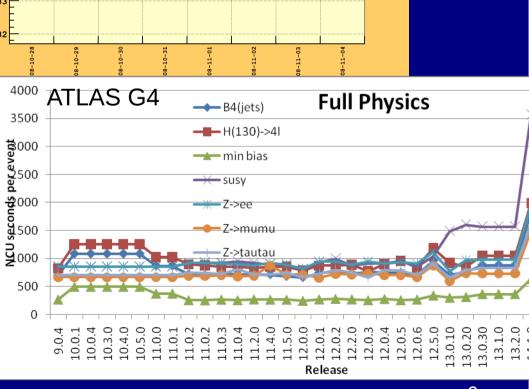


Track Performance Evolution

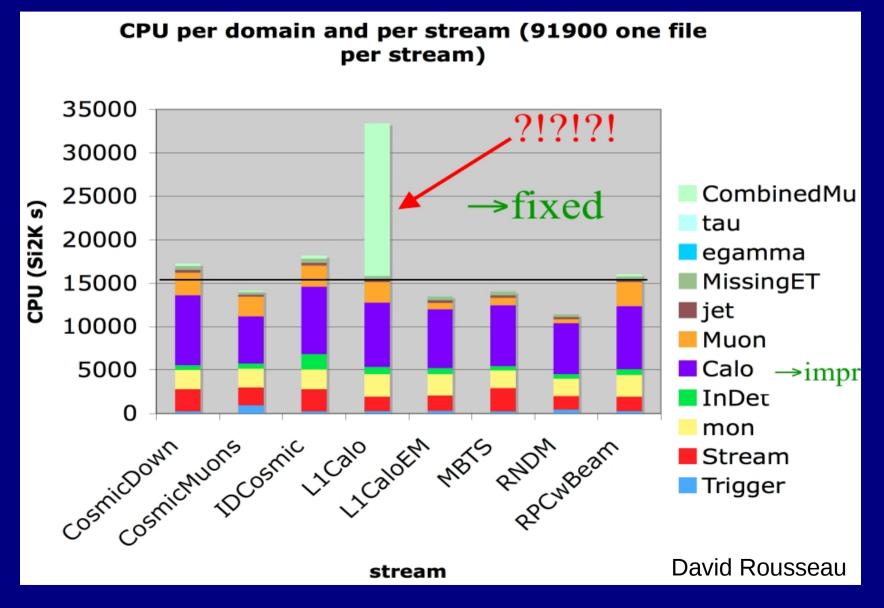


2008-1

Best tool for release coordinators to spot new issues



Actionable Information



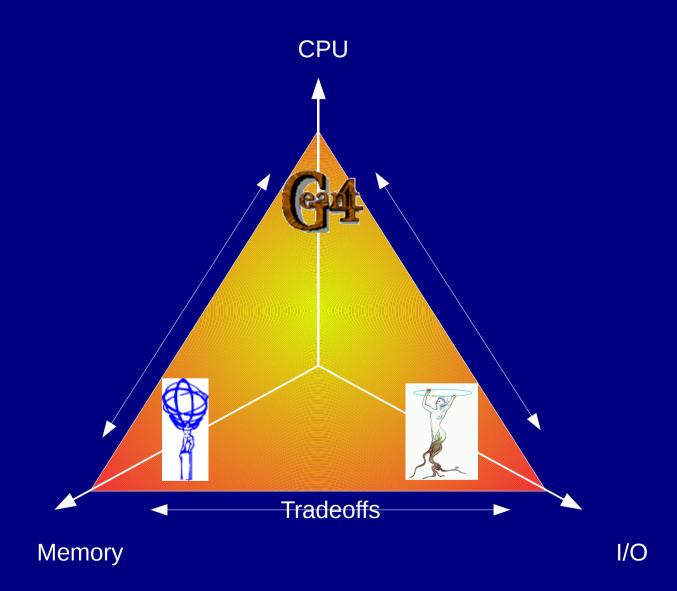
Performance Monitoring for Managers

Start early, never stop Monitor what runs in production, all of it Track evolution, intervene immediately Turn measurements into action items Raise profile of exercise, put physicists in charge

