

CERN and Computing and Storage

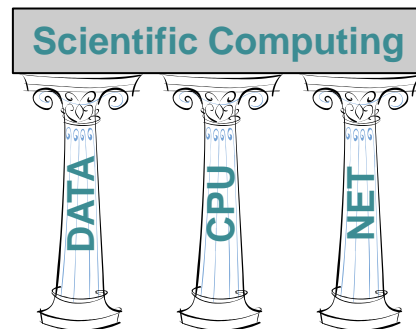
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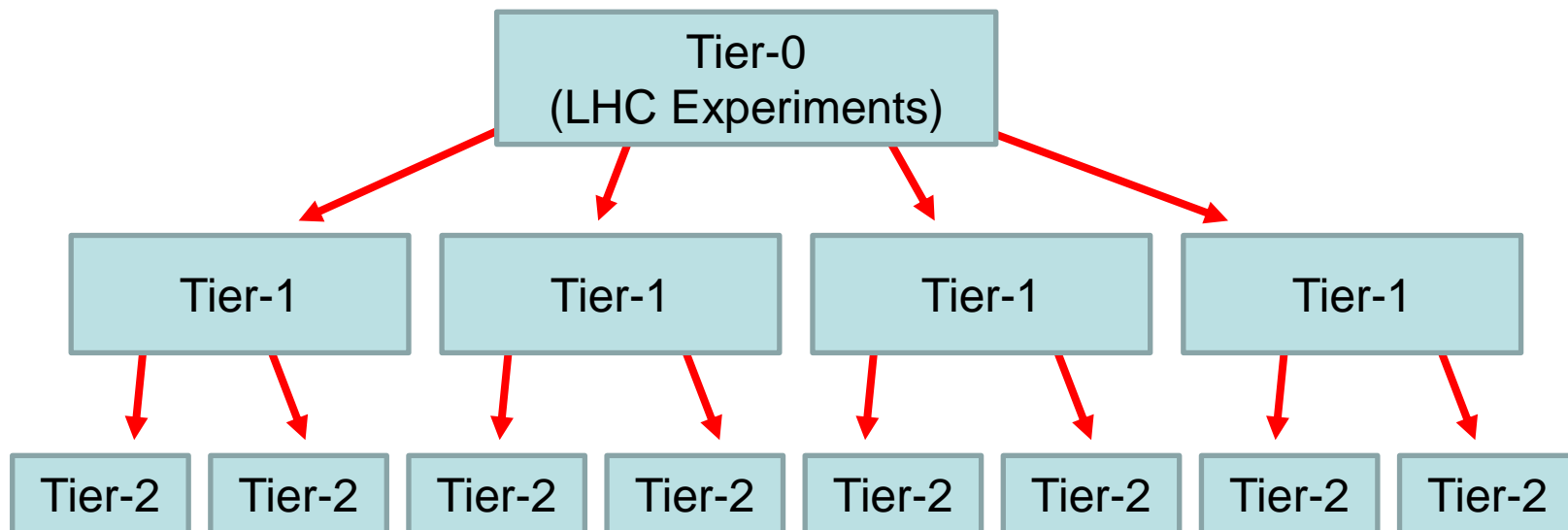
Big Data processing and analysis challenges in mega-science
experiments

**January 29-31, 2015,
Moscow, Dubna, Russia**

- Scientific computing for large experiments is typically based on a distributed infrastructure
- Storage is one of the main pillars
- A key parameter in distributed scientific computing is the efficiency
 - High efficiency requires to have the CPUs colocated with the data to analyze, using the network.



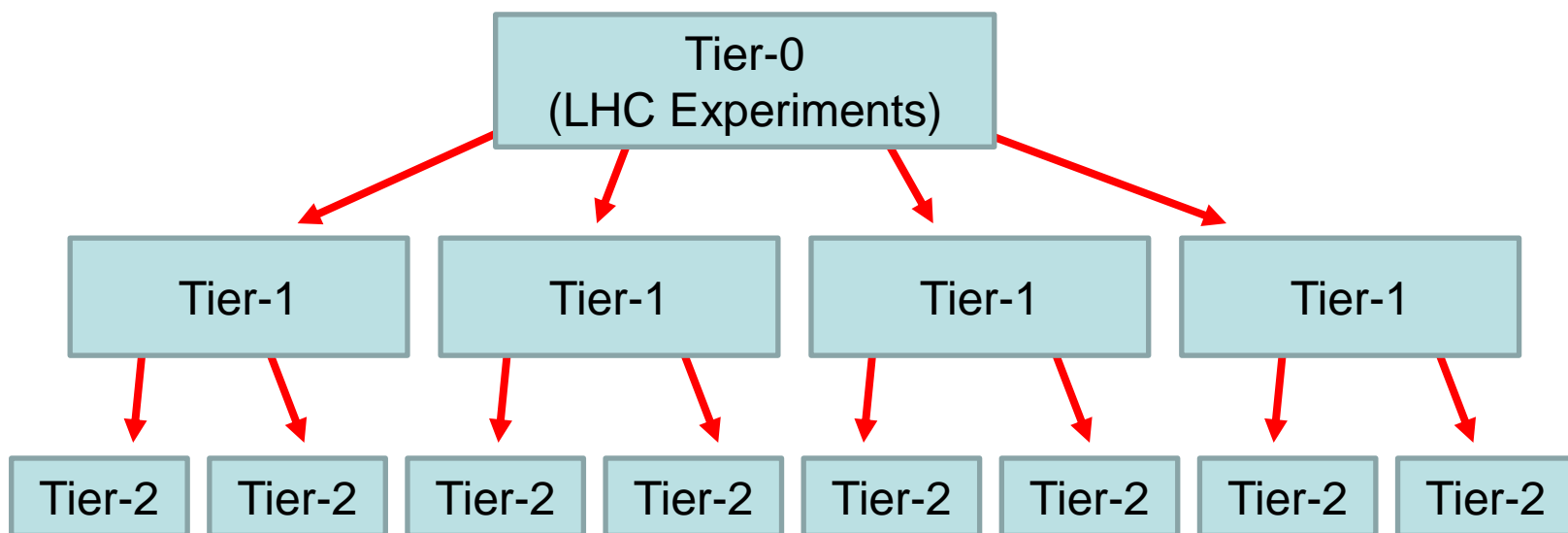
- Storage in scientific computing is distributed across multiple data centres
- Data flows from the experiments to all datacenters where there is CPU available to process the data



