

The Use and Integration of Distributed and Object-Based File-Systems at Brookhaven National Laboratory

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Robert Petkus, Maurice Askinazi, David Free, Bruce Gibbard, Jerome Lauret, Zhenping Liu, Ofer Rind, Thomas Throwe, Yingzi Wu

RHIC/USATLAS Computing Facility
Brookhaven National Laboratory

RCF Fileserver Overview

RHIC Computing Facility

- 2000+ dual-CPU analysis/reconstruction Linux farm
- 680 TB local disk storage (SCSI, SATA, and PATA)
 - 132 TB xrootd storage on 650 nodes utilized by STAR
 - 25 TB dCache storage on 128 nodes (pool servers) utilized by PHENIX
- 220 TB FC RAID5 NFS storage on 37 Sparc/Solaris servers, Veritas 4.0 (VxVM, VxFS)
- 100 TB object-based storage on 20 Panasas shelves
- 6 TB FC AFS storage on 3 Solaris servers

ACF Fileserver Overview

USATLAS Computing Facility

- 330+ dual-CPU Linux farm
 - 146 TB local SATA disk storage dedicated to dCache
 - 322 internal/external read pool nodes
 - 8 internal/external write pool nodes
- 20+ TB FC RAID5 NFS storage on 4 Sparc/Solaris servers, Veritas 4.0 (VxVM, VxFS)
- 1 TB FC AFS storage on 3 Solaris servers

Fileserver Criteria

- Fast, scalable, reliable, and fault tolerant
- Load balancing
- Security and centralized management
- Single, global namespace
- Transparent, uniform, “POSIX-like” data access

Different Implementation Philosophies:

- Hardware vs. software
- Central vs. distributed
- Existing protocol vs. new protocol
- Proprietary vs. open

NFS Centralized Storage

NFS, Solaris, Veritas

- NFS consistently performs reliably on Solaris 9
- Network data transfer rate 70-80 MB/sec with single 1 Gb NIC
- Ubiquitous, compatible, and mature

However,

- No load balancing, poor scalability, and insecure
- Veritas is expensive and unpredictable in the midst of a host of more able, free competitors; e.g., XFS

Usage at RHIC and USATLAS

- Home directories
- Container for reconstructed data, job output
- Predominantly used for read access or scratch space

AFS Centralized Storage

OpenAFS

- Secure, reliable distributed filesystem well suited as a repository for static data and WAN access
- Replication creates redundancy and increases read performance
- Bottleneck potential during writes

Usage at RHIC and USATLAS

- Central software repositories
- Home directories
- Web content
- Container for finished data

Panasas - Clustered Centralized Storage

- Integrated hardware-software solution
- Fast: direct and parallel data access
- Global namespace
- Security and centralized management
- Distributed metadata



Front

Components of each “shelf” at the RCF:

- 1 Director blade and 10 storage blades
- Up to 8TB raw storage
- 1 Gb of internal connectivity per blade (iSCSI)
- 4 Gb of external connectivity (etherchannel)

Panasas Architecture

ActiveScale Operating System

- Runs on Director Blade
- Divides files into data objects, which are arbitrary in size, and stripes them across storage blades
- Dynamically distributes workload across storage blades
- Each storage blade is filled to 90% capacity. The remaining 10% is reserved for rebuilding parity