Performance Optimization

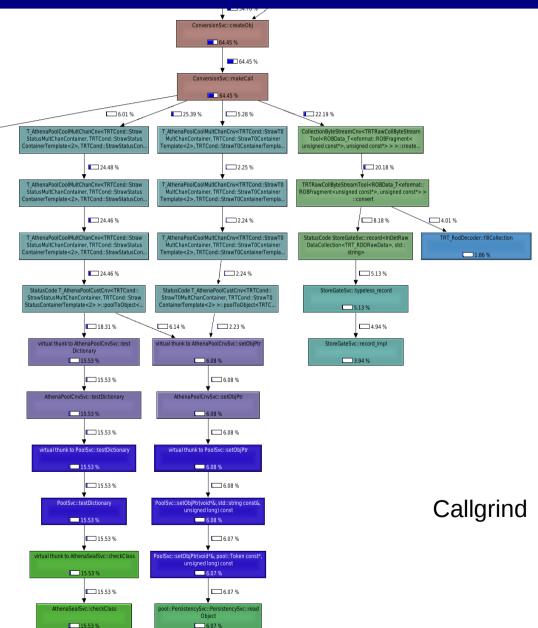
Monday Parallel Session 3 Pete Elmer – CMS Software Performance Giulio Eulisse – HEP C++ meets reality Thursday Parallel Session 12 Andrzej Nowak – An update on perfmon and the struggle to get it into the Linux kernel

Optimize What?



CPU Optimization

Tools provide Usage graph Line-by-line breakdown First 10% "easy" pass by value temporaries sub-optimal containers





CPU utilization



Three broad classes of changes which affect the CPU performance:

- 1 <u>Physics algorithmic</u> the program output changes
 - Extra cuts, simulating extra or more detailed effects, "Do we run the extra trackfinder?"
- 2 <u>Algorithmic, but technical</u> the program output does not change
 - Caching, lazy evaluation, removal of redundant calculations, data structures, etc.
- 3 <u>Purely technical</u> the program output does not change
 - Changes related to specific issues in C++, the memory management, the operating system, the compilers and the hardware where are applications run

By far the largest gains/losses come from #1, of course. Improvements from #2 and #3 are "free beer" in that there are no trade-offs with physics (though there may be trade-offs between CPU use and memory, etc.)

