

PetaCache: Data Access Unleashed

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Outline

- Motivation Is there a Problem?
- Economics of Solutions
- Practical Steps Hardware/Software
- Some Performance Measurements



Motivation

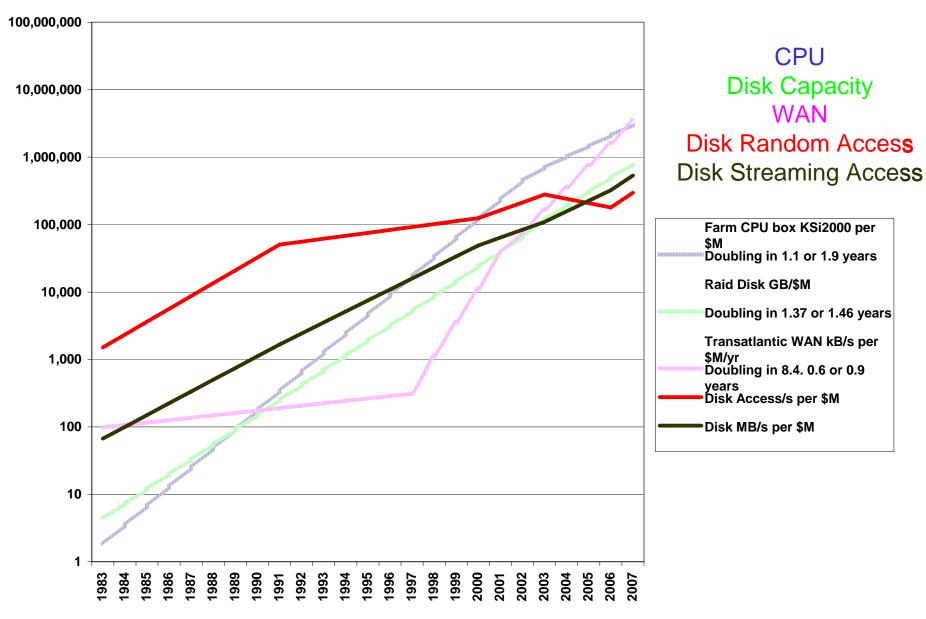
Storage In Research: Financial and Technical Observations

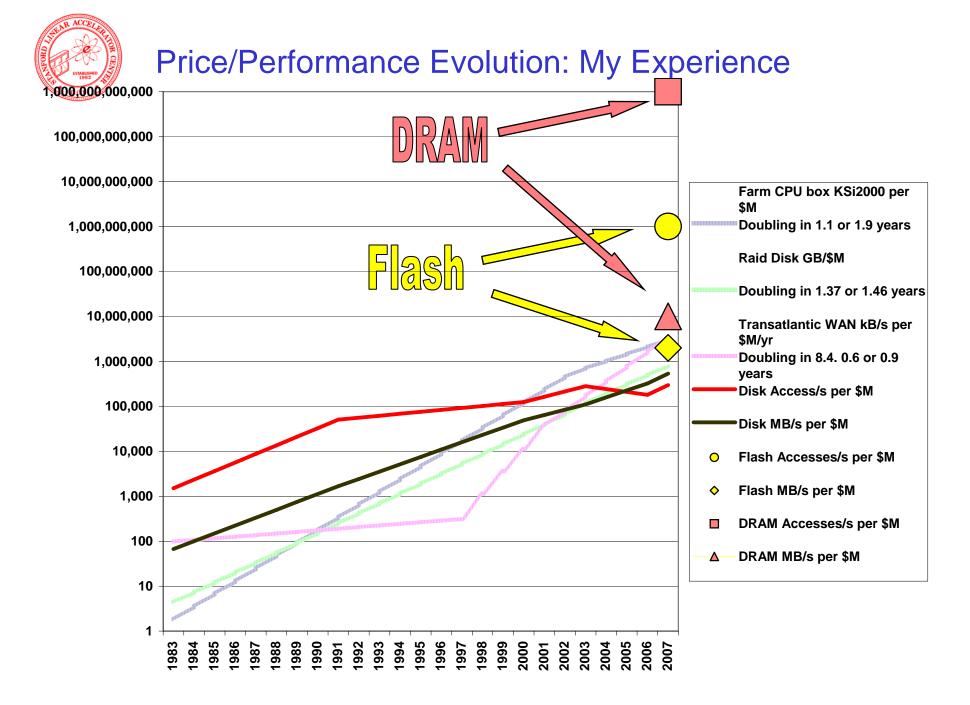
- Storage costs often dominate in research
 - CPU per \$ has fallen faster than disk space per \$ for most of the last 25 years
- Accessing data on disks is increasingly difficult
 - Transfer rates and access times (per \$) are improving more slowly than CPU capacity, storage capacity or network capacity.
- The following slides are based on equipment and services that I* have bought for dataintensive science

