

# Chains of functional primitives

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```
I = 0
J = 0
100 I = I + 1
    IF (I .GT. NEVENT) THEN
        GO TO 200
    EVENT = EVENTS(I)
    IF (CONDIT(EVENT)) THEN
        GO TO 100
    J = J + 1
    OUTPUT(J) = CALCUL(EVENT)
    GO TO 100
200 ! ...
```

```
for (i = 0; i < nEvents; i++) {
    event = events[i];
    if (condition(event))
        continue;
    output.push_back(
        calculation(event));
}
```

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val output = events.filter(condition).map(calculation)
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“condition” and “calculation” are user-defined functions passed as arguments to “filter” and “map.”

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```

The map/filter functional chain *says less* than the for loop.

“Step through the events *in order*, skip if the condition is met, and *incrementally grow the output list*.”

“Remove events for which the condition holds and apply the calculation to the remainder.”

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}
```

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val output = events.  
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```

- ▶ The for loop body could have operated on multiple events at once, but didn't in this case. The map functional cannot ever. Compilers and runtime environments can take advantage of this knowledge for vectorization or parallelization.

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- ▶ Left: `output.push_back` doesn't know how large the output can be and has to dynamically allocate.  
Right: `output.size <= events.size`; allocate and trim.



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- ▶ Left: `output.push_back` doesn't know how large the output can be and has to dynamically allocate.  
Right: `output.size <= events.size`; allocate and trim.
- ▶ To repartition the for loop, the user must be involved in the index arithmetic; the functionals are more abstract.

```
val output = events.filter(condition).map(calculation)
```

could mean

- ▶ Generate inline code for `condition` and `calculation` and vectorize the calculation.
- ▶ Evaluate them in a thread execution pool.
- ▶ Launch parallel jobs on a worldwide grid.
- ▶ Construct a CUDA kernel and pass `condition` and `calculation` to the GPU.
- ▶ Construct an intermediate list after `filter` and before `map`.
- ▶ Lazy-evaluate the `filter`, treating it like a Python iterator (no intermediate list).

Although these choices have significant performance consequences, they are secondary to the intention expressed in that line of code.

This style is fairly common among data analysts:

- ▶ R code is full of `apply/lapply/tapply`, and the R users I know try to avoid for loops whenever possible.
- ▶ “Map-reduce” launched an industry around Hadoop, and functional chains are the central paradigm of Spark.
- ▶ Functional primitives are hidden in the `SELECT`, `WHERE`, and `GROUP BY` clauses of SQL.
- ▶ LINQ, the data extraction sublanguage of .NET, is heavily functional.
- ▶ d3, a popular visualization library for Javascript, also manipulates data with functional chains.

Although it restricts flexibility, this paradigm seems to fit data analysis well.

Switching to this paradigm requires the user to become familiar with some functional primitives.

- ▶ Adopt the “there’s an app for that” mentality.

	input	function	output	operation
<b>map</b>	table of $A$	$f : A \rightarrow B$	table of $B$	apply $f$ to each row $A$ , get a table of the same number of rows $B$
a.k.a. “lapply” (R), “SELECT” (SQL), list comprehension (Python)				
<b>filter</b>	table of $A$	$f : A \rightarrow$ boolean	table of $A$	get a shorter table with the same type of rows
a.k.a. single brackets (R), “WHERE” (SQL), list comprehension (Python)				
<b>flatMap</b>	table of $A$	$f : A \rightarrow$ table of $B$	table of $B$	compose <b>map</b> and <b>flatten</b> , get a table of any length
a.k.a. “map” (Hadoop), “EXPLODE” (SQL), $>>=$ (Haskell)				

	input	function(s)	output	operation
reduce	table of $A$	$f : (A, A) \rightarrow A$	single value $A$	apply $f$ to the running sum and one more element
aggregate	table of $A$ , initial value $B$ ("zero")	$f : (A, B) \rightarrow B$ $f : (B, B) \rightarrow B$ (increment and combine)	single value $B$	accumulate a counter with a different data type from the input
aggregate by key	table of $\langle K, A \rangle$ , initial value $B$	$f : (A, B) \rightarrow B$ $f : (B, B) \rightarrow B$	pairs $\langle K, B \rangle$	aggregate independently for each key

a.k.a. "reduce" (Hadoop), "GROUP BY" (SQL)

More exotic functionals can handle specific cases.