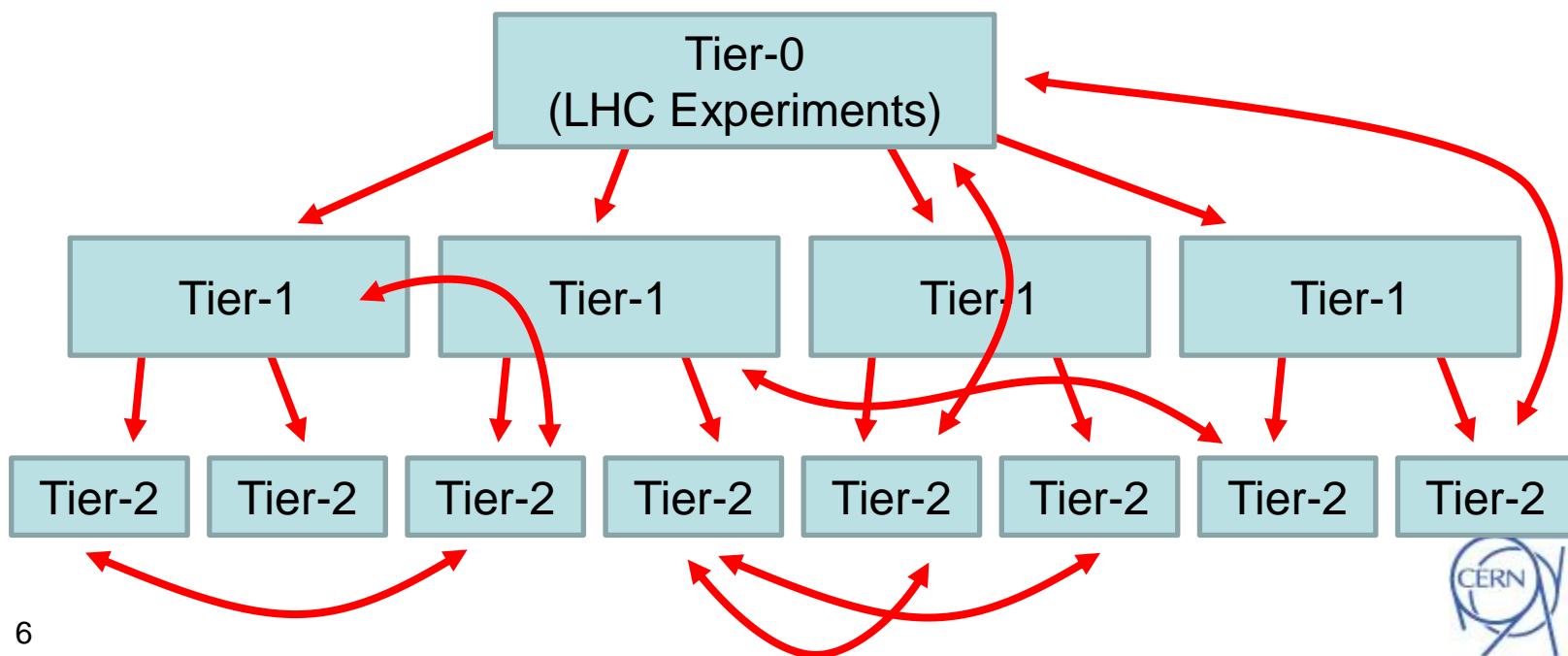
- Whenever a site has ...
 - idle CPUs (because no Data is available to process)
 - or excess of Data (because there is no CPU left for analysis)
- ... the efficiency drops



- Analysis made with high efficiency requires the data to be pre-located to where the CPUs are available

