Scalla/xrootd

Andrew Hanushevsky SLAC National Accelerator Laboratory Stanford University 19-August-2009 Atlas Tier 2/3 Meeting

http://xrootd.slac.stanford.edu/

Outline

System Component Summary **#** Recent Developments **#** Scalability, Stability, & Performance **#** ATLAS Specific Performance Issues **#** Faster I/O The SSD Option **#** Conclusions



The Components

■ xrootd

Provides actual data access

cmsd

Glues multiple xrootd's into a cluster

♯ cnsd

Glues multiple name spaces into one name space

♯ BeStMan

Provides SRM v2+ interface and functions

FUSE

Exports xrootd as a file system for BeStMan

GridFTP

Grid data access either via FUSE or POSIX Preload Library



Recent Developments

April, 2009 **T**orrent WAN transfers ■ May, 2009 # Auto-reporting summary monitoring data June, 2009 **#** Ephemeral files July, 2009 **#** Simple Server Inventory August, 2009



File Residency Manager (FRM)

Functional replacement for MPS scriptsCurrently, includes...

- Pre-staging daemon frm_pstgd and agent frm_pstga
 - Distributed copy-in prioritized queue of requests
 - Can copy from any source using any transfer agent
 - Used to interface to real and virtual MSS's
- frm_admin command
 - Audit, correct, obtain space information
 - Space token names, utilization, etc.
 - Can run on a live system



Torrent WAN Transfers

The xrootd already supports parallel TCP paths Significant improvement in WAN transfer rate Specified as xrdcp – S num ■ New Xtreme copy mode option Uses multiple data sources bit torrent-style ■ Specified as xrdcp –x Transfers to CERN; examples: 12MB/sec (1 stream) ■ 1 source (.de): \blacksquare 1 source (.us): 19MB/sec (15 streams) ■ 4 sources (3 x .de + .ru): 27MB/sec (1 stream each) ■ 4 sources + || streams: 42MB/Sec (15 streams each)

■ 5 sources (3 x .de + .it + .ro):

54MB/Sec (15 streams each)



Summary Monitoring

xrootd has built-in summary monitoring
In addition to full detailed monitoring
Can auto-report summary statistics
xrd.report configuration directive
Data sent to up to two central locations
Accommodates most current monitoring tools
Ganglia, GRIS, Nagios, MonALISA, and perhaps more

- Requires external xml-to-monitor data convertor
- Can use provided stream multiplexing and xml parsing tool
 - Outputs simple key-value pairs to feed a monitor script



Ephemeral Files

Files that persist only when successfully closed
 Excellent safeguard against leaving partial files

- Application, server, or network failures
 - E.g., GridFTP failures
- Server provides grace period after failure
 - Allows application to complete creating the file
 - Normal xrootd error recovery protocol
 - Clients asking for read access are delayed
 - Clients asking for write access are usually denied
 - Obviously, original creator is allowed write access

■ Enabled via xrdcp – P option or ofs.posc CGI element



Simple Server Inventory (SSI)

A central file inventory of each data server

- Does not replace PQ2 tools (Neng Xu, University of Wisconsin)
 - Good for uncomplicated sites needing a server inventory
- Inventory normally maintained on *each* redirector
 - But, can be centralized on a single server
 - Automatically recreated when lost
 - Updated using rolling log files
 - Effectively no performance impact
- Flat text file format
 - LFN, Mode, Physical partition, Size, Space token
 - "cns_ssi list" command provides formatted output



Stability & Scalability

xrootd has a 5+ year production history

- Numerous high-stress environments
 - BNL, FZK, IN2P3, INFN, RAL, SLAC
- Stability has been vetted
 - Changes are now very focused
 - Functionality improvements
 - Hardware/OS edge effect limitations
 - Esoteric bugs in low use paths

Scalability is already at the theoretical maximum

■ E.g., STAR/BNL runs a 400+ server production cluster



Performance I

Following figures are based on actual measurements These have also been observed by many production sites ■ E.G., BNL, IN2P3, INFN, FZK, RAL, SLAC CAVEAT! Figures apply only to the *reference* implementation Other implementations vary significantly Castor + xrootd protocol driver dCache + native xrootd protocol implementation DPM + xrootd protocol driver + cmsd XMI HDFS + xrootd protocol driver



