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	Year	Type	Bandwidth Mbit/s	Physical	Interf.	Protocol	
This three standards Did not have a large	1984	FDDI	100		fiber	TCP/IP, Dedicated	
influence on the development of Multi	1988	Ultranet	100	Dedicated	copper	TCP/IP	
Gbit/s Networks		ATM	155 – 625	Dedicated	fiber	ATM & encaps.	
	1989	НІРРІ	800	HIPPI-800 HIPPI-Ser.	copper fiber	Dedicated, TCP/IP. IPI3	
FIBRE LIMINNEL BSSOCIATION	1991	Fibre Channel	255 - 510,	FC-Phys	fiber	Dedicated	
	1999		1020 - 2040			TCP/IP, IPI3, SCSI	
	1995	Myrinet	1 Gbit/s	Dedicated		Dedicated, IP	
<u>Myrinet</u>	2000		2 Gbit/s	fiber		TCP/IP	
Gigabit Ethernet Alliance	1996	Gigabit Ethernet	1.25 Gbit/s	FC + IEEE 802.ae	copper fiber	TCP/IP	
Obsolete or Commodity now to							
		day					
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A company private solution to allow fast switched point to point connections



100 MByte/s channel speed but not full crossbar



Introduced in 1990 survived up to 1997



Not compliant with known network standards



In fact not a network device but a star-point connecting switch



Compliant with TCP/IP and some other protocols



Used in the CERN computer center from 1992 up to 1996 to connect the large main frames (Cray, IBM, Siemens/Fujitu





ligh Performance N etworking **ATM** = Asynchronous Transfer



Work on Amostaneed in 1984 and was released as a standard in 1988.



ATM is a flexible communication industry standard made for data, voice and video.



ATM is a packet switching system with small cells of 48 bytes payload + 5 bytes header.



ATM allows to set-up virtual channels.



Virtual channels rise problems at congestion, either data cells are lost or in secure mode result in overall high latency.



ATM had a certain popularity as a network with speeds of 155 Mbit/s and 622 Mbit/s.



Some people professed ATM as network and back plane replacement, connecting processor, memory and graphics parts aswell as the network. It never happened



ATM is still used at 2.5 Gbit/s OC48/SDH16 and 10 Gbit/s OC192/SDH48 but tendency is to go at the higher throughputs to the larger frames of POS and IP.



High Performance N etworking The HIPPI

With a need to conne<mark>ct supercom</mark>puters, Storage systems, and graphics I/O, LANL started to develop a new kind of Interconnect in 1987, guided by Don Tolmie.

Guidelines for this first High Performance Interface: 100 MByte/s, Simple and Cheap.

All big computer manufacturers participated in the development, under guidance of ANSI T11.3 and managed by HNF, the High-performance Networking Forum.

HIPPI-Phy: Physical Standard for point to point connections was accepted in October 1991, first equipment was demonstrated in 1990.

The standard was extended from point to point connections with the introduction of Data Switches and a switch protocol HIPPI-SW



Serial-HIPPI was proposed in 1991, and by CERN's influence accepted as a standard in 1995

Protocols:

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HIPPI framing protocol(not much used),TCP/IP(for communication)IPI 3(mainly used for storage).



In 1989 CERN started to work with HIPPI, with a great success in NA48 and in the computer center and was finally dismantled in 2002.

Let's have a quick look at the HIPPI technology

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In 1991 the Serial-HIPPI specification ware handed over to CERN.

From 1991 to 1993 Michael S. Haben did his PhD to serialize HIPPI signals in a collaboration with Birmingham University, H.P and CERN.



H.P made the chips and optical components that enabled Serial-HIPPI

In 1994 work started to connect the NA 48 experiment to the computer center using Serial HIPPI. With the result that ANSI accepted in 1995 as a standard HIPPI – Serial.



Serial HIPPI uses the same interface and the same protocols as HIPPI. Such transparent to the user.



Serial HIPPI has two physical standards:

Long Wavelength	1350 nm	Lasers	10 Km.
Short Wavelength	850 nm	Diodes	200 m.