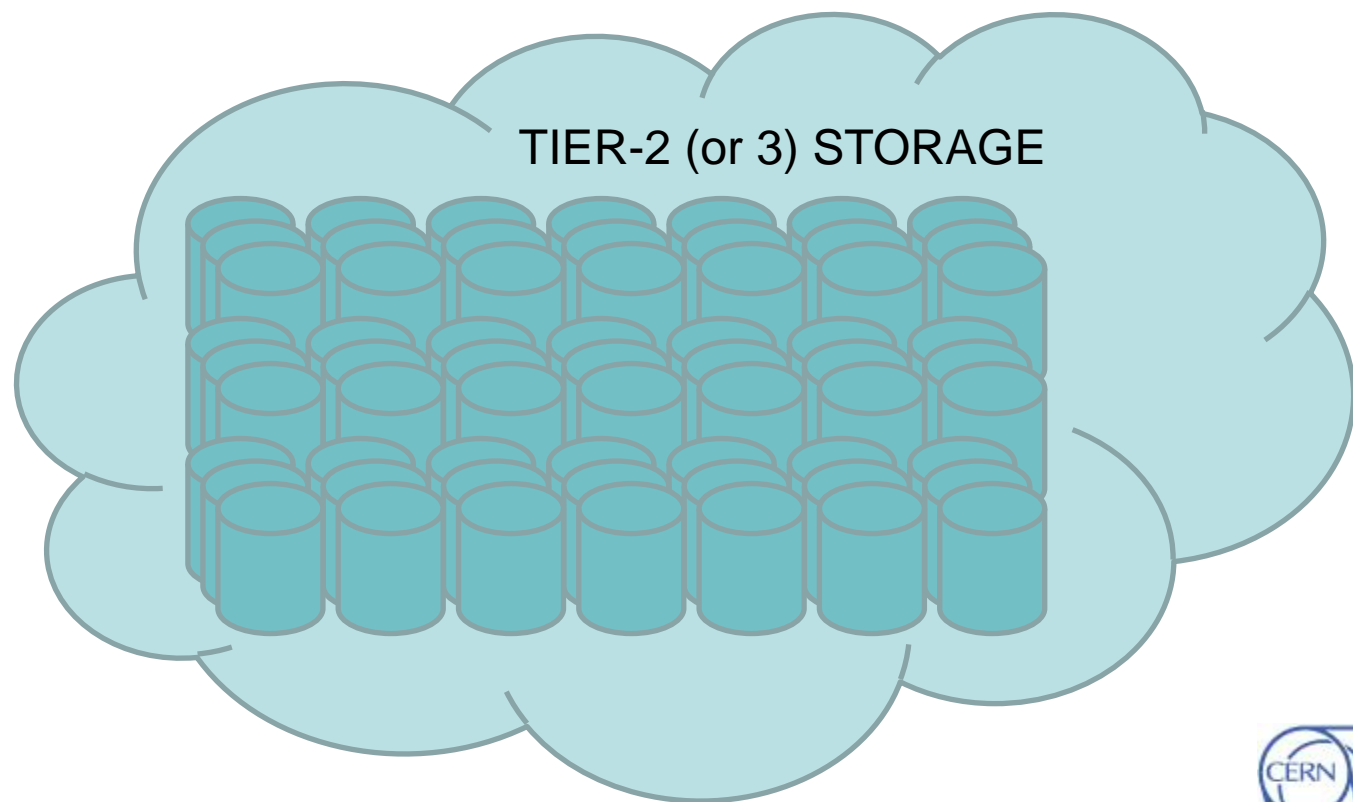


- Analysis made with high efficiency requires the data to be pre-located to where the CPUs are available
- Or to allow peer-to-peer data transfer
 - This allows sites with excess of CPU, to schedule the pre-fetching of data when missing locally or to **access it remotely** if the analysis application has been designed to cope with high latency

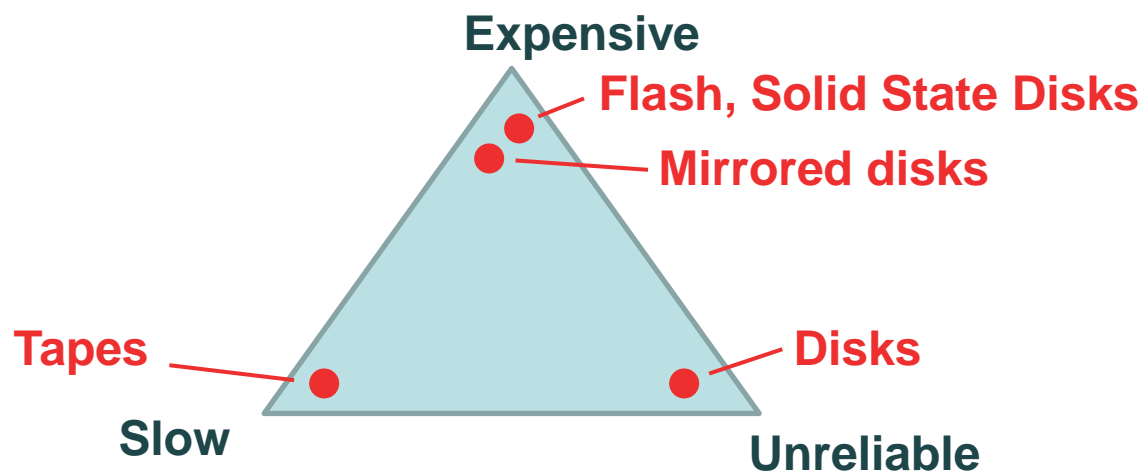


- Both approaches coexists in WLCG
- Data is pre-placed
 - This is the role of the experiments that plans the analysis
- Data is globally accessible and federated in a global namespace
 - The middleware always attempt to take the local data and uses an access protocol that redirects to the nearest remote copy when the local data is not available
 - All middleware and jobs are designed to minimize the impact of the additional latency that the redirection requires
- Using access protocols that allows global data federation is essential
 - http, xroot

