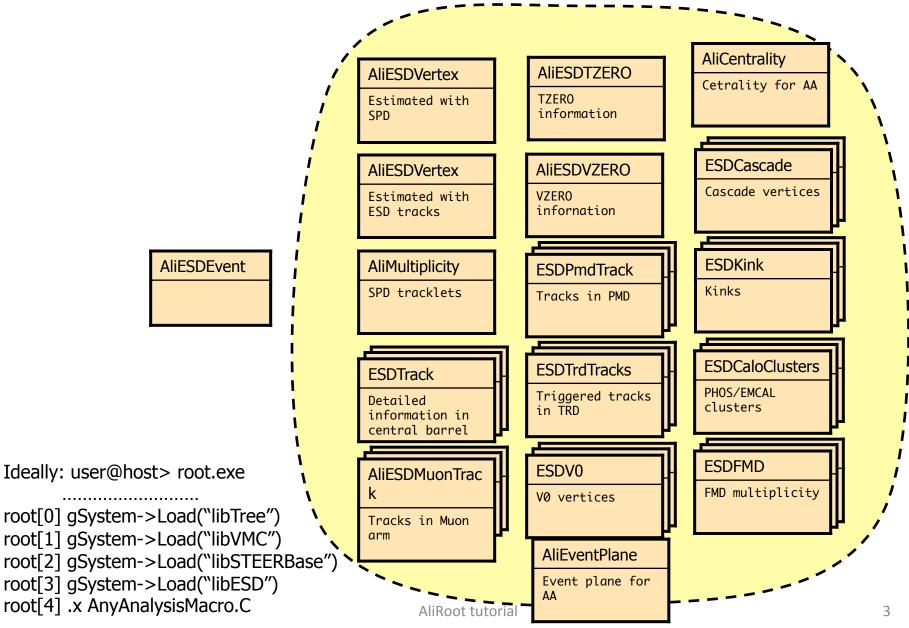
ALICE Event Data Model

P.Hristov 08/07/2018

What can we learn from the past? RUN1 & RUN2 EDM

The ESD



AliESDEvent and AliESDtrack classes

- Accumulation and exchange of tracking information among the barrel detectors
- Contained in the ESD and used for physics analysis

Class AliESDtrack: public AliExternalTrackParam

- final params
- reconstruction status flags
- length, time, combined PID
- vertex constrained params
- impact parameters & cov.matrix
- params at the outer TPC wall
- params at the inner TPC wall
- ...
- detector specific info (chi2, num.of clusters, PID...)

Common base classes for ESDs and AODs

AliVEvent

AliESDEvent

AliAODEvent

AliVHeader

AliESDHeader

AliAODHeader

- standard access to containers
- common getters and setters
- some differences in the interface (to be "cured"!)

AliVParticle

AliExternalTrackParam

AliAODTrack

. . .

