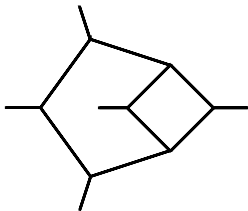


Integration-by-parts reductions via algebraic geometry

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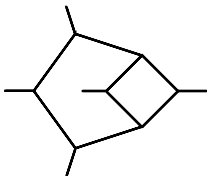
Amplitudes 2018

SLAC National Accelerator Laboratory, 19th of June 2018

Based on PRD **93**(2016)041701, 1712.09737 and 1805.01873
with J. Böhm, A. Georgoudis, H. Schönemann, M. Schulze, Y. Zhang

- 1 Motivation
- 2 IBP identities on unitarity cuts
- 3 Syzygy equations and their solution

- 4 Main example:



Integration-by-parts reductions

IBP identities arise from the vanishing integration of total derivatives,

[Chetyrkin, Tchakov, Nucl. Phys. B **192**, 159 (1981)]

$$\int \prod_{i=1}^L \frac{d^D \ell_i}{\pi^{D/2}} \sum_{j=1}^L \frac{\partial}{\partial \ell_j^\mu} \frac{v_j^\mu P}{D_1^{a_1} \dots D_k^{a_k}} = 0$$

where P and v_j^μ are polynomials in ℓ_i, p_j , and $a_i \in \mathbb{N}$.

Role in perturbative QFT calculations:

- **Reduction.** Reduce number of contributing loop integrals by factor of $\mathcal{O}(10^2) - \mathcal{O}(10^6)$ to basis.
- **Computing master integrals.** Enable setting up differential equations for basis integrals \mathcal{I}_j :

[Gehrmann and Remiddi, Nucl. Phys. B **580**, 485 (2000)]

$$\frac{\partial}{\partial x_m} \mathcal{I}(\mathbf{x}, \epsilon) = A_m(\mathbf{x}, \epsilon) \mathcal{I}(\mathbf{x}, \epsilon)$$

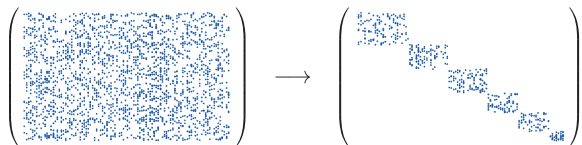
where x_m denotes a kinematical invariant.

IBP reductions on unitarity cuts

Standard approach: enumerate all linear relations and apply
Gauss-Jordan elimination to *large* linear systems

[Laporta, *Int.J.Mod.Phys. A* **15** (2000) 5087-5159]

Idea here: use *unitarity cuts* to block-diagonalize system



We use the Baikov representation ($k = \frac{L(L+1)}{2} + LE$),

$$I(N; a) \equiv \int \prod_{j=1}^L \frac{d^D \ell_j}{i\pi^{D/2}} \frac{N}{D_1^{a_1} \dots D_k^{a_k}} = \int \frac{dz_1 \dots dz_k}{z_1^{a_1} \dots z_k^{a_k}} \text{Gram}(z) \Big|_{(\vec{p}, \ell)}^{\frac{D-L-E-1}{2}} N$$

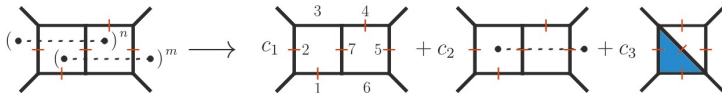
[Baikov, *Phys.Lett. B* **385** (1996) 404-410]

in which cuts are straightforward to apply,

$$\int \frac{dz_i}{z_i^{a_i}} \xrightarrow{\text{cut}} \oint_{\Gamma_{\epsilon(0)}} \frac{dz_i}{z_i^{a_i}} \quad i \in \mathcal{S}_{\text{cut}}$$

Example: Zurich-flag cut

Let us construct IBP identities on the Zurich-flag cut



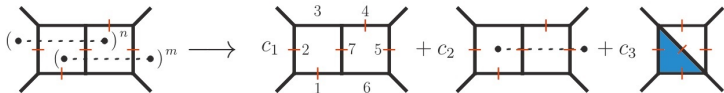
Define $S_{\text{cut}} = \{1, 2, 4, 5, 7\}$ and $G = \text{Gram}(\tilde{\mathbf{p}}, \ell)$.