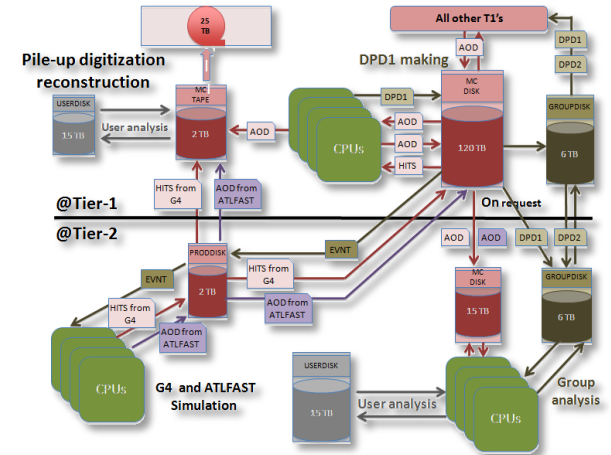
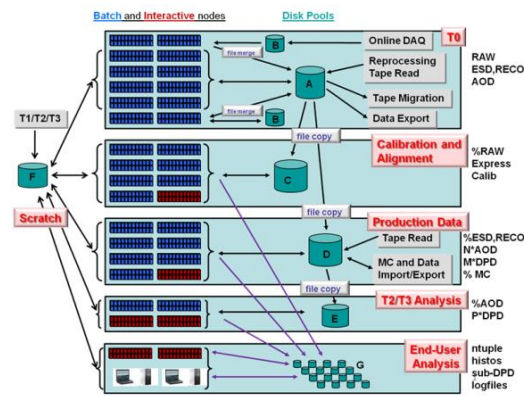
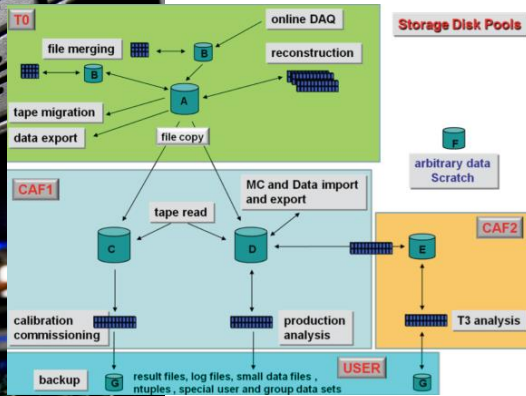
- A simple storage model: all data into the same storage
 - Uniform, simple, **easy to manage, no need to move data**
 - Can provide sufficient level of performance and reliability



- Different storage solutions can offer different quality of services
 - Three parameters: Performance, Reliability, Cost
 - You can have two but not three
- To deliver both performance and reliability you must deploy expensive solutions
 - Ok for small sites (you save because you have a simple infrastructure)
 - Difficult to justify for large sites

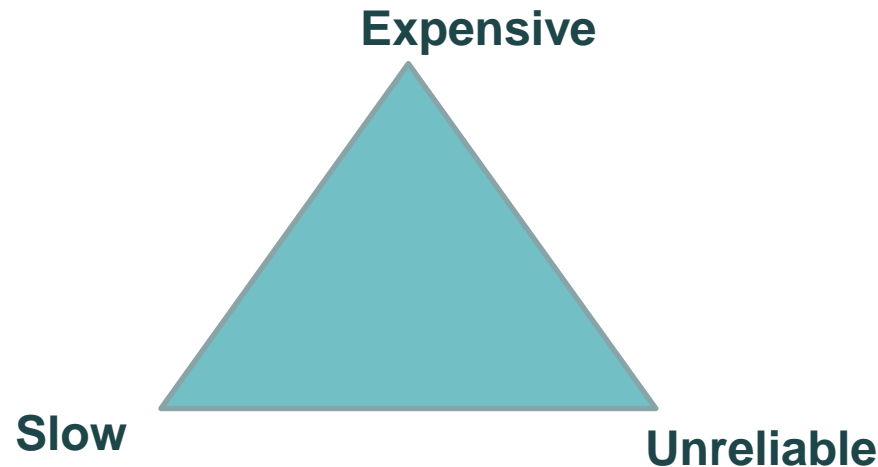


- Examples from LHC experiment data models



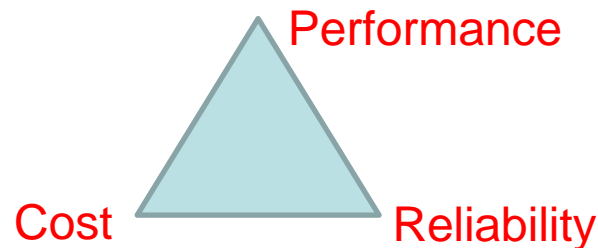
- Two building blocks to empower data analysis
 - Data pools with different quality of services
 - Tools for data transfer between pools

- Storage need to be able to adapt to the changing requirement of the experiments
 - Cover the whole area of the triangle and beyond

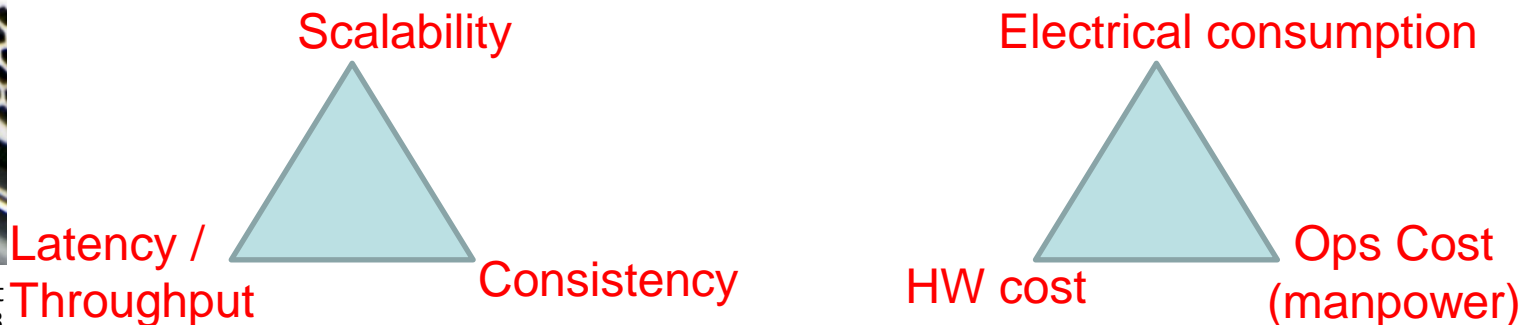


- **Why multiple pools with different quality of services are needed ?**
- Derived data used for analysis and accessed by thousands of nodes
 - Need high performance, Low cost, minimal reliability (derived data can be recalculated)
- Raw data that need to be analyzed
 - Need high performance, High reliability, can be expensive (small sizes)
- Raw data that has been analyzed and archived
 - Must be low cost (huge volumes), High reliability (must be preserved), performance not necessary

