Application of Many-core Accelerators for Problems in Astronomy and Physics

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Agenda

- Our Problems
- Recent Development of Many-core Accelerator Systems
- Our Approach to the problems
- Performance evaluation
- Summary

Particle Simulations

- Simulate evolution of the universe
 - As a collection of particles
 - Depending on scale, each particle represents
 - Galaxy
 - Star
 - Asteroid
 - Gas blob etc.
 - Particles are interacting
 - Mainly by gravity
 - Long-range force

Numerical Modeling

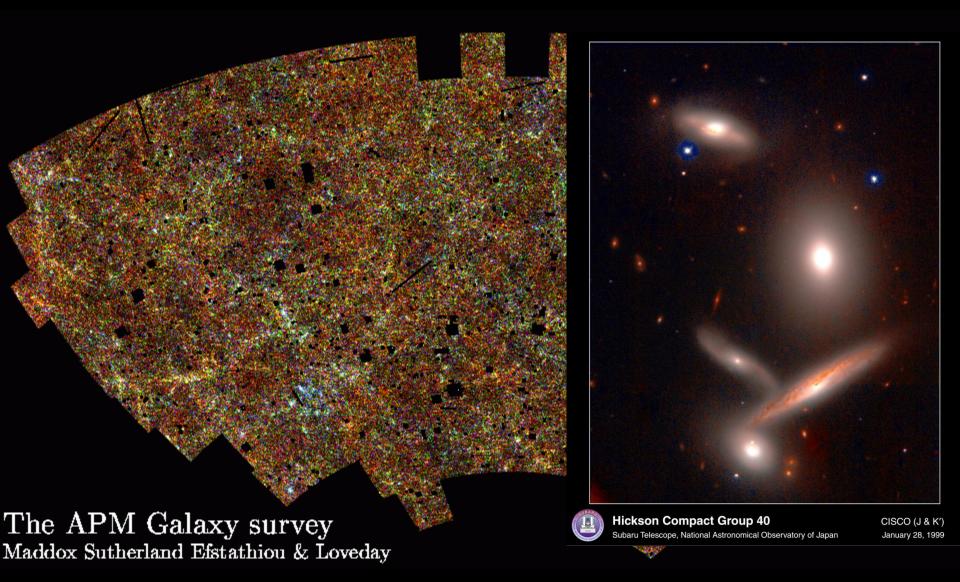
Solve ODE for many particles

$$\frac{d\vec{v_i}}{dt} = \sum_{j=1}^{N} \vec{f}(\vec{r_i} - \vec{r_j})$$

where f is gravity, hydro force etc...

- Two main problems
 - How to integrate the ODE?
 - How to compute RHS of ODE?
 - We will use accelerators for this part

Grand Challenge Problems



Grand Challenge Problems

- Simulations with very huge N
 - How is mass distributed in the Universe?
 - One big run with N $\sim 10^{9-12}$
 - Scalable on a simple big MPP system
 - Limited by memory size
- Modest N but complex physics
 - Precise modeling of formation of astronomical objects like galaxy, star, solar system.
 - Need many runs with N ~ 10⁶⁻⁷

Cluster Configuration

Cluster with accelerators for Modest N problems

Big MPP cluster for Large N problems

Number of nodes

Accelerator?

- A device that assist a main computer
 - for speeding a specific calculation
 - Cell, ClearSpeed, GPU etc.
- Many-core accelerator is
 - Parallel computer on a chip
 - Difficulties raised in parallel computing applies
 - Very high performance on specific tasks
 - Developing so fast
 - changes in mice year?

Many-core Accelerators

- Cell, ClearSpeed, GPU etc.
 - have FP units as many as 32 1000 or more
 - Number of FP units is continuously rising...
 - Driven by demand for high performance gaming!
 - 2 x growth with every generation (~1.5 yr or so)



Latest Cypress GPU (ATi)
1600 FP units (single precision)
Running at 850 MHz
1 GB
16x PCI-E gen2
Consume ~ 200W

