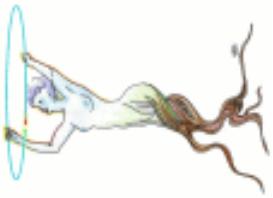
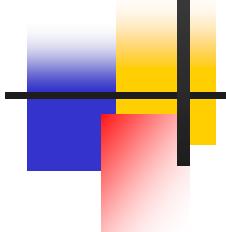


Bring the KB to the PB not the PB to the KB

Fons Rademakers

PROOF with AliEn and GLite



Parallel ROOT Facility

- The PROOF system allows:
 - Parallel analysis of trees in a set of files
 - Parallel analysis of objects in a set of files
 - Parallel execution of scripts
- On clusters of heterogeneous machines
- Its design goals are:
 - Transparency, scalability, adaptability

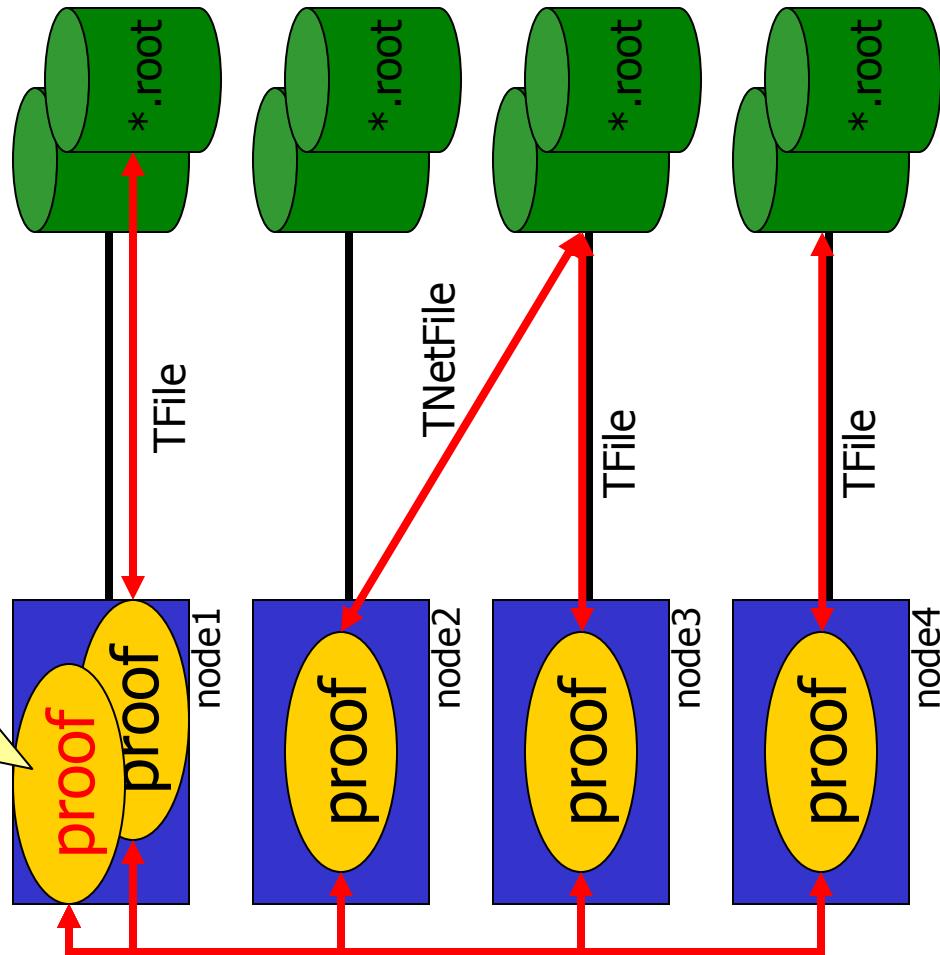
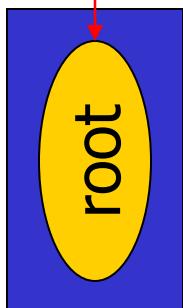


Parallel Script E

Local PC DF Cluster

```
#proof.conf  
slave node1  
slave node2  
slave node3  
slave node4
```

←Stdout/obj
ana.C→



```
$ root  
root [0] tree->Process("ana.C")  
root [1] gROOT->Proof("remote")  
root [2] chain->Process("ana.C")
```

root = master server
proof = slave server

Data Access Strategies

- Each slave get assigned, as much as possible, packets representing data in local files
 - If no (more) local data, get remote data via rootd and rfio (needs good LAN, like GB eth)
- In case of SAN/NAS just use round robin strategy



PROOF Transparency

- Make working on PROOF as similar as working on your local machine
- Return to the client all objects created on the PROOF slaves
 - The master server will try to add "partial" objects coming from the different slaves before sending them to the client

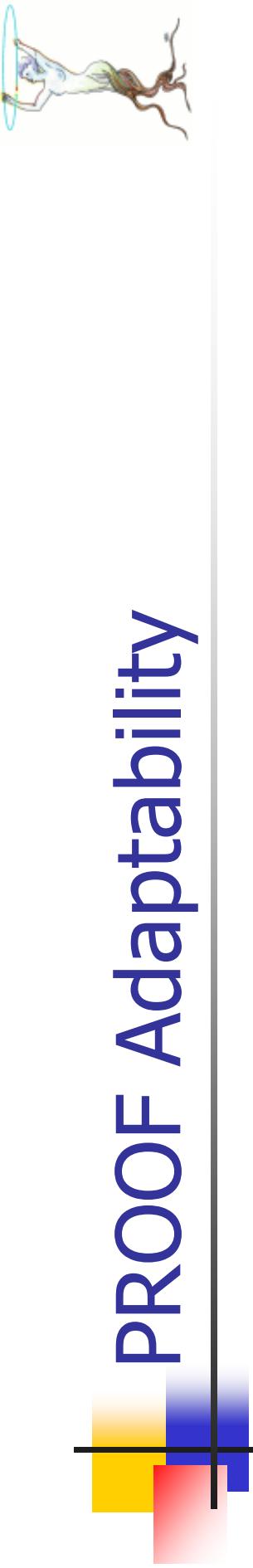


PROOF Scalability

- Scalability in parallel systems is determined by the amount of communication overhead (Amdahl's law)
- Varying the packet size allows one to tune the system. The larger the packets the less communications is needed, the better the scalability
- Disadvantage: less adaptive to varying conditions on slaves