- Many ways to split (performance, reliability, cost)



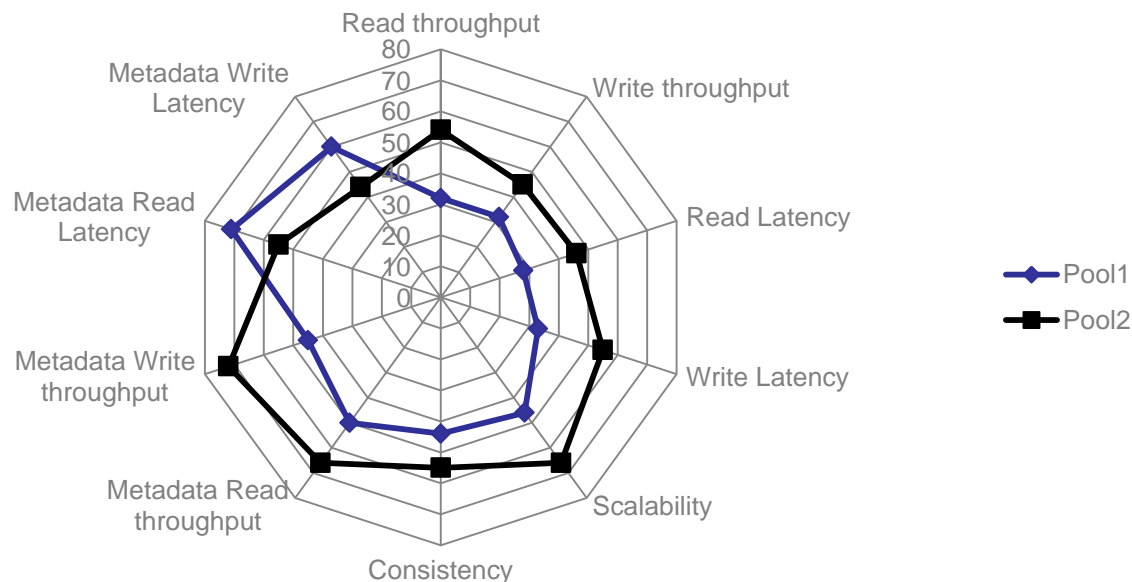
- Performance has many sub-parameters
- Cost has many sub-parameters
- Reliability has many sub-parameters





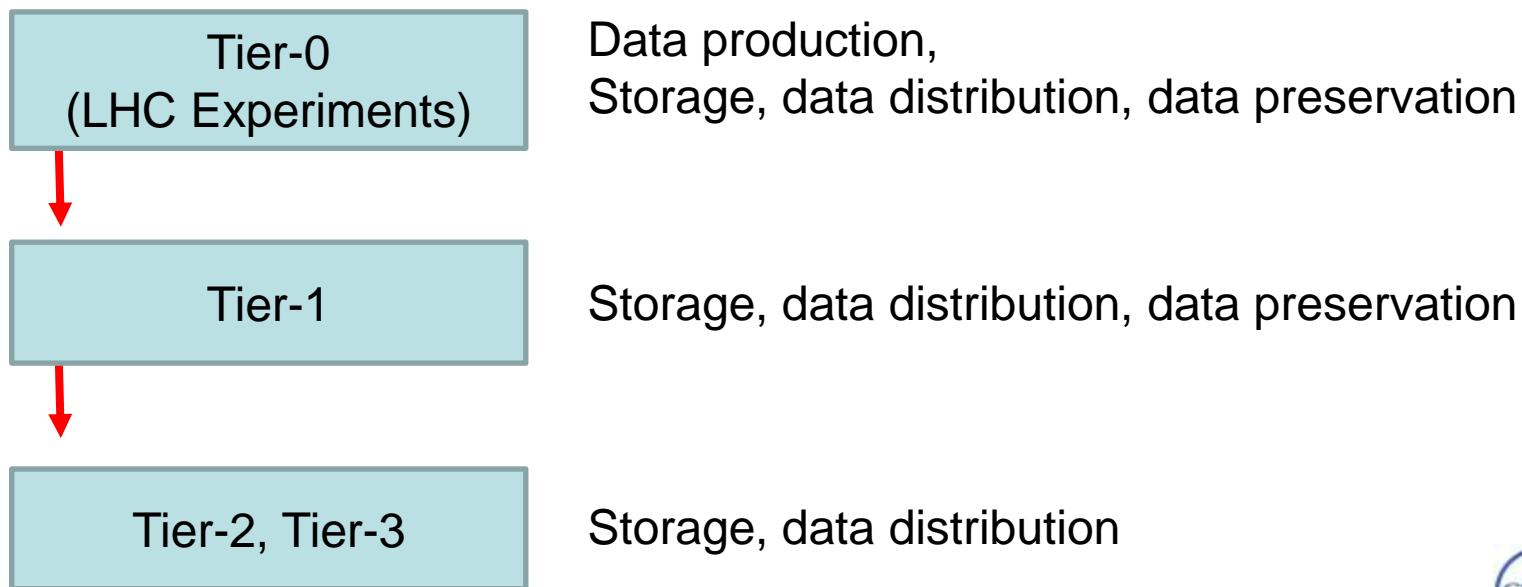
DSS And reality is complicated

- Key requirements: Simple, Scalable, Consistent, Reliable, Available, Manageable, Flexible, Performing, Cheap, Secure.
- Aiming for “à la carte” services (storage pools) with on-demand “quality of service”
- And where is scalability ?



