Jet SIFT-ing

A Scale-Invariant Jet Clustering Algorithm for the Substructure Era

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SIFT: Scale-Invariant Filter Tree Outline of Presentation

• SCALE INVARIANT

- Jet Clustering Background
- Motivation for Scale Invariance
- Algorithm Implementation
- Algorithm Visualization
- Algorithm Testing
- FILTER
 - Integrated Grooming
 - Remove Soft Co-Linear Radiation
 - A Natural Halting Condition

• TREE

- Fast Algorithms
- Multidimensional Trees

SIFT: SCALE-INVARIANT Filter Tree

- Traditional Jet Clustering imposes a fixed cone size, and thus a fixed scale on events
- Boosted objects tend to collimate and fall into a single jet radius
- Substructure techniques are essential for recovering information inside the jet
- However, these techniques are often complicated, with de- and re-clustering
- We propose as SCALE INVARIANT approach which is intrinsically suitable for tagging substructure AS the jet is being assembled

Collider Variables & Coordinates

- Transverse components (perpendicular to the beam) are very important (invariant under longitudinal boosts, P_{τ} total is zero)
- Differences in orientation characterized by ΔR , referring also to azimuth angle ϕ
- The pseudorapidity η is a proxy for the polar (beam) angle θ , defined such that differences $\Delta \eta$ are (almost) invariant under longitudinal boosts
- This invariance is exact for the rapidity *y* (difference is handling of MASS)

$$P_{\rm T} \equiv \sqrt{P_x^2 + P_y^2}$$
$$\Delta R \equiv \sqrt{(\Delta \eta)^2 + (\Delta \phi)^2}$$
$$\eta \equiv \frac{1}{2} \ln \left(\frac{|\vec{P}| + P_z}{|\vec{P}| - P_z} \right) \equiv -\ln \tan \left(\frac{\theta}{2} \right)$$
$$y \equiv \frac{1}{2} \ln \left(\frac{E + P_z}{E - P_z} \right) \equiv \ln \left(\frac{\sqrt{\cosh^2 \eta + \frac{M^2}{P_{\rm T}^2}} + \sinh \eta}{\sqrt{1 + \frac{M^2}{P_{\rm T}^2}}} \right)$$

FIG. 1: The pseudorapidity η (bold, orange) is plotted as a function of the polar angle θ . For comparison, the longitudinal rapidity y (fine, blue) is also shown for various values of $M/P_{\rm T}$, equal to $\{1/2, 1, 2, 5, 10, 20\}$ from top to bottom.



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Formation of Hadronic Jets

- The hard partonic event may result in the production of colored objects (at Feynman diagram level, e.g. MadGraph)
- These objects rapidly "shower", radiating quarks & gluons (e.g. Pythia)
- QCD confinement implies that strongly charged particles cannot exist as free objects at large separations; they must convert "hadronize" (e.g. Lund color strings in Pythia) into color-neutral particles such as pions, K mesons, etc.
- Color strings may convolve descendants of partonic objects with each other and even with the underlying beam; this is not still mitigated in a lepton collider

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Image: CMS

Image: Stefan Höche

Standard Jet Clustering Algorithms

- Hadronized objects need to be recombined in a manner that preserves correlation with the underlying hard (partonic) event
- 3 related algorithms reference an input angular width R₀ & differ by an index n
- Objects more widely separated than R₀ will never be clustered
- n = 0, or "Cambridge/Aachen" clusters objects with high angular adjacency
- n = +1, or "kT" additionally favors clustering of soft pairs first
- n = -1, or "Anti-kT" prioritizes clustering where one of the pair is hard
- Anti-kT is now the default jet clustering tool at LHC, with $R_0 = (0.4, 0.5)$
- It is robust against "soft" and "colinear" jet perturbations and has regular jet shapes which are favorable for calibration against pileup, etc.

$$\delta_{AB} \equiv \min\left[P_{\mathrm{T}A}^{2n}, P_{\mathrm{T}B}^{2n}\right] \times \left(\frac{\Delta R}{R_0}\right)^2$$

Jet Substructure

- Highly boosted mothers will tend to yield very collimated daughters
- In hadronic top quark decays t ⇒ W/b ⇒ u/d/b with COM energy above a TeV, the likelihood of resolving only 2 or even 1 discrete object increases
- For example, within, a "fat" (large $R_0 \gtrsim 1$), N-Subjettiness τ_N can characterize how well the event matches an N-prong hypothesis (axes chosen separately)
- The best discrimination comes from the ratio r_N , e.g. how much more 3-prong-like is the event than 2-prong like
- Variable cone sizes have also been considered to cope with loss of structure

