

Outline



Recent releases

Performance tests and optimization

- Query optimization on small data samples
- Scalability tests on large simulated samples
- Support of actual deployment with real data
- Work in progress



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COOL releases: experiments



- <u>COOL 2.4.0</u> (28 Feb 2008)
 - Current Atlas production release (2008 DB tests)

• <u>COOL 2.5.0</u> (10 Jun 2008)

- Used in LHCb, while Atlas has not picked this up
 - Requires user code changes (removal of SEAL)
- **<u>COOL 2.6.0</u>** (10 Nov 2008: yesterday!)
 - Will soon be integrated by both Atlas and LHCb

COOL 2.6.0a ? (Dec 2008 ?)

 Rebuild: new externals (ROOT 5.22, Boost 1.36) and new platforms (gcc43/SLC5?, VC9) for 2009



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COOL releases: performance

• COOL 2.4.0 (28 Feb 2008)

• <u>COOL 2.5.0</u> (10 Jun 2008)

- Major internal cleanup of queries and hints: same SQL, C++ and hints for all use cases
 - Faster and more stable queries for MV user tags
 - Using 240, Atlas had an issue fixed using SQL profiles
 - Faster and more stable queries for MV HEAD

COOL 2.6.0 (10 Nov 2008: yesterday!)

- Remove unnecessary COUNT(*) queries
- Faster IOV bulk insertion for MV closed IOVs
- Complete internal cleanup of query use cases
- Payload queries: should monitor performance



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- Proactive optimization on small samples
 - SQL query optimization of retrieval and insertion

Scalability tests: simulate large samples

- Atlas reference workload: store/read 10 years

React to requests from real deployment

- e.g. remove COUNT(*) queries in COOL 2.6.0
- Test the system with simultaneous connections



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Performance optimization

Main focus: performance for Oracle DBs
 Master T0 database for both Atlas and LHCb

Proactive performance test on small tables

- Test main use cases for retrieval and insertion
- Query times should be flat as tables grow larger
 - e.g. avoid full table scans
- Oracle performance optimization strategy
 - Basic SQL optimization (fix indices and joins)
 - Use hints to stabilize execution plan for given SQL
 - Instability from unreliable statistics, bind variable peeking
 - Determine best hints from analysis of "10053 trace" files



Systematic tests of known causes of instabilities

- Bind variable "peeking", missing or stale "statistics"
 - Instabilities observed in the Atlas 2007 tests (e.g. CNAF vs. Lyon)
- Stable performance after adding Oracle hints





(single IOV)

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COOL 2.5.0: MV user tags

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IOV retrieval from multi-version "user tags"
 Used by Atlas for all MV data (e.g. calibrations)

0.8 COOL-240--stat-peeklo COOL-240--stat-peekhi COOL-240--nost-peeklo COOL-240--nost-peekhi 0.7 COOL-HEAD-stat-peeklo COOL-HEAD-stat-peekhi COOL-HEAD-nost-peeklo COOL-HEAD-nost-peekhi 0.6 0.5 0.4 0.3 0.2 0.1 0 Øk 2k 4k 6k 8k 10k 12k 14k 16k 18k 20k Position of IOVs being queried

User Tag browseObjects()

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COOL 2.6.0: MV insertion



IOV insertion in multi-version case
 Problem was for [t1,t2], not for [t1,+inf]

Bug 17903 in bulk insertion (closedIOVs)



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COOL 2.6.0: COUNT(*)



 In COOL 2.5.0, IOVs were always counted before they were retrieved

- select count(*) from (select ... from IOVS)
- select ... from IOVS

In COOL 2.6.0, this is done on demand



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Scalability tests



- Proactive performance test on large tables
 - Stable insertion and retrieval rates (>1k rows/s)
 - Simulate data sets for 10 year of LHC operation



Test case: Atlas "DCS" data

- Measured voltages, currents...
- Largest Atlas data set
 - 1.5 GB (2M IOVS) / day

To do next: data partitioning

- Goal: ease data management
- Evaluating Oracle partitioning
 - Test possible performance impact



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Work in progress



Session sharing

- An Atlas job reads data from 10 COOL schemas
 - This now requires 10 sessions to the same DB
 - Will change COOL to use a single session

Schema improvements

- Store payload and metadata in 2 separate tables
 - Will reduce storage overhead (MV payload duplication)
- Partitioning