^{*} The WAN services from 1998 onwards were bought by Harvey Newman of Caltech



Price/Performance Evolution: My Experience







Another View

- In 1997 \$M bought me:
 - ~ 200-core CPU farm (~few x 108 ops/sec/core)

or

- ~ 1000-disk storage system (~2 x 10³ ops/sec/disk)
- Today \$1M buys me (you):
 - ~ 2500-core CPU farm (~few x 109 ops/sec/core)

or

- ~ 2500-disk storage system (~2 x 10³ ops/sec/disk)
- In 5 10 years?



Impact on Science

- Sparse or random access must be derandomized
- Define, in advance, the interesting subsets of the data
- Filter (skim, stream) the data to instantiate interest-rich subsets



Economics of Solutions

conomics of LHC Computing

- Difficult to get \$10M additional funding to improve analysis productivity
- Easy to re-purpose \$10M of computing funds if it would improve analysis productivity



Cost-Effectiveness

DRAM Memory:

- \$100/gigabyte
- SLAC spends ~12% of its hardware budget on DRAM
- Disks (including servers)
 - \$1/gigabyte
 - SLAC spends about 40% of its hardware budget on disk
- Flash-based storage (SLAC design)
 - \$10/gigabyte
 - If SLAC had been spending 20% of its hardware budget on Flash we would have over 100TB today.



Practical Steps

The PetaCache Project



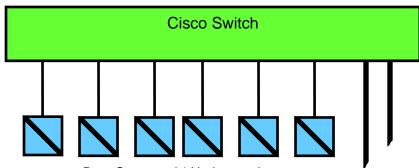
PetaCache Goals

- Demonstrate a revolutionary but cost effective new architecture for science data analysis
- Build and operate a machine that will be well matched to the challenges of SLAC/Stanford science

The PetaCache Story So Far

- We (BaBar, HEP) had data-access problems
- We thought and investigated
 - Underlying technical issues
 - Broader data-access problems in science
- We devised a hardware solution
 - We built a DRAM-based prototype
 - We validated the efficiency and scalability of our low-level dataaccess software, xrootd
 - We set up a collaboration with SLAC's electronics wizards (Mike Huffer and Gunther Haller) to develop a more cost-effective Flashbased prototype
- We saw early on that new strategies and software for data access would also be needed

DRAM-Based Prototype Machine (Operational early 2005)



Data-Servers 64 Nodes, each
Sun V20z, 2 Opteron CPU, 16 GB memory
1TB total Memory
Solaris or Linux (mix and match)