20b/24b up-coding is used for link security, incrementing bandwidth from 800 Mbit/s to 1 Gbit/s

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ligh Performance N etworking Upcoding; what is it, why?

$\frac{1}{2}$ of a $\frac{4b}{5b}$ (bad) example

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0 0	0	00 ┥	1 0	0	3 0
0 0	0	0 1	1 0	0	0 1
0 0	0	10	1 0	0	10
0 0	0	11	1 0	0	1
0 0	1	00	10	1	0 0
0 0	1	0 1 🦰	1 0	1	0 1
0 0	1	10	1 0	1	10
0 0	1	11	1 0	1	11
0 1	0	0 0	1 1	0	0 0
0 1	0	0 1	1 1	0	0 1
0 1	0	10	11	0	10
0 1	0	11	1 1	0	11
0 1	1	0 0	1 1	1	00
0 1	1	0 1	1 1	1	0 1
0 1	1	1 0	1 1	1	10
0 1	1	11	11	1	11

0000	0
0001	1
0010	2
0011	3
0100	4
0101	5
0110	6
0111	7
1000	8
1001	9
1010	10
1011	11
1100	12
1101	13
1110	14
1 1 11	15

Several ways of upcoding are possible, a number them standardized: of are 4b/5b. 10b/12b. 20b/24b. but sometimes patented.

This coding uses a even parity placed in the middle of the four bits + parity. (a bad example)

They all have a different influence on the total bandwidth with an average of + 20 %.

Data words are spread over a wider pattern, sometimes combined with bit shuffling, such to avoid cross talk.

Forbidden data words on non critical places in the coding are used for link synchronization.

The goal is to make link errors as low as possible BERR 10⁻¹² to 10⁻¹⁵ is normal.

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A point to point connection that can perform network functionality by means of FABRIC's.



Fibre Channel introduces "OPEN FIBER CONTROL" that shuts down the laser as the connector is open. **EYE Security**

Fibre Channel introduces a Layer structure from the Physical Level on.

Fibre Channel is also a standard handled at ANSI T 3.11







igh erformance Networking **Gigabit Ethernet 1000**

The success of Here and the partial success of Fibre Channel forced the Ethernet community in 1995 to react.

Based on 1 Gigabit payload the total bandwidth with upcoding and overhead should not be more than 1.2 Gigabit.

The Gigabit Ethernet community struggled a long time with technological problems on the physical link level.

Finally The Gigabit Ethernet group adapted the 1 Gigabit Fibre Channel Technology and over-clocked it by 20%.



This over-clocking used the technology at its limits and was in the beginning source of many link errors.



A certain company used illegally chip technology patented for Serial HIPPI, however the result ware the first reliable and successful PCI interfaces.



Finally the standards 802.3z describing the modified Fibre-Channel Physical link layer and 802.3ab which introduced new physical link layer with set-up algorithms that proofed successful are accepted in 1999.



Both standards forgot to solve the problem of the small 1500 byte Ethernet frames.

1000 Base-T Link Initiation

If a clever man goes to read a book



He first looks at the Index

As the link opens up, it looks what the other end contains: **AUTO-NEGOTIATION PAGES**

Contents:

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RSYNC	= Receive Word Synchronization
SD	= Signal Detect
LINK	= Link Detect
AN_NP	=AutoNegotiation Next page Status
AN_TX_NP	=AutoNegotiation TX Next page Status
AN_RX_NP	=AutoNegotiation RX Next page Statu
AN_RX_BP	=AutoNegotiation RX Base page Statu
AN_RMTRST	=AutoNegotiation Remote Restart Sta

Other Gigabit Ethernet Properties:

8B / 10B Upcoding

One of the non Data Codes is used for a PAUSE Frame with a value defining Pause time.

If PAUSE Frame value = 0 transmission restarts.

This Flow control send by the Receiver, at the end of a frame, stops the transmitter at the end of the frame.

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High Performance Networking 1000 Base-T and the frame size problem

Given: 1) A **16 MByte** file to be transferred over the network.

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2) Ethernet maintains the 1500 Byte Maximum Frame Size.

