



SUT TREE
2011KEE



ALICE
A JOURNEY OF DISCOVERY

Asian sites in ALICE – storage role in the context of Run 3

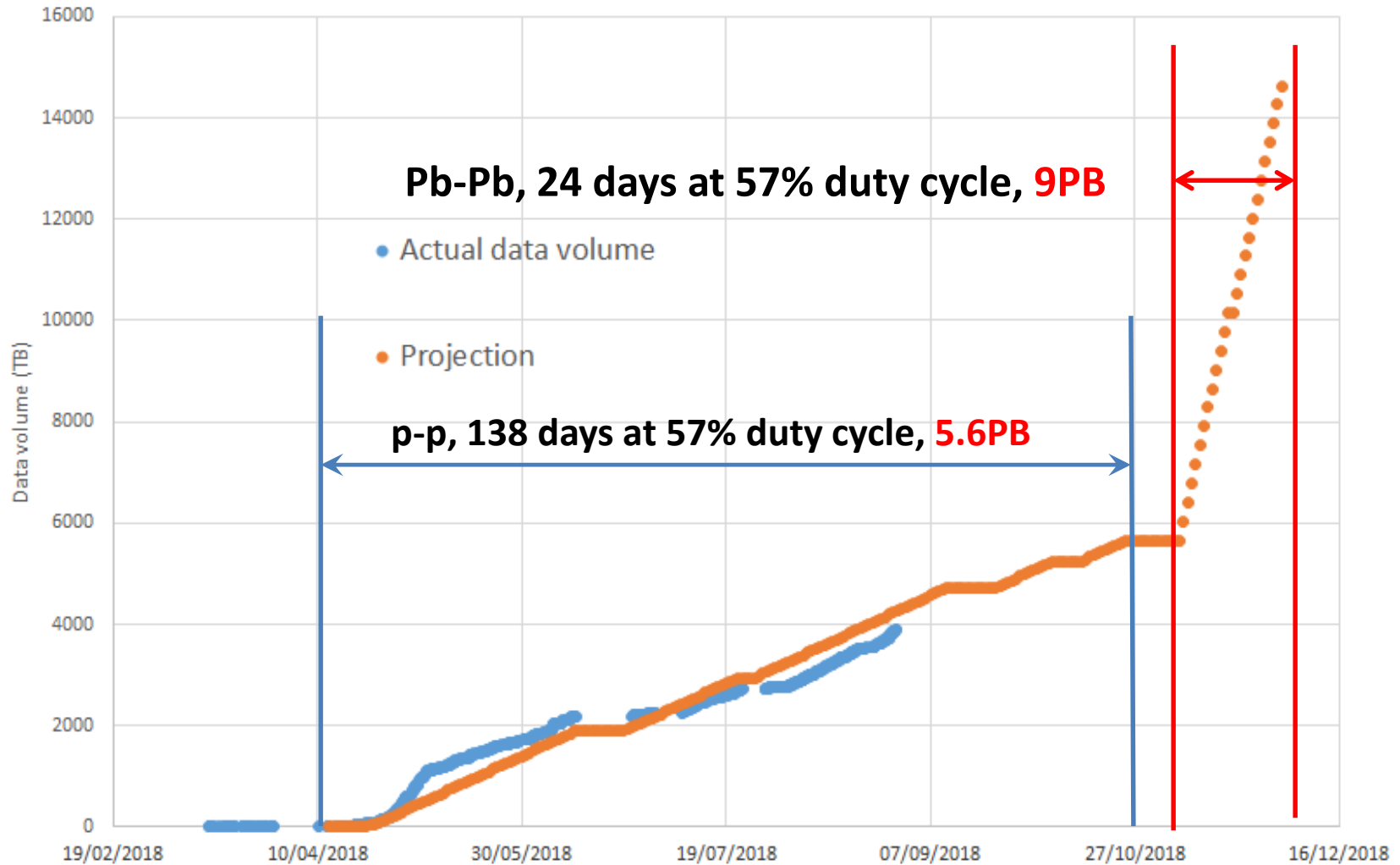
4-th Asia Tier Center Forum
Bangkok, Thailand

21 November 2018
Latchezar Betev

ALICE today

- Last data taking period of RUN2: Pb-Pb
- Expected to collect 9PB of data – 3x more than any other Pb-Pb period prior
- Possible due to a partial upgrade of the readout
- Will be followed up by data processing for ~4-5 months