AliESDtrack

Current content of the standard AOD

AliAODEvent

contains an (extendable) TList

AliAODHeader

event information

AliAODTrack

TClonesArray of tracks

AliAODVertex

TClonesArray of vertices

AliAODJet

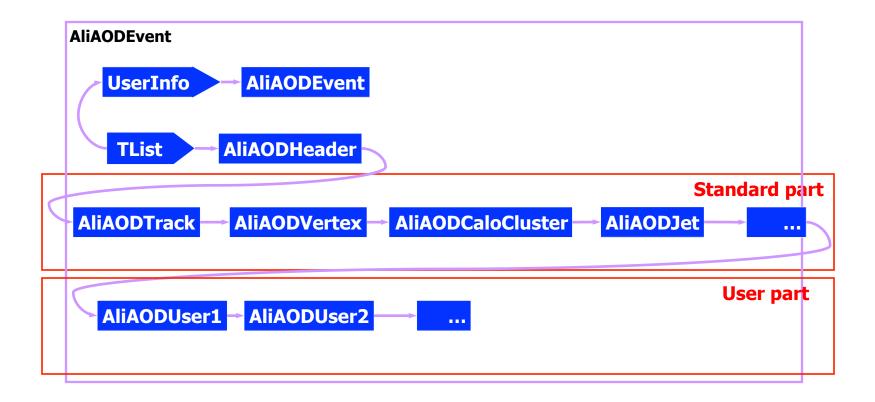
TClonesArray of jets

AliAODTracklets

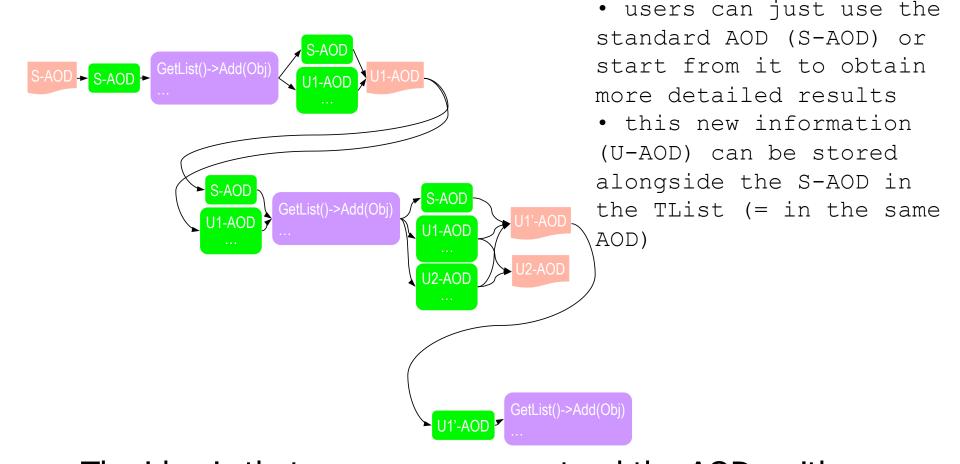
Container for SPD tracklets

General idea: ESD & AOD

Very simple... use a list



Extending the AOD



 The idea is that every user can extend the AODs with "non-standard" objects

AliRoot tutorial

Flexible and Extendable

Reading back

```
void AliAODEvent::GetStdContent() const
               // set pointers for standard content
                                                                     AliAODTrack* AliAODEvent::GetTrack(Int_t i) const
               fHeader = (AliAODHeader*)fAODObjects->At(0);
                                                                     { return (AliAODTrack*) fTracks->At(i); }
               fTracks = (TClonesArray*)fAODObjects->At(1);
               fVertices = (TClonesArray*)fAODObjects->At(2);
AliAODEvent
               fClusters = (TClonesArray*)fAODObjects->At(3);
                        = (TClonesArray*)fAODObjects->At(4);
               fJets
                                                                         ►fHeader
                                                                          Vertices
                                                                          Clusters
                                                                         fJets
                     AliAODEvent
    UserInfo
                     AliAODHeader
     TList
    AliAODTrack -
                                                            AliAODJet
                      AliAODVertex →
                                         AliAODCluster
    AliAODUser1
                        AliAODUser2
                                                            AliAODJet
```

•...and the same user also can access the "non-standard" objects from the previous slide

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Run1 & Run2 EDM: Summary

- Fully based on ROOT
- Event-oriented data model: trees of ESD & AOD/delta AOD, but also kinematics, ESD friends, track references, tags
 - Several ROOT files
 - Access to the different data via handlers
- Deep inheritance chains
- Containers of containers of objects

- Performance @ O2 prototype facility, GSI
 - Shared file system (Lustre)+ xrootd client plugin(0.6PB)
 - Job scheduler: Slurm, max.
 1000 job slots
 - Possibility for local or GRID jobs
- For one process
 - Simple copy: ~1.2 GB/s
 - Unzipping: ~100 MB/s
 - − P_T analysis: ~20 MB/s

Run1 & Run2: What did we learn?

