

# Increasing Tape Efficiency

HEPiX Fall 2008 Taipei

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- **Tape efficiency project**
- **Problem areas**
- **What has been done**
- **What is under development**
- **Roadmap**
- **Summary**



- All functionality dealing directly with storage on and management of tapes
  - Volume database
  - Migrations/recalls
  - Tape drive scheduling
  - Low-level tape positioning and read/write
- Team is from IT/DM
- Contributions from IT/FIO





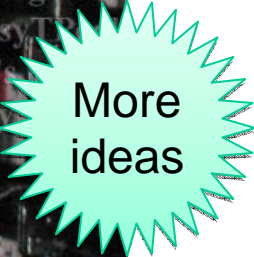
Work done

- Write more data per tape mount



Current work

- Use a more efficient tape format
  - The current tape format does not deal efficiently with small files



More ideas

- Improve read efficiency
  - Require modifications from disk to tape



DM

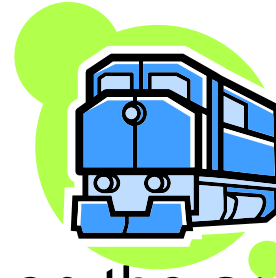
# What has been done



- Recall/migration policies

- Freight train approach

- Hold back requests based on the amount of data and elapsed time

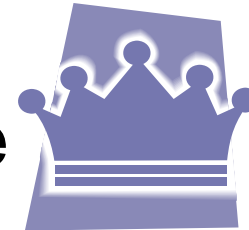


- Production managers rule

- Production managers plan relatively **large workloads** for CASTOR

- **Access control lists** give production managers a relatively larger percentage of resources

- **User and group based priorities** encourage users to work with their production managers





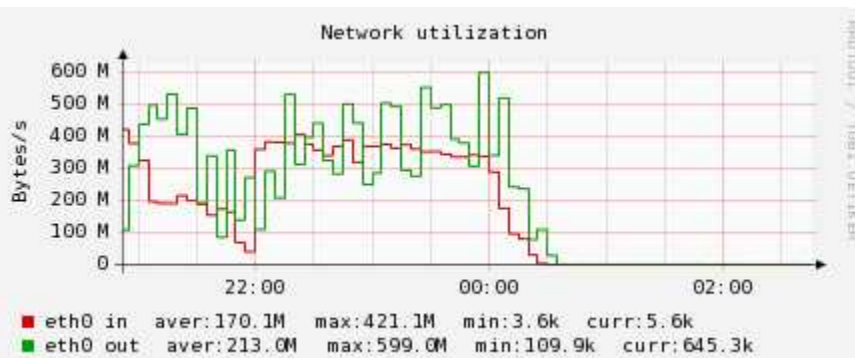
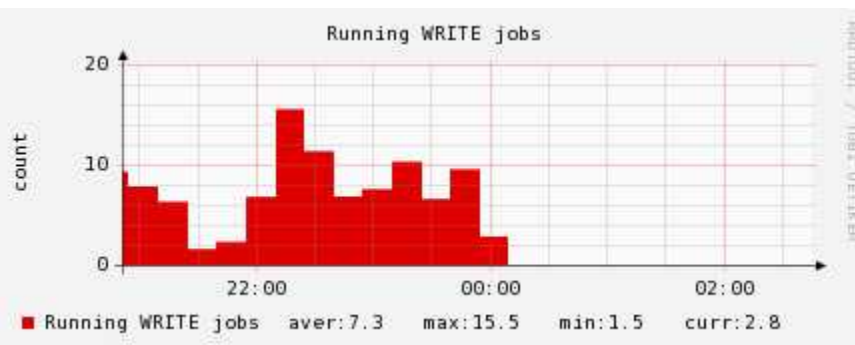
- Repacks the data from one set of tapes onto another set of tapes
- Repack is used for media migration
- Repack is used to defragment tapes