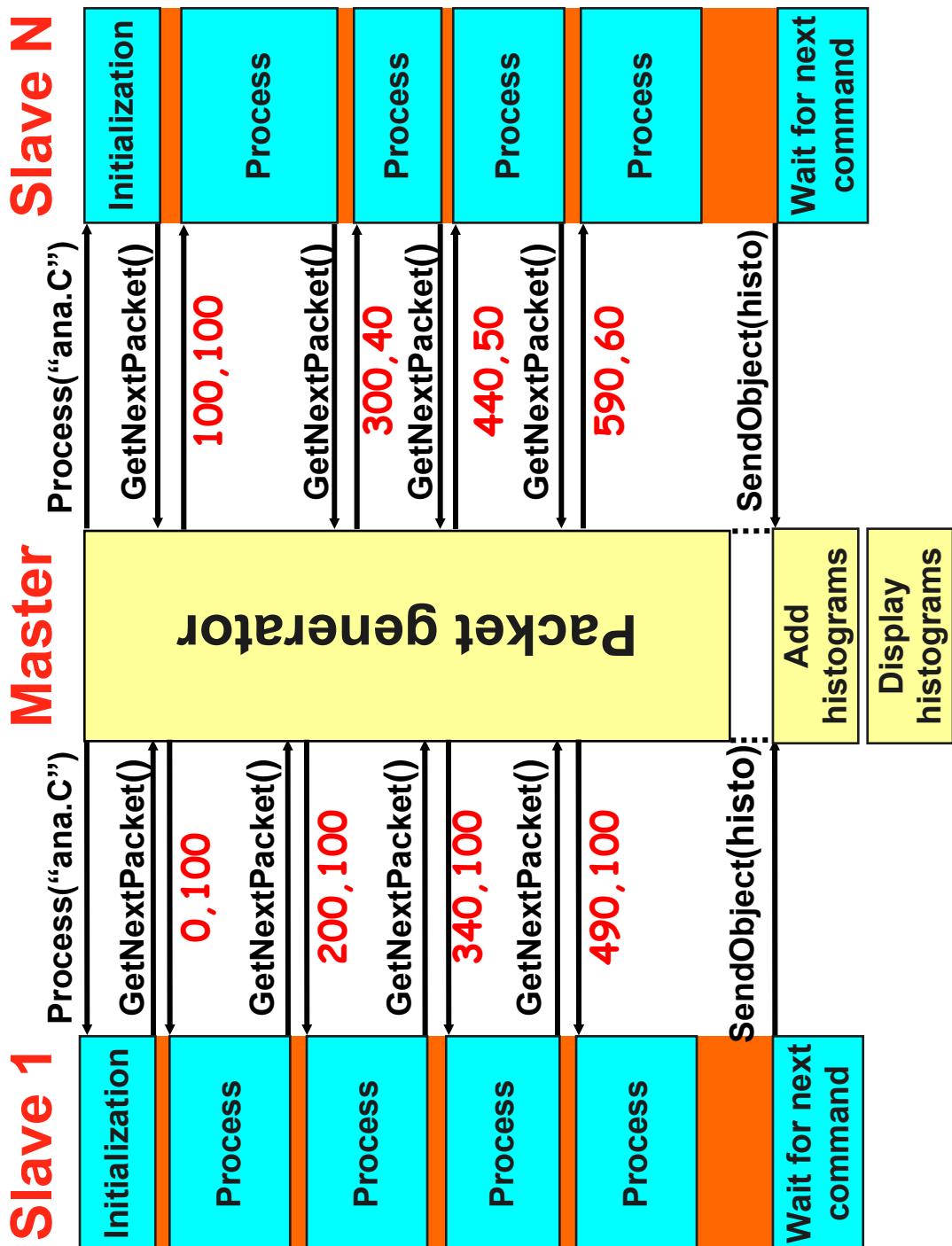


PROOF Adaptability



- Adaptability means to be able to adapt to varying conditions (load, disk activity) on slaves
- By using a "pull" architecture the slaves determine their own processing rate and allows the master to control the amount of work to hand out
 - Disadvantage: too fine grain packet size tuning hurts scalability

Workflow For Tree Analysis – Pull Architecture



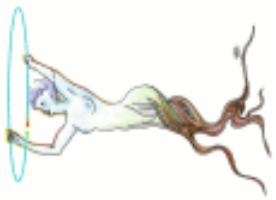
PROOF Error Handling

- Handling death of PROOF servers
 - Death of master
 - Fatal, need to reconnect
 - Death of slave
 - Master can resubmit packets of death slave to other slaves
- Handling of ctrl-C
 - OOB message is send to master, and forwarded to slaves, causing soft/hard interrupt

PROOF Authentication

- PROOF supports secure and un-secure authentication mechanisms
 - Same as for rootd
 - UsrPwd
 - SRP
 - Kerberos
 - Globus
 - SSH
 - UidGid

Architecture and Implementation

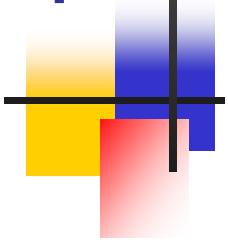




TSelector – The Algorithms

■ Basic ROOT TSelector

```
// Abbreviated version
class TSelector : public TObject {
protected:
    TList *fInput;
    TList *fOutput;
public
    void Init(TTree* );
    void Begin(TTree* );
    void SlaveBegin(TTree * );
    Bool_t Process(int entry);
    void SlaveTerminate();
    void Terminate();
};
```



TDSet – The Data



- Specify a collection of TTrees or files with objects

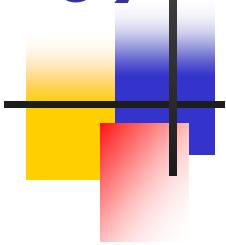
```
root[0] TDSet *d = new TDSet("TTree", "tracks", "/");
OR
root[0] TDSet *d = new TDSet("TEvent", "", "/objs");
root[1] d->Add("root://rcrs4001/a.root");
...
root[10] d->Print("a");
root[11] d->Process("mySelector.C", nentries, first);
```

- Returned by DB or File Catalog query etc.
- Use logical filenames ("fn:...")



Sandbox – The Environment

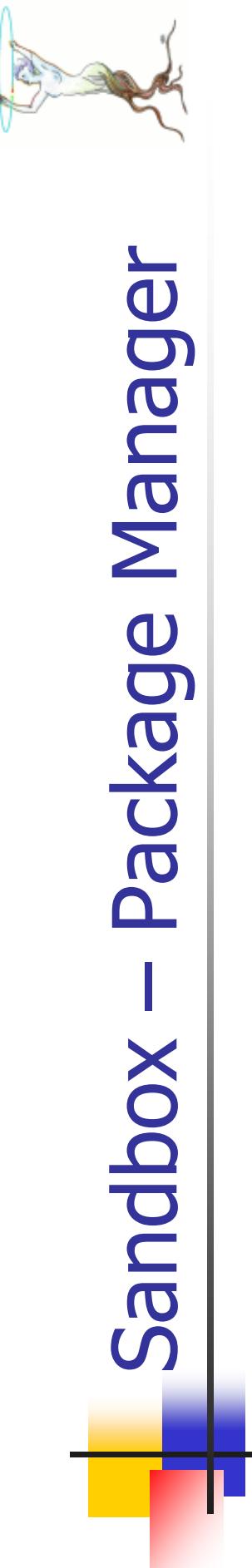
- Each slave runs in its own sandbox
 - Identical, but independent
- Multiple file spaces in a PROOF setup
 - Shared via NFS, AFS, shared nothing
- File transfers are minimized
 - Cache
 - Packages



Sandbox – The Cache

- Minimize the number of file transfers
 - One cache per file space
- Locking to guarantee consistency
- File identity and integrity ensured using
 - MD5 digest
 - Time stamps
- Transparent via `TProof::Sendfile()`