- ESD and AOD => the functionality diverges despite of the common base classes
- We have IO bottleneck because of
 - Slow storage not EDM issue
 - Decompression partially related to EDM
 - Deserialization purely EDM issue
- Flexibility without discipline: AOD easily becomes "trash bit" where everybody tries to store specific objects
- Complex ESD/AOD structure with nested objects, containers of containers, pointers, etc. causes deserialization overhead
 - ROOT cannot completely recover the losses
 - The flexibility comes with performance price

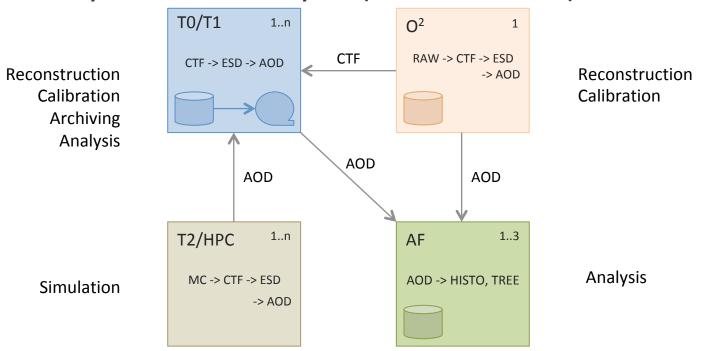
What do we plan? RUN3 EDM

O² Analysis Framework and Facilities



Run3: Important differences

- x100 more collisions (from 0.5-1 kHz to 50 kHz Pb-Pb)
- Continuous readout => raw data in compressed time frames (23 ms, 1000 MB Pb-Pb collisions, ~2GB)
- Only AOD for analysis (transient ESD)

