

# **Unified Storage Systems for Distributed Tier-2 Centres**

Greig A. Cowan, Graeme A. Stewart, Andrew Elwell University of Edinburgh & University of Glasgow





Greig A Cowan Distributed access to storage CHEP 2007

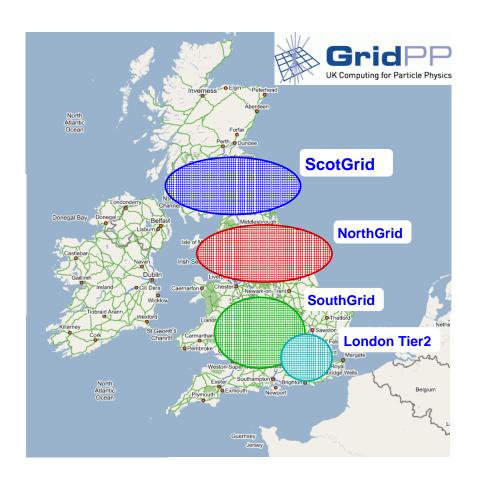


### **Overview**

- 1. Using Tier-2 resources efficiently.
  - Accessing data across the WAN
- 2. Testing setup
- 3. Results
  - (a) File open times, read times and rates
  - (b) Server response
- 4. Future work
- 5. Summary



## **Distributed Tier-2s**



- GridPP organised into four regional Tier-2s.
  - Helps with deployment and operations.
  - Cross-site support.
- Can we do better on a technical level?
  - Can we pool resources to use them more efficiently?
    - \* Storage at one site could be regarded as being "close" to CPU of another.

Greig A Cowan Distributed access to storage CHEP 2007



### **Access to data**

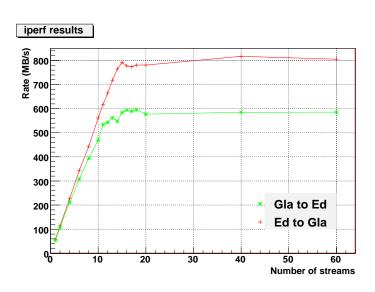
- Currently, jobs are sent to the compute element which is local to the data.
  - Users running a selection algorithm over a dataset.
- Often more efficient for the jobs to process the data directly on the SE.
  - Use POSIX-like protocols (rather than copying entire file to the WN).
    - \* rfio for CASTOR and DPM (with gsi)
    - \* (gsi) dcap for dCache
    - \* ROOT provides TGFALFile to allow access to these SEs on the grid.
- Problem: if batch farm where data is located is full, then jobs cannot run.
  - Other sites in Tier-2 may have spare capacity.
    - \* Inefficiency in system.



## Access to data across the WAN

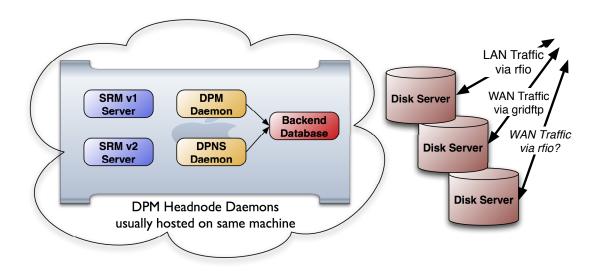
- Can we use the POSIX protocols to access storage across the wide area network?
  - Will this be transparent to users?, i.e.,
    - \* Can they access data at the same rate?
    - \* Does the efficiency of their jobs remain the same?
- ullet Production JANET-UK network between UKI-SCOTGRID-GLASGOW and ScotGrid-Edinburgh. RTT  $\sim 12$ s.







