### Managing Heterogeneous HTCondor Workloads

(or how I learned to stop worrying and love multicore)

William Strecker-Kellogg <willsk-at-bnl.gov>

## Support for Multicore Jobs in HTCondor

#### Paritionable Slots

- Each machine has one slot with all resources
- Jobs request portions of resources (CPUs / RAM) from these
  - Sliced into dynamic slots
  - Matching is done to "parent" slot
- See my Condor Week 2014 talk
- Working well since 2013

### Competition

- Starvation of multicore jobs in direct competition with smaller jobs
  - "Greedy" scheduling
- ► E.g. 3 cpus become free->3 single-core jobs start

### Defragmentation

Depth-first filling obviates need for a lot of defrag

## **ATLAS** Configuration

### All farm has one STARTD config

```
SLOT_TYPE_1=100%
NUM_SLOTS=1
NUM_SLOTS_TYPE_1=1
SLOT_TYPE_1_PARTITIONABLE=True
SlotWeight=Cpus
```

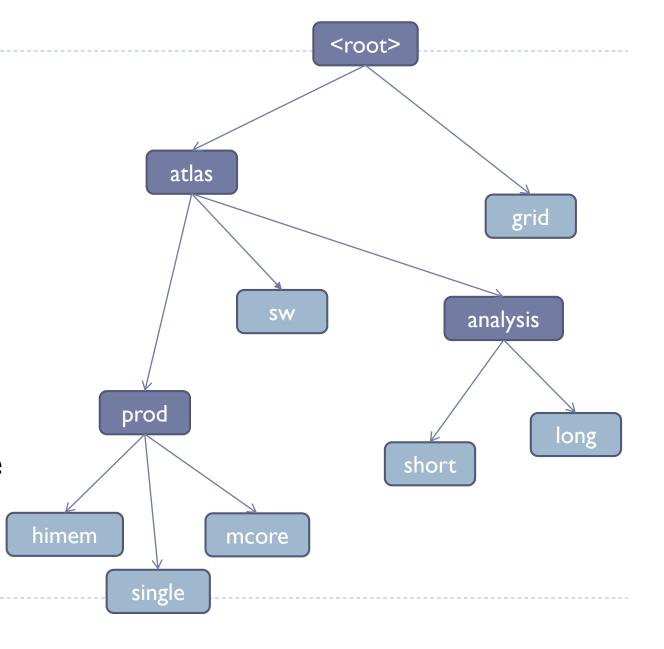
### Minimal Defragmentation

- Depth-first filling with NEGOTIATOR\_POST\_JOB\_RANK
- Start 2/hr, stop when 10 free CPUs appear

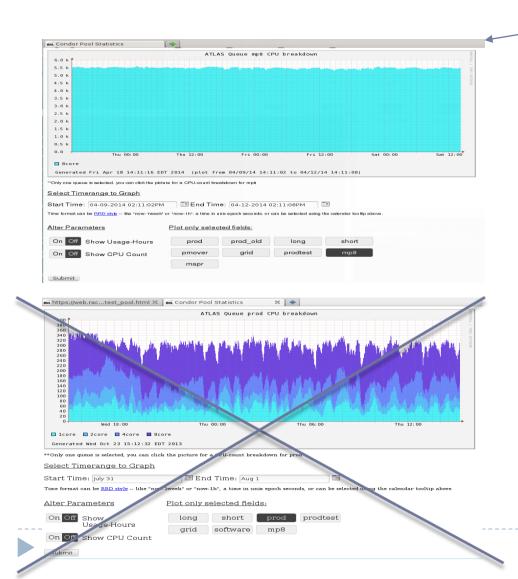
## Provisioning

- Use Hierarchical Group Quotas
  - Partition ATLAS into production / analysis
  - Partition jobs of different sizes into different groups at same level of tree
  - Logical divisions along other jobqualities (length)
  - Leaf nodes receive jobs

Up to site administrators to define



# Queues & Quotas



One size per-queue

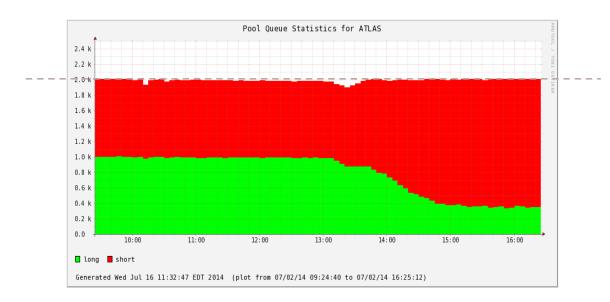
- Quotas on intermediate groups are sum(child-quotas)
  - Up the tree till the root-node has a quota the size of farm
- Jobs are segregated by resource-usage profile into different groups
- In ATLAS, Groups are 1-to-1 with Panda Queues