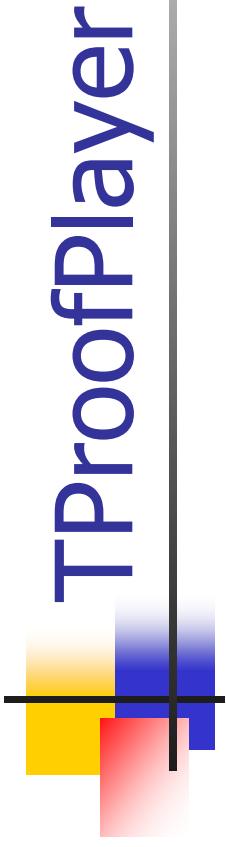


Sandbox – Package Manager

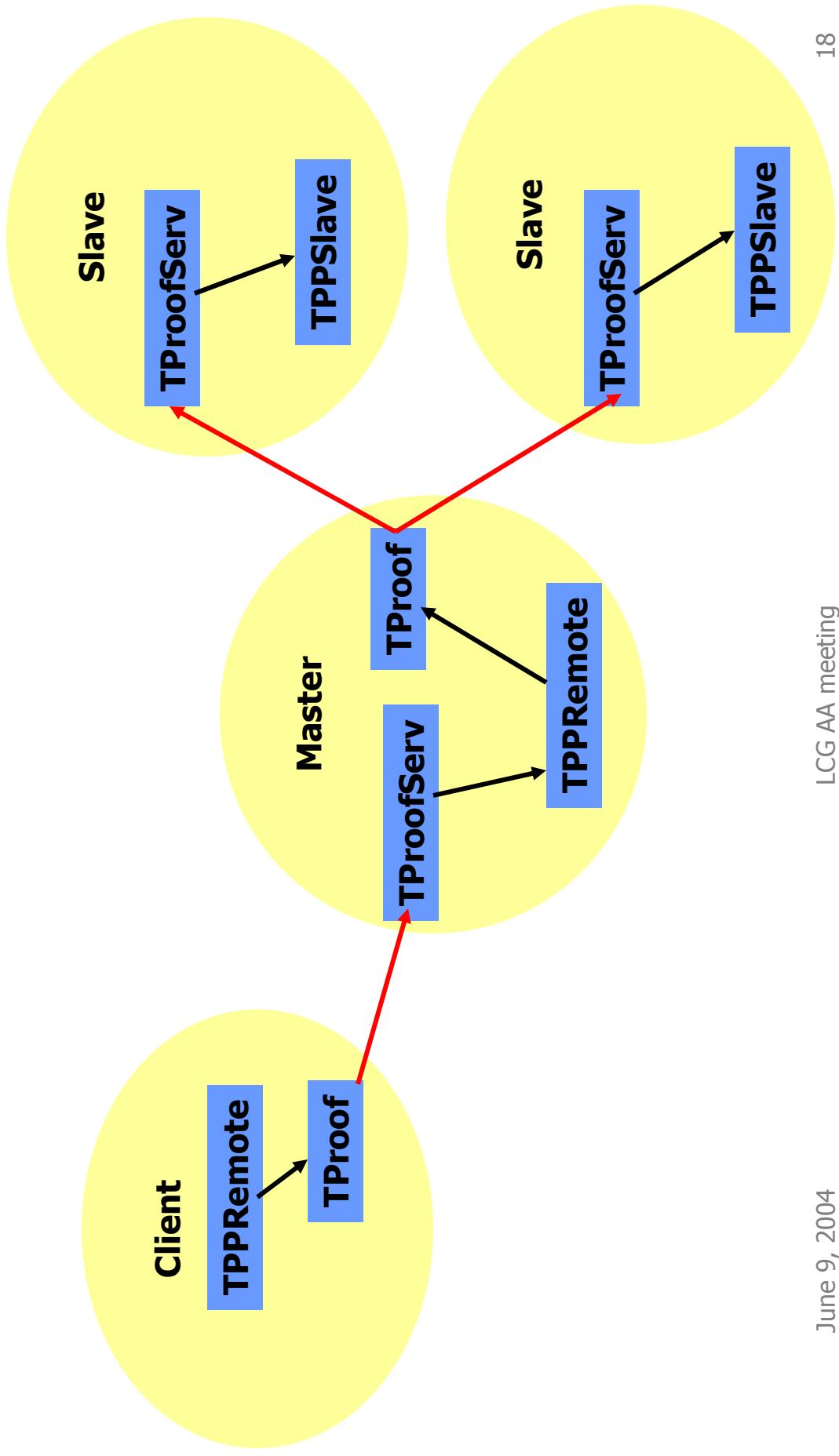
- Provide a collection of files in the sandbox
- Binary or source packages
- PAR files: PROOF ARchive. Like Java jar
 - Tar file, ROOT-INF directory
 - BUILD.sh
 - SETUP.C, per slave setting
- API to manage and activate packages

Implementation Highlights

- TProofPlayer class hierarchy
 - Basic API to process events in PROOF
 - Implement event loop
 - Implement proxy for remote execution
- TEventIter
 - Access to TTree or TObject derived collection
 - Cache file, directory, tree