- Reading the current ANSI AUL format is approximately twice as fast as writing
- Repack uses Castor as a cache
- Repack uses the cache to support asymmetric read/write drive allocation
- Repack is equivalent to one LHC experiment and as such is a good test run for Castor







- 4 drives reading
- 7 drives writing
- 400MBytes/s



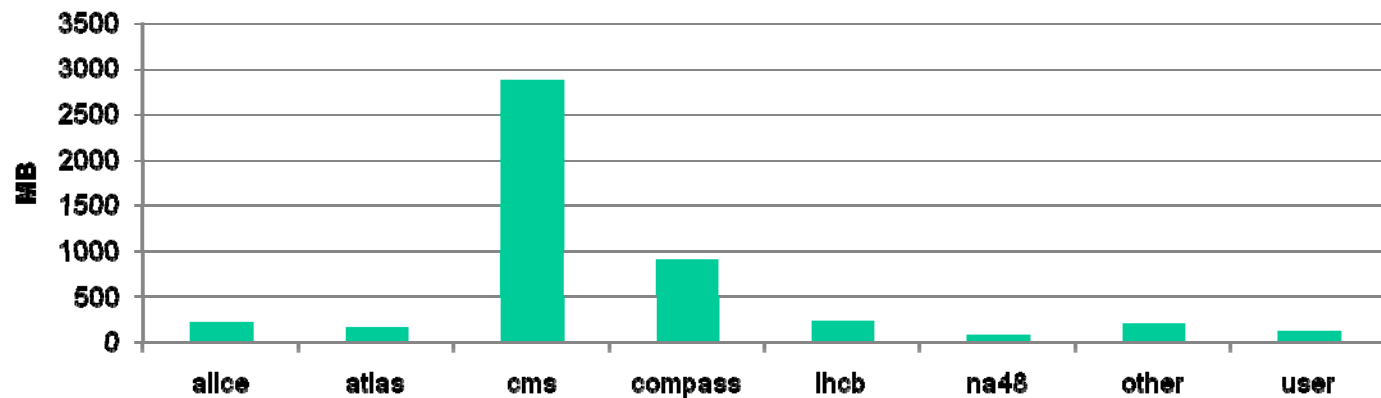
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# What is under development