Direct Flow Software

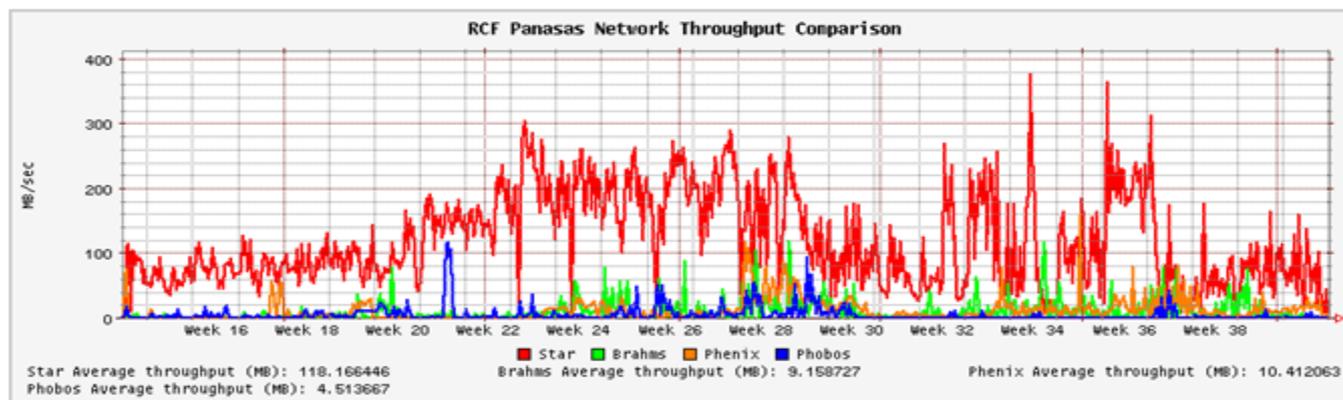
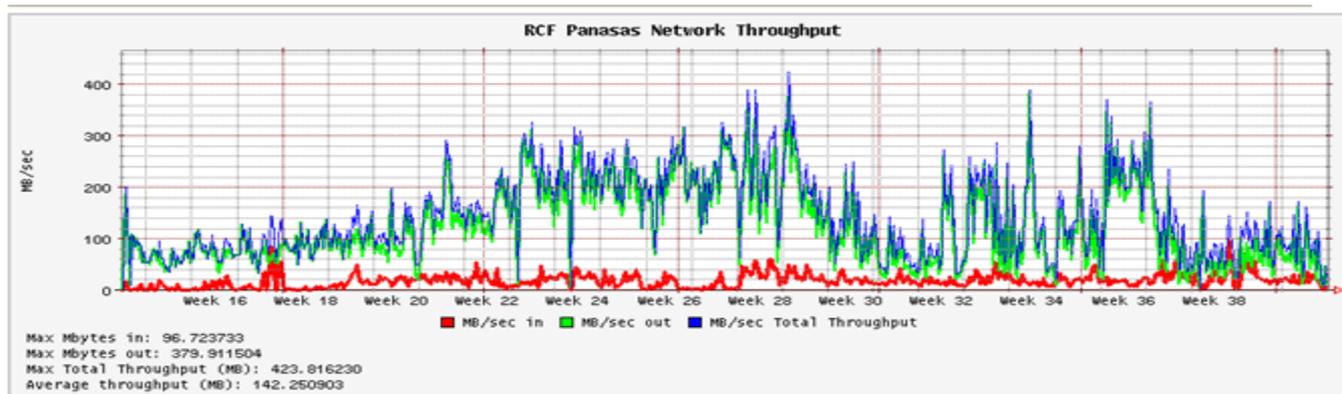
- Installed on the Linux compute node
- Direct data path from client to storage blades
- Optimizes data layout, caching and prefetching
- File is reconstructed at the compute node