TProofPlayer

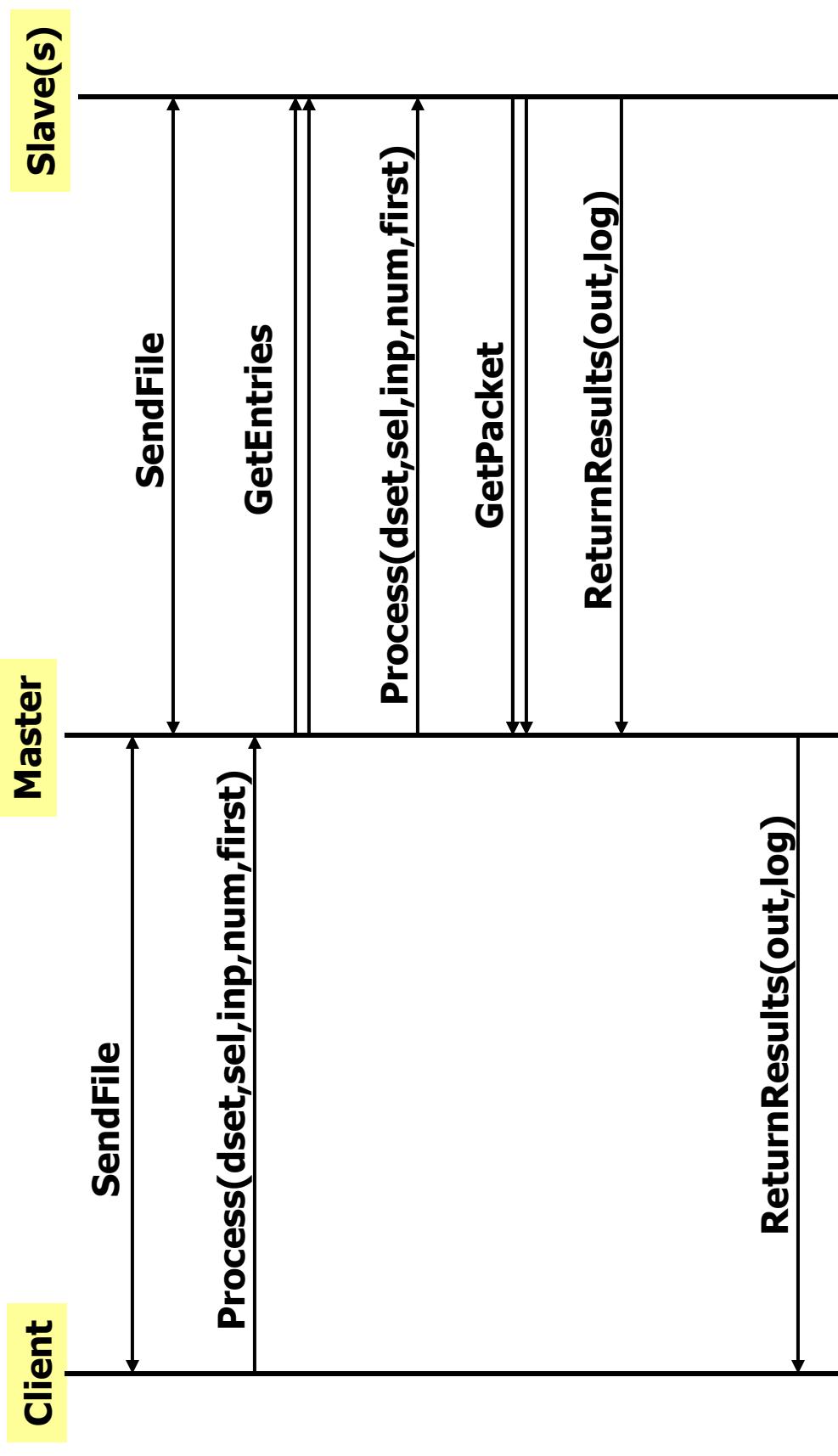


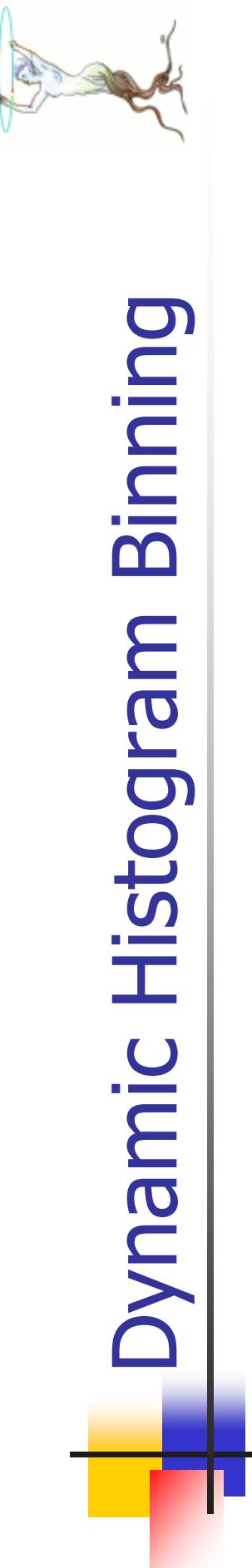
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Simplified Message Flow



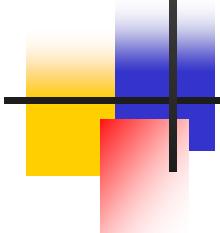


Dynamic Histogram Binning

- Implemented using TH1LimitsFinder class
- Avoid synchronization between slaves
- Keep score-board in master
 - Use histogram name as key
 - First slave posts limits
 - Master determines best bin size
 - Others use these values

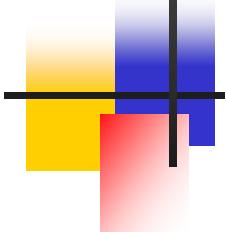
Merge API

- Collect output lists in master server
 - Objects are identified by name
 - Combine partial results
- Member function: Merge(TCollection *)
 - Executed via CINT, no inheritance required
 - Standard implementation for histograms and (in memory) trees
 - Otherwise return the individual objects





PROOF Scalability



8.8GB, 128 files

1 node: 325 s

32 nodes in parallel: 12 s

PROOF Scalability

Events/sec

5000000

4000000

3000000

2000000

1000000

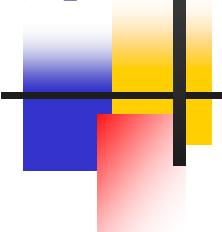
0

0 5 10 15 20 25 30 35 CPU's

32 nodes: dual Itanium II 1 GHz CPU's,
2 GB RAM, 2x75 GB 15K SCSI disk,
1 Fast Eth, 1 GB Eth nic (not used)

Each node has one copy of the data set
(4 files, total of 277 MB), 32 nodes:
8.8 Gbyte in 128 files, 9 million events

Setting Up PROOF



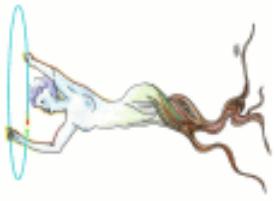
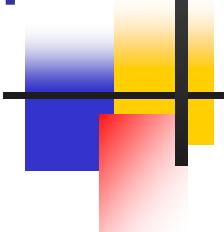
Setting Up PROOF

- Install ROOT system
- For automatic execution of daemons add proofd and rootd to /etc/inetd.conf (or in /etc/xinetd.d) and /etc/services (not mandatory, servers can be started by users)
 - The rootd (1094) and proofd (1093) port numbers have been officially assigned by IANA
- Setup proof.conf file describing cluster
- Setup authentication files (globally, users can override)

PROOF Configuration File

```
# PROOF config file. It has a very simple format:  
#  
# node <hostname> [image=<imagename>]  
# slave <hostname> [perf=<perfindex>]  
# [image=<imagename>] [port=<portnumber>]  
# [srp | krb5]  
# user <username> on <hostname>  
  
node csc02 image=nfs  
  
slave csc03 image=nfs  
slave csc04 image=nfs  
slave csc05 image=nfs  
slave csc06 image=nfs  
slave csc07 image=nfs  
slave csc08 image=nfs  
slave csc09 image=nfs  
slave csc10 image=nfs
```

The AliEn GRID



AliEn - A Lightweight GRID

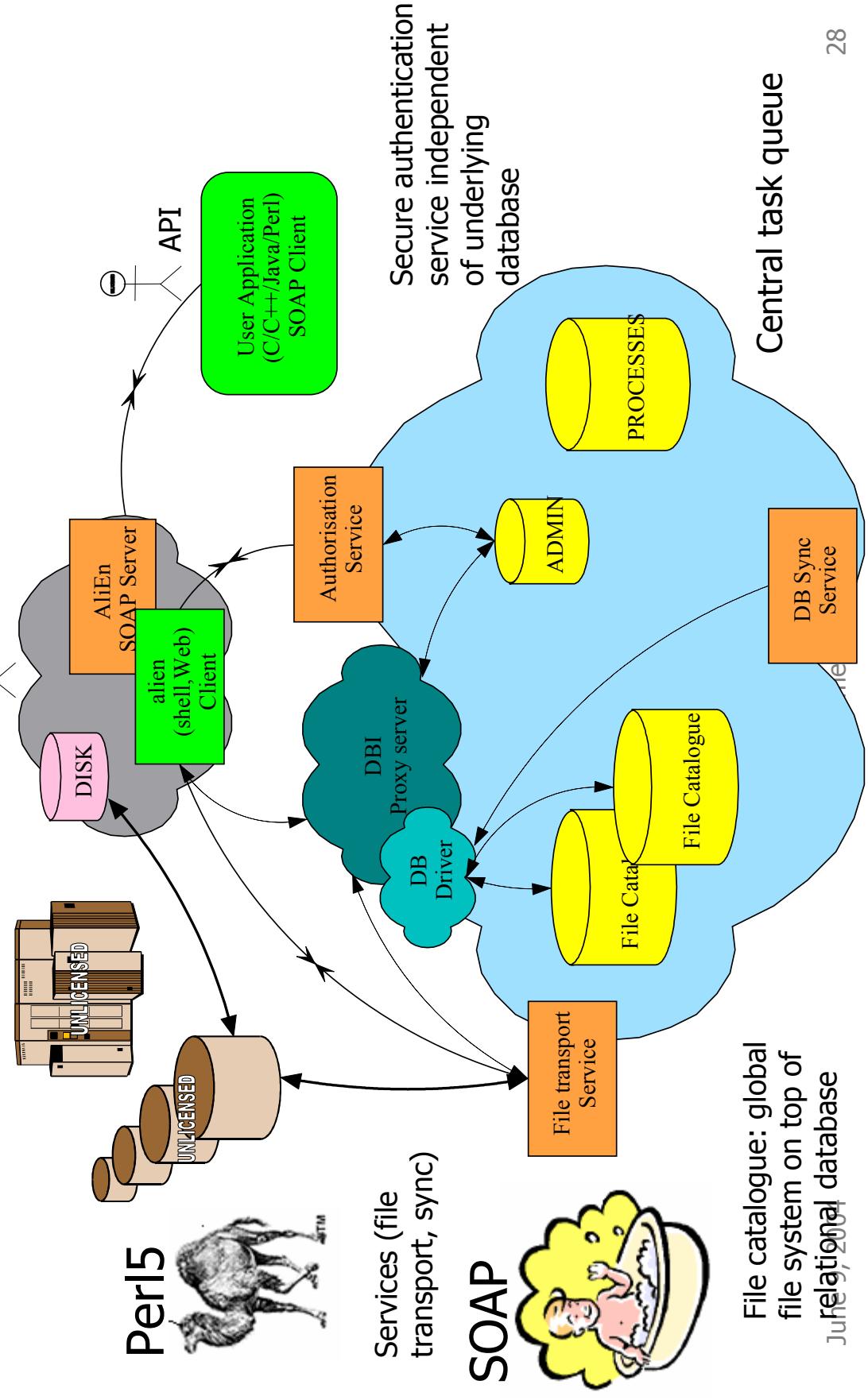
- AliEn (<http://alien.cern.ch>) is a lightweight alternative to full blown GRID
 - based on standard components (SOAP, Web services)
 - Distributed file catalogue as a global file system on a RDBMS
 - TAG catalogue, as extension
 - Secure authentication
 - Central queue manager ("pull" vs "push" model)
 - Monitoring infrastructure
 - C/C++/perl API
 - Automatic software installation with AliKit