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- See next talk



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Future deployment model





DB access via CORAL server

- Address secure authentication and connection multiplexing
- Development still in progress
 - Only minimal changes in COOL





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Manage conditions data of Atlas and LHCb

- <u>Time variation</u> (validity) and <u>versioning</u> (tags)
- Common project of Atlas, LHCb, CERN IT



Support for several relational databases

- Oracle, MySQL, SQLite, Frontier
- Access to SQL from C++ via the CORAL libraries



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COOL data model

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Modeling of conditions data objects

- System-managed common "metadata"

- Data items: many tables, each with many channels
- Interval of validity "IOV" [since, until]
- <u>Versioning information</u> with handling of interval overlaps
- User-defined schema for "data payload"
 - Support for fields of simple C++ types

Main use case: event reconstruction

- Lookup data payload valid at a given event time





COOL collaborators



Core development team

- Andrea Valassi (CERN IT-DM)
 - 80% FTE (core development, project coordination, release mgmt)
- Marco Clemencic (CERN LHCb)
 - 20% FTE (core development, release mgmt)
- Sven A. Schmidt (Mainz ATLAS)
 - 20% FTE (core development)
- Martin Wache (Mainz ATLAS)
 - 80% FTE (core development)
- Romain Basset (CERN IT-DM)
 - 50% FTE (performance optimization) + 50% FTE (scalability tests)
- On average, around 2 FTE in total for development since 2004

Collaboration with users and other projects

- Richard Hawkings and other Atlas users and DBAs
- The CORAL, ROOT, SPI and 3D teams

Former collaborators

• G. Pucciani, D. Front, K. Dahl, U. Moosbrugger



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- **COOL basics** (only what is needed to understand the rest...)
 - Data model basics
 - Use case for this talk: MV tags (relational schema and SQL query)
 - Performance plots (how we define 'good' performance)
- Oracle performance optimization strategy
 - Basic SQL optimization (fix indexes and joins)
 - Execution plan instabilities (same SQL, different plans)
 - Observe (causes: unreliable statistics, bind variable peeking)
 - Analyze (<u>10053 trace files</u> and the BEGIN_OUTLINE block)
 - Fix (rewrite queries to please the Optimizer; then add <u>hints</u>)







Conditions data

- Detector data that vary in time and may be versioned
- Several use cases (different schemas and SQL queries to optimize)
 - Temperatures, voltages (measured single version, SV)
 - Calibration, alignment (computed multiple versions, MV)
- **COOL conditions objects** ("IOV"s interval of validity)
 - Metadata: channel (c), IOV (t_{since} , t_{until}), version or tag (v)
 - Data: user-defined "payload" (x1,x2,...)
 - Typical query: retrieve the condition data payload X that was valid at time T in channel C for tag V

• COOL relational implementation (based on CORAL)

Several backends (Oracle, MySQL...); C++ only (no PL/SQL)



DM Test case: MV tag retrieval



Query: fetch all IOVs in [T1,T2] in tag PROD in all channels

2. For each channel C, select IOVs in tag PROD in [T1, T2]

(this is a very large table – and the most delicate part of the query to optimize)









Is query time the same for all values of parameters T1, T2?

– It was not in the initial COOL releases (≤ COOL 2.3.0)

• Query time is higher for more recent IOVs than for older IOVs

"tagId=PROD AND chId=C AND ((since \leq T1< until) OR (T1 < since \leq T2))"



DM Basic optimization: better SQL

In tag PROD, in each channel at most one IOV is valid at T1

- Build a better SQL strategy from this constraint (unknown to Oracle)
 - The constraint is enforced in the C++ code, not in the database







- So, we thought the job was done...
 - Query time used to increase, we managed to make it flat
- But... every now and then our tests or our users reported performance issues again (...?...)
 - Example: different performance in ATLAS tests at CNAF and LYON
- Symptoms: same SQL, different execution plan
 - In time, we identified two possible causes for this:
 - Bind variable peeking
 - Optimal exec plan for finding old IOVs and recent IOVs are different
 - Problem if optimal plan for old IOVs is used for finding recent IOVs
 - Missing or unreliable statistics
 - Optimal exec plan is computed starting from wrong assumptions



Execution plan instabilities: plots

- Systematic study of 6 (2x3) cases
 - 2 cases for b.v. peeking: peek "low" (old IOVs) or "high" (recent IOVs)

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- 3 cases for statistics: none, full, unreliable (empty tables)