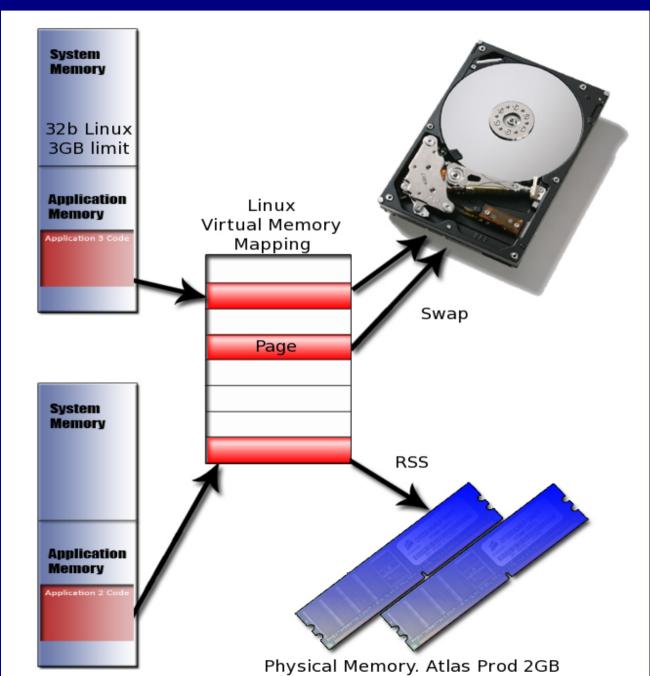
P. Elmer

CHEP 09 Collaboration Meeting - Prague

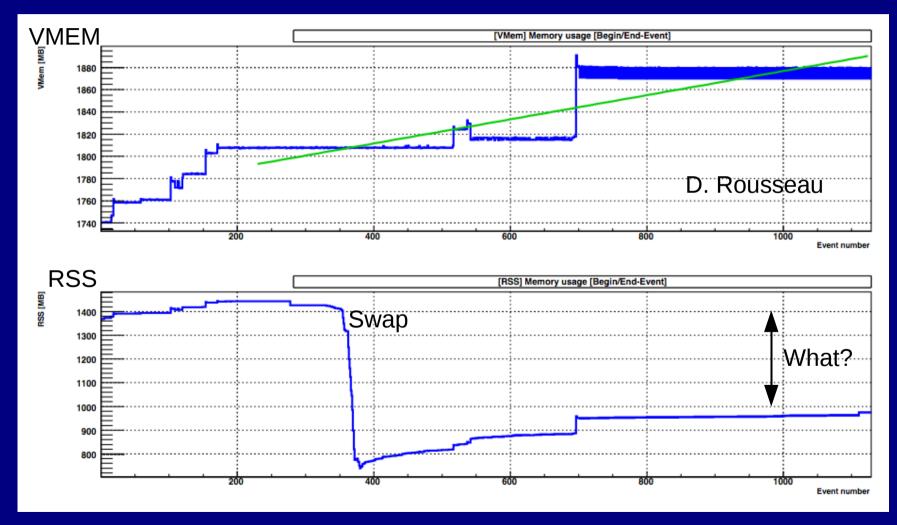
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Memory Optimization

Linux Memory Management

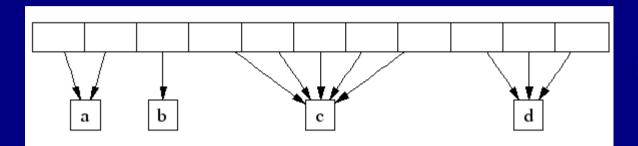


Swap in Action

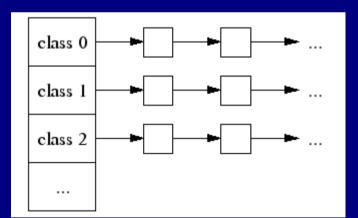


Event Number

Heap Allocation (tcmalloc)



Heap set of pages

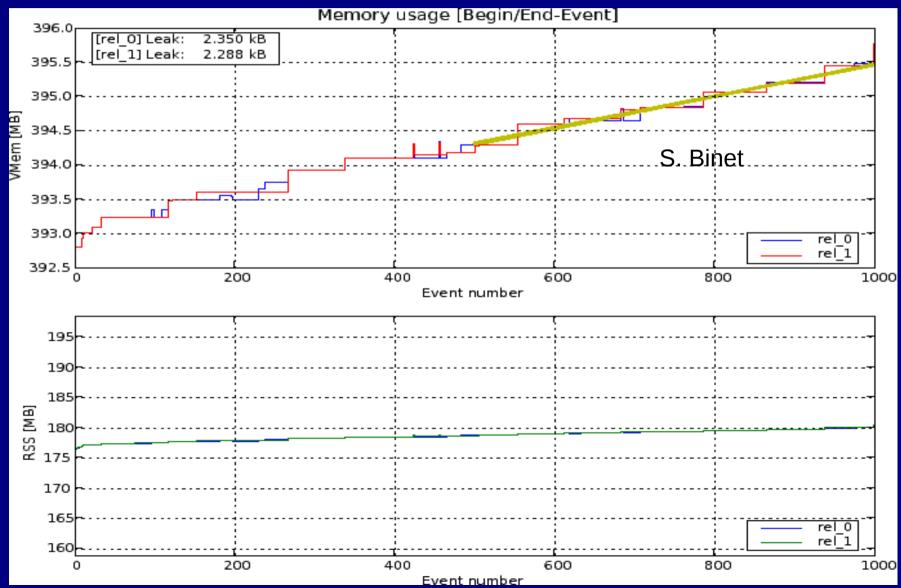


Size object pools

Allocation: – Grab free entry from list Deallocation

- Return slot to list
- Return memory "lazily"