The Core GRID Functionality !!

- AliEn is routinely used in the different ALICE data challenges
- AliEn has been released as the EGEE GLite prototype

AliEn Components

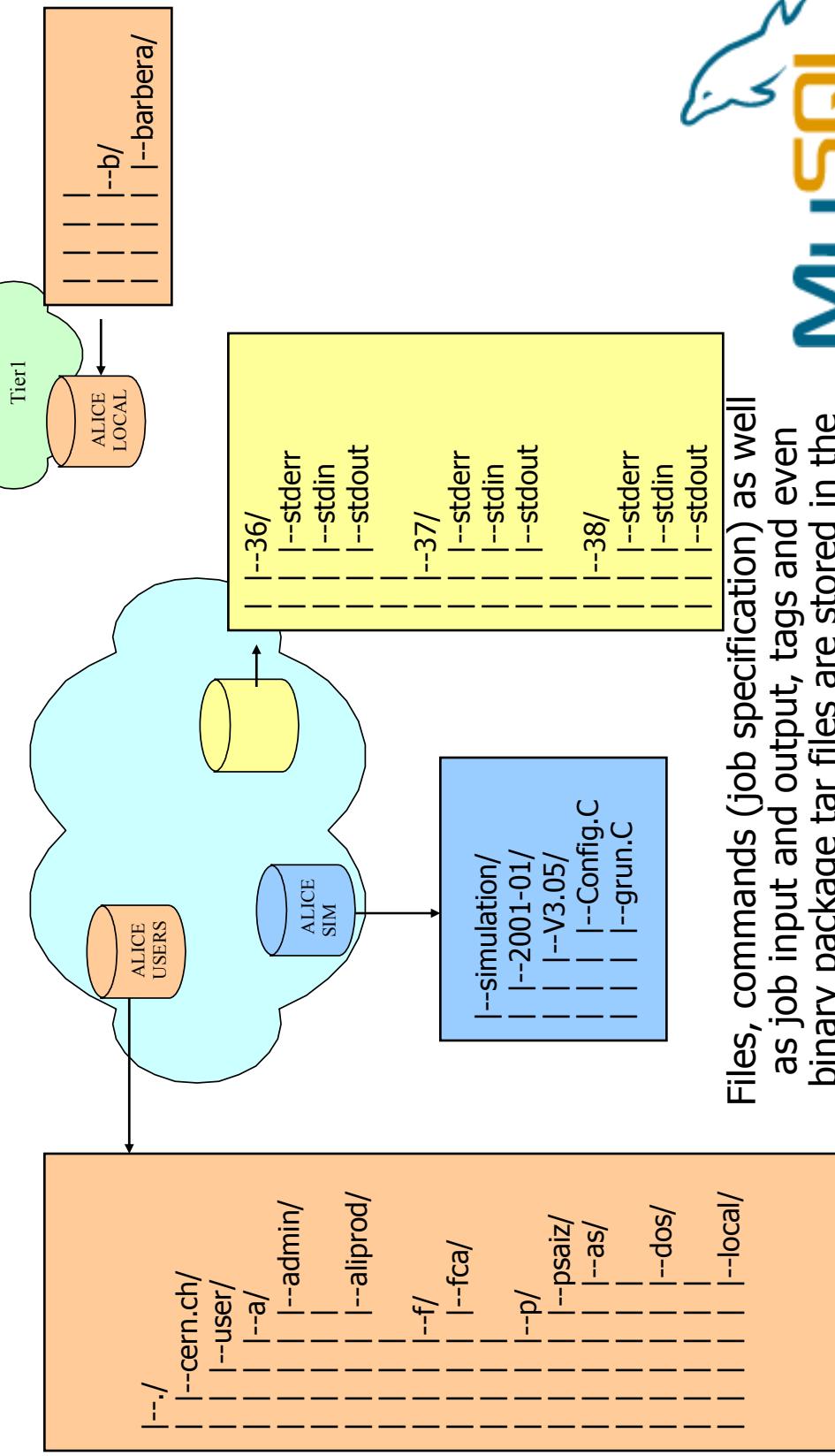
Architecture



AliEn Components



File catalogue

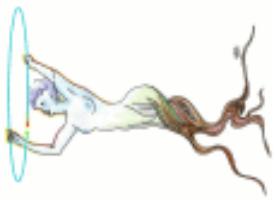
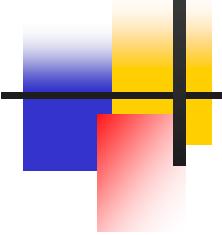


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PROOF and the GRID



PROOF Grid Interface

- PROOF can use a Grid Resource Broker to detect which nodes in a cluster can be used in the parallel session
- PROOF can use Grid File Catalogue and Replication Manager to map LFN's to PFN's
- PROOF daemons can be started by Grid job scheduler
- PROOF can use Grid Monitoring Services
- Access via abstract Grid interface