Analyze plans: 10053 trace



- Look at the plan that was used for your query execution
 - More reliable than 'explain plan', 'set autotrace' and other methods...
- Look at how and why the Optimizer chose this plan
 - Bind variable values
 - Alternative plans attempted
 - Were user-supplied hints understood and used?
 - The "Dumping Hints" section at the end

• Look at the Optimizer's outline for the chosen plan

- Get inspiration from the outline to prepare your user-supplied hints
 - The "BEGIN_OUTLINE_DATA" section towards the end





- This is an iterative process! In summary:
 - 1. Execute your query for many cases (peek high/low...)
 - 2. Get plan and outline for a case with good performance
 - You want your plan to look like this in the end for all cases
 - 3. Do you need some query rewrite?
 - Are query blocks not named? Add QB_NAME and go to 1.
 - Is Oracle rewriting your query? Change SQL and go to 1.
 - Is Oracle using a different join order? Change SQL and go to 1.
 - 4. Is there a case with bad performance? Get its outline.
 - What is different in 'good' outline? Add as a hint and go to 1.
 - Was your hint not used or not useful? Try another and go to 1.
 - 5. Do all cases have good performance? You made it!







- Generate a 10053 trace file 'myfile.trc'
 - From SQL*Plus
 - ALTER SESSION SET EVENTS
 '10053 TRACE NAME CONTEXT FOREVER, LEVEL 1';
 - ALTER SESSION SET tracefile_identifier='myfile'
 - From CORAL:
 - export CORAL_ORA_SQL_TRACE_ON="10053"
 - export CORAL_ORA_SQL_TRACE_IDENTIFIER="myfile"
- Retrieve the trace file
 - Ask your friendly DBA to get it from the server's udump...
 - But please avoid generating (and asking for) trace files unless you need them...;-)







- You should invalidate existing exec plans between tests
 - To remove the effect of bind variable peeking (e.g. when testing the effect of different bind variable values)
 - To make sure that execution plans are recomputed and ORA-10053 trace files are as complete as possible
- To invalidate existing execution plans you may:
 - Flush the shared pool (DBA only affects the whole DB)
 - Simpler hack: alter a relevant table in a dummy way
 - e.g. "ALTER TABLE mytable LOGGING;"



Query rewrite – are you in control? Department

Master your query blocks

- Name your query blocks syntax is "/*+ QB_NAME(xxx) */"
 - Else the Optimizer will name them for you (e.g. "SEL\$1")
- The Optimizer rewrites your query blocks? Do it yourself!
 - Symptoms: query block names like "SEL\$3F979EFD", keywords like "MERGE" (remove inline views) or "CONCAT" (expand as union all)
 - Solution: do what the Optimizer would do (e.g. remove MERGE by expanding subqueries in WHERE clause into normal joins)

• Master the order of your joins

- The Optimizer reorders your joins? Do it yourself!
 - Copy the Optimizer's favorite order from the "LEADING" keyword







Default hints added in COOL 2.3.1 release

Standard HEAD tag MC browseObjects()

- Stable good plans in all 6 cases (2 bind var peeking x 3 statistics)









- No support for hints
 - Implemented in COOL queries using SQL injection
 - Prepend the hint "/*+...*/" to the 1^{st} item in the SELECT list
 - This hack does not work for UPDATE, INSERT, DELETE
 - CORAL support request <u>sr #103420</u>
- No support for subqueries in WHERE clause
 - Implemented in COOL queries using SQL injection
 - CORAL receives a WHERE clause that explicitly contains a fully qualified "(SELECT ... FROM ...) " subquery
 - COOL needs to know if it is talking to Oracle or MySQL (quotes)
 - CORAL support request <u>sr #103547</u>



M Performance – in progress



- Handle all use cases consistently in C++ code
 - SV, MV tags (~CVS tags) or 'user tags' (~CVS branches)
 - Goal: same performance optimization in all use cases
 - Share a single C++ method to define the general SQL strategy (with internal switches for use-case-dependent SQL fragments)
 - So far each use case was optimized separately

Evaluate Oracle partitioning

- Goal: ease data management (long-term archiving)
 - Partitioned tables with partitioned (local) indexes
- Evaluate impact (benefits?) for performance too

Performance for non-Oracle backends

- Using the same SQL is not always possible
 - MySQL performance is bad with subqueries
- Lower priority