On S_{cut} , the double-box integral takes the form

$$I_{\text{cut}}^{\text{DB}}[P] = \prod_{i \in S_{\text{cut}}} \oint_{\Gamma_\epsilon(0)} \frac{d\tilde{z}_i}{\tilde{z}_i} \int \prod_{j \notin S_{\text{cut}}} d\tilde{z}_j \frac{G(\tilde{\mathbf{z}})^{\frac{D-6}{2}}}{\tilde{z}_3 \tilde{z}_6} P(\tilde{\mathbf{z}})$$

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Relabeling $z_{\{1,2,3,4\}} = \tilde{z}_{\{3,6,8,9\}}$, this becomes

$$I_{\text{cut}}^{\text{DB}}[P] = \int \frac{dz_1 dz_2 dz_3 dz_4}{z_1 z_2} G(\mathbf{z})^{\frac{D-6}{2}} P(\mathbf{z})$$

Generic total derivative on cut

Need to find IBP identities which involve

$$I_{\text{cut}}^{\text{DB}}[P] = \int \frac{dz_1 dz_2 dz_3 dz_4}{z_1 z_2} G(\mathbf{z})^{\frac{D-6}{2}} P(\mathbf{z})$$

Total derivatives \longrightarrow **IBP identities**. Generic total derivative on cut:

$$\begin{aligned} 0 &= \int \left[\sum_{i=1}^4 \frac{\partial}{\partial z_i} \left(\frac{a_i(\mathbf{z}) G(\mathbf{z})^{\frac{D-6}{2}}}{z_1 z_2} \right) \right] dz_1 \cdots dz_4 \\ &= \int \left[\sum_{i=1}^4 \left(\frac{\partial a_i}{\partial z_i} + \frac{D-6}{2G} a_i \frac{\partial G}{\partial z_i} \right) - \sum_{j=1,2} \frac{a_j}{z_j} \right] \frac{G(\mathbf{z})^{\frac{D-6}{2}}}{z_1 z_2} dz_1 \cdots dz_4 \end{aligned}$$

The **red term** corresponds to an integral in $(D-2)$ dimensions, and the **purple term** in general produces doubled propagators.

IBP identities from syzygies

To avoid **dimension shifts** and **doubled propagators** in

$$0 = \int \left[\sum_{i=1}^4 \left(\frac{\partial a_i}{\partial z_i} + \frac{D-6}{2G} a_i \frac{\partial G}{\partial z_i} \right) - \sum_{j=1,2} \frac{a_j}{z_j} \right] \frac{G(\mathbf{z})^{\frac{D-6}{2}}}{z_1 z_2} dz_1 \cdots dz_4$$

we demand that each term is *polynomial*,

$$\sum_{i=1}^4 a_i \frac{\partial G}{\partial z_i} + bG = 0$$
$$a_j + b_j z_j = 0$$

with a_i, b_i, b polynomials in z . Such eqs. are known as **syzygy equations**.

[Gluzza, Kajda, Kosower, PRD83(2011)045012], [Schabinger, JHEP01(2012)077], [Ita, PRD94(2016)116015]

Obtain IBPs by plugging (a_i, b) into the top equation.

Note: (qa_i, qb) is also a solution, for polynomial q .

Strategy to solve syzygy equations

Solve syzygy equations with c cuts

$$a_j + b_j z_j = 0, \quad j = 1, \dots, k-c \quad (1)$$

$$\sum_{j=1}^{m-c} a_j \frac{\partial G}{\partial z_k} + bG = 0 \quad (2)$$

as follows.

- 1) The generators of (1) are trivial:

$$\mathcal{M}_2 = \langle z_1 \mathbf{e}_1, \dots, z_k \mathbf{e}_k, \mathbf{e}_{k+1}, \dots, \mathbf{e}_m \rangle$$

- 2) Generators $\mathcal{M}_1 = \langle (\mathbf{a}_1, \dots, \mathbf{a}_m, b), \dots \rangle$ of (2) for the *off-shell* case $c = 0$ **can be explicitly found:**

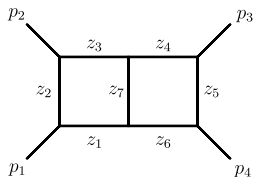
$$(\mathbf{a}_\alpha, b) = \left(\sum_{k=1}^{E+L} (1 + \delta_{ik}) x_{jk} \frac{\partial z_\alpha}{\partial x_{ik}}, 2\delta_{ij} \right)$$

where $x_{ij} = v_i \cdot v_j$ with $v_{i,j} \in \{p_1, \dots, p_E, \ell_1, \dots, \ell_L\}$.

[Böhm, Georgoudis, KJL, Schulze, Zhang, 1712.09737]

- 3) Take module intersection $\mathcal{M}_1|_{\text{cut}} \cap \mathcal{M}_2|_{\text{cut}}$

Example 1: syzygies of planar double box



Set $P_{12} = p_1 + p_2$ and

$$z_1 = \ell_1^2, \quad z_2 = (\ell_1 - p_1)^2, \quad z_3 = (\ell_1 - P_{12})^2$$

$$z_4 = (\ell_2 + P_{12})^2, \quad z_5 = (\ell_2 - p_4)^2, \quad z_6 = \ell_2^2$$

$$z_7 = (\ell_1 + \ell_2)^2, \quad z_8 = (\ell_1 + p_4)^2, \quad z_9 = (\ell_2 + p_1)^2$$