Panasas performance snapshot

RCF Panasas Report: Last 6month

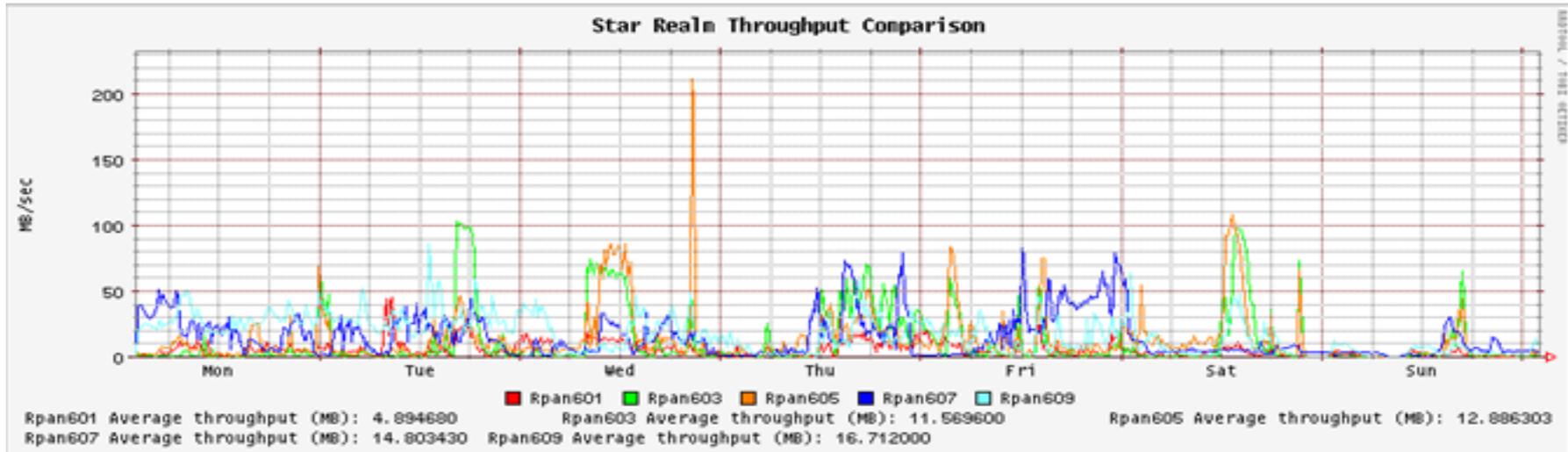
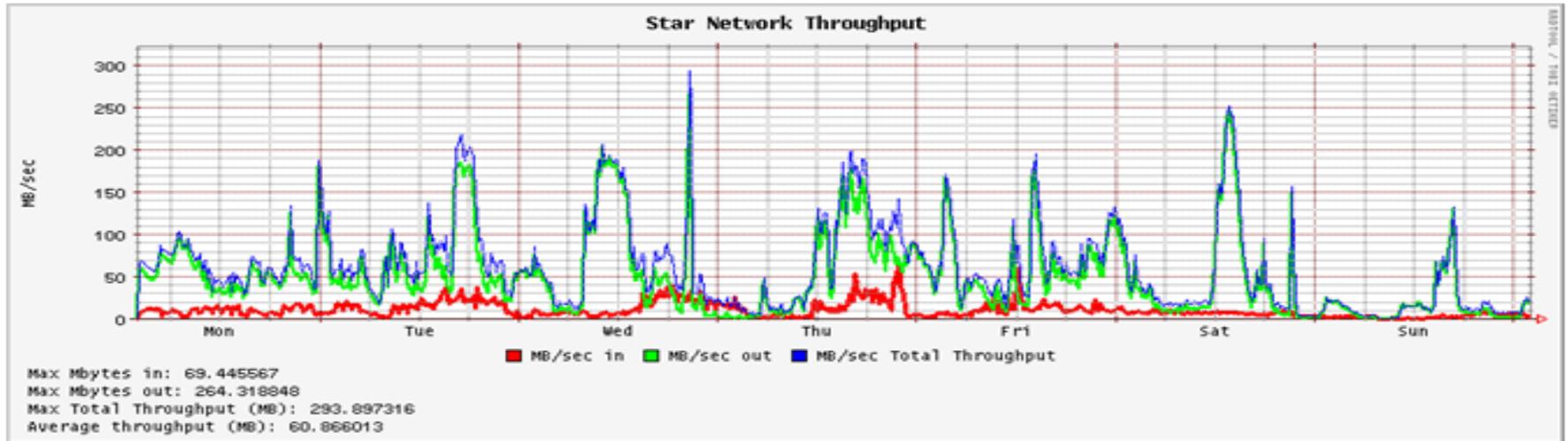
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Panasas performance snapshot



Panasas in Practice

Sure it's fast and the technology is innovative and promising,

But...

- Proven to be unreliable and finicky
- Expensive, but utilizing low-rent hardware – you wouldn't or couldn't recycle this for something else
- Maxed-out volumes are unstable requiring the use of hard quotas which further diminish usable capacity
- A storage blade failure frequently equates to lengthy <filesystem offline> reconstruction and recovery
- The Linux clients are buggy kernel modules that often crash the node
- No LDAP support.