PetaCache MICS + HEP-BaBar Funding



DRAM-Based Prototype





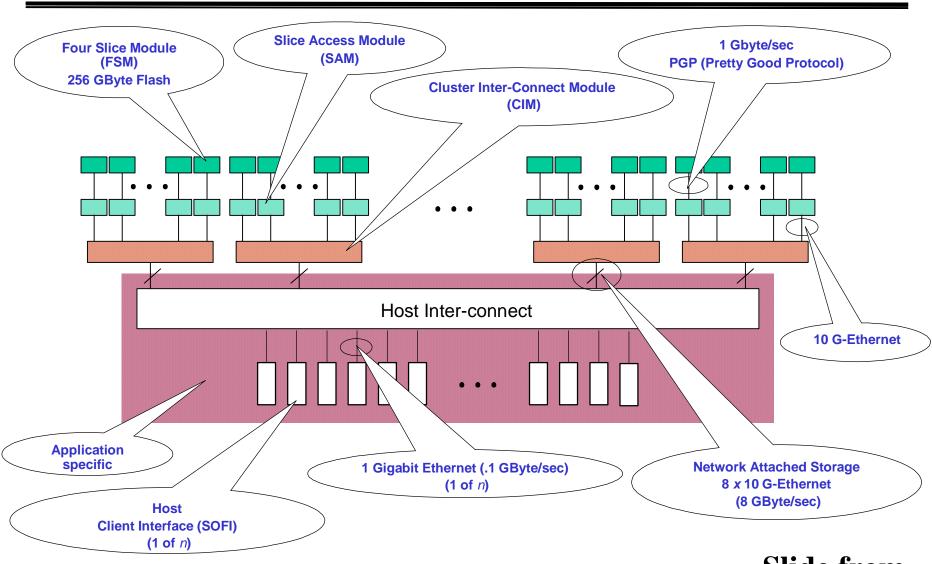
FLASH-Based Prototype Operational Real Soon Now