Only need to find explicit relation $z = Ax$. Here $A = \begin{pmatrix} 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ -2 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ -2 & 0 & -2 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 2 & 0 & 2 & 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 2 & 0 & 2 & 0 & 2 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 2 & 1 \\ -2 & 0 & -2 & 0 & -2 & 0 & 1 & 0 & 0 & 0 \\ 0 & 2 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{pmatrix}$

Set $t_{i,j} = (a_\alpha, b)$. The syzygy generators are *linear* in the z_k

$$t_{4,1} = (z_1 - z_2, z_1 - z_2, -s + z_1 - z_2, 0, 0, 0, z_1 - z_2 - z_6 + z_9, t + z_1 - z_2, 0, 0)$$

$$t_{4,2} = (s + z_2 - z_3, z_2 - z_3, z_2 - z_3, 0, 0, 0, z_2 - z_3 + z_4 - z_9, -t + z_2 - z_3, 0, 0)$$

$$t_{4,3} = (-s + z_3 - z_8, t + z_3 - z_8, z_3 - z_8, 0, 0, 0, z_3 - z_4 + z_5 - z_8, z_3 - z_8, 0, 0)$$

$$t_{4,4} = (2z_1, z_1 + z_2, -s + z_1 + z_3, 0, 0, 0, z_1 - z_6 + z_7, z_1 + z_8, 0, -2)$$

$$t_{4,5} = (-z_1 - z_6 + z_7, -z_1 + z_7 - z_9, s - z_1 - z_4 + z_7, 0, 0, 0, -z_1 + z_6 + z_7, -z_1 - z_5 + z_7, 0, 0)$$

$$t_{5,1} = (0, 0, 0, s - z_6 + z_9, -t - z_6 + z_9, z_9 - z_6, z_1 - z_2 - z_6 + z_9, 0, z_9 - z_6, 0)$$

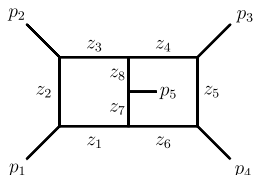
$$t_{5,2} = (0, 0, 0, z_4 - z_9, t + z_4 - z_9, -s + z_4 - z_9, z_2 - z_3 + z_4 - z_9, 0, z_4 - z_9, 0)$$

$$t_{5,3} = (0, 0, 0, z_5 - z_4, z_5 - z_4, s - z_4 + z_5, z_3 - z_4 + z_5 - z_8, 0, -t - z_4 + z_5, 0)$$

$$t_{5,4} = (0, 0, 0, s - z_3 - z_6 + z_7, -z_6 + z_7 - z_8, -z_1 - z_6 + z_7, z_1 - z_6 + z_7, 0, -z_2 - z_6 + z_7, 0)$$

$$t_{5,5} = (0, 0, 0, -s + z_4 + z_6, z_5 + z_6, 2z_6, -z_1 + z_6 + z_7, 0, z_6 + z_9, -2)$$

Example 2: syzygies of non-planar double pentagon



Set $P_{i,j} \equiv p_i + p_j$ and

$$\begin{aligned} z_1 &= \ell_1^2, & z_2 &= (\ell_1 - p_1)^2, & z_3 &= (\ell_1 - P_{1,2})^2, \\ z_4 &= (\ell_2 - P_{3,4})^2, & z_5 &= (\ell_2 - p_4)^2, & z_6 &= \ell_2^2, \\ z_7 &= (\ell_1 + \ell_2)^2, & z_8 &= (\ell_1 + \ell_2 + p_5)^2, & z_9 &= (\ell_1 + p_3)^2, \\ z_{10} &= (\ell_1 + p_4)^2, & z_{11} &= (\ell_2 + p_1)^2 \end{aligned}$$

Here $z = Ax$ with

$$A = \begin{pmatrix} 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ -2 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ -2 & 0 & 0 & -2 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & -2 & 0 & -2 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & -2 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 2 \\ -2 & -2 & -2 & -2 & -2 & -2 & -2 & -2 & 1 & 2 & 1 \\ 0 & 0 & 0 & 0 & 2 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 2 & 0 & 0 & 1 & 0 & 0 \\ 0 & 2 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 \end{pmatrix}$$

and the syzygy generators are again compact:

$$\begin{aligned} f_{5,1} &= (z_1 - z_2, z_1 - z_2, -z_1 z_2 + z_1 - z_2, 0, 0, 0, \\ & z_1 - z_2 - z_6 + z_{11}, -z_1 z_2 - z_1 z_3 - z_1 z_4 + z_1 - z_2 - z_6 + z_{11}, \\ & z_1 z_3 + z_1 - z_2, z_1 z_4 - z_1 - z_2, 0, 0), \\ f_{5,2} &= (z_1 z_2 + z_2 - z_3, z_2 - z_3, z_2 - z_3, 0, 0, 0, \\ & -z_1 z_4 + z_1 + z_2 + z_4 + z_7 - z_6 - z_9 - z_{10} - z_{11}, \\ & z_1 z_3 + z_1 z_4 + z_1 + z_2 + z_4 + z_7 - z_6 - z_9 - z_{10} - z_{11}, \\ & z_1 z_2 + z_2 z_3 + z_2 - z_3, -z_1 z_3 - z_1 z_4 - z_2 z_3 - z_2 - z_3, 0, 0, 0), \\ f_{5,3} &= (z_9 - z_1, -z_1 z_3 - z_1 + z_9, -z_1 z_3 - z_2 z_3 - z_1 + z_9, 0, 0, 0, \\ & z_3 z_4 - z_1 - z_4 + z_5 + z_9, -z_1 z_3 - z_2 z_3 - z_1 - z_4 + z_5 + z_9, \\ & z_9 - z_1, z_3 z_4 - z_1 + z_5, 0, 0), \\ f_{5,4} &= (z_{10} - z_1, -z_1 z_4 - z_1 + z_{10}, -z_1 z_4 - z_2 z_4 - z_1 + z_{10}, \\ & 0, 0, 0, -z_1 - z_3 + z_6 + z_{10}, \\ & z_1 z_2 + z_1 z_3 + z_2 z_3 - z_1 - z_3 + z_6 + z_{10}, \\ & z_3 z_4 - z_1 + z_{10}, z_{10} - z_1, 0, 0), \\ f_{5,5} &= (2z_1, z_1 + z_2, -z_1 z_2 + z_1 + z_3, 0, 0, 0, \\ & z_1 - z_6 + z_7, -z_1 z_2 + 2z_1 + z_3 - z_6 + z_7 - z_9 - z_{10}, \\ & z_1 + z_9, z_1 + z_{10}, 0, -2), \\ f_{5,6} &= (-z_1 - z_6 + z_7, -z_1 - z_7 - z_{11}, \\ & z_1 z_2 + z_3 z_4 - 2z_1 - z_3 - z_4 + z_6 + z_9 + z_{10}, 0, 0, 0, \\ & -z_1 + z_6 + z_7, z_1 z_2 - 2z_1 - z_3 + z_6 + z_9 + z_{10}, \\ & z_3 z_4 - z_1 - z_4 + z_5 - z_6 + z_7, -z_1 - z_3 + z_7, 0, 0), \quad (6.11) \\ f_{6,1} &= (0, 0, 0, -z_1 z_3 - z_1 z_4 - z_6 + z_{11}, -z_1 z_3 - z_6 + z_{11}, \\ & z_1 - z_6, z_1 - z_2 - z_6 + z_{11}, \\ & -z_1 z_3 - z_1 z_4 - z_1 - z_2 - z_6 + z_{11}, 0, 0, z_1 - z_6, 0), \\ f_{6,2} &= (0, 0, 0, z_1 z_3 + z_1 z_4 + z_1 + z_3 + z_4 + z_7 - z_6 - z_9 - z_{10} - z_{11}, \\ & z_1 z_3 + z_1 z_4 + z_2 z_3 + z_1 + z_3 + z_4 + z_7 - z_6 - z_9 - z_{10} - z_{11}, \\ & -z_1 z_3 - z_1 z_4 + z_1 + z_3 + z_4 + z_7 - z_6 - z_9 - z_{10} - z_{11}, \\ & -z_1 z_4 + z_1 + z_2 + z_4 + z_7 - z_6 - z_9 - z_{10} - z_{11}, \\ & z_1 z_3 + z_1 z_4 + z_1 + z_2 + z_4 + z_7 - z_6 - z_9 - z_{10} - z_{11}, 0, 0, \\ & -z_1 z_4 + z_1 + z_3 + z_4 + z_7 - z_6 - z_9 - z_{10} - z_{11}, 0), \\ f_{6,3} &= (0, 0, 0, z_3 - z_4, z_3 - z_4, z_3 z_4 - z_4 + z_5, \\ & z_3 z_4 - z_1 - z_4 + z_5 + z_9, -z_1 z_3 - z_2 z_3 - z_1 - z_4 + z_5 + z_9, 0, 0, \\ & z_1 z_3 + z_3 z_4 - z_4 + z_5, 0), \\ f_{6,4} &= (0, 0, 0, -z_3 z_4 - z_5 + z_6, z_6 - z_5, z_6 - z_5, \\ & -z_1 - z_2 + z_6 + z_{10}, z_1 z_2 + z_1 z_3 + z_2 z_3 - z_1 - z_3 + z_6 + z_{10}, 0, 0, \\ & z_1 z_4 - z_5 + z_6, 0), \\ f_{6,5} &= (0, 0, 0, z_1 - z_6 + z_7 - z_9 - z_{10}, -z_6 + z_7 - z_{10}, \\ & -z_1 - z_6 + z_7, z_1 - z_6 + z_7, -z_1 z_2 + z_1 z_3 + z_5 - z_6 + z_7 - z_9 - z_{10}, \\ & 0, 0, -z_2 - z_6 + z_7, 0), \\ f_{6,6} &= (0, 0, 0, -z_3 z_4 + z_4 + z_6, z_3 + z_6, 2z_6 - z_1 + z_6 + z_7, \\ & z_1 z_2 - 2z_1 - z_3 + z_6 + z_9 + z_{10}, 0, 0, z_6 + z_{11}, -2). \end{aligned}$$