Technology	Frames to be handle	Move Instructions (32 bit words)	Transmitter Interrupts	Receiver Interrupts	Transfer Time without latency or interrupts
10Base-T	<u>+</u> 10 000	4 000 000	1 (end of file)	<u>+</u> 10 000	1.5 seconds
100Base-T	<u>+</u> 10 000	4 000 000	1 (end of file)	<u>+</u> 10 000	0.15 seconds
1000Base-T	<u>+</u> 10 000	4 000 000	1 (end of file)	<u>+</u> 10 000	0.015 seconds
10000Base-T	<u>+</u> 10 000	4 000 000	1 (end of file)	<u>+</u> 10 000	0.0015 seconds

If we only look at the number of Instructions/sec** to be handled (ignoring the Interrupts) with:

Ethernet	or	10Base-T needs	5	Mega	Instructions/second to fill the Pipe	A 5 MHz processor is enough	but in	1990
Fast Etherne	t or	100Base-T needs	50	Mega	Instructions/second to fill the Pipe	A 50 MHz processor is enough	but in	1995
GigE	or	1000Base-T needs	0.5	Giga	Instructions/second to fill the Pipe	A 0.5 GHz processor just now	but in	1998
10GE	or	10000Base-T needs	5	Giga	Instructions/second to fill the Pipe	No 5 GHz processor available		



High Performance N etworking Solutions ?

The first interfaces did or	nly 30 to 40 hbyte/s	Why ?	Look at the receiving end
Frame:	A 1500 Byte frame at 10 b	it a Byte = 15 բ	usec
Interrupts:	5 μsec for a 300 MHz pro	cessor = 1500	Clock Cycles.
Move and Check frames:	minimum 2 X 1500 cycles	= 3000 Clock	Cycles = 10 μsec
Transfer Time:	15 + 2.5 +10 = 27.5 μsec/fr	rame	
Bandwidth:	1 / 275.10 ⁻⁵ (1500) = 37 №	IByte/s	

Solutions:

Use faster processors

Hardware CRC Generation = The interface generates the CRC code at the fly and keeps it available to be read by the processor

Interrupt Aggregation = collecting all the frames of a file in local memory in the Interface and generate an interrupt (end of file) as the transfer is complete (or if local memory is almost full).

Traffic Offload Engines (also called TCP/IP offload Engine or TOE) = handling all TCP manipulations in a local processing engine at the interface level and transfer the Data by DMA.







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This Statistics come from Mellanox: are they reliable ?? are they honest ??

Alacritech can never do 1700 Mbit/s on a 1000Mbit/s link. It is a 2.5 Gbit network and uses a network DMA in stat of TCP/IP

Intel PRO/1000T is an early pre-production module that can do better.

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igh erformance Networking SCI = Scalable Coherent

interface

First ideas on SCI started in 1987 resulting in an accepted standard: ANSI/IEEE 1596/1992.



SCI first of all an OPEN distributed bus, with network capabilities.



Physical connections are based on parallel serial links with ring oriented interconnects able to connect up to 64 K nodes.



Including switches, multiple rings can be combined in a network like structure.



Dedicated frame format and dedicated protocol, but active IP encapsulation possible.



Algorithms for Cache Coherency and Memory Locks are foreseen.



Not foreseen: Low latency, Fast priority handling, But transfers use DMA.



Link speeds: depending IC technology from 125 Mbit/s (1993) to 3 Gbit/s (2000).





High Performance N etworking References:

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HIPPI-SC, High-Performance Parallel Interface -Switch Control, ANSI X3.222-1996, Rev 3.2, ISO/IEC 11518-6, April 9, 1997.

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For Fibre Channel information

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http://www.hnf.org/

http://hsi.web.cern.ch/HSI/hippi/

http://www.fibrechannel.org/index.html http://data.fibrechannel-europe.com/index.html

http://www.10gea.org/GEA1000BASET1197_rev-wp.pdf

SEEQ 8101 Gigabit Ethernet Controller (Data Sheet), April 27, 1998

The Scalable Coherent Interface (SCI) is an approved ISO/ANSI/IEEE Standard, 1596-1992