For instance,

- ▶ `collection.skip(n)` to offset a collection by `n`
- ▶ `zip(collections*)` to walk through collections in lock-step

can be combined to compare an event with the previous event:

```
zip(events, events.skip(1)).map(operation_on_pairs)
```

Or perform nested loops (SQL JOIN):

- ▶ `cartesian` to loop over all pairs  $i, j$  of a collection
- ▶ `triangular` to loop over pairs  $i, j \geq i$  of a collection

Different functional names because the user thinks of them differently; each would have to be optimized differently, anyway.

This is very similar to what I'm doing with Histogrammar (<https://github.com/diana-hep/histogrammar>), which introduces a dozen functional primitives that are all variations on aggregate.

```
histogram = Bin(100, 0, 20, fill_rule, Count())

hist2d = Bin(binsX, lowX, highX, fillX,
             Bin(binsY, lowY, highY, fillY, Count()))

profile = Bin(binsX, lowX, highX, fillX, Deviate(fillY))

box_whiskers = Bin(binsX, lowX, highX, fillX, Branch(
    Minimize(fillY), Quantile(0.25, fillY), Quantile(0.5, fillY),
    Quantile(0.75, fillY), Maximize(fillY)))

unknown_support = SparselyBin(binWidth, fillX, Count())

efficiency = Fraction(cut, Bin(100, 0, 20, fill_rule, Count()))
```

where all `fill_rules` are lambda functions.



To be fluent, one needs a good syntax for lambdas.

C++ `[] (Datum d) { return sqrt(d.px*d.px + d.py*d.py); }`

Scala `{d: Datum => Math.sqrt(d.px*d.px + d.py*d.py)}`

Python `lambda d: math.sqrt(d.px**2 + d.py**2)`

R `function (d) { sqrt(d.px^2 + d.py^2) }`

Unlike all the rest, Python lambdas are fundamentally limited to one-line expressions (no statements, such as local assignment).

I have some code (written in Python) that extends Python's grammar to include multiline assignments like this:

```
lambda d: sqrt(d.px**2 + d.py**2)
def(d -> sqrt(d.px**2 + d.py**2))
```

```
def(x -> y = sqrt(x),
      z = 2*y,
      z**2)
```

or even

```
def(x -> y = sqrt(x),
      z = 2*y if y < 1,
          3*y if y < 2,
          4*y else,
      z**2)
```

which might be useful for extended functionals in Python.

I also have a suite of functional primitives with naïve implementations in the attached `functional_chains.py`.

```
class Data(object):
    def __init__(self, generator):
        self.generator = generator

    def map(self, fcn):
        return Data(fcn(x) for x in self.generator)

    def filter(self, fcn):
        return Data(x for x in self.generator if fcn(x))

    def flatMap(self, fcn):
        return Data(itertools.chain.from_iterable(fcn(x) for x
            in self.generator))

    ...
```

...with comparisons to TTree.Draw strings in PyROOT.

```
ttree.Draw("fTracks.fPx >> hPx")
assert Data.source().flatMap(lambda event: event.fTracks).
    map(lambda track: track.fPx).verify("hPx")
```

```
ttree.Draw("fMatrix[][0] >> hMatrix0")
assert Data.source().flatMap(lambda _: _.fMatrix[:,0]).
    verify("hMatrix0")
```

```
ttree.Draw("fTemperature - 20 * Alt$(fClosestDistance[9], 0) " +
    ">> hClosestDistanceAlt")
assert Data.source().filterMap(lambda _: _.fTemperature - 20 *
    _.fClosestDistance.getOrElse(9, 0.0)).
    verify("hClosestDistanceAlt")
```