Computing module intersections

Given $\mathcal{M}_1 = \langle v_1, \dots, v_p \rangle$ and $\mathcal{M}_2 = \langle w_1, \dots, w_q \rangle$ with v_i, w_j m -tuples of polynomials. Let Q denote the $m \times (p+q)$ matrix

$$Q = \begin{pmatrix} \vdots & & \vdots & \vdots & & \vdots \\ v_1 & \cdots & v_p & w_1 & \cdots & w_q \\ \vdots & & \vdots & \vdots & & \vdots \end{pmatrix}$$

Then compute wrt. POT and variable order $[z_1, \dots, z_m] \succ [s_{ij}]$

$$\langle h_1, \dots, h_t \rangle \equiv \text{Gröbner basis of column space of } \begin{pmatrix} & Q & \\ 1 & & 0 \\ & \ddots & \\ 0 & & 1 \end{pmatrix}$$

Selecting $h_i = (\overbrace{0, \dots, 0}^m, x_1, \dots, x_p, y_1, \dots, y_q)$, we have

$$0 = \sum_{j=1}^p x_j v_j + \sum_{k=1}^q y_k w_k$$

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$$Q = \begin{pmatrix} \vdots & & \vdots & \vdots & & \vdots \\ v_1 & \cdots & v_p & w_1 & \cdots & w_q \\ \vdots & & \vdots & \vdots & & \vdots \end{pmatrix}$$

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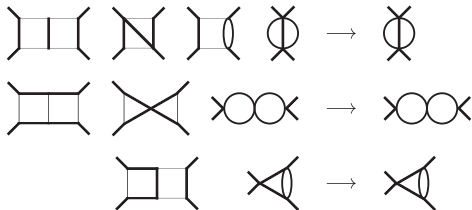
Selecting $h_i = (\overbrace{0, \dots, 0}^m, x_1, \dots, x_p, y_1, \dots, y_q)$, we have

$$0 = \sum_{j=1}^p x_j v_j + \sum_{k=1}^q y_k w_k \implies \sum_{j=1}^p x_j v_j = - \sum_{k=1}^q y_k w_k \in \mathcal{M}_1 \cap \mathcal{M}_2$$

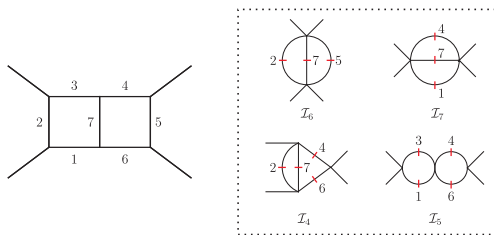
Hence $\sum_{j=1}^p x_j v_j$ generate $\mathcal{M}_1 \cap \mathcal{M}_2$, taking (x_1, \dots, x_p) from each h_i .

Spanning set of cuts for IBPs

To find the complete IBP reduction, we must consider the cuts associated with “uncollapsible” masters:



A bit more explicitly, the cuts we need to consider are



Main example: non-planar hexagon box

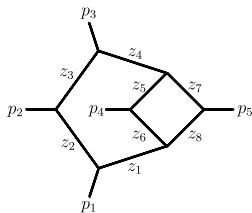
Task: IBP reduce non-planar hexagon box with numerator insertions of degree four in the z_i

[Chicherin, Henn, Mitev JHEP **05**(2018)164]

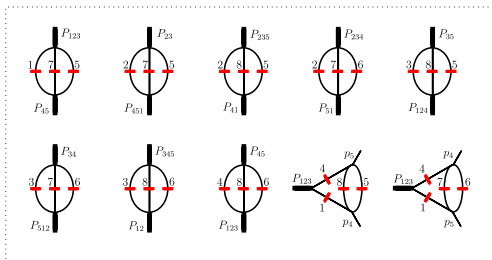
[Badger, Brønnum-Hansen, Hartanto, Peraro, PRL **120**(2018)092001]

[Abreu, Cordero, Ita, Page, Zeng, 1712.03946]

[Chawdhry, Lim, Mitov, 1805.09182]

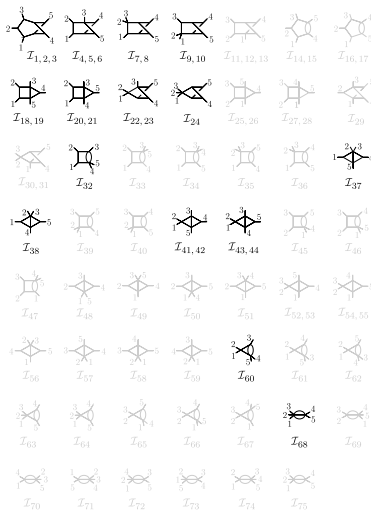


There are 10 cuts to consider:



Non-planar hexagon box: spanning set of cuts

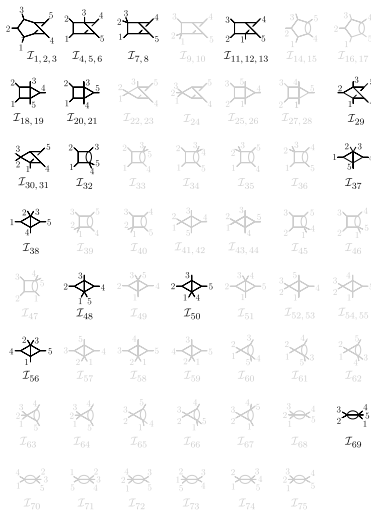
Construct and solve IBP identities on a spanning set of cuts.



Cut $\{1, 5, 7\}$

Non-planar hexagon box: spanning set of cuts

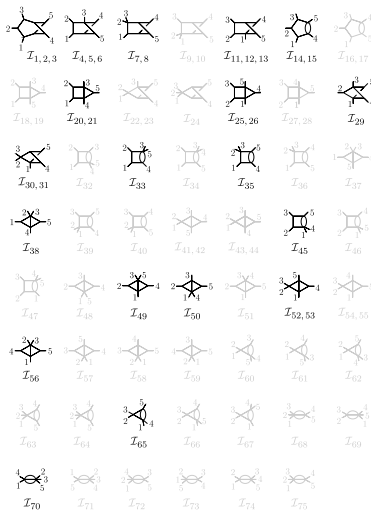
Construct and solve IBP identities on a spanning set of cuts.



Cut $\{2, 5, 7\}$

Non-planar hexagon box: spanning set of cuts

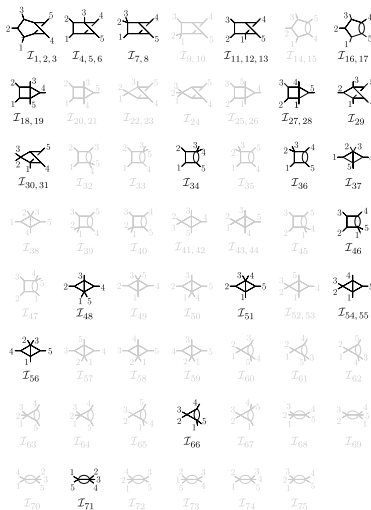
Construct and solve IBP identities on a spanning set of cuts.



Cut {2, 5, 8}

Non-planar hexagon box: spanning set of cuts

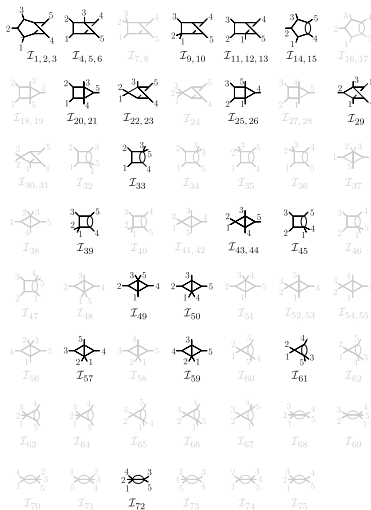
Construct and solve IBP identities on a spanning set of cuts.



Cut {2, 6, 7}

Non-planar hexagon box: spanning set of cuts

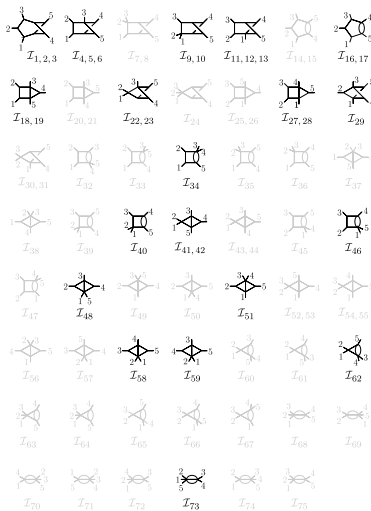
Construct and solve IBP identities on a spanning set of cuts.



Cut $\{3, 5, 8\}$

Non-planar hexagon box: spanning set of cuts

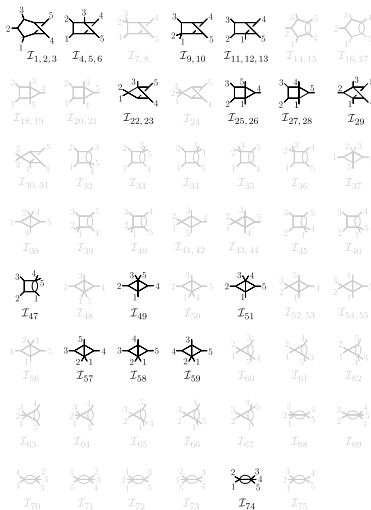
Construct and solve IBP identities on a spanning set of cuts.