Performance II

Latency

Capacity vs. Load



xrootd latency < 10µs → network or disk latency dominates
Practically, at least ≈10,000 Ops/Second with linear scaling
xrootd+cmsd latency (not shown) 350µs →»1000 opens/second_</pre>



Performance & Bottlenecks

High performance + linear scaling Makes client/server software virtually transparent A 50% faster xrootd yields 3% overall improvement

- Disk subsystem and network become determinants
 - This is actually excellent for planning and funding

HOWEVER

- Transparency makes other bottlenecks apparent
 Hardware, Network, Filesystem, or Application
 - Requires deft trade-off between CPU & Storage resources
 - But, bottlenecks usually due to unruly applications
 - Such as ATLAS analysis



ATLAS Data Access Pattern





ATLAS Data Access Problem

Atlas analysis is fundamentally indulgent

While xrootd can sustain the request load the H/W cannot
 Replication?

- Except for some files this is not a universal solution
 - The experiment is already disk space insufficient
- **#** Copy files to local node for analysis?
 - Inefficient, high impact, and may overload the LAN
 - Job will still run slowly and no better than local disk
- **♯** Faster hardware (e.g., SSD)?
 - This appears to be generally cost-prohibitive
 - That said, we are experimenting with smart SSD handling



Faster Scalla I/O (The SSD Option)

Latency only as good as the hardware (xrootd adds < 10µs latency)

- **Scalla** component architecture fosters experimentation
 - Research on intelligently using SSD devices





The ZFS SSD Option

Decided against this option (for now)

- Too narrow
 - OpenSolaris now or Solaris 10 Update 8 (likely 12/09)
 - Linux support requires ZFS adoption
 - Licensing issues stand in the way
- Current caching algorithm is a bad fit for HEP
 - Optimized for small SSD's
 - Assumes large hot/cold differential
 - Not the HEP analysis data access profile



The xrootd SSD Option

Currently architecting appropriate solution

- Fast track is to use staging infrastructure
 - Whole files are cached
 - Hierarchy: SSD, Disk, Real MSS, Virtual MSS
- Slower track is more elegant
 - Parts of files are cached
 - Can provide parallel mixed mode (SSD/Disk) access
 - Basic code already present
 - But needs to be expanded
- **#** Will it be effective?



Disk vs SSD With 323 Clients



Disk I/O

SSD I/O



What Does This Mean?

Well tuned disk can equal SSD Performance ■ True when number of well-behaved clients < *small* **n** Either 343 Fermi/GLAST clients not enough or Hitting some undiscovered bottleneck **#** Huh? What about ATLAS clients? Difficult if not impossible to get Current grid scheme prevents local tuning & analysis Desperately need a "send *n* test jobs" button • We used what we could easily get Fermi read size about 1K and somewhat CPU intensive



Conclusion

Xrootd is a lightweight data access system

- Suitable for resource constrained environments
 - Human as well as hardware
- Rugged enough to scale to large installations
 - CERN analysis & reconstruction farms
- Flexible enough to make good use of new H/W

Smart SSD

Available in OSG VDT & CERN root package
 Visit the web site for more information
 http://xrootd.slac.stanford.edu/



Acknowledgements

Software Contributors

- Alice: Derek Feichtinger
- CERN: Fabrizio Furano, Andreas Peters
- Fermi/GLAST: Tony Johnson (Java)
- Root: Gerri Ganis, Beterand Bellenet, Fons Rademakers
- SLAC: Tofigh Azemoon, Jacek Becla, Andrew Hanushevsky, Wilko Kroeger
- LBNL: Alex Sim, Junmin Gu, Vijaya Natarajan (BestMan team)
- **#** Operational Collaborators
 - BNL, CERN, FZK, IN2P3, RAL, SLAC, UVIC, UTA
- **#** Partial Funding
 - US Department of Energy
 - Contract DE-AC02-76SF00515 with Stanford University