Memory Leaks



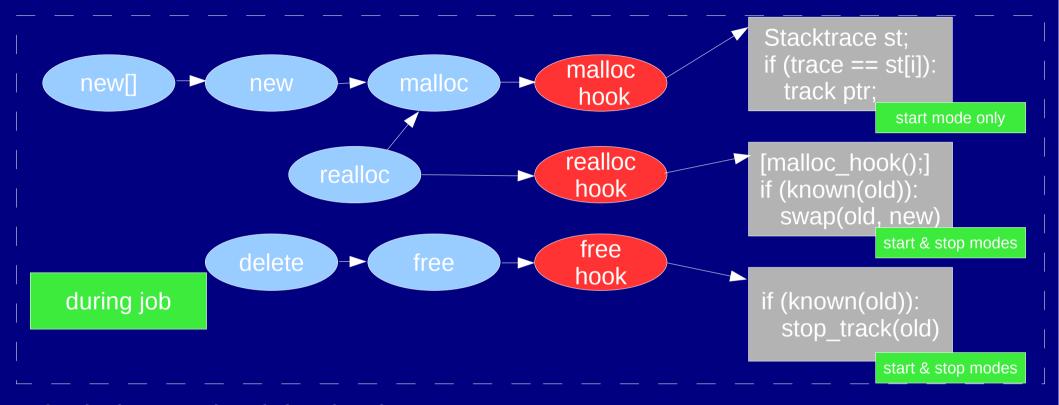
Fighting Leaks: Valgrind

Interpreter tracks mallocs/frees Reports leak candidates

==29199== 300 bytes in 15 blocks are definitely lost in loss record 1006 of 1301 ==29199== at 0x3414B6D6: operator new(unsigned) (vg_replace_malloc.c:133) ==29199== by 0x3F79D4F0: Trk::TrackSummaryTool::createSummary(Trk::Track const&) (in /afs/ dist/nightlies/rel/atlrel_1/Tracking/TrkTools/TrkTrackSummaryTool/TrkTrackSummaryTool-00-11-0 opt/libTrkTrackSummaryToolLib.so) ==29199== by 0x3F7BBDA0: InDet::PriVxTopAlg::m_preselect(Trk::Track const*) (in /afs/cern.ch /nightlies/rel/atlrel_1/InnerDetector/InDetRecAlgs/InDetPriVxFinder/InDetPriVxFinder-01-00-01/i6 opt/libInDetPriVxFinder.so)

Slow(~10x), memory hungry(2-3x)

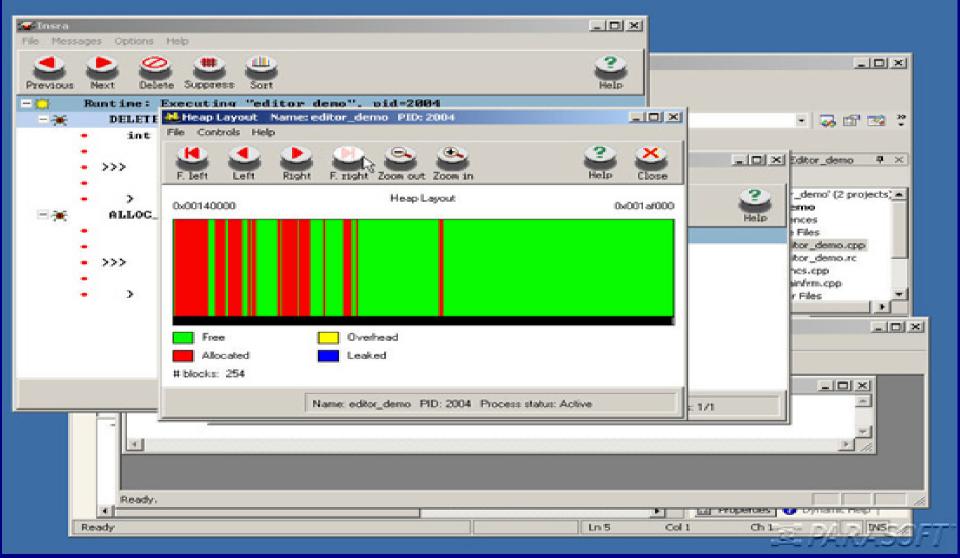
Lightweight: Hephaestus Wim Lavrijsen



leak detected, originating in: CaloClusterBuilderSE::CreateImpactInCalo (Trk::Track const*) (28 bytes) 0x8272ef2 CaloClusterBuilderSE::CreateImpactInCalo(Trk::Track const*)