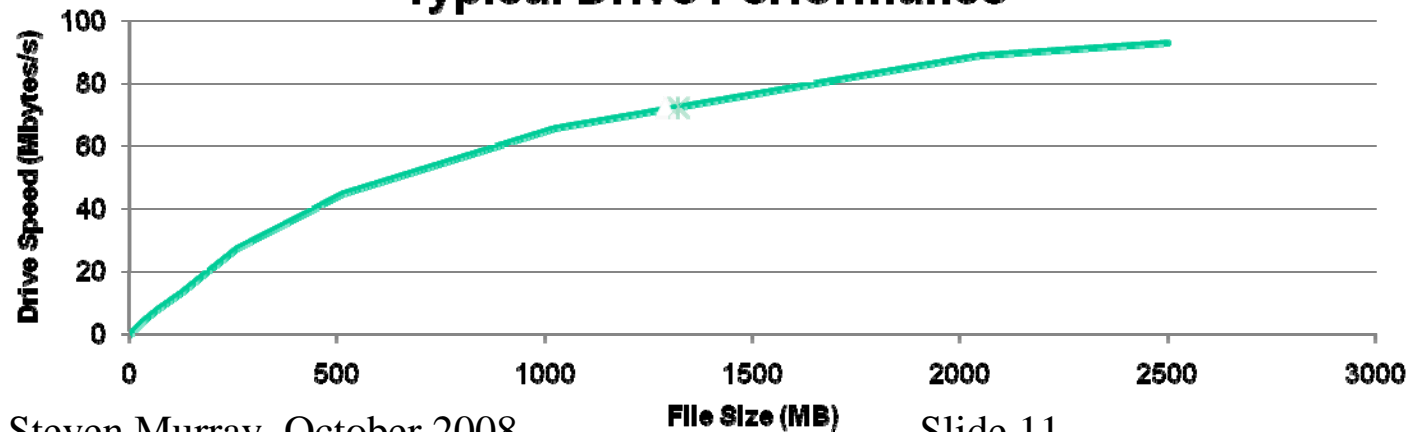


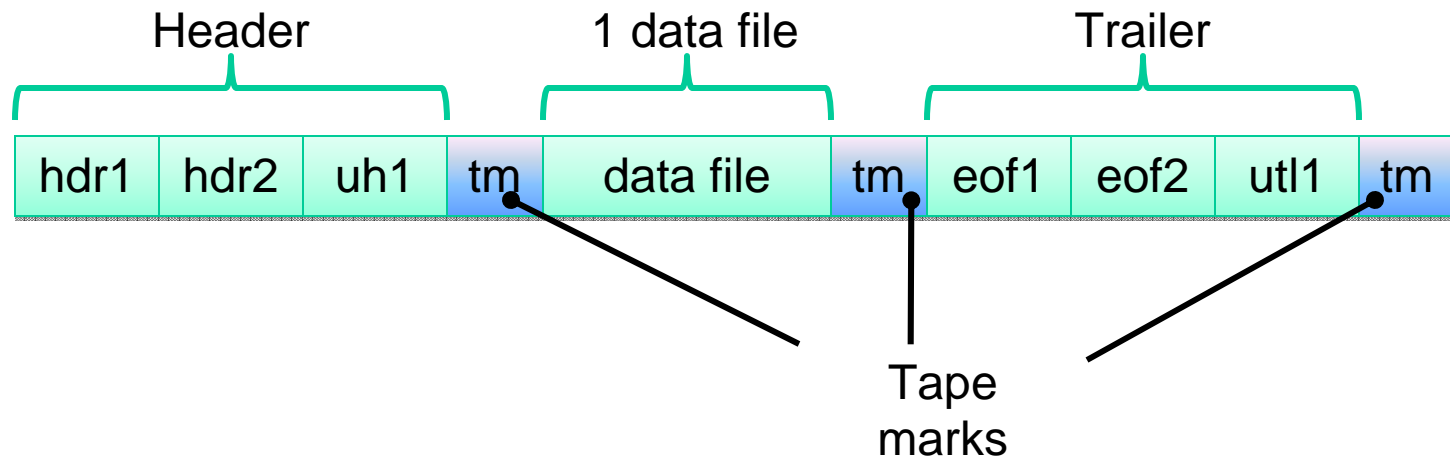
- Users were encouraged to store large files in Castor
- Unfortunately Castor contains many small files

Average Filesize per VO



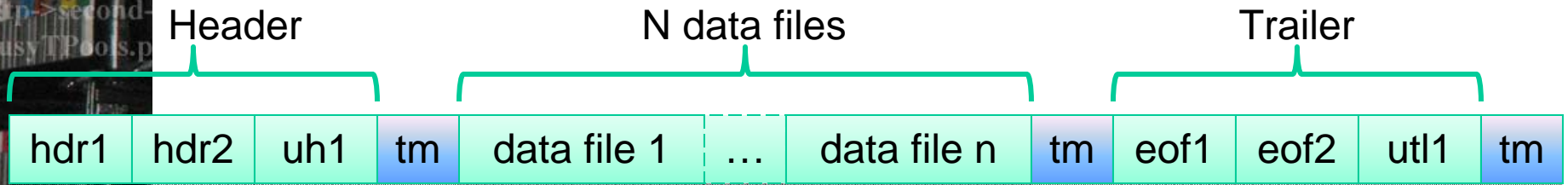
Typical Drive Performance





- ANSI AUL format
- 3 tape marks per file
- 2 to 3 second per tape mark
- 9 seconds per data file independent of its size