TOP500 List

Two systems use accelerators out of top 5 systems

Rank	Site	Computer/Year Vendor	Cores	R _{max}	R _{peak}
1	Oak Ridge National Laboratory United States	Jaguar - Cray XT5-HE Opteron Six Core 2.6 GHz / 2009 Cray Inc.	224162	1759.00	2331.00
2	DOE/NNSA/LANL United States	Roadrunner - BladeCenter QS22/LS21 Cluster, PowerXCell 8i 3.2 Ghz / Opteron DC 1.8 GHz, Voltaire Infiniband / 2009 IBM PowerX	122400 Cell 8	1042.00	1375.78
3	National Institute for Computational Sciences/University of Tennessee United States	Kraken XT5 - Cray XT5-HE Opteron Six Core 2.6 GHz / 2009 Cray Inc.	98928	831.70	1028.85
4	Forschungszentrum Juelich (FZJ) Germany	JUGENE - Blue Gene/P Solution / 2009 IBM	294912	825.50	1002.70
5	National SuperComputer Center in Tianjin/NUDT China	Tianhe-1 - NUDT TH-1 Cluster, Xeon E5540/E5450, ATI Radeon HD 4870 2, Infiniband / 2009 NUDT Radeon	⁷¹⁶⁸⁰ HD48	563.10 70	1206.19

Green500 List

All top systems use accelerators

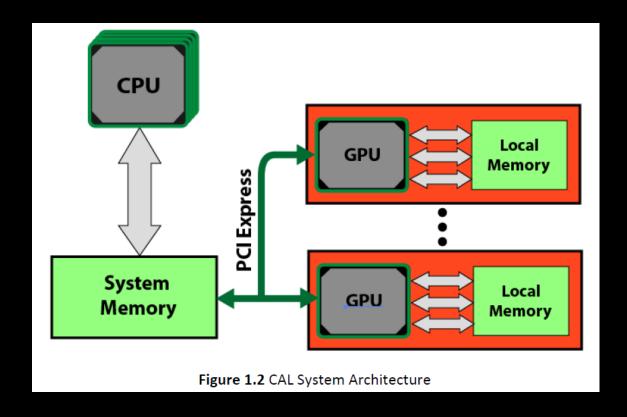
Green500 Rank	MFLOPS/W	Site*	Computer*	Total Power (kW)	TOP500 Rank*
1	722.98	Forschungszentrum Juelich (FZJ)	Q ACE SFB TR Cluster, PowerXCell 8i, 3.2 Glz, 3D-Torus	59.49	110
1	722.98	Universitaet Regensburg	QLACE SFB TR Cluster, PowerXCell 8i, 3.2 Glz, 3D-Torus	59.49	111
1	722.98	Universitaet Wuppertal	QLACE SFB TR Cluster, PowerXCell 8i, 3.2 GLz, 3D-Torus	59.49	112
4	458.33	DOE/NNSA/LANL	PowerXCell 8i	276	29
4	458.33	IBM Poughkeepsie Benchmarking Center	BladeCenter QS22/LS21 Cluster, PowerXCell 8i 3.2 Ghz / Opteron DC 1.8 Gliz, Infiniband	138	78
6	444.25	DOE/NNSA/LANL	BladeCenter QS22/LS21 Cluster, PowerXCell 8i 3.2 Ghz / Opteron DC 1.8 Glz, Voltaire Infiniband	2345.5	2
7	428.91	National Astronomical Observatory of Japan	GRAPE-DR	51.2	445
8	379.24	National SuperComputer Center in Tianjin/NUDT	Radeon HD4870	1484.8	5
9	378.77	King Abdullah University of Science and Technology	Blue Gene/P Solution	504	18

Using GPU is easy if...

- Use the existing library
 - LINPACK relies on DGEMM
 - DGEMM performance of GPU > 100 Gflops
 - FFT on GPU ~ 50 Gflops (SP)
 - N-body on GPU ~ 100 Gflops (DP)
- For more general problems
 - Rewriting the existing code base
 - Rewriting itself is not so difficult
 - Optimizing it is the problem depending on a given architecture