## Surplus Sharing

- Surplus sharing is controlled by boolean accept\_surplus flag on each queue
  - Quotas are normalized in units of SlotWeight (CPUs)
- Groups with flag set to True can take unused slots from their siblings
  - Parent groups with flag allow surplus to "flow down" the tree from their siblings to their children
  - Parent groups without accept\_surplus flag constrain surplus-sharing to among their children

# Surplus Sharing

- Scenario: analysis has quota of 2000 and no accept\_surplus; short and long have a quota of 1000 each and accept\_surplus on
  - ▶ short=1600, long=400...possible
  - short=1500, long=700...impossible (violates analysis quota)



## Management of Queues

#### We have a solution

- 1. Database to hold group-tree information and parameters
- 2. Flask website to manage quotas easily
- 3. Scripts to inject changes from database into your condor pool

### ▶ Available on my github

- Installable as python PIP package or RPM
- Site in production at BNL for the last 3 months

Demo...

## Sharing with Multicore

- Multicore queues can't have accept-surplus set alongside single-core queues
  - Scheduling is greedy, and would require much more defragmentation to work

#### Solutions:

- Factor your tree so this never happens?
  - This is nice but we can do better!
- Automatically rebalance surplus based on demand
  - Need a way to read demand... this is site-specific
  - 2. Assign parameters to determine "weight" of queues

## Weight, Demand & Threshold

- ▶ Each queue's "weight" is the size of its jobs
  - ▶ 8 for mcore...
  - 2 for high-memory...
- Demand is read with a script from PANDA
  - This is the only site-specific code that needs to be written
  - Could be a simple condor\_q lookup of Idle jobs
- Threshold is what average demand for last hour must be above in order for queue to be considered "full"
  - Adjust to best fit how "fast" each queue moves vs. how big the jobs are