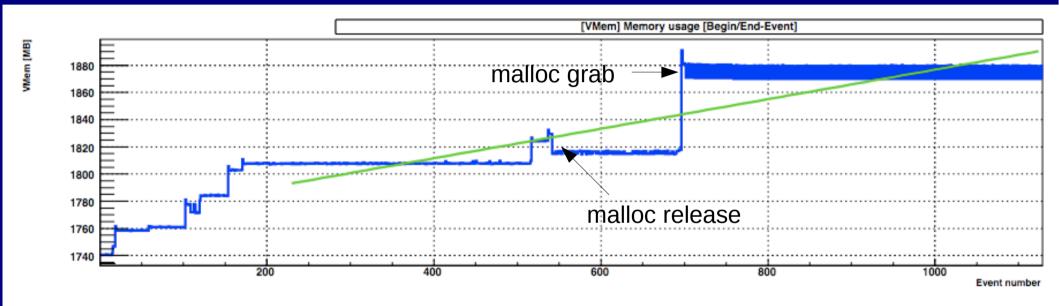
Requires preloading

Deluxe: insure++



Requires special compilation

VMEM on the Rise



Heap Fragmentation

Possibly with different lifetimes

Visualizing Fragmentation (Totalview Memory Debugger)

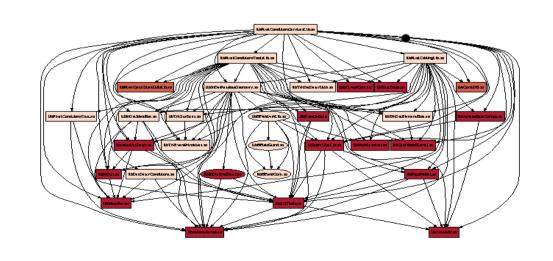
| Configuration Leak Detecti Options Detect Leaks Detective | | |
|--|---|---|
| > filterapp(11574) | | Í |
| > 0x0804b934 - 0x0809c000 (3 | (321.70KB) | |
| Heap Information Backtra | Memory block:TypeAllocatedFilteredNoSize5.10KBStart Address0x0804c968 | |
| Category Category Category Category Corrupted Gua Corrupted Gua Corrupted Gua Corrupted Gua Contect Co | End Address 0x0804ddcf Backtrace ID 9 Point of allocation: File stl_alloc.h Methoddefault_alloc_template::_S_chunk_alloc Line 490 Guard Blocks: None | |
| S Contractions of the second s | Backtrace 13 Backtrace 13 Backtrace | |

Fighting Fragmentation

ATLAS ~10% memory lost to fragmentation Always reserve memory for containers Do not new small objects If you have to (polymorphism), use a segregated allocator (e.g. memory pool) Do not mix long-lived and short-lived allocations If you have to, consider using a private heaps

Fighting Library Bloat

Play with gcc symbols visibility Reduce symbol duplications templates, statics Optimize packaging studying dependencies



Summary

Monitor Core Performance

- Always
- Using the code that runs in production
 Optimize Core Performance
 - When you must
 - Guided by tools
 - Be aware of trade offs
 - Know when to stop

Thanks

D. Brown, J. Lauret, B. Simmons, K. Ciba,
M. Clemencic, M. Cattaneo, J. Apostolakis,
F. Carminati, R. Brun, F. Rademakers,
J. Kowalkowski, M. Paterno, S. Snyder,
L. Sexton-Kennedy, D. Olson, C. Tull,
S. Binet, D. Rousseau, W. Lavrijsen,
M. Tatarkhanov, S. Jarp, A. Nowak,
P. Elmer, Y. Yao, F. Winklmeier

Fork, COW, and vmem

Production systems must adapt to advanced memory optimization

- Athena/Gaudi MP use fork (and COW) to dramatically reduce real memory usage by aggressively sharing pages among evt workers
- Linux KSM promises to do this automatically
- We can run 4 RecExCommon instances in a 4GB machine without swapping

No use if jobs get killed based on (wrong) vmem statistics

Optimization in the Multi/Many-core Era Scaling currently limited by I/O Tilt CPU/Memory balance further Cache conflicts increasingly important Sharing memory pages crucial Production systems must adapt