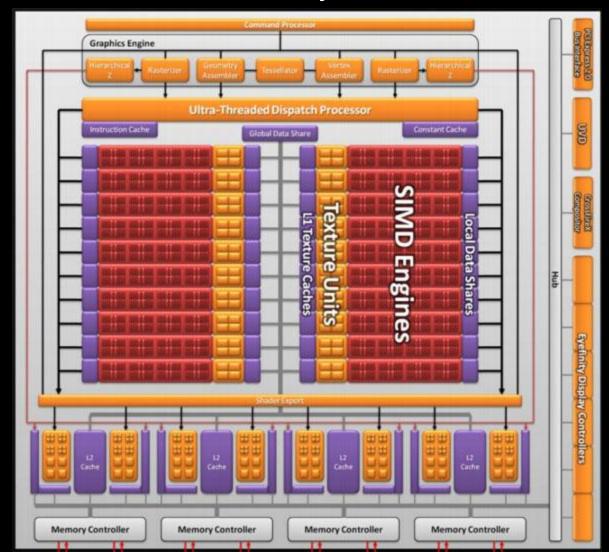
Architecture of Accelerators (1)

- CPU controls GPU
 - Application running on CPU
 - kernel running on GPU



Architecture of Accelerators (2)

GPU consists of many FP units



Challenges

- How to program many-core systems?
 - Like a vector-processor but not exactly same
 - Many programming models/APIs for rapidly changing architectures

- Memory wall
 - at the local memory
 - 2.7 Tflops vs. 153 GB s⁻¹
 - at I/O the accelerators
 - Only 16 GB s⁻¹
 - External I/O in cluster configuration is more severe

Programming Many-core Accelerators

- To use accelerators, need two programs
 - A program running on host
 - A program running on accelerators
 - Compute kernel
- Example
 - C for CUDA / Brook+
 - Host program in C++
 - Compute kernel in extended C
 - Function with appropriate keyword
 - Separate source code

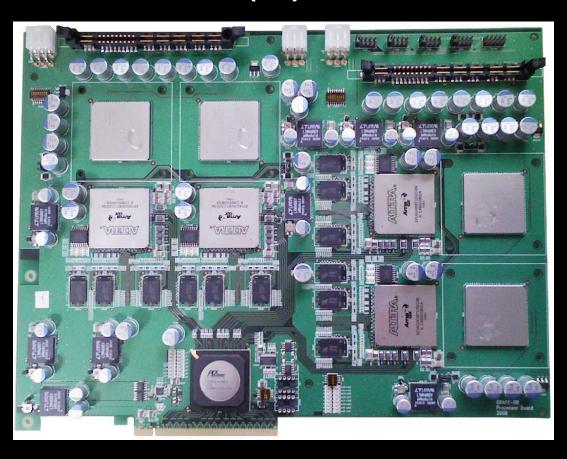
Programming efforts require

- on how we I/O to/from accelerators
 - Mainly programming for CPU
 - relatively easy
- on how we use FP units
- on how we use internal memories
 - Programming for GPU
 - strongly dependent on a given architecture
 - where we need to optimize
- on how we program a cluster of GPU
 - no definitive answer

GRAPE-DR (1)



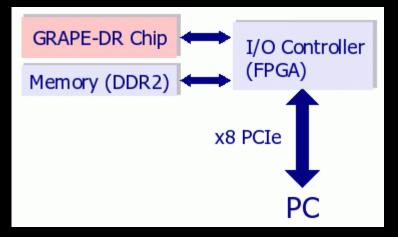
One Chip: 512 PEs Running at 400 MHz 8x PCI-E gen1 288 MB Consume ~ 50 W

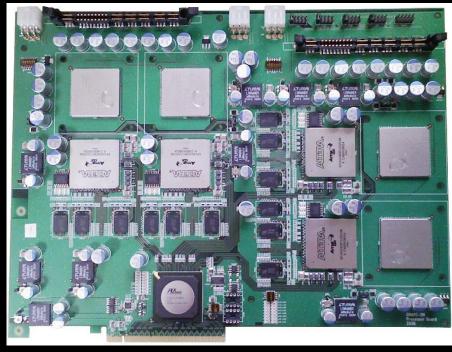


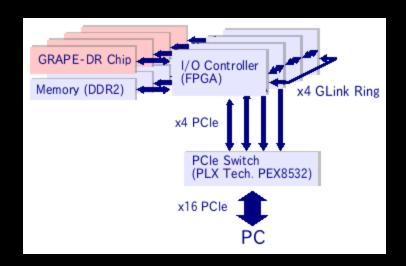
Ranked at 445th on TOP500 Ranked at 7th on Green500