Each 256 KB data file block written to tape includes a 1 KB header

- Multi-file block format within the ANSI AUL format
- Header per block for “self description”
- 3 tape marks per n files
- n will take into account:
  - A configurable maximum number of files
  - A configurable maximum size
  - A configurable maximum amount of time to wait



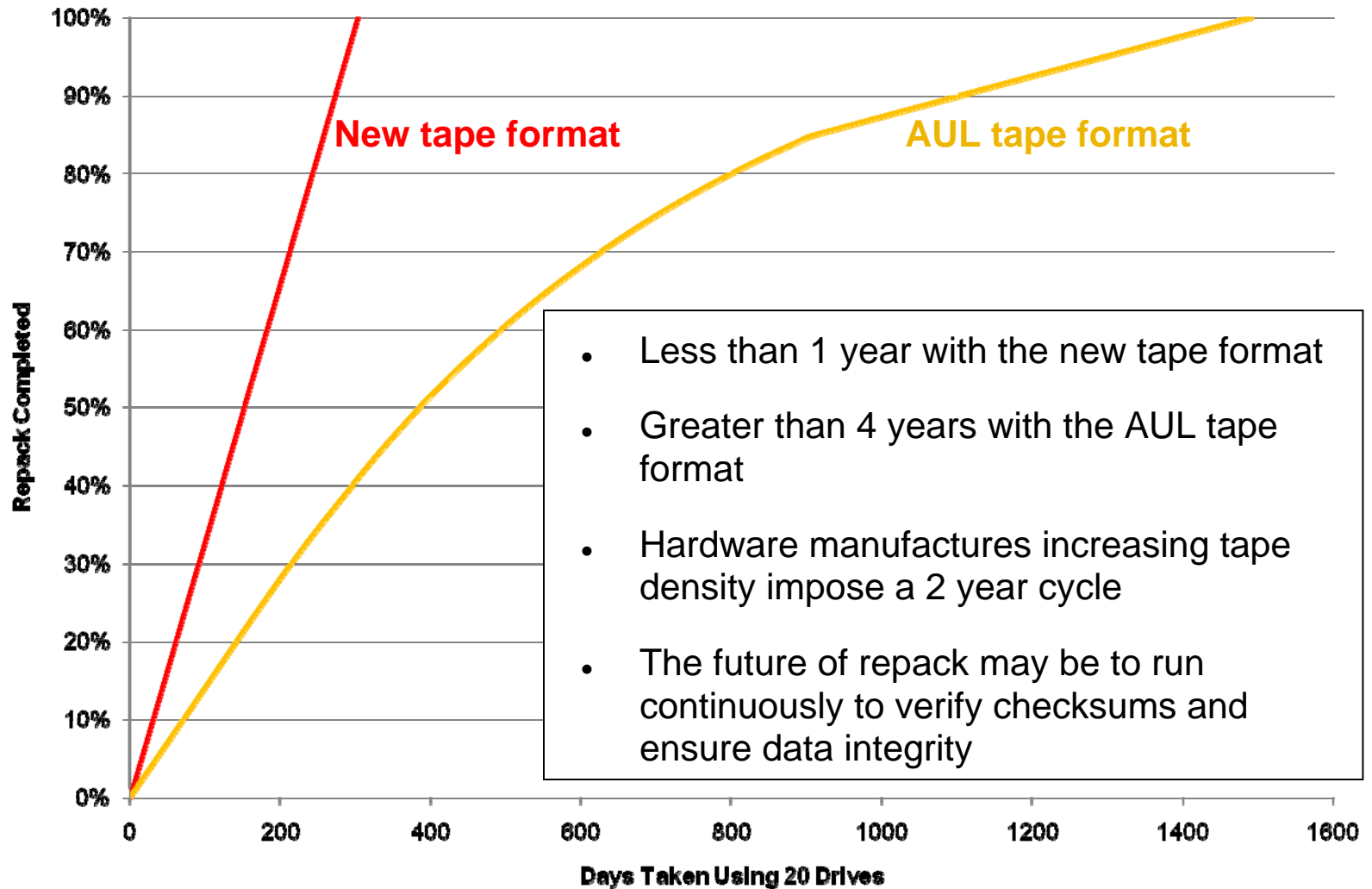


#	Meta-data name	Explanation	Examples for Data	Bytes
1	VERSION_NUMBER	The version of the block format	09.13	5
2	HEADER_SIZE	Header size in bytes	01024	5
3	CHECKSUM_ALGORITHM	Name of the checksum algorithm	Adler-32	10
4	HEADER_CHECKSUM	Adler-32 checksum	4146884724 00000000000000012	10
5	TAPE_MARK_COUNT	Sequential number addressing the migration-files on the tape	345	20
6	BLOCK_SIZE	Block size in bytes inclusive of header	0000262144	10
7	BLOCK_COUNT	Block offset from the beginning of the tape. Tape marks and labels are included in the count	00000000000000012 345	20
8	BLOCK_TIME_STAMP	Time since the Epoch (00:00:00 UTC, January 1, 1970), measured in seconds	1222332810	10
9	STAGER_VERSION	The version of the stager software	2.1.7.18	15
10	STAGER_HOST	The DNS name of the stager host including the domain	c2cms2stager.cern.c	30
11	DRIVE_NAME	Will be provided by a local configuration file	h 0003592028	10
12	DRIVE_SERIAL	Will be provided by a local configuration file	00000000456000001 642	20
13	DRIVE_FIRMWARE	Will be provided by a local configuration file	D3I0_C90	10