Cut $\{3, 6, 7\}$

Non-planar hexagon box: spanning set of cuts

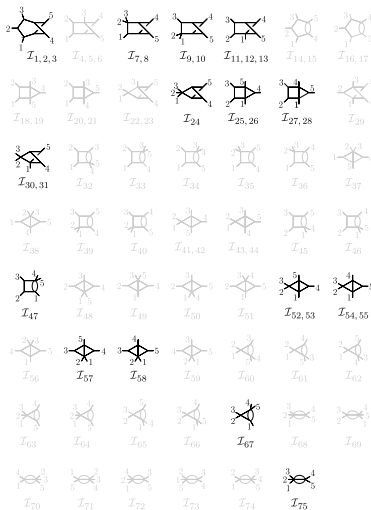
Construct and solve IBP identities on a spanning set of cuts.



Cut $\{3, 6, 8\}$

Non-planar hexagon box: spanning set of cuts

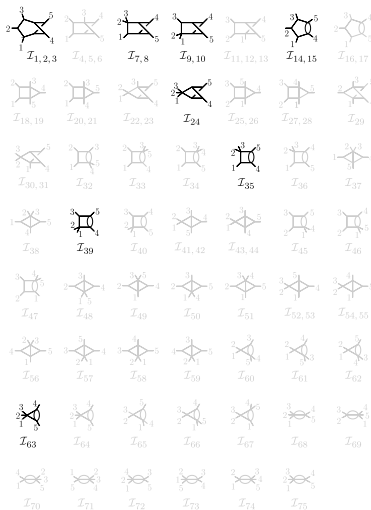
Construct and solve IBP identities on a spanning set of cuts.



Cut $\{4, 6, 8\}$

Non-planar hexagon box: spanning set of cuts

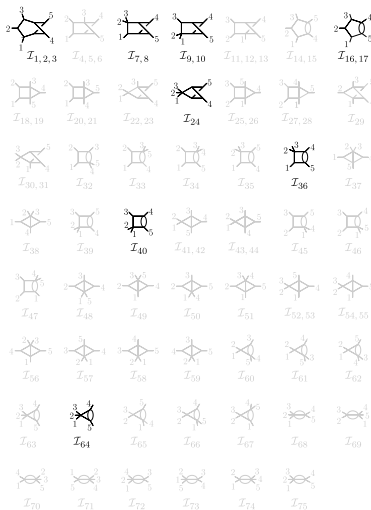
Construct and solve IBP identities on a spanning set of cuts.



Cut $\{1, 4, 5, 8\}$

Non-planar hexagon box: spanning set of cuts

Construct and solve IBP identities on a spanning set of cuts.



Cut $\{1, 4, 6, 7\}$

- Timing to compute $M_1|_{\text{cut}} \cap M_2|_{\text{cut}}$: **25-800 s**
(on 24 cores, 3.40 GHz)

- Size of generating systems after trimming: **1.5-10 MB**

Plug **resulting generators** into ansatz for total derivative:

$$0 = \int \left[\sum_{i=1}^{m-c} \left(\frac{\partial a_{r_i}}{\partial z_{r_i}} + \frac{D-L-E-1}{2G(z)} a_{r_i} \frac{\partial G}{\partial z_{r_i}} \right) - \sum_{i=1}^{k-c} \frac{a_{r_i}}{z_{r_i}} \right] \frac{G(z)^{\frac{D-L-E-1}{2}}}{z_{r_1} \cdots z_{r_{k-c}}} dz_{r_1} \cdots dz_{r_{m-c}}$$

- Resulting linear systems to solve:
700-1200 equations, size **1 MB**, density **1.5%**

Gauss-Jordan elimination of IBP systems

To find the IBP reductions, Gauss-Jordan eliminate IBP systems.

Some remarks:

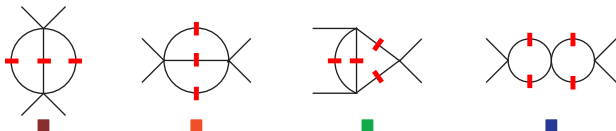
- To preserve sparsity, use a *total pivoting* strategy (i.e., allow column swaps)
- For cut $\{1, 4, 6, 7\}$, the RREF can be performed fully analytically, requiring 31 minutes on one core and 1.5 GB RAM.
- \vdots
- For $\{3, 6, 7\}$, assigned numerical values to two s_{ij} .
Ran 440 points on cluster (2.5 h and 1.8 GB RAM per job).
Used interpolation code to get analytical results (23 min and 15 GB RAM on one core).

[von Manteuffel and Schabinger, PLB **744**(2015)101]

[Peraro, JHEP12(2016)030]

Merging on-shell IBP reductions

By solving the IBP identities on the following cuts



we reconstruct the *complete IBP reductions* by merging the partial results.