- Needed:
 - Tools to transfer data effectively across pools of different quality
 - Storage elements that can be reconfigured “on the fly” to meet new requirements without moving the data
- Examples
 - Moving petabytes of data from a multiuser disk pool into a reliable tape back-end
 - Increasing the replica factor to 5 on a pool containing condition data requiring access from thousands of simultaneous users
 - Deploying petabytes of additional storage in few days

- Three main roles
 - Storage (store the data) *Size in PB*
 - Distribution (ensure that data is accessible) *Availability*
 - Preservation (ensure that data is not lost) *Reliability*



- You can achieve high reliability using standard disk pools
 - Multiple replicas
 - Erasure codes
- Here is where tapes can play a role
 - Tapes ??



- Tapes have a bad reputation in some use cases
 - Slow in random access mode
 - high latency in mounting process and when seeking data (F-FWD, REW)
 - Inefficient for small files (in some cases)
 - Comparable cost per (peta)byte as hard disks



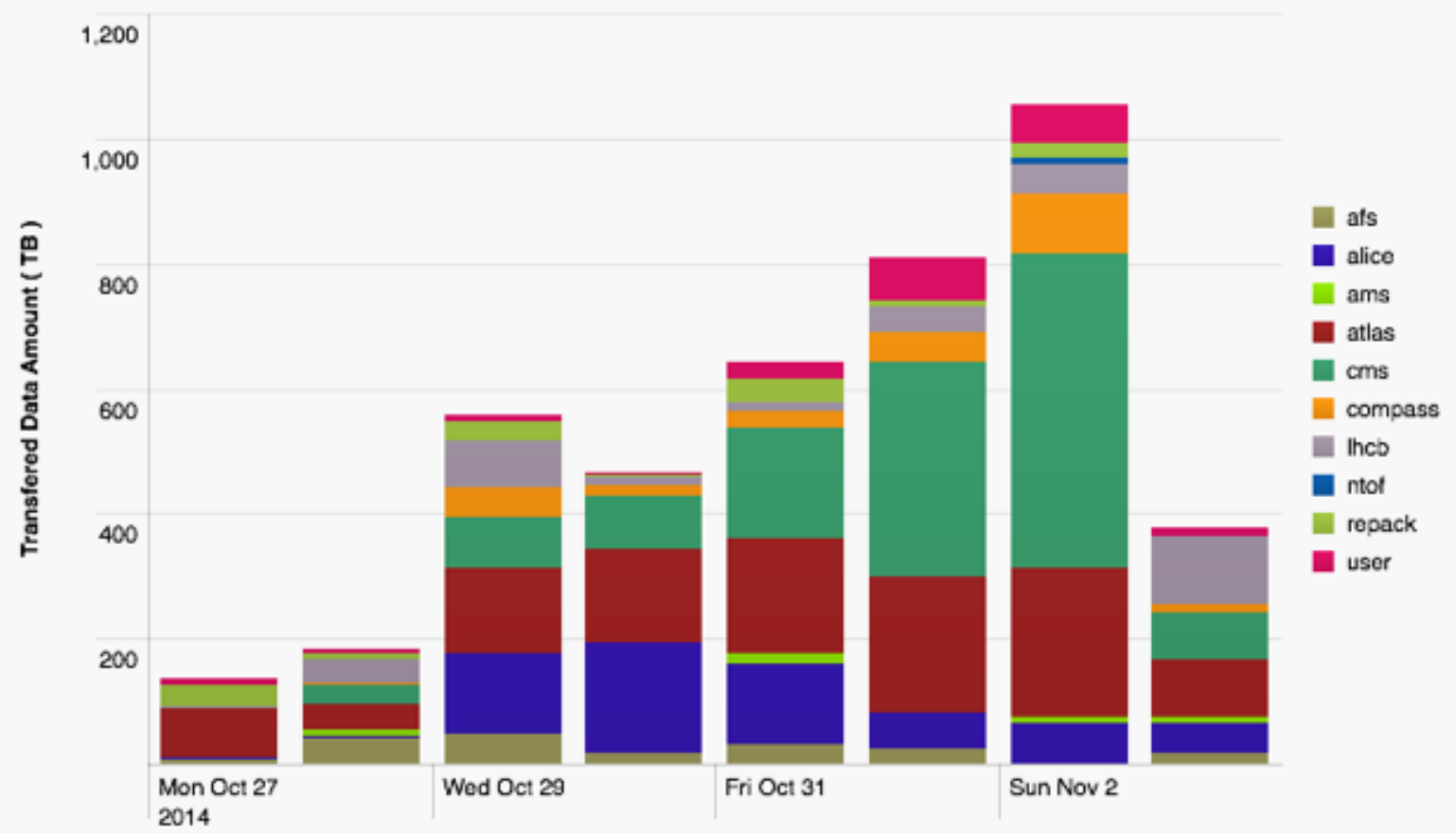
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 - Slow in random access mode
 - high latency in mounting process and when seeking data (F-FWD, REW)
 - Inefficient for small files (in some cases)
 - Comparable cost per (peta)byte as hard disks
- Tapes have also some advantages
 - Fast in sequential access mode
 - > 2x faster than disk, with physical read after write verification (4x faster)
 - Several orders of magnitude more reliable than disks
 - Few hundreds GB loss per year on 80 PB raw tape repository
 - Few hundreds TB loss per year on 50 PB raw disk repository
 - No power required to preserve the data
 - Less physical volume required per (peta)byte
 - Inefficiency for small files issue resolved by recent developments
 - Nobody can delete hundreds of PB in minutes
- Bottom line: if not used for random access, tapes have a clear role in the architecture





