Scalla/xrootd

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http://xrootd.slac.stanford.edu/

Outline

System Overview

- What's it made of and how it works
- **#** Opportunistic Clustering
 - Batch nodes as data providers
- # Expansive Clustering
 - Federation for speed and fault tolerance
 - The Virtual Mass Storage System
- **#** Fullness vs Simplification



Full Scalla/xrootd Overview





The Components

- **x**rootd
 - Provides actual data access
- **#** cmsd
 - Glues multiple xrootd's into a cluster
- # cnsd
 - This might not be needed Glues multiple name spaces into one name space
- **#** BeStMan
 - Provides SRM v2+ interface and functions
- **#** FUSE
 - Exports xrootd as a file system for BeStMan
- **#** GridFTP
 - Grid data access either via FUSE or POSIX Preload Library



Getting to xrootd hosted data

 Automatic when files named root:// 	
 Manually, use TXNetFile() object 	
Note: identical TFile() object will not work with xrootd!	
# xrdcp	Native Set
 The native copy command 	
# POSIX preload library	这些法国合议
 Allows POSIX compliant applications to use xrootd 	
# gridFTP	Simple Add
BeStMan (SRM add-on)	
 srmcp for srm-to-srm copies 	Intensive
# FUSE	Full Grid Set
 Linux only: xrootd as a mounted file system 	



Cluster Maneuvering





Corresponding Configuration File

```
# General section that applies to all servers
#
all.export /atlas
if redirector.slac.stanford.edu
all.role manager
else
all.role server
fi
all.manager redirector.slac.stanford.edu 3121
# Cluster management specific configuration
#
cms.allow *.slac.stanford.edu
# xrootd specific configuration
#
xrootd.fslib /opt/xrootd/prod/lib/libXrdOfs.so
xrootd.port 1094
```



File Discovery Considerations

The redirector does not have a catalog of files

- It always asks each server, and
- Caches the answers in memory for a "while"
 - So, it won't ask again when asked about a past lookup
- # Allows real-time configuration changes
 - Clients never see the disruption
- **#** Does have some side-effects
 - The lookup takes less than a millisecond when files exist
 - Much longer when a requested file does not <u>exist!</u>



Handling Missing Files



Missing File Considerations

System optimized for "file exists" case! This is the classic analysis situation Penalty for going after missing files **#** Aren't *new* files, by definition, missing? Yes, but that involves writing data! The system is optimized for reading data • So, creating a new file *will* suffer a 5 second delay • Can minimize the delay by using the **xprep** command Primes the redirector's file memory cache ahead of time



Why Do It This Way?

Simple, lightweight, and ultra-scalable Ideal for opportunistic clustering E.g., leveraging batch worker disk space Ideal fit with PROOF analysis Has the R³ property (Real-Time Reality Representation) Allows for ad hoc changes Add and remove servers and files without fussing

Restart anything in any order at any time

- Ideal for expansive clustering
 - E.g., cluster federation & globalization
 - Virtual mass storage systems and torrent transfers



Opportunistic Clustering

■ Xrootd *extremely* efficient of machine resources ■ Ultra low CPU usage with a memory footprint 20 ≈ 80MB

Ideal to cluster just about anything





Opportunistic Clustering Caveats

Using batch worker node storage is problematic

- Storage services must compete with actual batch jobs
 - At best, may lead to highly variable response time
 - At worst, may lead to erroneous redirector responses
- # Additional tuning will be required
 - Normally need to renice the cmsd and xrootd
 - As root: renice –n -10 –p *cmsd_pid*
 - As root: renice –n -5 –p *xroot_pid*
- You must not overload the batch worker node
 Especially true if exporting local work space



Opportunistic Clustering & PROOF

■ Parallel Root Facility layered on xrootd Good architecture for "map/reduce" processing **#** Batch-nodes provide PROOF infrastructure Reserve and use for interactive PROOF Batch scheduler must have a drain/reserve feature Use nodes as a parallel batch facility Good for co-locating application with data Use nodes as data providers for other purposes



PROOF Analysis Results



Akira's talk about "Panda oriented" ROOT analysis comparison at the Jamboree

http://indico.cern.ch/getFile.py/access?contribId=10&sessionId=0&resId=0&materialId=slides&confId=38991



Sergey Panitkin

Expansive Clustering

Xrootd can create ad hoc cross domain clusters
Good for easily federating multiple sites
This is the ALICE model of data management
Provides a mechanism for "regional" data sharing
Get missing data from close by before using dq2get
Architecture allows this to be automated & demand driven
This implements a Virtual Mass Storage System



Virtual Mass Storage System





What's Good About This?

Fetch missing files in a timely manner Revert to dq2get when file not in regional cluster **#** Sites can participate in an ad hoc manner The cluster manager sorts out what's available **The Can use R/T WAN** access when appropriate **The Can significantly increase WAN xfer rate** Using torrent-style copying



Torrents & Federated Clusters





Improved WAN Transfer

The xrootd already supports parallel TCP paths Significant improvement in WAN transfer rate ■ Specified as xrdcp – S num **X**treme copy mode uses multiple data sources ■ Specified as xrdcp –x Transfers to CERN; examples: 12MB/sec (1 stream) ■ 1 source (.de): $\blacksquare 1$ source (.us): 19MB/sec (15 streams) ■ 4 sources (3 x .de + .ru): 27MB/sec (1 stream each) ■ 4 sources + || streams: 42MB/Sec (15 streams each) 54MB/Sec (15 streams each) **5** sources $(3 \times .de + .it + .ro)$:



Expansive Clustering Caveats

Federation & Globalization are easy if
Federated servers are *not* blocked by a firewall
No ALICE xroot servers are behind a firewall
There are alternatives
Implement firewall exceptions

Need to fix all server ports
Use proxy mechanisms

Easy for some services, more difficult for others
 All of these have been tried in various forms
 Site's specific situation dictates appropriate approach



Summary Monitoring

Needed information in almost any setting **#** Xrootd can auto-report summary statistics Specify xrd.report configuration directive **#** Data sent to one or two locations Use provided mpxstats as the feeder program Multiplexes streams and parses xml into key-value pairs Pair it with any existing monitoring framework ■ Ganglia, GRIS, Nagios, MonALISA, and perhaps more



Summary Monitoring Setup





Putting It All Together





Can't We Simplify This?

The **cnsd** present for **XrootdFS** support Provide composite name space for "Is" command **# FUSE** present for **XrootdFS** support **# XrootdFS & FUSE** for **BeSTMan** support **# BeSTMan** for **SRM** support **I SRM** for push-type grid data management dq2get is a pull function and only needs gridFTP **#** Answer: **Yes!** This can be simplified.



Tearing It All Apart





In Conclusion...

■ Xrootd is a lightweight data access system

- Suitable for resource constrained environments
 - Human as well as hardware
- Geared specifically for efficient data analysis
- **#** Supports various clustering models
 - E.g., PROOF, batch node clustering and WAN clustering
 - Has potential to greatly simplify Tier 3 deployments
- **#** Distributed as part of the OSG VDT
 - Also part of the CERN root distribution
- Visit http://xrootd.slac.stanford.edu/



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