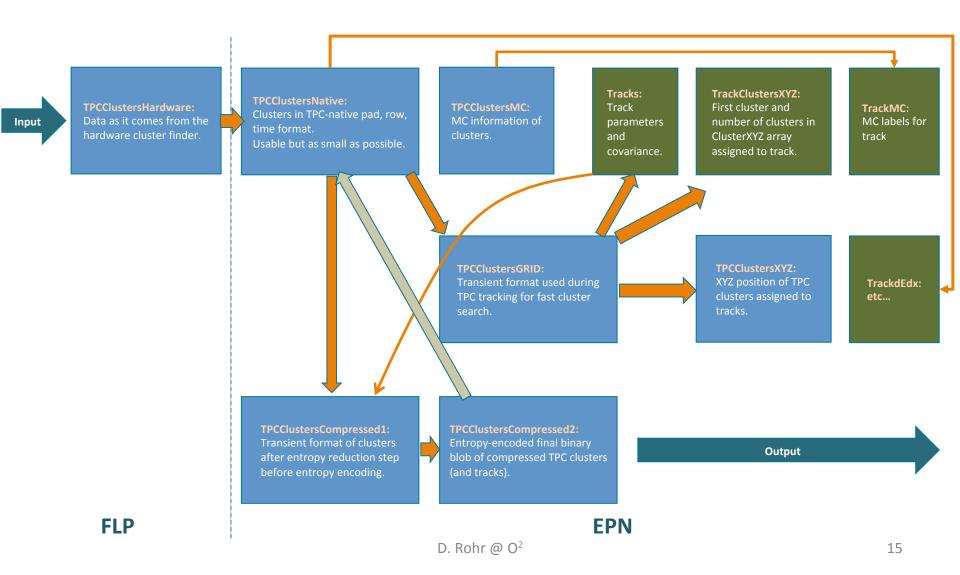


P.Buncic @ O^2

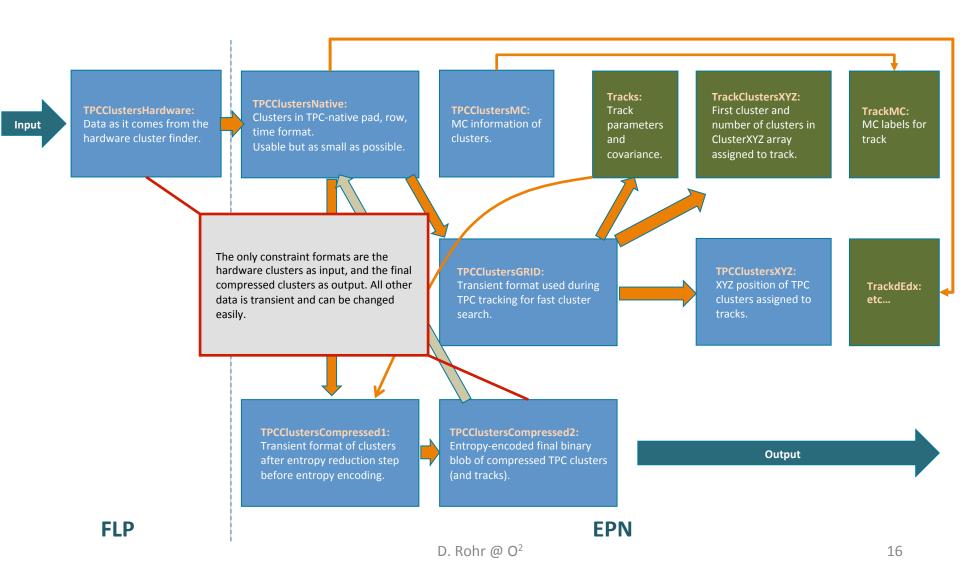
Run3 – important differences wrt Run1/2

- No "events" in general continuous readout without trigger
- The collision is defined by its primary vertex
 - It is an association (index) of primary tracks
 - The secondary objects: vertexes, kinks, tracks, etc. are associated with the primary vertex, but can be reassociated
- Many collisions in the same time frame
- Some collisions are triggered (they are "events" in Run1/2 sense)

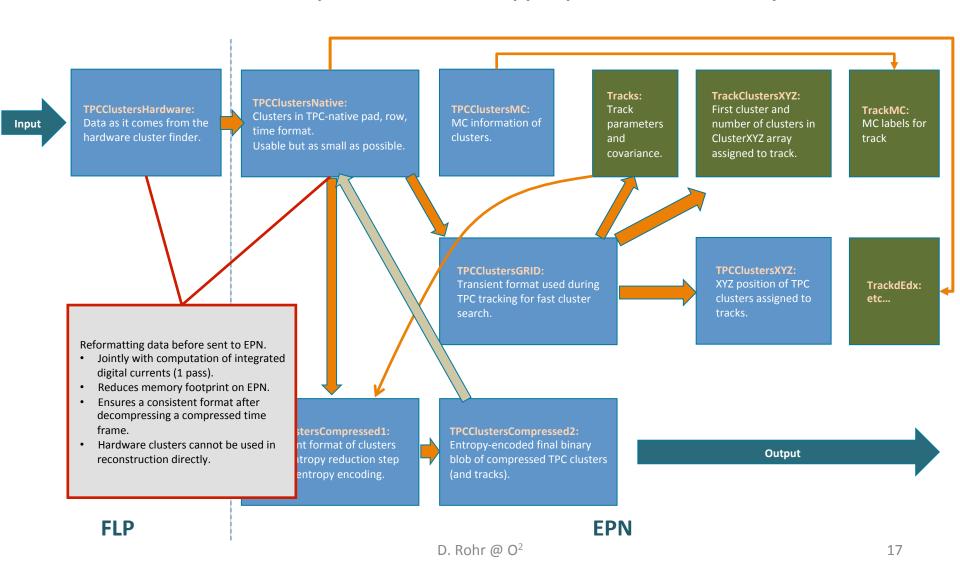
- TPC produces the bulk of data, so a format that minimizes memory consumption is needed.
 - We should avoid data duplication, but some copy steps seem to be necessary.