Tape Verification Read (lxbsq1332, lxfsrd0301) (TB) - in the last 7 days

Export



Source: G. Cancio / CERN



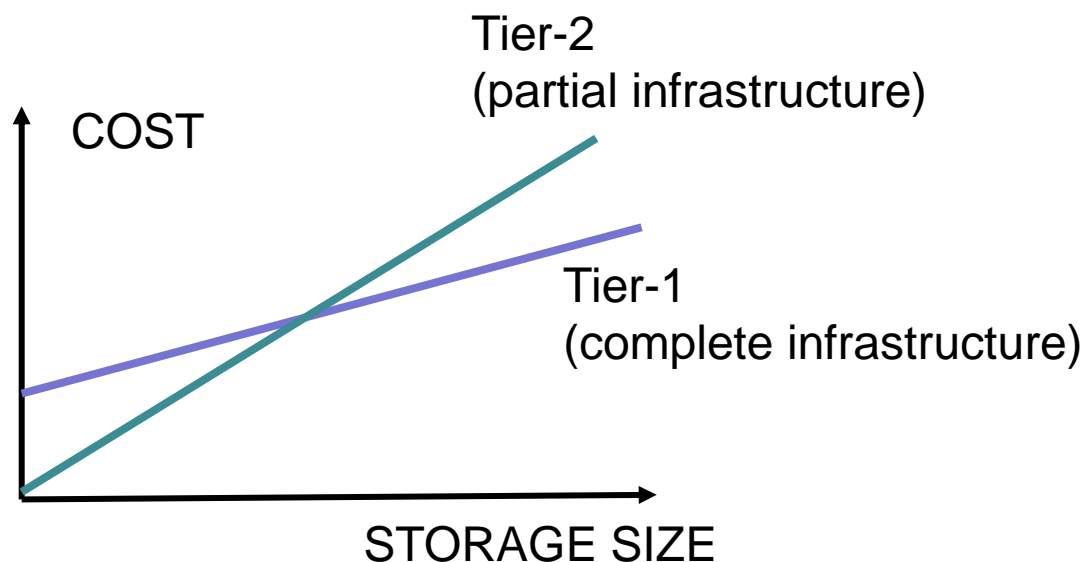


DSS

Data & Storage Services

Practically ... which solution ?

- There are fixed costs for the skilled manpower that is required to manage the high efficiency but complex solutions
 - Disk technology, Database technology, Tape technology, Networks, Databases...



- Many options available
 - interoperability is a key requirement for WLCG
 - DPM, dCache
- AT CERN
 - Separation between HOT (disks) and COLD (tapes) storage
 - CASTOR (used for the tape storage backend)
 - EOS (used for the disk storage frontend)
 - guaranteed low-latency file access for analysis
 - collaborative storage platform with space and access regulations
 - expand/exchange/shrink during operation
 - semi-automatic disk ejection & healing – reduced maintenance

- Example:
 - DPM is a consolidated product which is fully supported by CERN. It meets all requirements of a small site
 - Integrates well within both xroot and http data federations
 - EOS is the new design for a scalable infrastructure designed to minimize management cost. It is ready to meet the requirements for larger sites and multi-site environments
 - Larger functionalities

- Storage for High Energy Physics or Scientific computing can be
 - Too simplistic
 - Simple
 - Typical for Small Tier-2
 - Scalable and Reliable
 - Typical for a Tier-1 / Large Tier-2
 - Too complex
- Behind the software you use, it is always the people that makes the difference



DSS

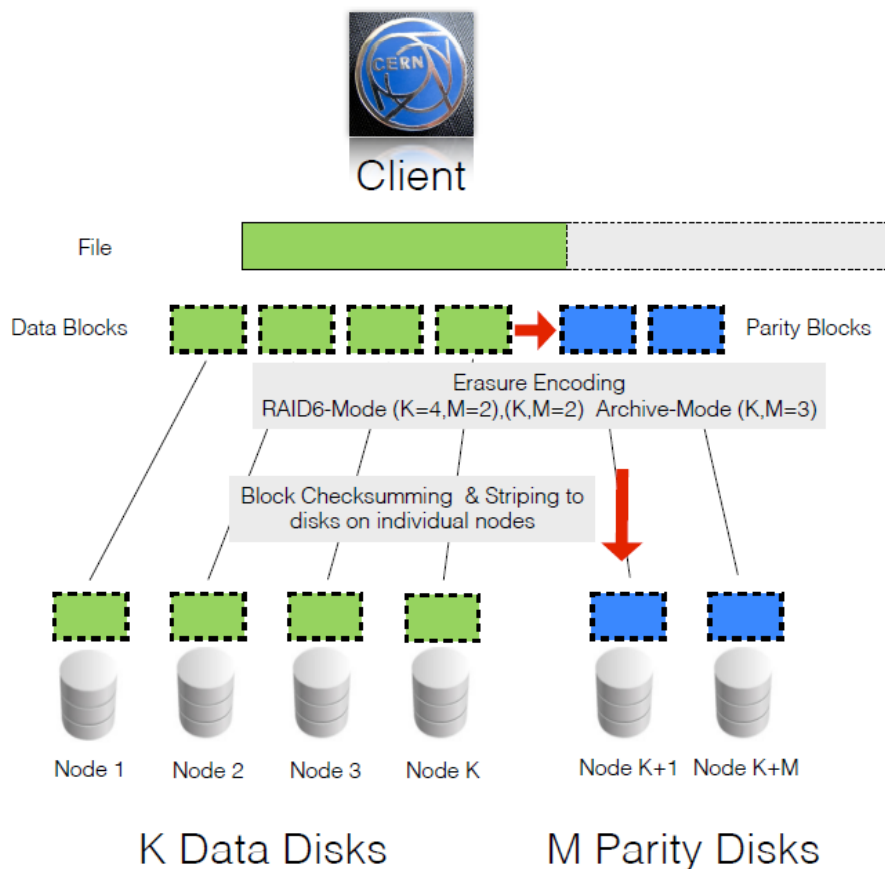
Areas of developments in storage at CERN