## Tier-2 storage

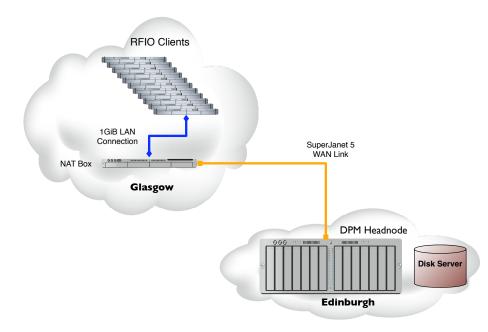


- DPM developed by EGEE as a lightweight solution for disk storage management at Tier-2 institutes.
- See other talks/posters at CHEP07 for further details.



# **Testing method**

- We wrote our own RFIO client application.
  - Reading data appears to be the main use case.
  - Configurable to meet needs of our study, i.e.,
    - \* RFIO mode
    - \* read block size
    - \* reading pattern (sequential, skipping, random)
    - \* Allows us to stress the SE.
  - "Skipping" means that we read a block of data, then skip ahead  ${\cal M}$  blocks and read again, until EOF.
- ullet Seed client onto N nodes and simultaneously start reading 1GB files from ScotGRID-Edinburgh DPM.





## **RFIO** protocol

- GSI-enabled protocol which allows POSIX file operations, permitting byte-level access to files.
  - clients require a X.509 Grid certificate signed by a trusted CA.
  - can use RFIO over the wide area network.
  - Ports must be opened in site firewall.
- RFIO library allows the client to choose from four modes of operation (see rfiosetopt() man page):
- 1. NORMAL: one call per read.
- 2. RFIO\_READBUF: fills internal buffer to service requests.
- 3. RFIO\_READAHEAD: uses internal buffer and reads until EOF.
- 4. RFIO\_STREAM: separate TCP streams for control and data.



## **RESULTS**

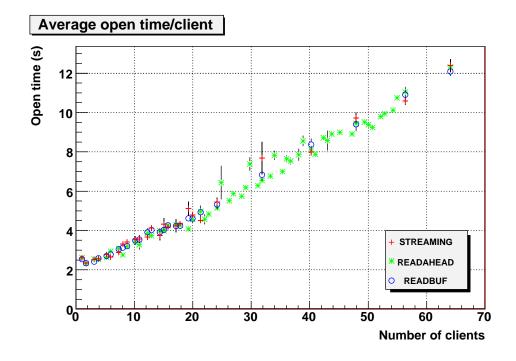
Greig A Cowan Distributed access to storage CHEP 2007



# File open times

### Sequential reading

- Linear increase in the open time with client number.
- ullet Large number of clients can increase open times up to  $>12\mathrm{s}$ .