GRAPE-DR (2)





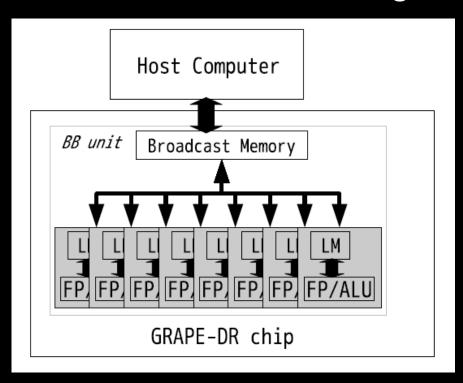




http://kfcr.jp/

Many-core Accelerators

- Both GRAPE-DR and R700 GPU
 - DP performance > 200 GFLOPS
 - Have many local registers: 72/256 words
 - Resource sharing in SP and DP units



But different in

- R700 has more complex VLIW stream cores
- **R700** has no BM
- R700 has faster memory I/O
- •DR has reduction network for efficient summation

Numerical Modeling

Solve ODE for many particles

$$\frac{d\vec{v_i}}{dt} = \sum_{j=1}^{N} \vec{f}(\vec{r_i} - \vec{r_j})$$

where f is gravity, hydro force etc...

- Two main problems
 - How to integrate the ODE?
 - How to compute RHS of ODE?
 - We will use accelerators for this part

A simple way to compute RHS

Compute force summation as

```
for i = 0 to N-1
   s[i] = 0
  for j = 0 to N-1
    s[i] += f(x[i], x[j])
Fig. 1. A simple nested loop to computer a general force calculation.
```

- Each s[i] can be computed independently
 - Massively parallel if N is large
 - Given i & j, each f(x[i],x[j]) can be computed independently if f() is complex

Unrolling (vectrization)

 Parallel nature enable us to unroll the outer-loop in n-ways

```
for i = 0 to N-1 each 4

s[i] = s[i+1] = s[i+2] = s[i+3] = 0

for j = 0 to N-1

s[i] += f(x[i], x[j])

s[i+1] += f(x[i+1], x[j])

s[i+2] += f(x[i+2], x[j])

s[i+3] += f(x[i+3], x[j])
```

- Two types of variables
 - x[i] and s[i] are unchanged during j-loop
 - x[j] is shared at each iteration
- Map computation for each x[i] to PE on accelerators

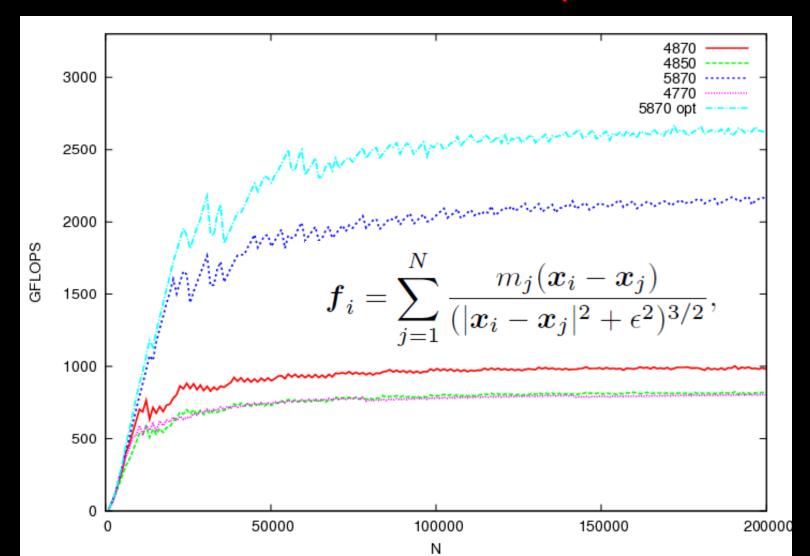
Optimization on GPU

```
for i = 0 to N-1
  acc[i] = 0
                             ~ 300 Gflops
  for j = 0 to N-1
    acc[i] += f(x[i], x[j])
for i = 0 to N-1 each 4
  acc[i] = acc[i+1] = acc[i+2] = acc[i+3] = 0
  for j = 0 to N-1
   acc[i] += f(x[i], x[j])
                                              ~ 500 Gflops
   acc[i+1] += f(x[i+1], x[j])
    acc[i+2] += f(x[i+2], x[j])
   acc[i+3] += f(x[i+3], x[j])
for i = 0 to N-1 each 4
  acc[i] = acc[i+1] = acc[i+2] = acc[i+3] = 0
  for j = 0 to N-1 each 4
    for k = 0 to 3
                                              ~ 700 Gflops
      acc[i] += f(x[i], x[j+k])
      acc[i+1] += f(x[i+1], x[j+k])
      acc[i+2] += f(x[i+2], x[j+k])
```