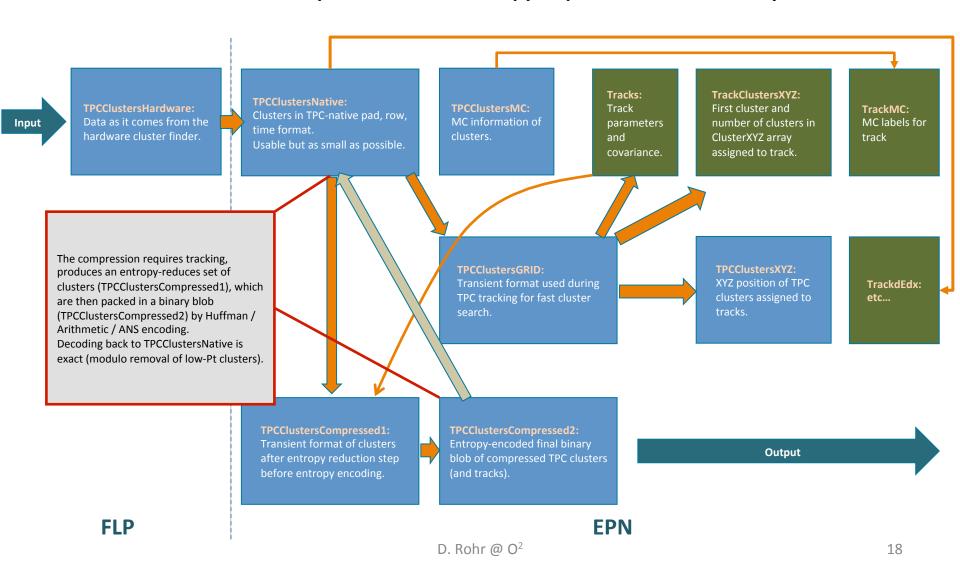
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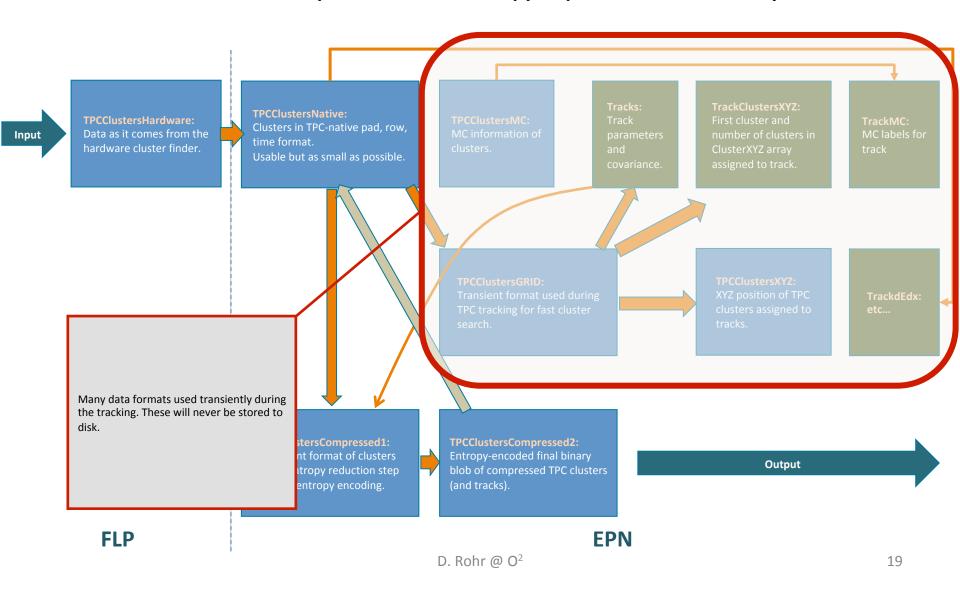
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Changes in the software concept

- From sequential ROOT-based processing and format to message based multiprocessing (see the ALFA talk @ CHEP2018)
- The message serialization is very important
 - Multiple languages are allowed
 - Heterogeneous architectures (different OS, CPU, GPU, FPGA,...)
- We are investigating different options also for analysis

Analysis data format: requirements

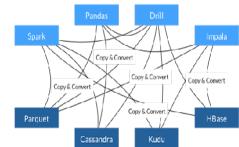
New data format should reduce as much as possible the cost of deserialization: some generality will be lost for the sake of improved speed

- Simple, flat: numbers only (no classes), use tables, cross-reference via numeric indices
- Columnar: SoA in-memory structure for better growing/shrinking and vectorization
- Extensible: base format is immutable, but easily extensible because it's SoA
- Chunked: a single timeframe can be divided in smaller units processable in parallel
- Zero size for null objects: filtered-out fields do not use RAM memory
- Recompute, don't store: do not store everything because recomputing may be cheaper
- No data restructuring: disk → memory → network should use similar representations
 D. Berzano @ CHEP2018