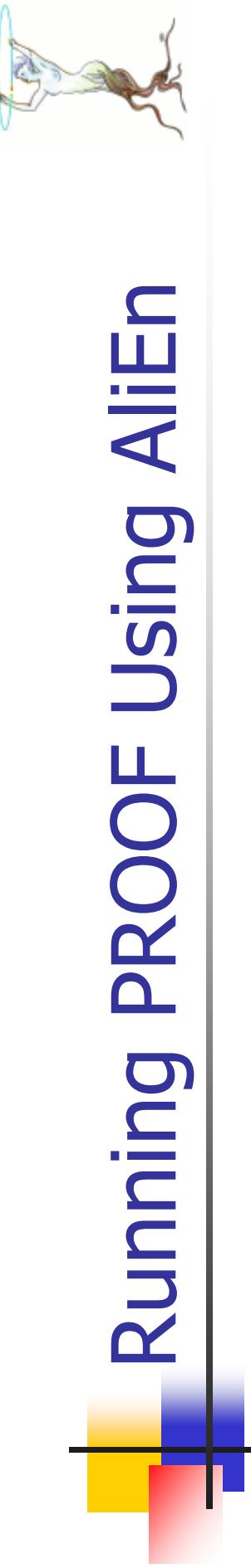
TGrid Class –

Abstract Interface to AliEn

```
class TGrid : public TObject {  
public:  
    virtual Int_t AddFile(const char *lfn, const char *pfn) = 0;  
    virtual Int_t DeleteFile(const char *lfn) = 0;  
    virtual TGridResult *GetPhysicalFileNames(const char *lfn) = 0;  
    virtual Int_t AddAttribute(const char *lfn,  
                               const char *attrname,  
                               const char *attrval) = 0;  
    virtual Int_t DeleteAttribute(const char *lfn,  
                               const char *attrname) = 0;  
    virtual TGridResult *GetAttributes(const char *lfn) = 0;  
    virtual void Close(Option_t *option="") = 0;  
  
    virtual TGridResult *Query(const char *query) = 0;  
  
    static TGrid *Connect(const char *grid, const char *uid = 0,  
                         const char *pw = 0);  
  
    ClassDef(TGrid,0) // ABC defining interface to GRID services  
};
```



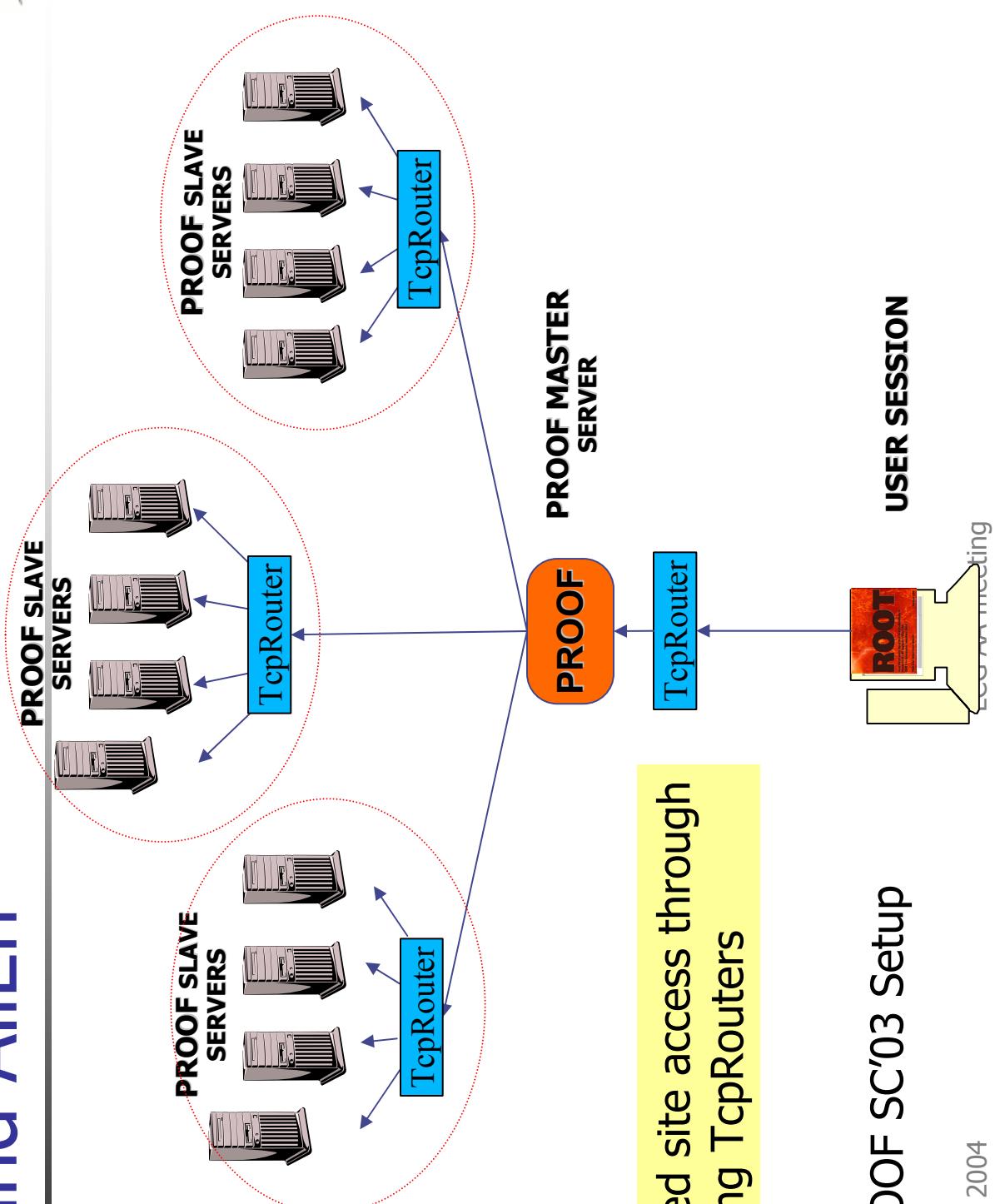
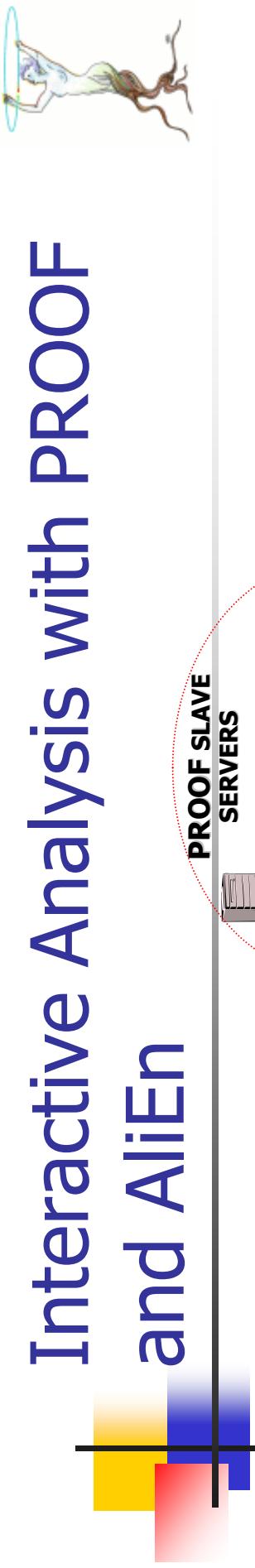
Running PROOF Using Alien



```
TGrid *alien = TGrid::Connect ("alien") ;  
  
TGridResult *res;  
res = alien->Query ("lfn://alice/simulation/2001-04/v0.6*.root") ;  
  
TDSet *treeset = new TDSet ("TTree", "AOD") ;  
treeset->Add (res) ;  
  
gROOT->Proof (res) ; // use files in result set to find remote nodes  
treeset->Process ("myselector.C") ;  
  
// plot/save objects produced in myselector.C  
. . .
```

This scenario was demonstrated by ALICE at SC'03 in Phoenix

Interactive Analysis with PROOF and AliEn



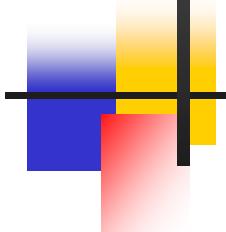
Guaranteed site access through
Multiplexing TcpRouters