acc[i+3] += f(x[i+3], x[j+k])

Performance of O(N²) algorithm

On a recent GPU ~ 1.3 Tflops



Our Compiler

- Accelerates force summation loop
- Support two accelerators
 - R700/R800 architecture GPU
 - GRAPE-DR
 - Developed by J.Makino etal.
- Precision controllable
 - Single, Double, & Quadruple precision
 - QP through DD emulation techniques
 - Partially support mixed precision

Our programming model

User write a source in DSL such as

```
LMEM xi, yi, zi, e2;
BMEM xj, yj, zj, mj;
RMEM ax, ay, az;
dx = xj - xi;
dy = yj - yi;
dz = zj - zi;
r1i = rsqrt(dx**2 + dy**2 + dz**2 + e2);
af = mj*r1i**3;
ax += af*dx;
ay += af*dy;
az += af*dz;
```

 Our compiler generates optimized machine code for GPU / GRAPE-DR

Comparison

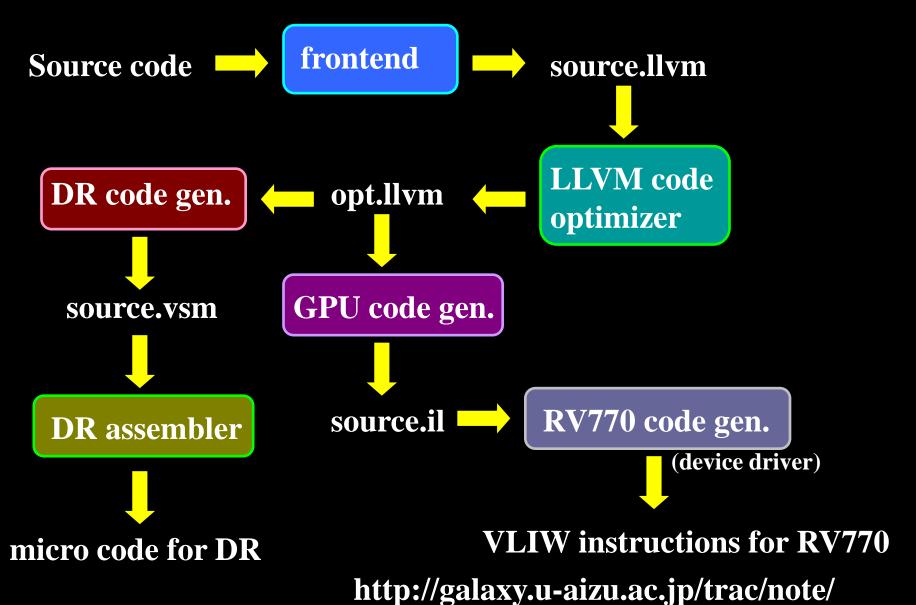
- Our approach is in between two conventional approaches
 - Automatic parallel compiler
 - A user just feed an existing source code
 - But not effective in general
 - Let-users-do-everything-type compiler
 - C for CUDA, OpenCL, Brook+ etc.
 - A user have to specify every details of
 - Memory layout and its movement
 - SIMD operations
 - Threads management on GPU

Details of our compiler

- Written in C++
 - Prototype was developed in Ruby

- We use following software/library
 - Boost sprit for the parser
 - Low Level Virtual Machine for the optimizer
 - Google template library for the code generators