14	DRIVE_HOST	The DNS name of the host including the domain	tpsrv250.cern.ch	30
15	VOL_DENSITY	The storage capacity of the tape	700.00GB	10
16	VOL_ID	Site specific numbering system (the sticker on a tape)	T02694	20
17	VOL_SERIAL	Volume Serial Number	T02694	20
18	DEVICE_GROUP_NAME	The device group name that linked the tape to the drive	3592B1 00000001099511627	10
19	FILE_SIZE	The size of the data file in bytes	776	20
20	FILE_CHECKSUM	Adler-32 checksum	1926860616	10
21	FILE_NS_HOST	The DNS name of the host including the domain	castorns.cern.ch	30
22	FILE_NS_ID	The name server ID of the data file	226994274	20
23	FILE_PROGESSIVE_CHECKSUM	Adler-32. Progressive checksum of all the blocks written to tape so far for the current data file	1234567890 00000000000000012	10
24	FILE_BLOCK_COUNT	Block offset from the beginning of the data file	345	20
<b>Header size before file_name :</b>				<b>375</b>
25	FILE_NAME	Last "x" bytes of the filename from the name server. This field acts as a padding to the nearest KiB.		649
<b>Header size :</b>				<b>1024</b>

VERSION\_NUMBER  
 HEADER\_SIZE  
 CHECKSUM\_ALGORITHM  
 HEADER\_CHECKSUM  
 TAPE\_MARK\_COUNT  
 BLOCK\_SIZE  
 BLOCK\_COUNT  
 BLOCK\_TIME\_STAMP  
 STAGER\_VERSION  
 STAGER\_HOST  
 DRIVE\_NAME  
 DRIVE\_SERIAL  
 DRIVE\_FIRMWARE  
 DRIVE\_HOST  
 VOL\_DENSITY  
 VOL\_ID  
 VOL\_SERIAL  
 DEVICE\_GROUP\_NAME  
 FILE\_SIZE  
 FILE\_CHECKSUM  
 FILE\_NS\_HOST  
 FILE\_NS\_ID  
 FILE\_PROGESSIVE\_CHECKSUM  
 FILE\_BLOCK\_COUNT  
 FILE\_NAME