- 5 TB of Flash memory
- Fine-grained, high bandwidth access





Building Blocks

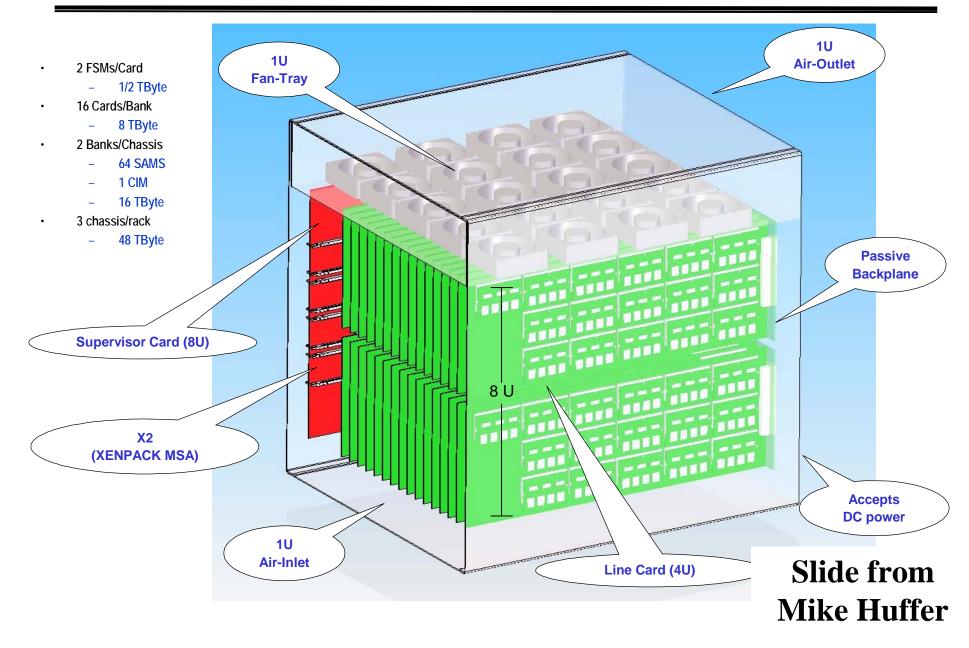


Slide from Mike Huffer





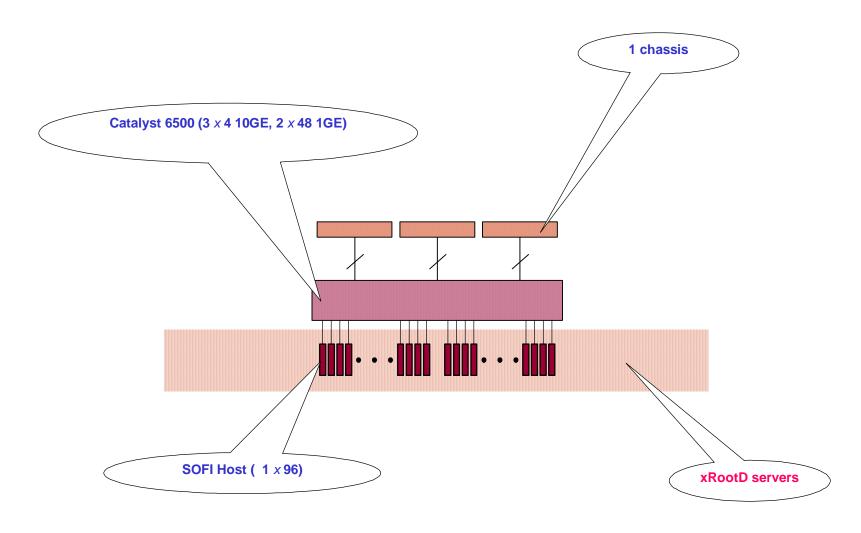
The "Chassis"







48 TByte facility



Slide from Mike Huffer



Commercial Product

Violin Technologies

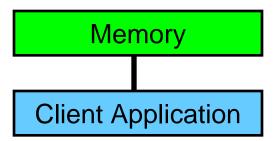
- 100s of GB of DRAM per box (available now)
- TB of Flash per box (available real soon now)
- PCle hardware interface
- Simple block-level device interface
- DRAM prototype tested at SLAC