Compiler work flow



Example 1: N-body

Simple softened gravity

$$f_i = \sum_{j=1}^{N} \frac{m_j(x_i - x_j)}{(|x_i - x_j|^2 + \epsilon^2)^{3/2}},$$



```
LMEM xi, yi, zi, e2;
BMEM xj, yj, zj, mj;
RMEM ax, ay, az;

dx = xj - xi;
dy = yj - yi;
dz = zj - zi;

rli = rsqrt(dx**2 + dy**2 + dz**2 + e2);
af = mj*rli**3;

ax += af*dx;
ay += af*dy;
az += af*dz;
```

Example 2: Feynman-loop integral

$$I = \int_{0}^{1} dx \int_{0}^{1-x} dy \int_{0}^{1-x-y} dz F(x, y, z),$$

$$F(x, y, z) = D(x, y, z)^{-2}$$

$$D = -xys - tz(1 - x - y - z) + (x + y)\lambda^{2} + (1 - x - y - z)(1 - x - y)m_{e}^{2} + z(1 - x - y)m_{f}^{2}.$$
(2)

LMEM xx, yy, cnt4;

QD operations on GPU

- We have implemented so-called DD emulation scheme on GPU&GRAPE-DR
 - QD variable is expressed as summation of two double precision variables
 - QD operations are emulated with DP operations
 - At least 20 times slower performance
 - Practical performance is more than 30 times slower on Core i7 CPU

Performance of QP operations

- Computation of Feynman-loop integral
 - elapsed time in QP operations

	N = 256	N = 512	N = 1024	N = 2048	clock
GRAPE-DR	0.21	1.21	7.83	55.1	380
RV770	0.09	0.66	5.03	39.7	750
Core i7	7.39	59.0	472		2670

- − CPU ~ 80 Mflops
- $R700 GPU \sim 6.43 7.57 Gflops$
- GRAPE-DR ~ 2.67 5.46 Gflops
- Tow reasons why QP is so fast
 - High compute density
 - DR & R700 are register rich

Development of QP arithmetic units

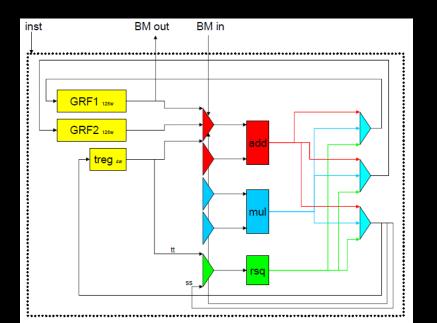
- QP emulation is not efficient
 - A factor of 20 performance penalty
 - Power consumption

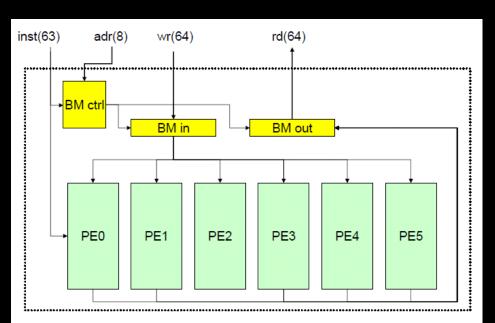
- If we have a dedicated QP unit
 - should be faster and energy efficient
 - but no commercial demand (yet)

We investigated a prototype of accelerator with QP arithmetic units

Status of Project

- We have implemented QP arithmetic units
 - Designed for Feynman integrals
 - 116 bit for mantissa, 11 bit for exponent
 - Add & Mul & inverse sqrt units
 - Implemented by VHDL





Summary

- Is a many-core accelerator is effective for
 - Massively parallel problems : YES
 - Monte-calro on million phase space points
 - O(N²) problems : YES
 - Gravity, Feynman integrals
 - $-O(N^{1.5})$ problems : Yes
 - Matrix multiply (DGEMM)
 - O(N log N) & O(N) problems
 - Generally it is not easy to optimize...
 - High precision operations : Yes
- Key is data reuse = high compute density

Conclusion

- Many-core accelerators are effective in problems in astronomy and physics
 - But how to program it effectively?
- We have constructed a compiler for manycore accelerators
 - That accelerate force-calculation-loop
 - Features simplicity and controllable precision
- Planed Extension
 - Support O(N log N) method on GPU