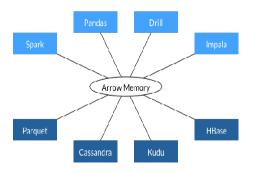
Apache Arrow Development

We have been experimenting with Apache Arrow: in-memory, language independent columnar data format. Has an on-disk companion too (Apache Parquet)

- Leverages vectorization and fits our other requirements
- **Units:** data organized in Tables, made of immutable Columns. Columns shared among tables (no copy)
- Memory management: Columns backed by Buffers, which allow for custom Memory Pools
- Meant for interoperability: allows for data exchange within the Apache ecosystem and outside, widely supported
- Fits the ALICE Run 3 data model based on message passing

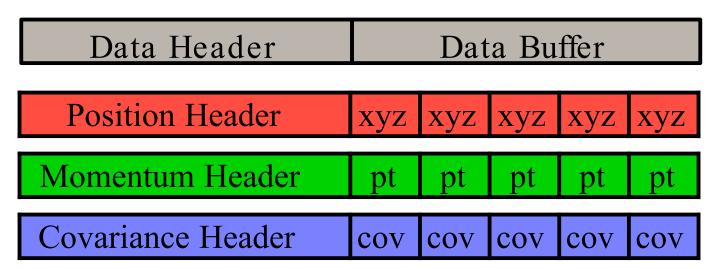






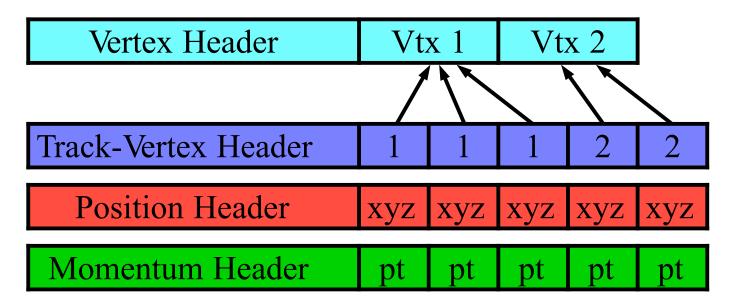
Data Message Development

- Use over the network message queuing (=buffers)
- Headers carry metadata used to uniquely identify the message payload
 - Data buffer type
 - Data buffer size
- Data element have to be flat (self-contained)
 - Any data type, arrays, structs, ...
 - No pointers



References in Data Message

- References between data buffers via the same index
- Separate data buffers can have different entry number
 - Use separate buffer for the reference
 - Track-Vertex contains the index of the vertex related to the track



More investigations: SOAContainer (LHCb)

- Code available in GitLab
- Tutorial <u>exists</u>
- Aims
 - exploiting modern CPUs/GPUs is hard enough
 - SOAContainer should help you try SOA code quickly, with low barrier of entry
 - · code should be readable, and do what it says writable by experts and non-experts alike
 - produce reasonable code to start with (profile, optimize later)
 - will not magically solve your performance problems
 - brain not outdated yet: need good idea and right algorithm
 - may still want to look at assembly to judge if compiler needs help (simpler loops, break dependencies, ...)
 - compilers may not autovectorise the way you want them to
- Use a skin to describe
 - which fields make up an SOA object
 - give the underlying tuple a nice interface gets all the methods defined in the fields
 - can define methods that access more than one field

More investigations: Libflatarray

- Who should use it?
 - Basically anyone who wants to store objects or structs in an array with 2 or more dimensions and needs to iterate through them.

Features

- object-oriented interface, but SoA memory layout
- good for vectorization (SSE, AVX, CUDA)
- zero address generation overhead
- fast iterators, yield access to neighboring elements
- efficient copy in/out
- Small development team (5 contributors)

Conclusions and Plans

- Variety of data formats in the quasi-online O² processing
- ALICE is designing a new flat event (collision) data model for the Run3 analysis
- We have gained a lot of experience from Run1/2 (positive: organized analysis in trains; negative: design of EDM)
- Several options are in different stages of investigation
- The plan is to have detailed comparison by the end of the year and finalize the Run3 EDM design