Given N axes
$$\hat{n}_k$$
, $\tau_N = \frac{\sum_{i \in J} p_{T,i} \min(\Delta R_{ik})}{\sum_{i \in J} p_{T,i} R_0}$
 $r_N = \frac{\tau_N}{\tau_{N-1}}$

A Scale-Invariant Jet Algorithm

- It may be worth asking whether alternative techniques could provide intrinsic resiliency to boosted event structure; this requires dropping the input scale R₀
- It would be good to "asymptotically" recover the favorable behavior of Anti-kT
- Numerator should favor angular collimation; we propose ΔM^2 , similar to JADE
- Denominator should suppress soft pair clustering; we propose a sum of E_T
- Result is dimensionless, Lorentz invariant (longitudinally in the denominator), and free from references to external / arbitrary scales

$$\delta_{AB} \equiv \frac{\Delta M_{AB}^2}{E_{TA}^2 + E_{TB}^2} \qquad \qquad M^{A,B} \equiv \sqrt{\left(P_{\mu}^A + P_{\mu}^B\right)\left(P_A^\mu + P_B^\mu\right)} \\ = \sqrt{M_A^2 + M_B^2 + 2\left(E^A E^B - \vec{P}^A \cdot \vec{P}^B\right)} \\ \lim_{M_A = M_B = 0} \Rightarrow \sqrt{2|\vec{P}^A||\vec{P}^B|\left(1 - \cos\Delta\varphi^{B,A}\right)} \\ E_T \equiv \sqrt{M^2 + \vec{P}_T \cdot \vec{P}_T} = \sqrt{E^2 - P_z^2} \\ \lim_{M = 0} \Rightarrow |\vec{P}_T|$$

Hadronic TTbar Scale-Invariant Clustering

https://youtu.be/u9Z4qDuXL84



Test of Pre/Post Merger Statistic for Di-jets



- 95% of pairs reconstructed prior to 0.1
- 95% of final final mergers are after 2.0
- Results are invariant wrt beam energy

Visualization of Statistic Jump at Clustering



• The event jettiness count is intrinsically imprinted on the clustering history

Matching of final 6 objects with Truth-Level Quarks





Lepton to TTbar 2.5 TeV Anti-KT 0.5 with Ghosts

https://youtu.be/1fhbhlDrORA



Lepton to TTbar 2.5 TeV Scale Invariant Clustering with Ghosts

https://youtu.be/kxUmgv1HHMs



SIFT: Scale-Invariant FILTER Tree

- Running to termination can lead to merging of stray radiation
- Take a cue from "Soft Drop" (2014 Larkoski, Marzani, Soyez, Thaler)
- This procedure "Grooms" a jet by removing soft, wide-angle radiation to mitigate contamination from ISR, UE, and pileup
- SD iteratively DECLUSTERS C/A, dropping softer object unless & until:

$$\frac{\min(P_{TA}, P_{TB})}{P_{TA} + P_{TB}} > z_{\text{cut}} \left(\frac{\Delta R_{AB}}{R_0}\right)^{\beta}$$

- Typically, $z_{\rm cut}$ is $\mathcal{O}(0.1)$, and $\beta > 0$ for grooming
- We propose a scale-invariant analog which is applied within the original clustering itself.

$$\frac{E_{TA}E_{TB}}{E_{TA}^2 + E_{TB}^2} > \frac{\Delta M_{AB}^2}{2E_{TA}E_{TB}} \implies \delta_{AB} \equiv \frac{\Delta M_{AB}^2}{E_{TA}^2 + E_{TB}^2} < \frac{2E_{TA}^2 E_{TB}^2}{(E_{TA}^2 + E_{TB}^2)^2}$$

- The softer object is considered isolated unless it passes this FILTER
- This provides a natural halting condition to prevent total assimilation
- Curiously, the dynamic threshold is symmetric under $E_T \rightarrow 1/E_T$

Hadronic TTbar Scale-Invariant Clustering with Filtering

https://youtu.be/rDsBeEBTimw

Lepton to TTbar 2.5 TeV SIFT Filtered Clustering with Ghosts

https://youtu.be/G1XB5sQaolk

SIFT: Scale-Invariant Filter TREE

- A jet clustering algorithm is USELESS practically unless it is FAST
- Critical issue is the scaling dimension with number N of constituents
- A naïve implementation is CUBIC $O(N^3)$ because there are N mergers with a scan over N x N possible pairings at each stage. TOO SLOW!
- Why is FastJet (Cacciari, Salam, Soyez) FAST?
- FJ Lemma trims to $\mathcal{O}(N^2)$ by scanning only GEOMETRIC nearest neighbors
- How? The magic of "min of a min" facilitates factorization
- GLOBAL min of δ_{AB} has the property that *B* minimizes ΔR_{AB} if $P_{TA}^{2n} < P_{TB}^{2n}$

$$\delta_{AB} \equiv \min(P_{TA}^{2n}, P_{TB}^{2n}) \times \left(\frac{\Delta R_{AB}}{R_0}\right)^2$$