An example of an IBP relation produced by our method ($\chi \equiv t/s$):

$$\begin{aligned}
 & \left(\text{Diagram 1} \right)^2 = \frac{(D-4)s^2\chi}{8(D-3)} \text{Diagram 2} - \frac{(3D-2\chi-12)s}{4(D-3)} \text{Diagram 3} + \frac{(4-D)(9\chi+7)}{4(D-3)} \text{Diagram 4} \\
 & + 2 \text{Diagram 5} + \frac{(10-3D)(2\chi-13)}{8(D-4)s} \text{Diagram 6} + \frac{2D(\chi+1)-8\chi-7}{2(D-4)s} \text{Diagram 7} \\
 & + \frac{9(3D-10)(3D-8)}{4(D-4)^2s^2\chi} \text{Diagram 8} + \frac{(3D-10)(3D-8)(2\chi+1)}{2(D-4)^2(D-3)s^2} \text{Diagram 9}
 \end{aligned}$$

- Fully analytic IBP reductions of the 32 hexagon boxes

$$\begin{aligned} & \{I(1, 1, 1, 1, 1, 1, 1, 1, 0, 0, -4), & I(1, 1, 1, 1, 1, 1, 1, 1, 0, -1, -3), & I(1, 1, 1, 1, 1, 1, 1, 1, 0, -2, -2) \\ & I(1, 1, 1, 1, 1, 1, 1, 1, 0, -3, -1), & I(1, 1, 1, 1, 1, 1, 1, 1, 0, -4, 0), & I(1, 1, 1, 1, 1, 1, 1, 1, -1, 0, -3) \\ & I(1, 1, 1, 1, 1, 1, 1, 1, -1, -1, -2), & I(1, 1, 1, 1, 1, 1, 1, 1, -1, -2, -1), & I(1, 1, 1, 1, 1, 1, 1, 1, -1, -3, 0) \\ & I(1, 1, 1, 1, 1, 1, 1, 1, -2, 0, -2), & I(1, 1, 1, 1, 1, 1, 1, 1, -2, -1, -1), & I(1, 1, 1, 1, 1, 1, 1, 1, -2, -2, 0) \\ & I(1, 1, 1, 1, 1, 1, 1, 1, -3, 0, -1), & I(1, 1, 1, 1, 1, 1, 1, 1, -3, -1, 0), & I(1, 1, 1, 1, 1, 1, 1, 1, -4, 0, 0) \\ & I(1, 1, 1, 1, 1, 1, 1, 1, 0, 0, -3), & I(1, 1, 1, 1, 1, 1, 1, 1, 0, -1, -2), & I(1, 1, 1, 1, 1, 1, 1, 1, 0, -2, -1) \\ & I(1, 1, 1, 1, 1, 1, 1, 1, 0, -3, 0), & I(1, 1, 1, 1, 1, 1, 1, 1, -1, 0, -2), & I(1, 1, 1, 1, 1, 1, 1, 1, -1, -1, -1) \\ & I(1, 1, 1, 1, 1, 1, 1, 1, -1, -2, 0), & I(1, 1, 1, 1, 1, 1, 1, 1, -2, 0, -1), & I(1, 1, 1, 1, 1, 1, 1, 1, -2, -1, 0) \\ & I(1, 1, 1, 1, 1, 1, 1, 1, -3, 0, 0), & I(1, 1, 1, 1, 1, 1, 1, 1, 0, 0, -2), & I(1, 1, 1, 1, 1, 1, 1, 1, 0, -1, -1) \\ & I(1, 1, 1, 1, 1, 1, 1, 1, 0, -2, 0), & I(1, 1, 1, 1, 1, 1, 1, 1, -1, 0, -1), & I(1, 1, 1, 1, 1, 1, 1, 1, -1, -1, 0) \\ & I(1, 1, 1, 1, 1, 1, 1, 1, 0, 0, -1), & I(1, 1, 1, 1, 1, 1, 1, 1, 0, -1, 0) \} \end{aligned}$$

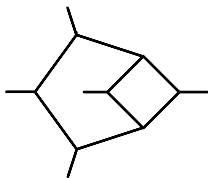
can be downloaded from (268 MB compressed / 790 MB uncompressed)

https://github.com/yzhphy/hexagonbox_reduction/releases/download/1.0.0/hexagon_box_degree_4_Final.zip

- Our results agree with fully numerical results from FIRE5 C++ (6 hours per point).

[A. Smirnov, CPC 189(2015)182]

- New formalism for IBP reductions. Main ideas: cuts, IBP identities from syzygies, total pivoting, rational reconstruction
- Obtained the **fully analytic** IBP reductions of



with numerator insertions up to degree 4 in the z_i .

- Powerful framework. IBP reductions for further $2 \rightarrow 3$ two-loop processes **seem well within reach**.