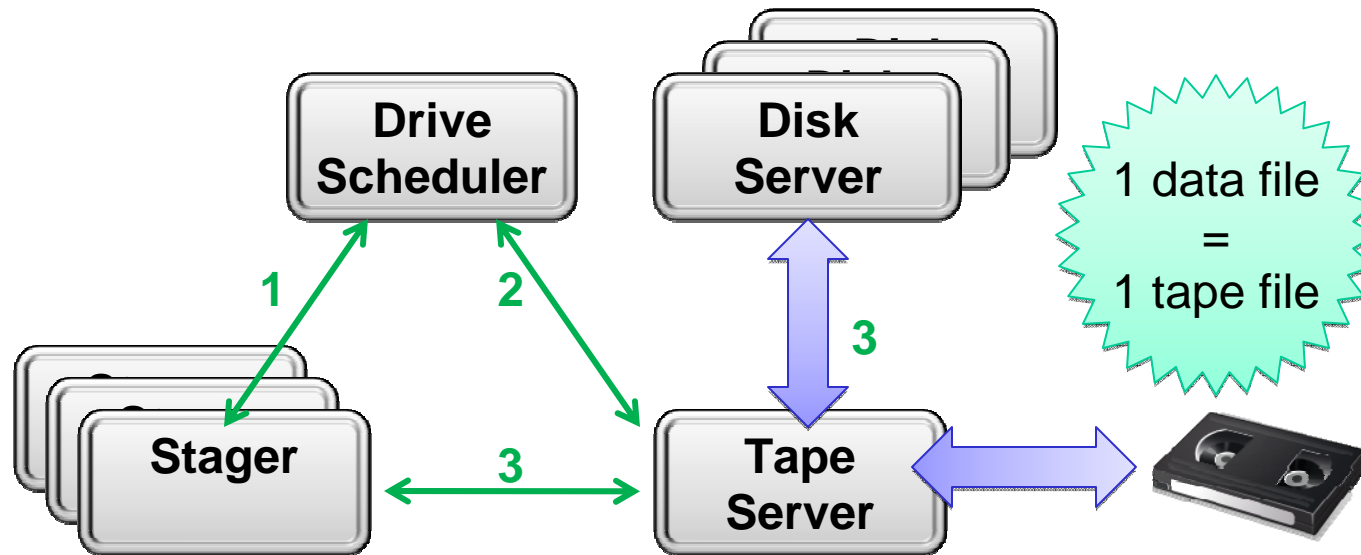




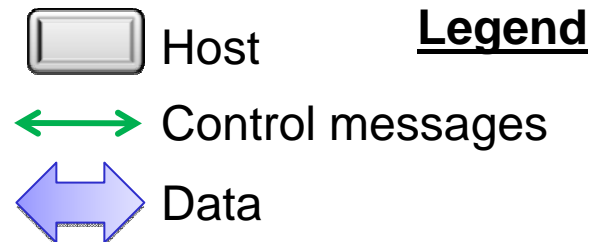


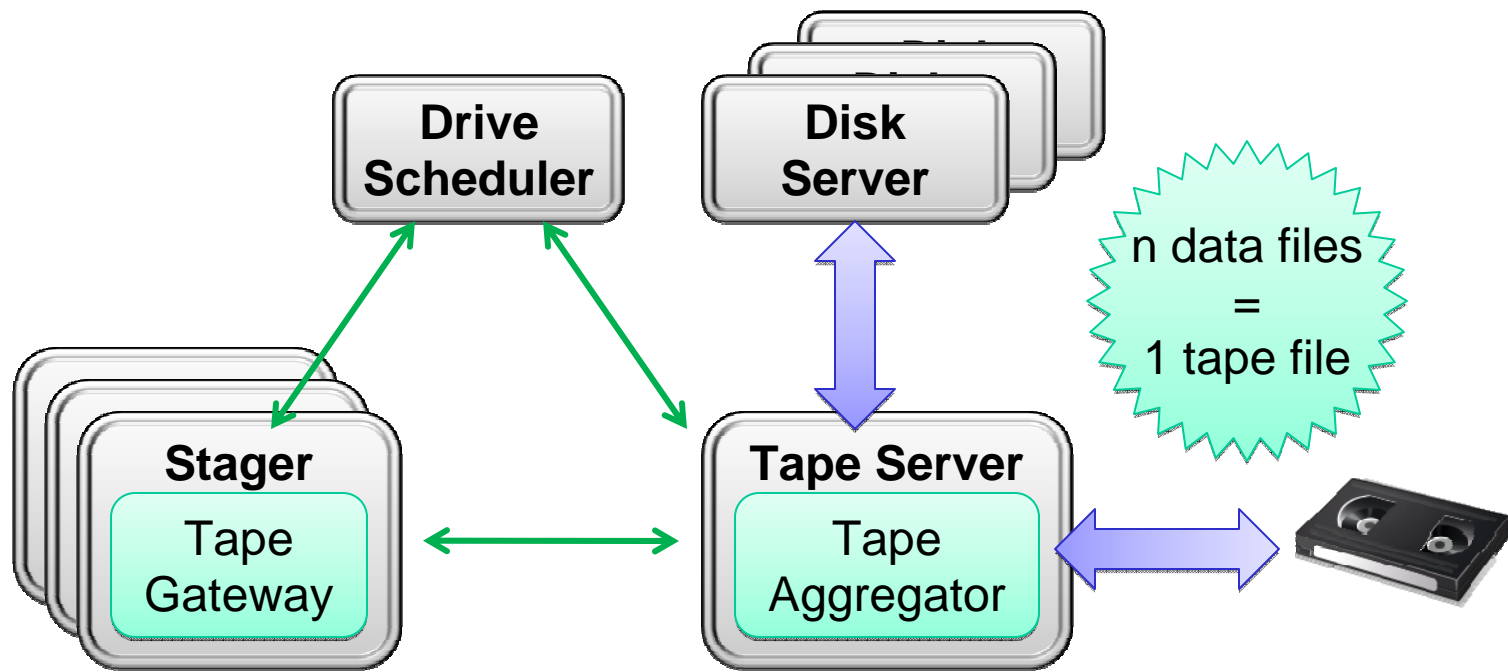
- The new tape format is only half of the story
- An aggregator needs to be inserted into the disk ↔ tape data streams
- Anything old that is replaced is an opportunity for code re-use and increased maintainability via the Castor framework



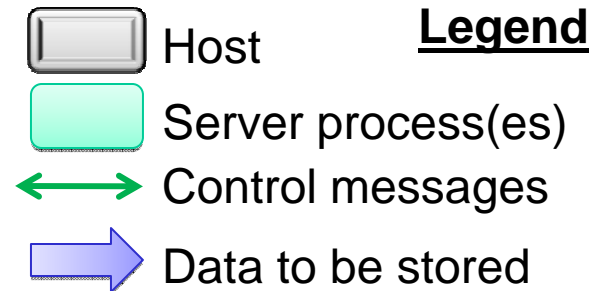


1. Stager requests a drive
2. Drive is allocated
3. Data is transferred to/from disk/tape based on file list given by stager





- The tape gateway will replace RTCPCliendD
- The tape gateway will be stateless
- The tape aggregator will wrap RTCPD



Date	Actions
Beginning Q4 2008	Put repack into full production will at least 20 drives. Expecting 700 MB/s.  Conclude new tape format architecture.
End Q1 2009	Release first functional prototype of new tape format.
End Q2 2009	Write new tape format with repack only.  Read new tape format everywhere.
End Q3 2009	Read and write everywhere
Beginning Q1 2010	Replace RTCPD with tape aggregator





- We have improved the efficiency of tape by increasing the amount of data we write per mount
- Repack uses Castor as cache to support asymmetric drive read/write allocation
- We are currently developing a new tape format to increase write performance
- The future of repack may be to run continuously to constantly verify data integrity in addition to media migration and tape defragmentation
- We will continue to identify the greatest efficiency improvements that require the least effort