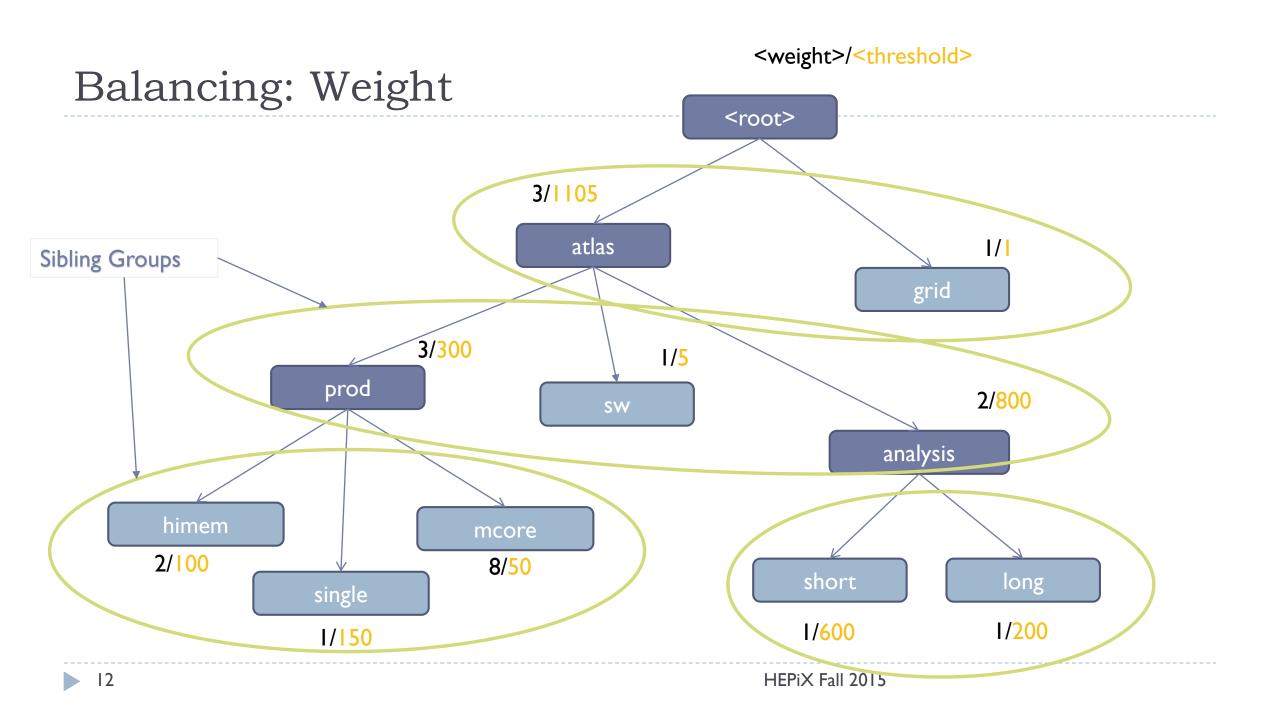
# Balancing Algorithm

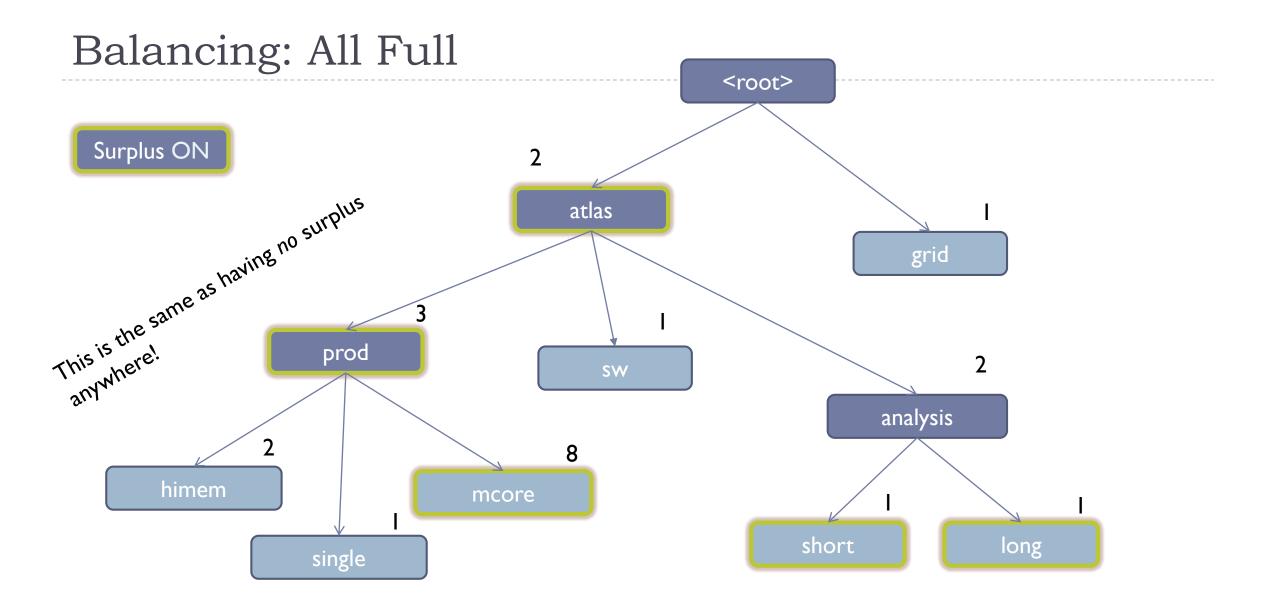
### **Assign Variables**