Some Performance Measurements

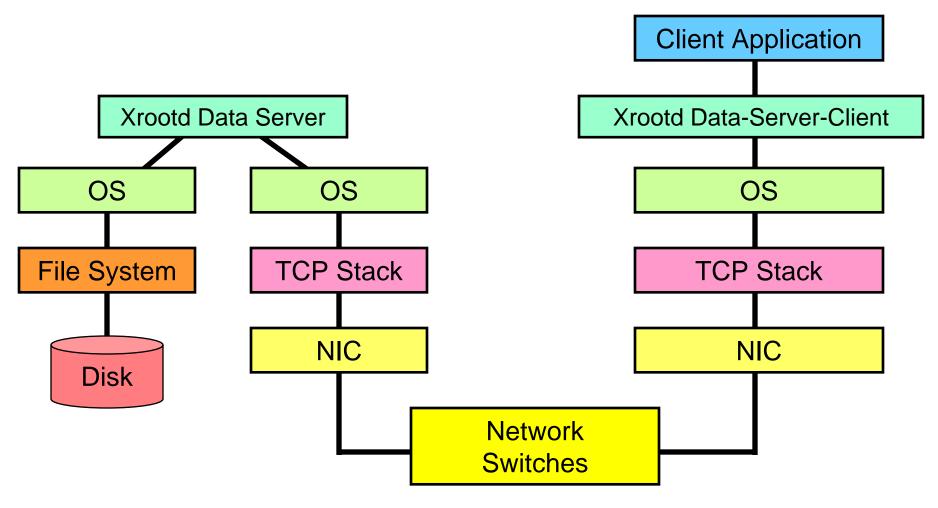


Latency (1) Ideal





Latency (2) Current reality

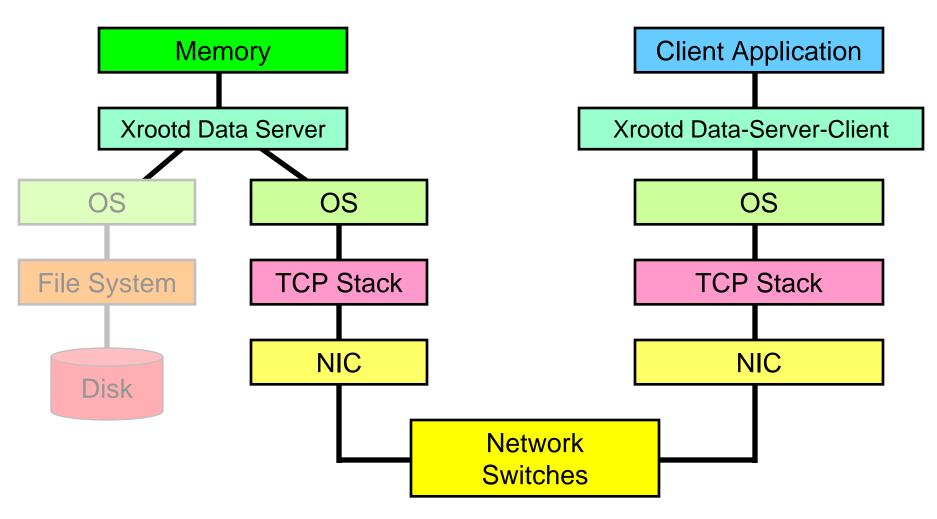


Richard P. Mount, SLAC

March 7, 2007



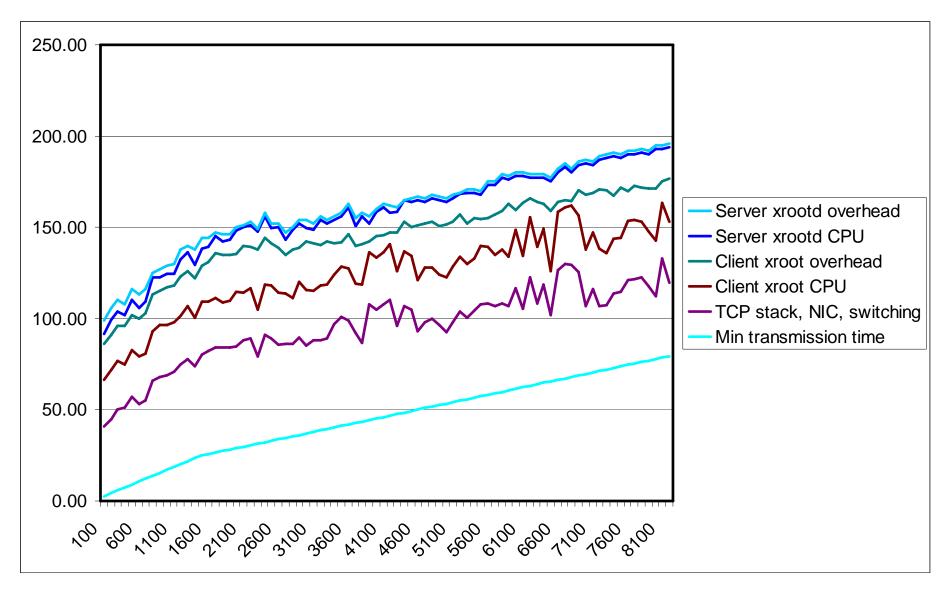
Latency (3) Immediately Practical Goal



Richard P. Mount, SLAC

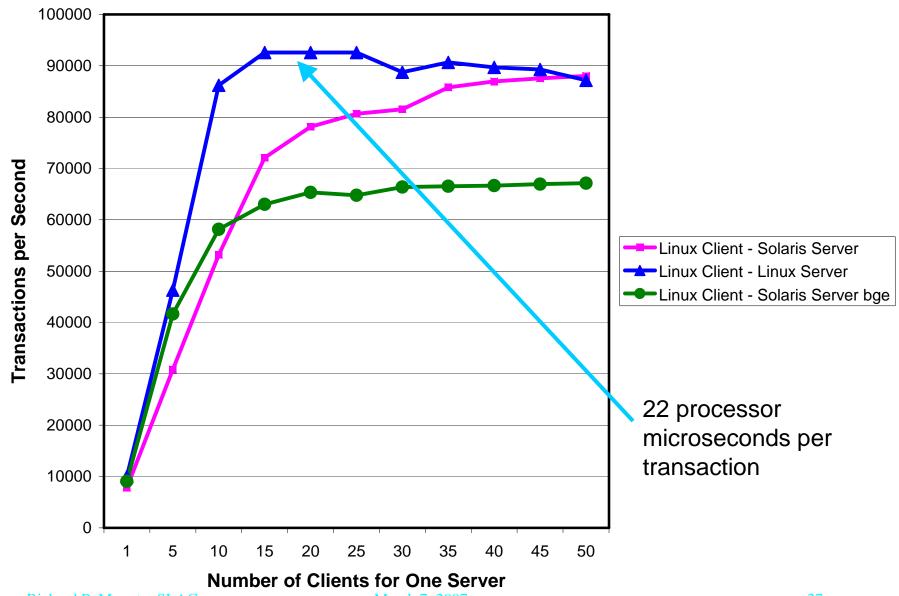
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DRAM-Based Prototype Latency (microseconds) versus data retrieved (bytes)





DRAM-Based Prototype Throughput Measurements



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27



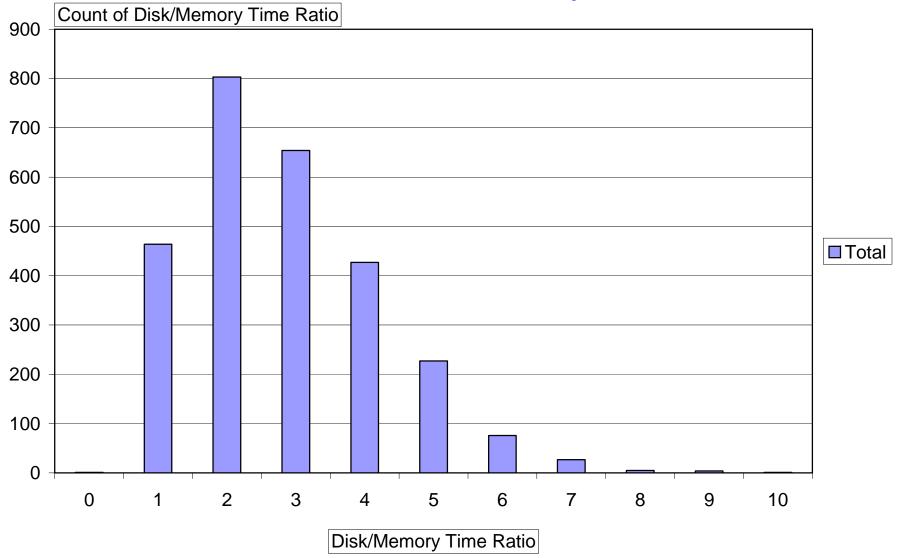
Throughput Tests

ATLAS AOD Analysis

- 1 GB file size (a disk can hold 500 1000 of these)
- 59 xrootd client machines (up to 118 cores)
 performing top analysis getting data from 1 server.
- The individual analysis jobs perform sequential access.
- Compare time to completion when server uses its disk, compared with time taken when server uses its memory.



DRAM-based Protoype ATLAS AOD Analysis





Comments and Outlook

- Significant, but not revolutionary, benefits for high-load sequential data analysis – as expected.
- Revolutionary benefits expected for pointerbased data analysis – but not yet tested.
- The need to access storage in serial mode has become part of the culture of dataintensive science – why design a pointerbased analysis when its performance is certain to be abysmal?
- TAG database driven analysis?