

## Georgian Technical University

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# New Tweakable Block Cipher 



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## WHAT IS CRYPTOGRAPHY?

Cryptography (From Greek means "secret writing") is the practice and study of techniques for secure communication in the presence of third parties.


## ENCRYPTION ALGORITHMS

- Classified as symmetric and asymmetric classes.
- Symmetric algorithms are a class of algorithms for cryptography that use the same cryptographic keys for both encryption of plaintext and decryption of ciphertext. The keys may be identical or there may be a simple transformation to go between the two keys. The keys, in practice, represent a shared secret between two or more parties that can be used to maintain a private information link.


## ENCRYPTION ALGORITHMS

- Asymmetric algorithms, is a class of cryptographic algorithms which requires two separate keys, one of which is secret (or private) and one of which is public. Although difference, the two parts of this key pair are mathematically linked. The public key is used to encrypt plaintext or to verify a digital signature, whereas the private key is used to decrypt ciphertext or to create a digital signature.


## ENCRYPTION ALGORITHMS

- As it is widely known for protection of the information generally symmetric block algorithms shall be applied, as the open key systems speed is quite low.


## SYMMETRICAL CRYPTOSYSTEM



## SYMMETRICAL CRYPTOSYSTEM

- In order to cover the open text structure the most effective way is to apply for two transformations: confusion and diffusion.
- Confusion is the transformation, the goal of which is to cover the connection among the keys and the ciphertext, and the goal of the diffusion is to render each symbol of the ciphertext dependent onto all the symbols of the open text, which would enable us to cover the open text structure.


## SYMMETRICAL CRYPTOSYSTEM

- As in symmetric algorithms it is impossible to use the complex mathematical transformations (that diminishes the fast action of the algorithm), in order to achieve such goals in the modern symmetric cryptography replacement and displacement operations are applied for with the multiple iterations.


## SYMMETRICAL CRYPTOSYSTEM

- Unfortunately, the block ciphers have significant fallback. That is their determined nature, which is expressed in the fact that the same text by means of the same keys is always transferred into the same cipher text. This fallback is tried to be suppressed by means of the encrypting regimes, in which the initialization vector is applied for, which enables to transfer the same text with the same keys into various cipher texts.


## TBC

- In 2002 the Article by M. Liskov, R. Rivest and D. Wagner was published, the idea stated in which that, initialization vector might be used not in the regime of encrypting, but in the algorithm itself. Such ciphers are called tweakable block ciphers.


## TBC

Tweakable block cipher overview


## HILL ALGORITHM

- In 1929 American mathematician Lester S.Hill by means of utilization of the linear algebra created n gram encrypting algorithm, which enables to make one outgoing symbol of the ciphertext dependet onto the n number of the incoming symbols.



## MODIFIED HILL ALGORITHM

- Our goal is to construct tweakable block cipher algorithm, in which in order to cover efficiently the open text structure, by us modified Hill's algorithm shall be used. In crypto algorithm 256 bits block is encrypted with the confidential keys. Upon entrance into the algorithm, the block to be encrypted shall be represented by means of the matrix, which is called the standing matrix, where each is the binary byte. Binary line to be encrypted shall be recorded in the matrix from the left to the right horizontally.


## MODIFIED HILL ALGORITHM

$$
M=\left(\begin{array}{llll}
a_{11} & a_{12} & a_{13} & a_{14} \\
a_{21} & a_{22} & a_{23} & a_{24} \\
a_{31} & a_{32} & a_{33} & a_{34} \\
a_{41} & a_{42} & a_{43} & a_{44}
\end{array}\right)
$$

- All the operations, which are completed for the text to be encrypted into the algorithm, are completed on this matrix. We will deal with one operation only, which provides the open text structure effective covering into the ciphertext.


## MODIFIED HILL ALGORITHM

- This operation mathematically might be recorded quite simply: $\mathrm{M} \times \mathrm{A}(\bmod 256)$. Where A is the matrix and that matrix shall by all means have the reverse matrix.



## ENCRYPTION

- Let us suppose we have open text: Evariste Galois was a French mathematician. We take the starting 16 symbols, transform them into ASCII code, the obtained result into binary system and adding to the first half of the initial key. Then finding multiplicative inverse for each byte in $\mathrm{GF}\left(2^{8}\right)$ field:

| E | v | a | r | i | s | t | e |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 69 | 118 | 97 | 114 | 105 | 115 | 116 | 101 |
| space | G | a | 1 | o | i | s | space |
| 32 | 71 | 97 | 108 | 111 | 105 | 115 | 32 |

## ENCRYPTION

| 00111101 | 00011011 | 11001011 | 10100100 |
| :--- | :--- | :--- | :--- |
| 11110110 | 11110011 | 11100000 | 11011001 |
| 11011100 | 01000110 | 01000010 | 10111110 |
| 11001011 | 10100110 | 01000010 | 11000100 |

The obtained binary string we transfer into the decimal system and represent 4 x 4 dimensional A matrix:

## ENCRYPTION

| 61 | 27 | 203 | 164 |
| :---: | :---: | :---: | :---: |
| 246 | 243 | 224 | 217 |
| 220 | 70 | 66 | 190 |
| 203 | 166 | 66 | 196 |

N matrix calculated by us in advance:

## ENCRYPTION

N Matrix:

| -1 | -2 | -2 | -2 |
| :---: | :---: | :---: | :---: |
| 2 | -1 | -2 | 2 |
| 1 | 1 | 1 | 2 |
| -1 | 1 | 2 | -1 |

## ENCRYPTION

A matrix is multiplied for N matrix, as the result of which $4 \times 4$ dimensional $A_{1}$ matrix is received again:

| 32 | 218 | 355 | 174 |
| :---: | :---: | :---: | :---: |
| 247 | -294 | -320 | 225 |
| -204 | -254 | -134 | -358 |
| -1 | -310 | -280 | -138 | transferred into the binary system:


| 32 | 218 | 99 | 174 |
| :---: | :---: | :---: | :---: |
| 247 | 218 | 192 | 225 |
| 52 | 2 | 122 | 154 |
| 755 | 207 | 237 | 118 |

## ENCRYPTION

| 32 | 218 | 99 | 174 | 247 | 218 | 192 | 225 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 00100000 | 11011010 | 1100011 | 10101110 | 11110111 | 11011010 | 11000000 | 11100001 |
| 52 | 2 | 122 | 154 | 255 | 202 | 232 | 118 |
|  |  |  |  |  |  |  |  |
| 00110100 | 00000010 | 1111010 | 10011010 | 11111111 | 11001010 | 11101000 | 01110110 |

## ENCRYPTION

- Then binary string is summarized with formed round first key by means of XOR,
- Key 1.1.

| 00110110 | 00100011 | 10001100 | 01101000 | 00111001 | 10010010 | 11001110 | 11010010 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 01100001 | 11111100 | 11001101 | 01111110 | 01000111 | 11111110 | 00010011 | 11010100 |

Binary string and Key 1.1. XOR summarization result:

| 00010110 | 11111001 | 11101111 | 11000110 | 11001110 | 01001000 | 00001110 | 00110011 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 01010101 | 11111110 | 10110111 | 11100100 | 10111000 | 00110100 | 11111011 | 10100010 |

## ENCRYPTION

- The Result is summarized with Tweak entrance by means of XOR:
-Tweak 1.1.:

| 01001010 | 10011110 | 10110001 | 10011100 | 00110101 | 0010101 | 01001111 | 01110010 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 11101111 | 10110010 | 01000111 | 10010010 | 10110100 | 01101110 | 01111001 | 00011000 |

In result we get this binary string:

| 01011100 | 01100111 | 01011110 | 01011010 | 11111011 | 01011101 | 01000001 | 01000001 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 10111010 | 01001100 | 11110000 | 01110110 | 00001100 | 01011010 | 10000010 | 10111010 |

## ENCRYPTION

- Then we transfer it to the decimal system and get next round first data:

| 92 | 103 | 94 | 90 | 251 | 93 | 65 | 65 |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| 186 | 76 | 240 | 118 | 12 | 90 | 130 | 186 |

## ENCRYPTION

- With the analogue method we act at B matrix only instead of N matrix we use M matrix and result is:
- M Matrix

| -2 | -1 | 2 | 2 |
| :---: | :---: | :---: | :---: |
| -2 | -2 | -1 | -2 |
| 1 | 1 | 1 | 2 |
| 2 | 1 | -1 | -1 |

## DECRYPTION

| $\mathbf{w}$ | a | s | space | a | space | F | r |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 119 | 97 | 115 | 32 | 97 | 32 | 70 | 114 |
| e | n | c | h | space | m | a | t |
| 101 | 110 | 99 | 104 | 32 | 109 | 97 | 116 |


| 01111000 | 01100011 | 11010000 | 11111011 |
| :--- | :--- | :--- | :--- |
| 10001011 | 00111000 | 01001011 | 01011001 |
| 10101110 | 11101101 | 00010111 | 01100110 |
| 10010111 | 01011011 | 10111101 | 10101011 |

## ENCRYPTION

- And at the end we get next round second data:

| 200 | 227 | 246 | 129 | 41 | 177 | 31 | 49 |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| 87 | 34 | 111 | 62 | 109 | 208 | 226 | 110 |

## DECRYPTION

- Decryption is the reversed process of encryption with the insignificant differences. While encrypting instead of the applied N and M matrixes we use 256 module reversed $\mathrm{N}^{-1}$ and $\mathrm{M}^{-1}$ matrixes accordingly.
$\mathrm{N}^{-1}$ matrix:

| -1 | 2 | -2 | 2 |
| :---: | :---: | :---: | :---: |
| -2 | -1 | -2 | -2 |
| 1 | 1 | 1 | 2 |
| 1 | -1 | 2 | -1 |

## DECRYPTION

- $\mathrm{M}^{-1}$ matrix:

| -2 | -1 | 2 | 2 |
| :---: | :---: | :---: | :---: |
| -2 | -2 | -1 | -2 |
| 1 | 1 | 1 | 2 |
| 2 | 1 | -1 | -1 |

## CONCLUSIONS

- In our work is done an attempt to build a new tweakable block cipher, which in our opinion, as a result of our work has been achieved. The algorithm is very fast and can easily be realized, both in hardware and software. The algorithm must be checked and analyzed with a variety of tests. At this point, we can only assess in advance it's cryptoresistant against cryptographic attacks. The current algorithm is cryptoresistant against to all currently known cryptanalysis. We are working on algorithm, if possible, patent algorithm and continue research in this direction.


## THANK $\mathbf{H} Y O U$