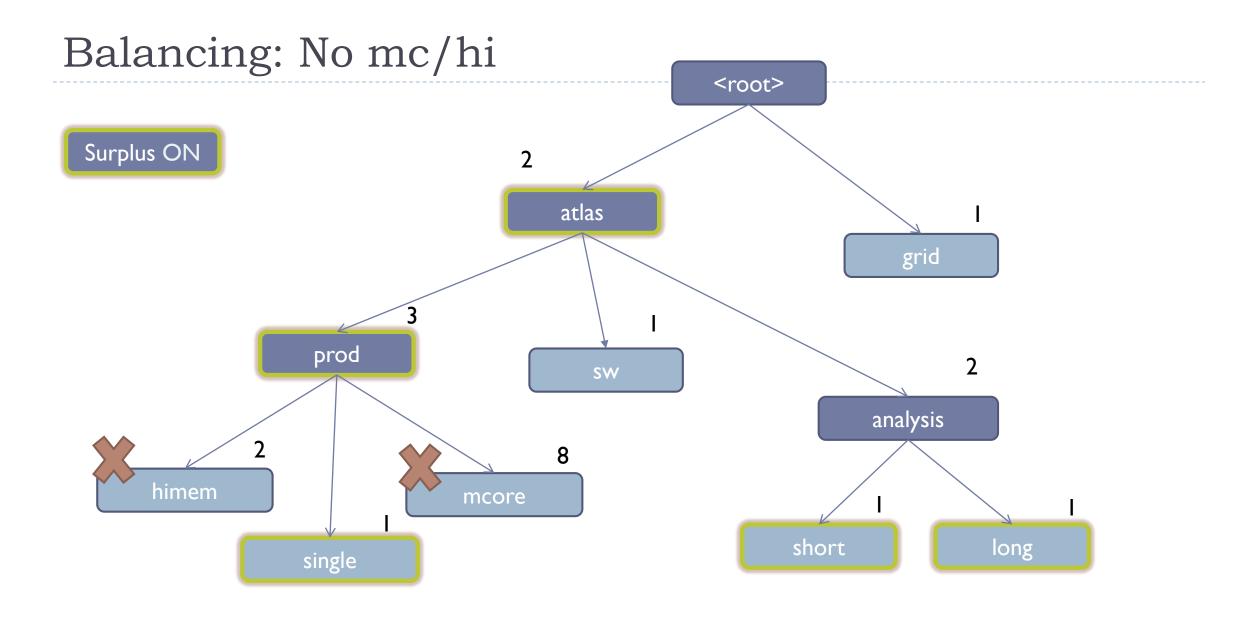
- Weights of intermediate groups
  - avg(w<sub>children</sub>)
- Demand of intermediate groups
  - sum(d<sub>children</sub>)