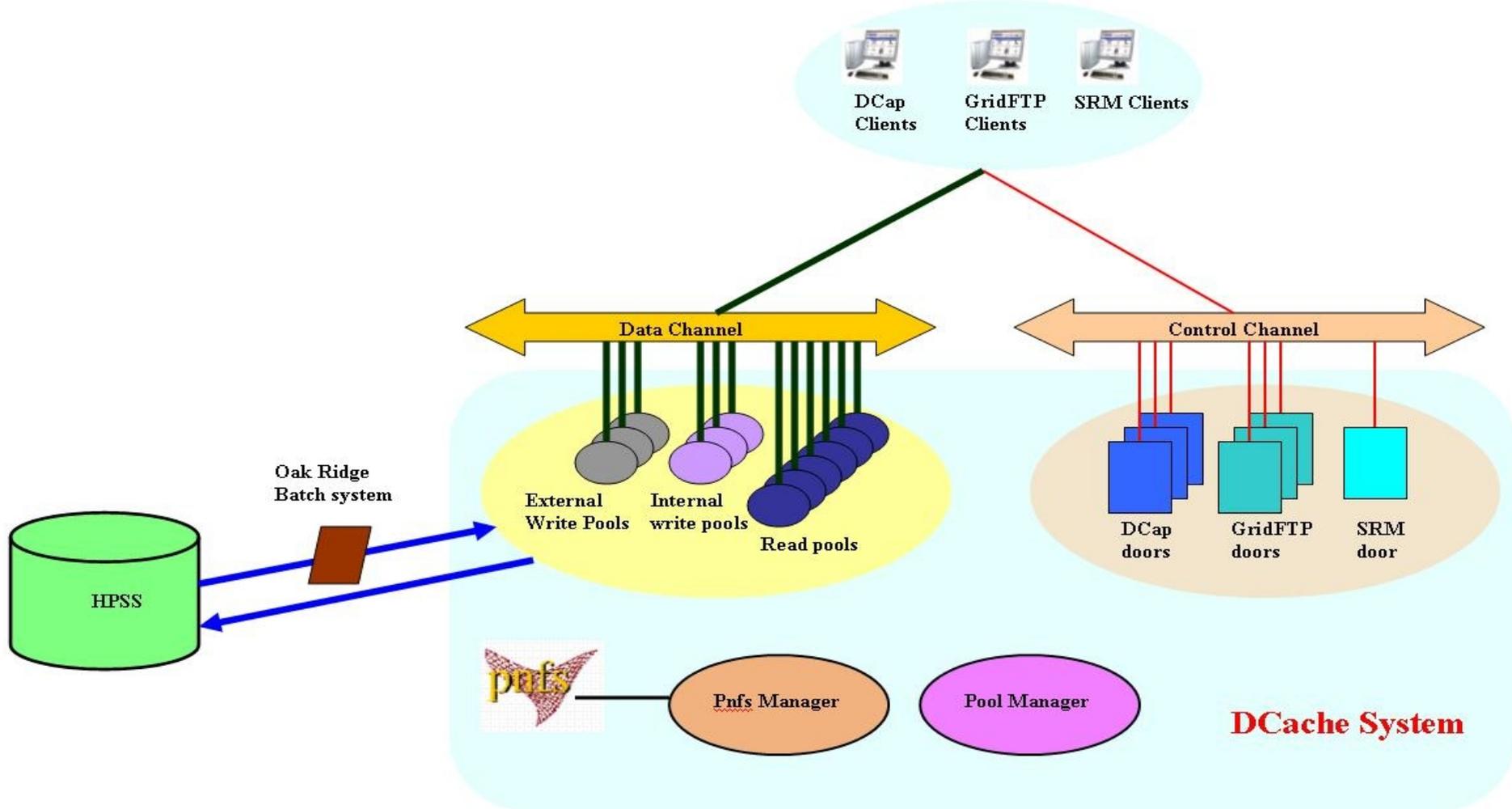
Usage at RHIC

- A replacement for NFS

dCache Distributed Disk Caching System

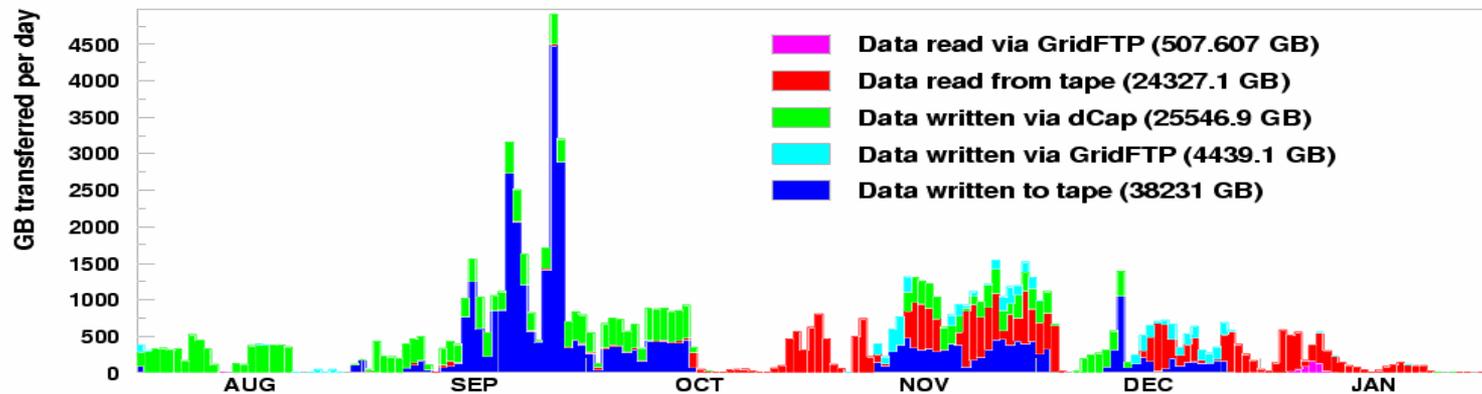
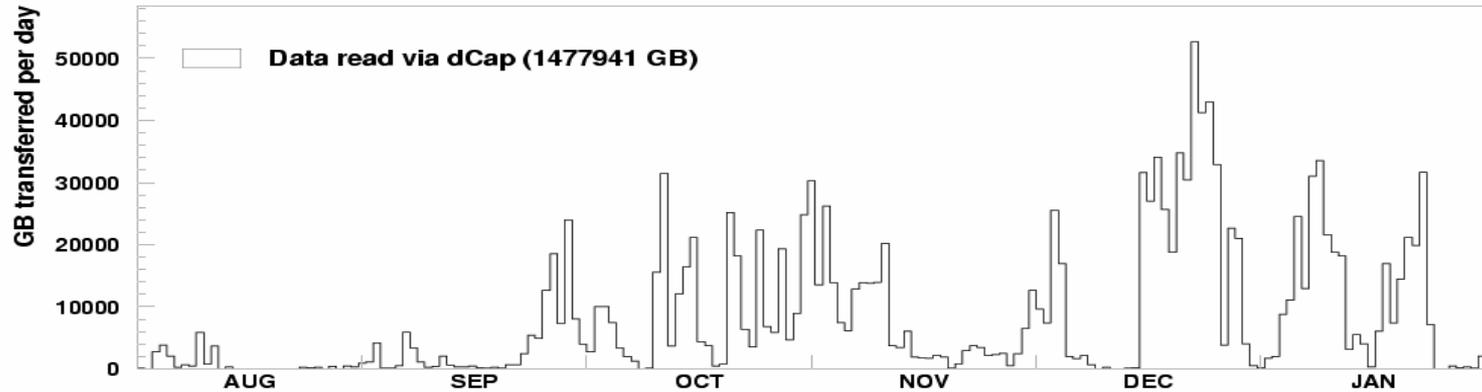
- Heterogeneous data access – data is stored in disk pools or in HPSS.
- Global namespace
- Secure (GSI, Kerberos), scalable, redundant, and fault-tolerant
- Load balancing: Detection of a “hot spot” will dynamically replicate files using cost metric algorithms; e.g., lowest latency, least load, disk space, internodal data transfer vs. HPSS refetch
- Transparent access to data in disk pools or HPSS
- Multiple data access methods (doors):
 - Local POSIX-like access via DCAP
 - SRM: Protocol negotiation, checksumming, space allocation
 - GsiFTP: secure WAN data transport
- General UNIX filesystem access via PNFS, a database and exportable filesystem made available via NFS that serves and stores metadata.
*PNFS does not store the actual files
- Improved administration and monitoring

USATLAS dCache Architecture



PHENIX dCache Usage

Phenix dCache Statistics, Aug. 2005 - Jan. 2006



dCache Issues

- Write pool nodes:
 - Need dedicated write pool nodes – frequent crashing when sharing computing resources
 - Better performance using the XFS filesystem
 - Need reliable disks
- PNFS database is single point of failure and potential bottleneck. Need multiple, fault tolerant metadata dbs
- If the PNFS server crashes on transfer (very rare), files written into the system might be improperly registered and/or invalid requiring a manual clean-up of residual data
- Read or write access only (not RW)
- No mechanism or policy for user/group based quotas on transfer limits

dCache at the RACF

USATLAS Usage

- Hybrid model: The majority of compute nodes are also read pool nodes thus maximizing the potential of each system
- Data analysis: write raw data into HPSS, analyse on farm, write back to HPSS
- Grid production: On-site, dCap is the only source and destination for data. Off-site, data is accessible via GridFTP and SRM clients
- Oak Ridge Batch System for backend tape prestage