- Then, with a FAST O(log N) algorithm for caching neighbors, the combined runtime can be "linearithmic" O(N log N). GOLD STANDARD!
- Signature of $O(\log N)$ algorithms is halving of problem size with each cycle
- Example is "bisection" method of traversing a sorted list
- The FAST approach to finding nearest neighbors can use a TREE

Can SIFT be FAST?

- If yes, there needs to be something like a "GEOMETRIC" measure
- As originally expressed, the metric is not even written in terms of coordinates
- For massless A & B, $\Delta M_{AB}^2 = 2P_A^{\mu}P_{\mu}^B \Rightarrow 2P_A P_B(1 \cos\Delta\theta) \approx P_A P_B(\Delta\theta^2 \Delta\theta^4/12)$
- But, we need to refer to the collider coordinates of A & B directly ($\Delta \eta_{AB}$, $\Delta \phi_{AB}$, etc.)
- Conjecture: for massive A & B, it will actually be Δy_{AB} that is relevant
- Boost from the $P_z = 0$ frame into the lab:

$$\begin{pmatrix} E \\ P_Z \end{pmatrix} = \begin{pmatrix} \cosh y & \sinh y \\ \sinh y & \cosh y \end{pmatrix} \begin{pmatrix} E_T \\ 0 \end{pmatrix} = \begin{pmatrix} E_T \cosh y \\ E_T \sinh y \end{pmatrix}$$
$$2P_A^{\mu}P_{\mu}^{B} = 2(E_A E_B - P_Z^{A} P_Z^{B} - P_T^{A} P_T^{B} \cos \Delta \theta_{AB})$$
$$= 2(E_T^{A} E_T^{B} [\cosh y^A \cosh y^B - \sinh y^A \sinh y^B] - P_T^{A} P_T^{B} \cos \Delta \theta_{AB})$$
$$= 2(E_T^{A} E_T^{B} \cosh \Delta y^{AB} - P_T^{A} P_T^{B} \cos \Delta \theta_{AB})$$

- We are getting WARM. BUT the difference between $E_T \& P_T$ (i.e. MASS) means that we CANNOT perfectly factorize kinematics from geometrics
- Nevertheless, we can proceed. BUT, we must seek neighbors in a 3D or 4D space
- The FastJet engine (Voronoi Tesselation) is 2D. We need a custom engine.
- NOTE: hyperbolic cosine differs from cosine in that all Taylor terms are POSITIVE

Building an D-Dimensional Tree

- "Balanced KD-Tree" framework (2003 Procopiuc, Agarwal, Arge, Vitter) is suitable
- The forking property of a tree allows $O(\log N)$ traversal
- Each descending "row" of the tree sorts on the next cyclic coordinate index
- To stay "balanced" we never add objects to a tree after initial construction
- We maintain a "forest" of trees of doubling size, as needed
- Protocols for pruning, grafting, and merging leaves must be built in
- Be sure to not reinject $O(N^2)$ scaling in these updates. Non-Trivial!
- Protocols for neighbor finding under a user defined metric must be built in
- Use "templating" to allow input from user-defined data structures
- Cyclic indices: extend by half principal domain either way & build "image" leaves
- Status: working D-dimensional $\mathcal{O}(N \log N)$ implementation exists / tested on Anti-kt
- Currently, this is being ported to C++ for increased speed in the "coefficient"

Conclusions and Ongoing Work

- SIFT is a SCALE INVARIANT clustering algorithm designed specifically for substructure
- FILTER-ing of soft and co-linear radiation can be done as the jet is clustered
- Organization of the data structure in a balanced TREE can make clustering fast
- The clustering history holds *information* it may be better to not halt at fixed radius.
- Could the algorithm be applied to existing fat jets for exclusive clustering?
- What is the jet-energy resolution width, and does it vary with P_T ?
- How does SIFT fare with pileup subtraction?
- How does the absolute mass of reconstructed particles connect?
- Is the distilled clustering history amenable to machine learning applications?
- Can SIFT intrinsically confront the problem of tagging boosted objects?

Thank You

(movie notebook available upon request to jwalker@shsu.edu)