For EOS (Reserve slides)

Source: A.J.Peters/ CERN

- No need for a central relational database
- Designed for infinite scalability and arbitrary reliability
- expand/exchange/shrink during operation
semi-automatic disk ejection & healing
- Multi site federation support
 - Wigner / Geneva > 1000 Km distance

- Arbitrary Reliability – RAIN
 - Redundant Array of Independent Nodes



RAIN

- Block Striping allows parallel IO
 - boost single file performance
 - e.g. with $K=4$ 280 MB/s streaming write & 400 MB/s streaming write
- Erasure Encoding + Block Checksumming allows
 - error correction on the fly
 - up to M concurrent disk failures without data loss
- disk space savings compared to file replication
 - e.g. 25% less space needed for (4+2) vs. 2 replica

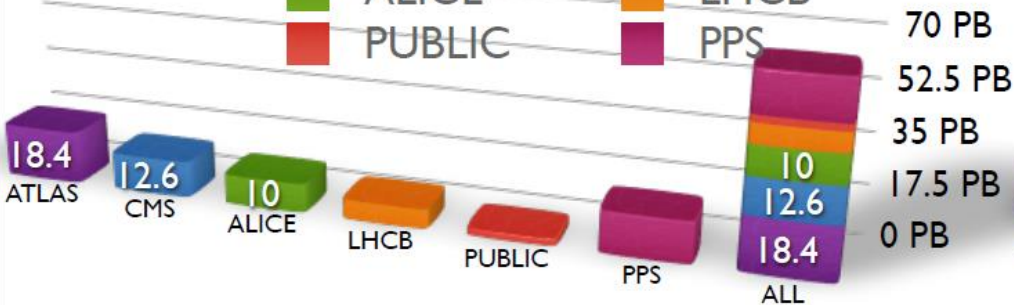


EOS in production at CERN

Raw Space

64.36 PB

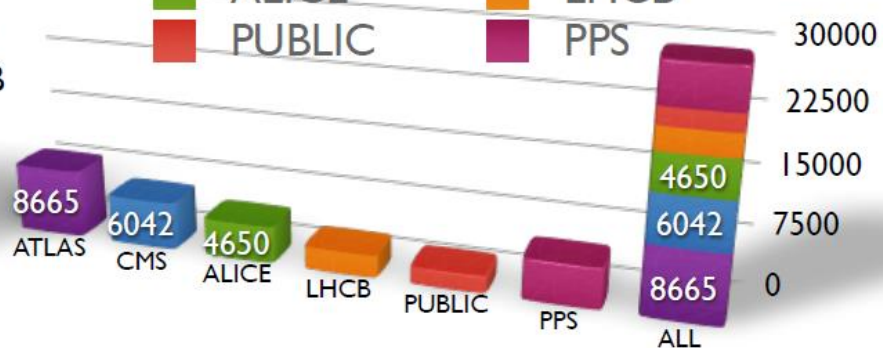
- ATLAS
- ALICE
- PUBLIC
- CMS
- LHCb
- PPS



Harddisks

29.4k

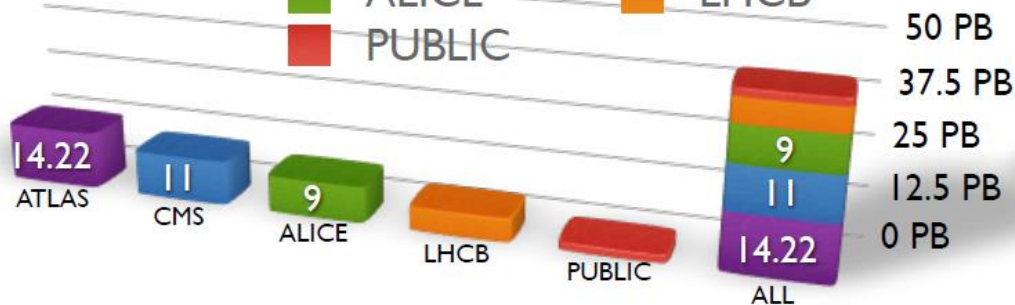
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Used Space

42.1 PB

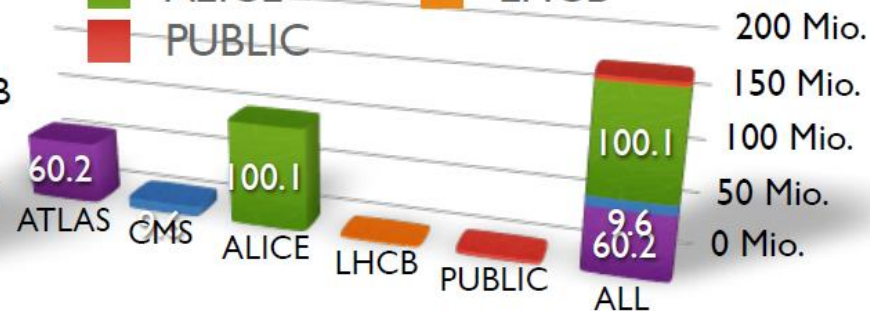
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Stored Files

174 Mio.

- ATLAS
- ALICE
- PUBLIC
- CMS
- LHCb



*EOSPUBLIC started 4.6.2013





- Five multi-PB installations
 - few thousand disks per instance
- Simplified life-cycle management workflows for on-going replacement/repair of hardware
- JBOD disks (no RAID controller) using software RAIN
- low-latency file access with in-memory namespace - 174M files
- fine-grained access control and quota management
- GRID storage element



- Accessed from thousands of CERN-local and remote batch nodes with Kerberos and GSI authentication
- Full support for XRootD, GridFTP, HTTP(S) & WebDav protocol - and – mountable with FUSE as a remote file system