RHIC Usage

- Hybrid model for a subset of compute nodes in the PHENIX experiment
- dCache metadata (PNFS) exported/imported directly to/from PHENIX file catalogue
- Data retrieval from HPSS via Carousel (developed at RHIC) which allows for policy-based request throttling

Xrootd: Accessing and Managing Distributed Data

- High performance, multiplexed data access per client
- Global namespace
- Secure (Kerberos and GSI plugins)
- P2P architecture with distinct data/control flow akin to GnuTella
- Low CPU overhead
- Load balancing (open load balancers (OLB)) – conceptually similar to dCache “hot spot” detection wherein data is dynamically replicated on demand via definable cost metrics
- No single POF – each element in the system is completely redundant and fault-tolerant
 - Data clients (also data servers as configured in the STAR experiment) have transparent access to data
 - Data Servers – A master copy of the data in HPSS will be retrieved to another server in the event of a failure
 - Redirectors – traffic managers that direct clients to data
 - OLB Managers (monitors load and availability of data) and Servers (maps data location)

Xrootd continued

Issues

- Limited data access methods – however, the merging of xrootd with SRM is in development to provide a complete grid solution
- No access via a general UNIX filesystem-like interface – psychological convenience?
- Read or write access only (not RW)
- No mechanism or policy for user/group based quotas on transfer limits

Usage at RHIC

- The STAR experiment with 650 xrootd dataservers in deployment is the largest in the world
- Hybrid deployment – compute nodes double as file servers

Other Solutions

- **GPFS:** data block striping across disks. Parallel reads/writes. Fail-over, replication, distributed metadata. Extremely fast and reliable but cost prohibitive
- **BlueArc:** ASICs dedicated to NFS, network, and filesystem. Feature-rich, fast and robust NFS solution – but still NFS...
- **Ibrix:** Meta servers assigned to segments in a disk pool. Traditional NFS or “Fusion”, a proprietary protocol
- **Lustre:** Object-based, distributed storage solution. Software only. Active development. Older versions are free
- **NFSv4.1:** One protocol for file access, locking, and mounting. Kerberos integration, client caching and delegation policies. When? ?
- **ZFS:** Combined volume manager and 128-bit filesystem. Transactional writes, 64-bit data checksums. Proprietary, Solaris-only and still not ready for general consumption

Feature Comparison of RHIC/USATLAS “Fileservers”

	NFS	AFS	Panasas	dCache	xrootd
Low Cost	✓	✓		✓	✓
Reliability	✓	✓		✓	✓
Speed			✓	✓	✓
Scalability		✓	✓	✓	✓
Fault Tolerance		✓ ¹	✓	✓ ²	✓
Load Balancing		✓ ¹	✓	✓	✓
Security		✓	✓	✓	✓
Centralized Mngmt		✓	✓	✓	✓
Global Namespace		✓	✓	✓	✓
POSIX Access	✓	✓	✓	✓ ³	✓
Open Source	✓	✓			✓
User Policy Mngmt	✓ ⁴	✓	✓ ⁵		

1: When using replication

2: Single PNFS DB is a POF

3: POSIX-like access

4: Dependent on the underlying filesystem being exported

5: No group-based quotas

Conclusion

- No single implementation is a panacea providing all solutions
- New file servers are constantly being developed/perfected and will continue to be tested
- Heterogeneous solutions will remain in place but the following trends are clear:
 - A move from centralized toward distributed storage
 - USATLAS and PHENIX will continue to deploy as well as expand and enhance their utilization of dCache
 - STAR, with the largest xrootd deployment in the world, is focused on xrootd, specifically a grid implementation using SRM
- Panasas robustness and reliability is ultimately disappointing
- NFS will hobble along for some time functioning as a home directory and scratch storage server
- AFS, popular at CERN and BNL, will remain the repository of choice for static software