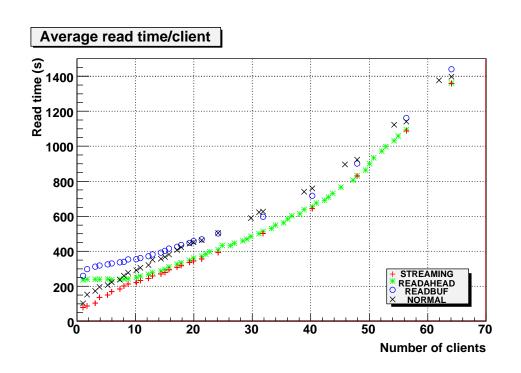


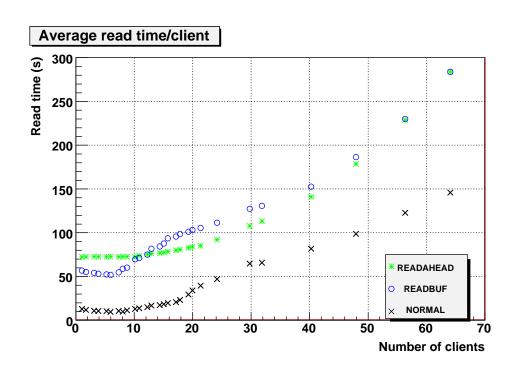


## **Read times**

#### **Sequential reading**

### Skipping through 10% of the file





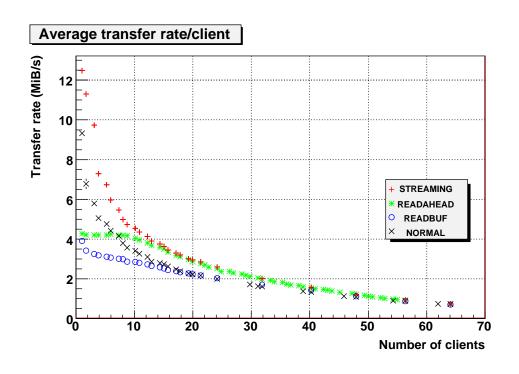
- LHS: STREAMING comes out on top for small number of clients. Not much difference for large number.
- RHS: NORMAL mode leads to optimal access.

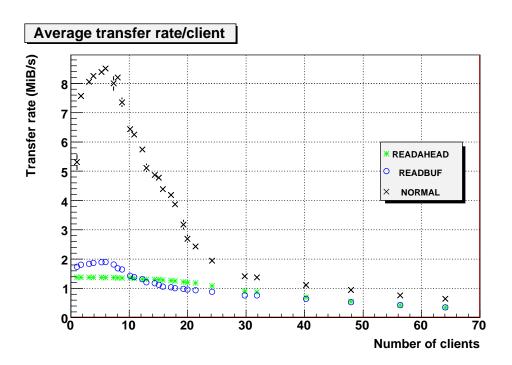


## **Read rates**

#### Sequential reading

### Skipping through 10% of the file

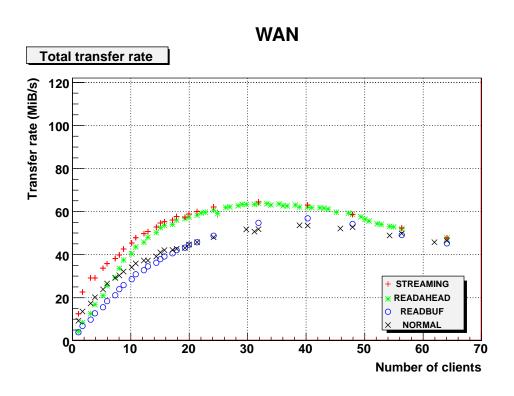


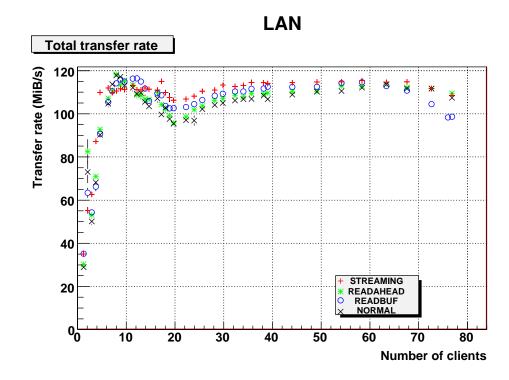


- Large number of clients, rates down to  $\sim$ 1MiB/s per job (NB single DPM server).
  - ATLAS software expects rates of O(10)MiB/s per job in 2008.



# **Comparison with LAN access**



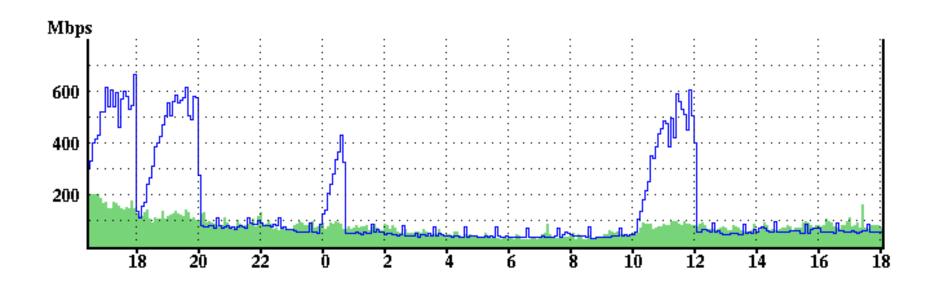


- Peak total rate across WAN  $\sim$ 65MiB/s.
  - Contention on the network. Max expected 100MiB/s.
- Peak total rate across LAN  $\sim$ 110MiB/s.
  - Single server. Dedicated bendwidth.