AliEn/PROOF SC'03 Setup

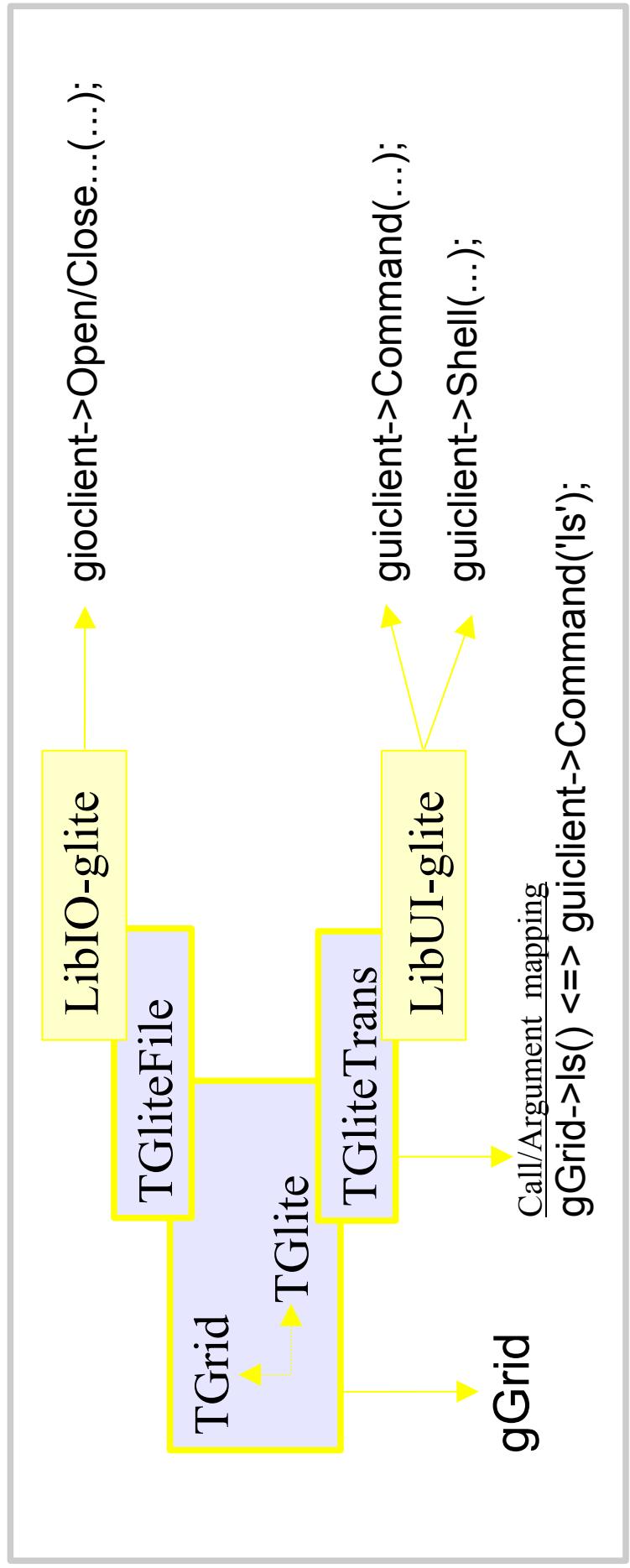
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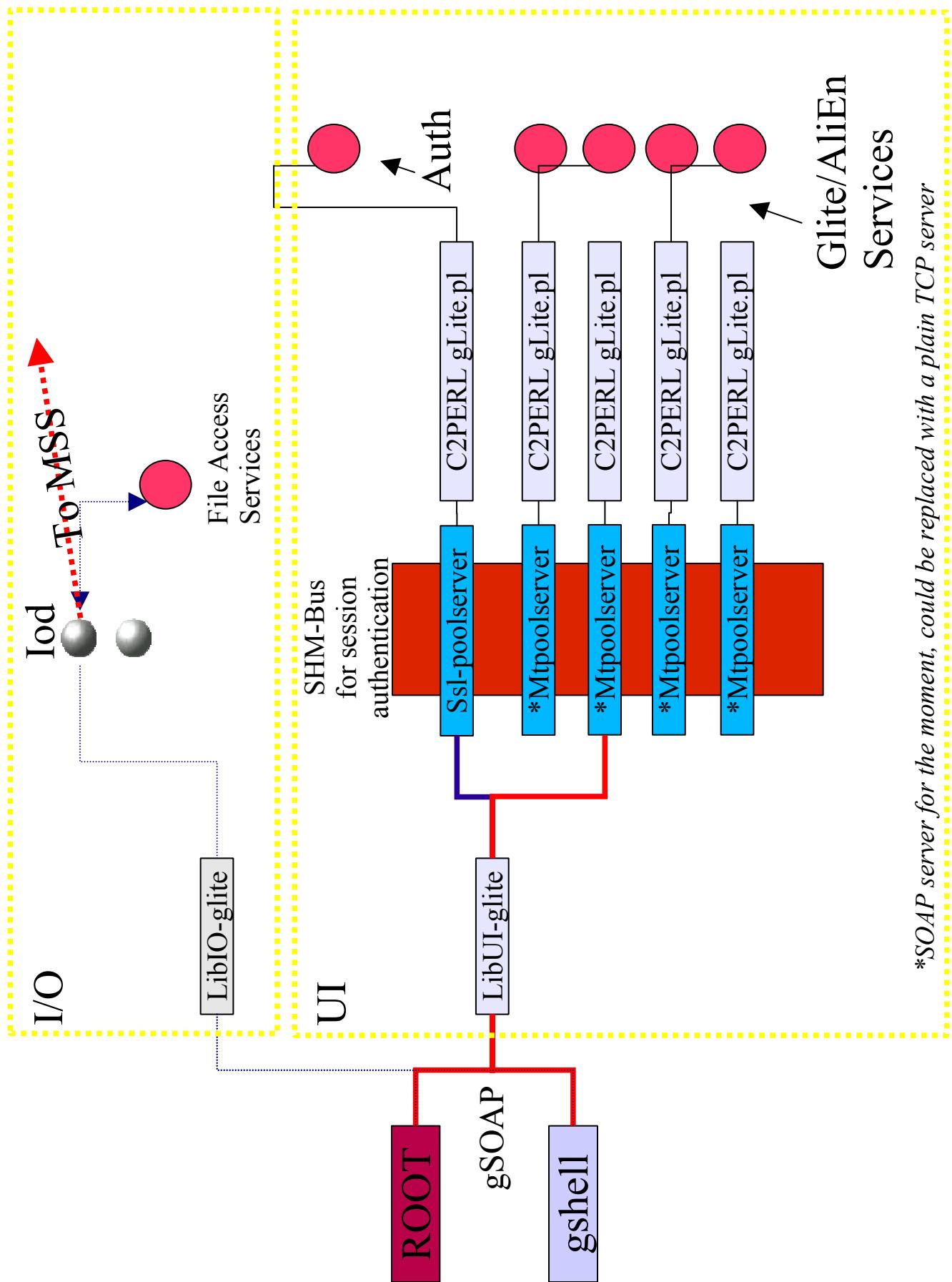
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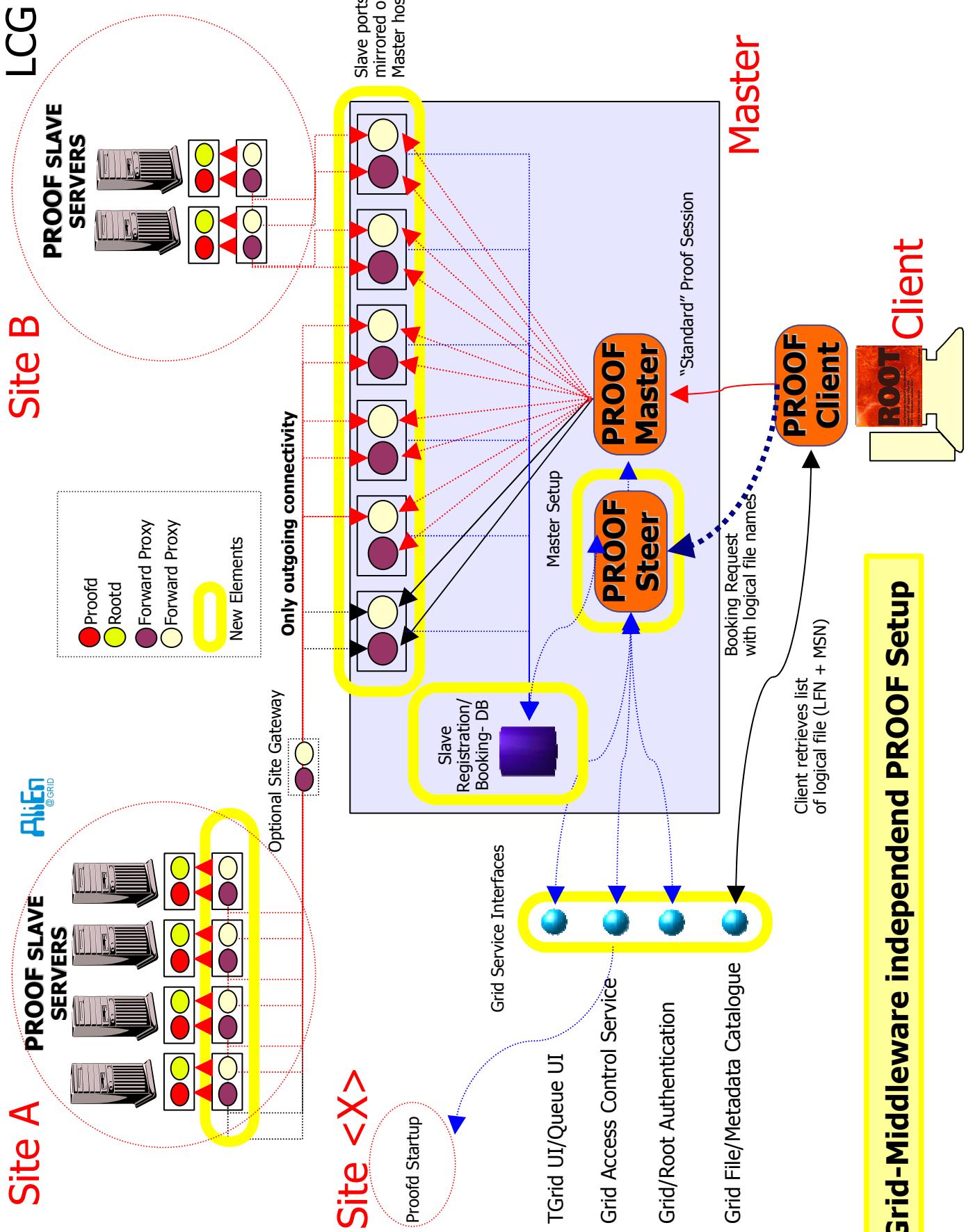
PROOF and GLite