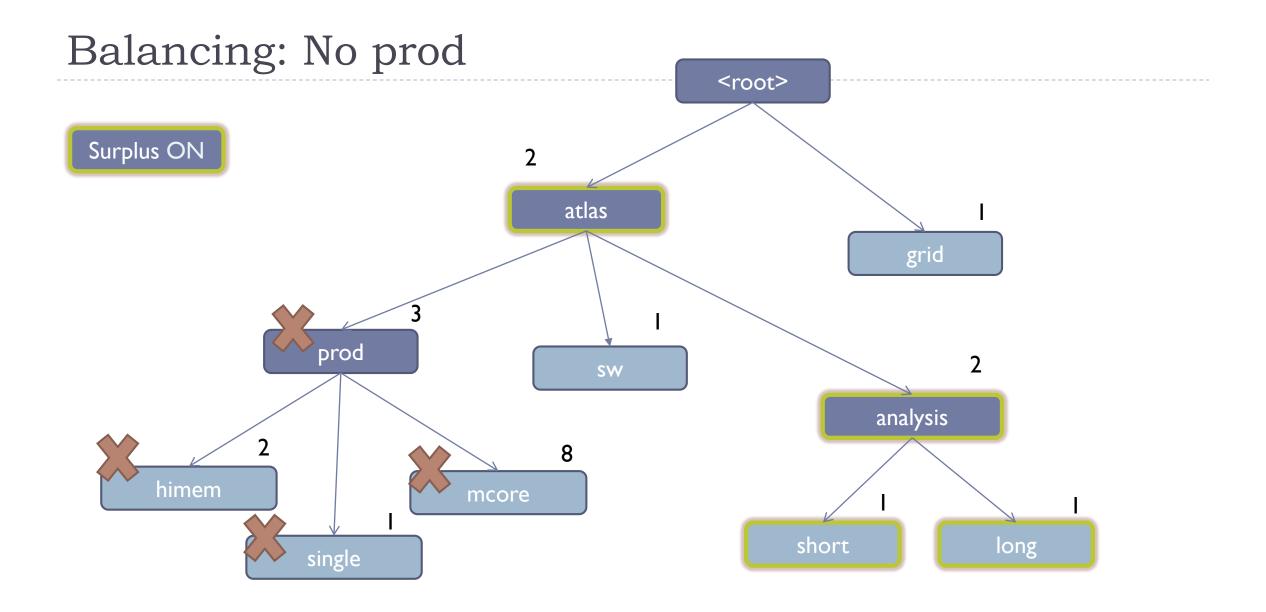
### **Algorithm**

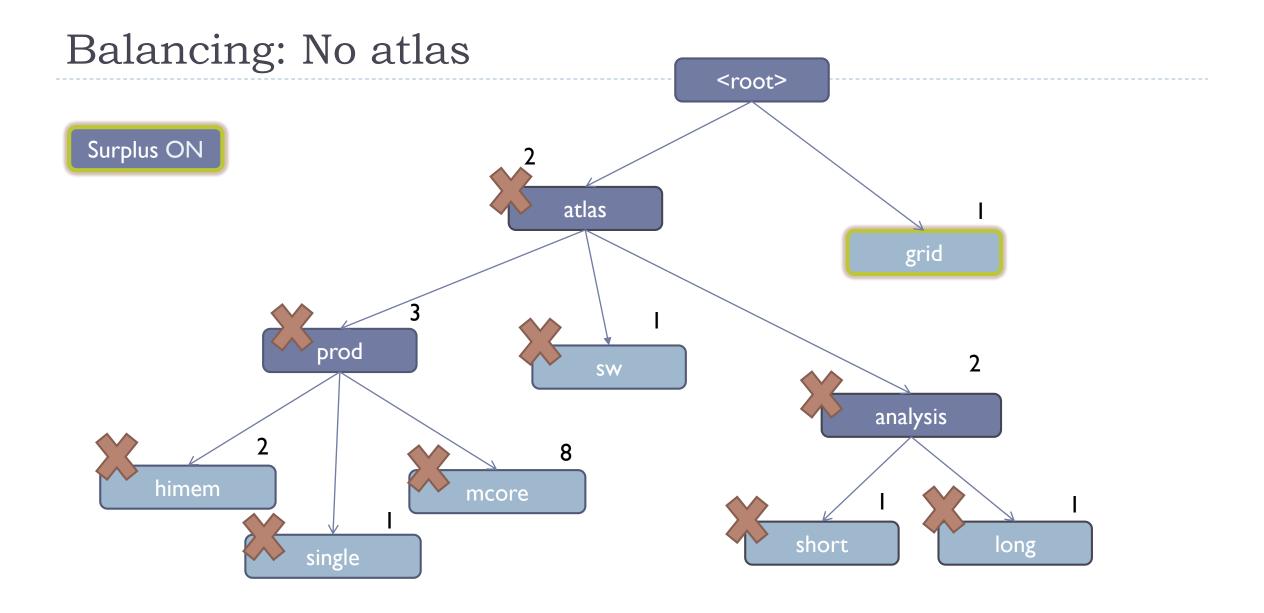
- Depth-first traversal
  - For each equal-level sibling-group
    - Set accept\_surplus to TRUE for only all the highest-weighted queues that have demand



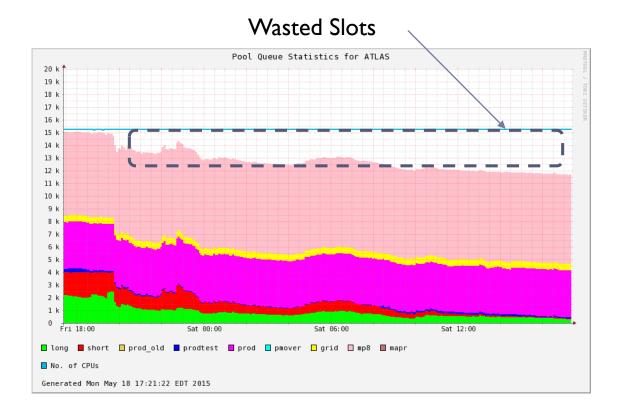


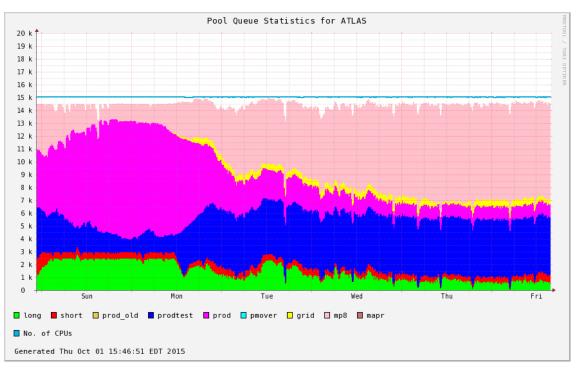






### Results





7 HEPiX Fall 2015

### Results & Context

- Multicore is ready to take slack from other production queues
- Spotty analysis-demand the past few months has allowed many millions of CPU-hours to go unwasted
- If all ATLAS has a lull in demand, OSG jobs can fill the farm

#### Who is this useful for?

- Algorithm works for any tree
  - Extensible beyond ATLAS where work is structured outside of batch system
- A multi-tenant service provide with a hierarchy of priorities
  - Really a problem of efficient provisioning, not scheduling
- Constraints
  - Workflow defined outside of HTCondor
  - Must map job-species to groups

## The End

Thank You!

Questions? Comments?

19 HEPiX Fall 2015