• Becomes **IO-bound** at a large number of clients, rate begins to decrease.



# Impact on network



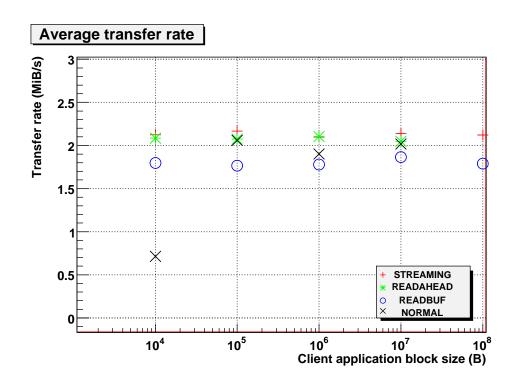
• Background traffic on the production network <100Mbps.

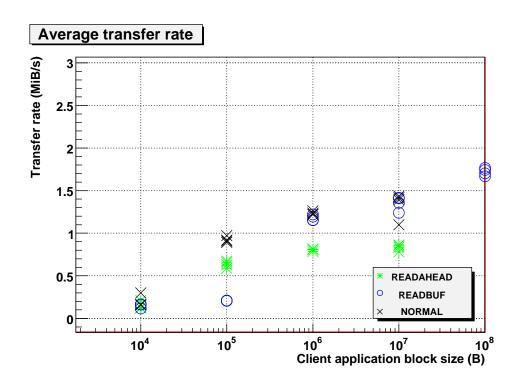


# **Application block size**

### Sequential reading

#### Skipping through 10% of the file





- No change when sequentially reading the file.
- When skipping through the file, higher rates achieved with larger block sizes, particularly for READBUF mode.



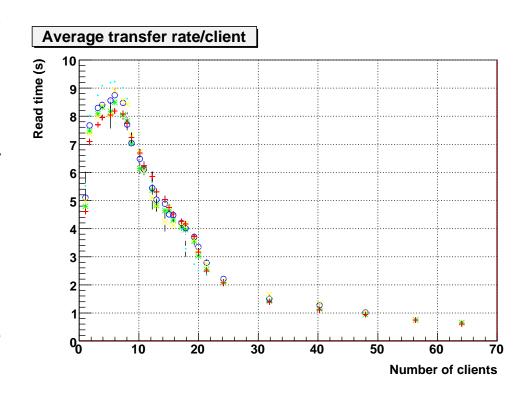
## **Client TCP parameters**

- Since we are moving data across the WAN, TCP kernel parameters **could** have a impact on the data throughput.
- Initial work looked at increasing the maximum TCP window size.
- e.g., in /etc/sysctl.conf we varied parameters such as, net.ipv4.tcp\_rmem and net.core.rmem\_max.
- Looked at increasing window sizes from 0.5MB up to 16MB.



# Variation with client TCP parameters

- Different colours correspond to different TCP window sizes.
- Very little difference.
  - Probably expected when such a large number of clients are simultaneously reading data.
  - Slight improvement at small client numbers with a larger window.
- Application optimisations probably required before tuning the networking parameters.