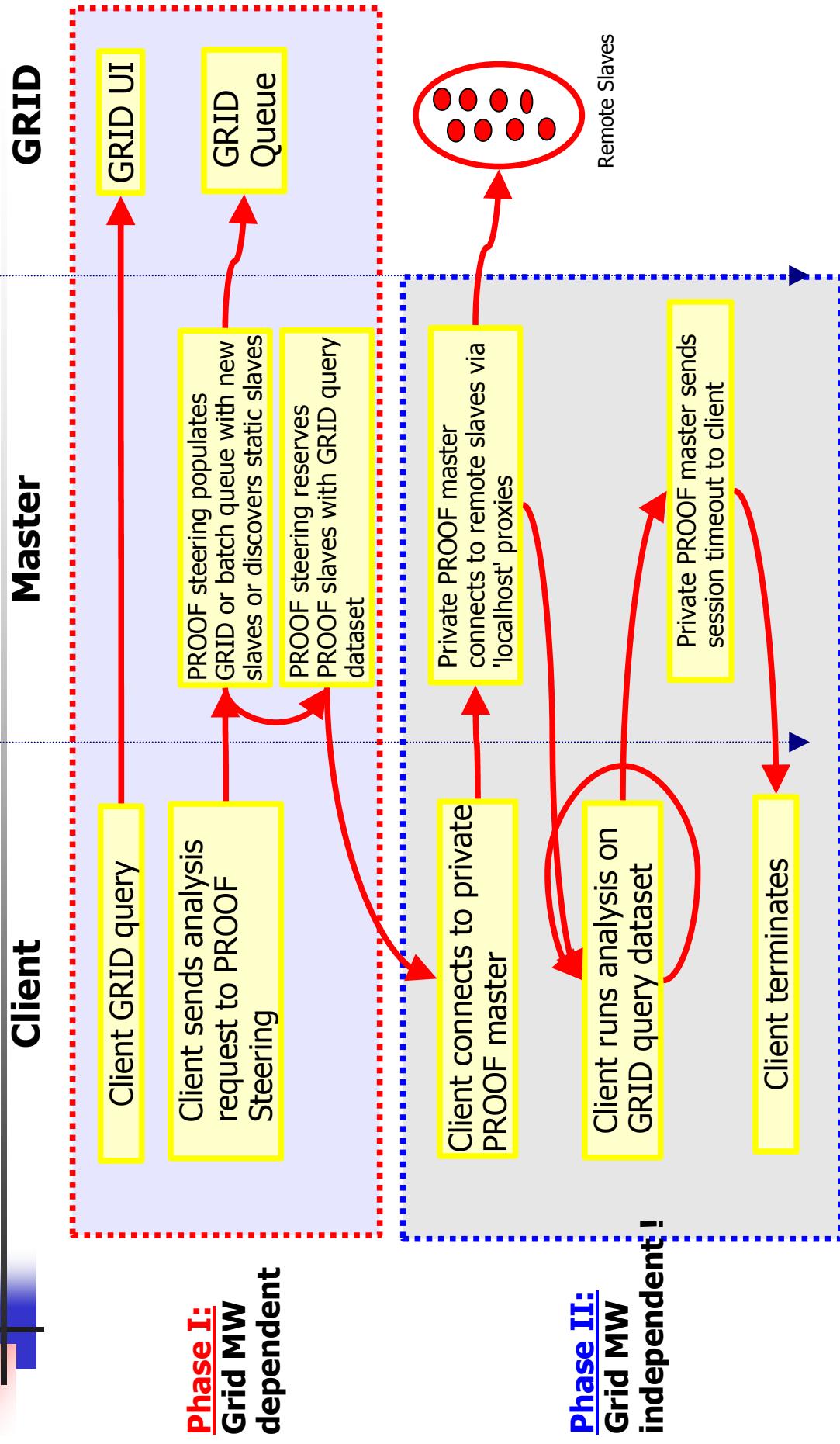
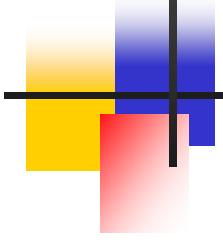
TGlite Interface







PROOF Session Diagram



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Conclusions

- The PROOF system on local clusters provides efficient parallel performance on up to O(100) nodes
- Combined with Grid middleware it becomes a powerful environment for “interactive” parallel analysis of globally distributed data
- ALICE plans to use the PROOF/GLite combo for phase 3 of its PDC’04