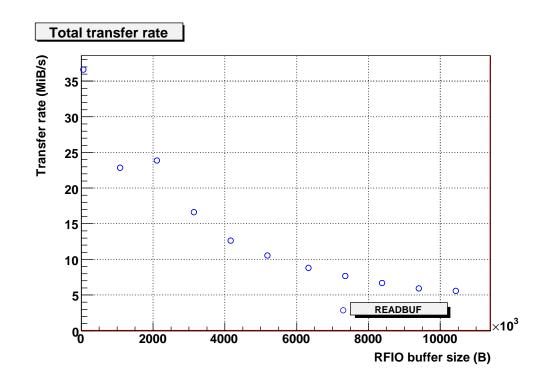




## Transfer rate vs. RFIO buffersize

#### Skipping through 10% of the file

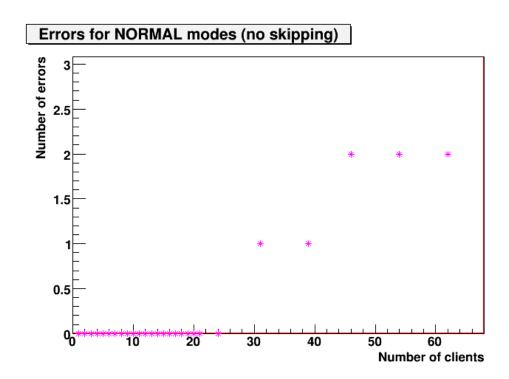
- RFIO READBUF mode uses a fixed size client side buffer for data transfer.
  - Parameter is RFIO IOBUFSIZE in /etc/shift.conf.
  - Can we see any dependence on the size of the buffer?
  - Plot shows that for a constant block size of 1MB, increasing the RFIO buffer leads to a reduced total transfer rate.





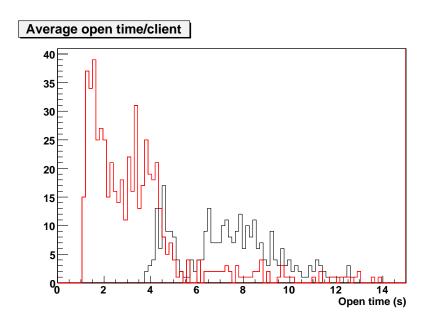
### File access errors

- Server performance degrades slightly when many clients simultaneously attempt to open files.
  - We are intentionally stressing the system.
- Substantial improvement over versions of DPM <</li>
  1.6.5, which could not support more than ~40 opens per second.

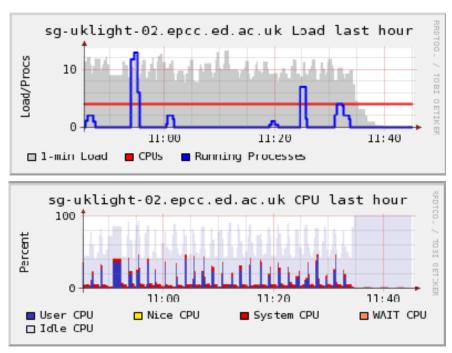




## **Server load**



Difference in the open time for < 20 clients (red) and  $\geq 20$  clients (black).



Load on single DPM server when 30 clients are simultaneously reading using NORMAL RFIO mode.



### **Future Work**

- Would like to repeat tests using lightpath as this gives:
  - Dedicated bandwidth that will not impact on other users.
  - Smaller RTT, of order 2ms.
- Alternative data access patterns. Extreme cases.
  - i.e., Use 1 client to open 1000 files on the SE and then send them to sleep (3000).
- Run some real analysis jobs.
  - ROOT TTreeCache will allow efficient data access across WAN.
    - \* See talk 284 at CHEP07
- Create a single DPM that spans both Glasgow and Edinburgh sites.



## **Summary**

- Using DPM and RFIO, our study has shown that it is possible to access storage across the WAN.
- This opens up possibilities for optimising storage and CPU usage within ditributed Tier-2s.
- Principle could be more widely applied to the Grid.
  - Rather than having many replicas of files spread over the Grid, closely linked sites could access a replica within their geographical region.
    - \* ATLAS already have a cloud model for data management...
- Studied the ideal cases where clients were behaving as expected.
  - How does the system respond in non-optimal cases?
- We saw good utilisation (60%) of the production network, but these rates may not be sufficient when large numbers of clients are running.
  - Need to investigate the potential of dedicated lightpaths.