Data taking projection



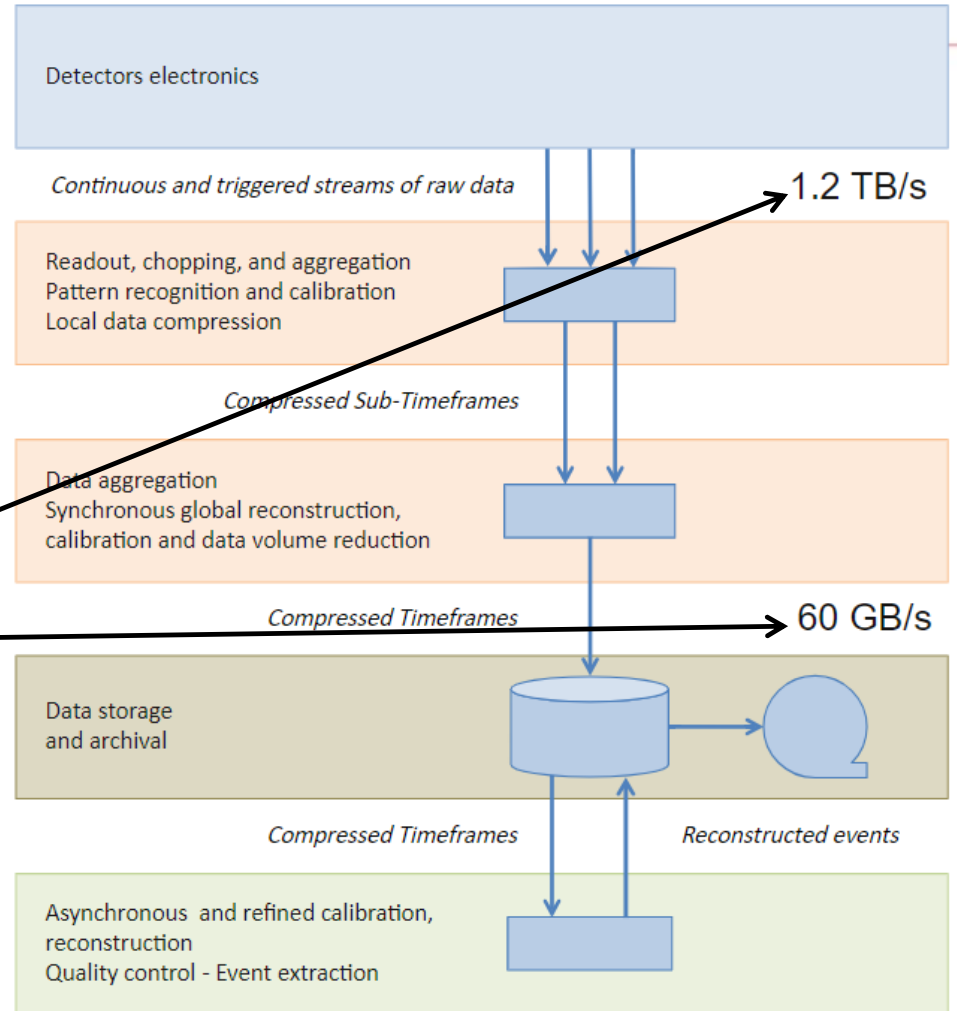
What follows?

- Intensive upgrade period – detector, readout and software
- Run3 – ALICE will run with upgraded everything and new physics programme
 - Focused on charm physics
 - Substantially higher (x50) data rates from detectors
 - Compression facility O2 to bring the data back into a manageable volume for offline storing/processing
 - A 20% growth in computing resources + the O2 facility will be enough to process the Run3+ data

O2 in a nutshell

- Custom-build cluster (CPU/GPU/FPGA) for efficient data compression

~20x



What changes in the computing model

- High reliance on the synchronous data compression – without it, the numbers will not match
- The role of the Grid computing remains largely familiar with some changes to
 - Additional analysis facility as a separate entity
 - More load on storage and redefined function (details to follow)

The upgrade challenge

- Order of magnitude more storage and files to manage
- Higher client-server (tasks) interaction rates
 - Estimated to be ~ 200 kHz (from about 20kHz today)
 - Affects the load on existing services – not expected to scale
- O2 facility addition to the system will bring new resources
 - Integration is a challenging problem
- Preserve the data model as much as possible
 - Storage provisioning and management is one of the toughest, most expensive and most time-consuming items on the Grid

The file catalogue

- All files on the Grid are annotated in the catalogue
 - Logical File Names (LFN) namespace with ACLs

`-rwxr-xr-x alidaq alidaq 264403565 Sep 09 22:10 /alice/data/2016/LHC16n/000261088/raw/16000261088034.205.root`

- Pointer to file location and GUIDs for SE filenames

`root://alice-tape-se.gridka.de:1094//10/33903/76cebd12-76a0-11e6-9717-0d38a10abeef`

`root://voalice10.cern.ch//castor/cern.ch/.../16000261088034.205.root`

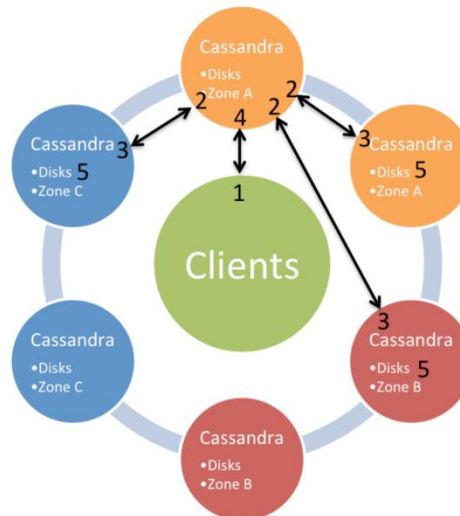
- 6 billion LFNs, 3.5 billion physical files (archives of multiple logically combined LFNs)
- Running on MySQL (master-slave) DB with cache
 - Somewhat outdated technology by today's standards
- For 150k running jobs – 20KHz read / 1KHz change/delete

The file catalogue (2)

- Will continue (together with the common task queue) to be the backbone of the ALICE distributed computing
- The basic functionality will remain unchanged
- The backend will be changed
 - To be able to cope with increased load

The new backend

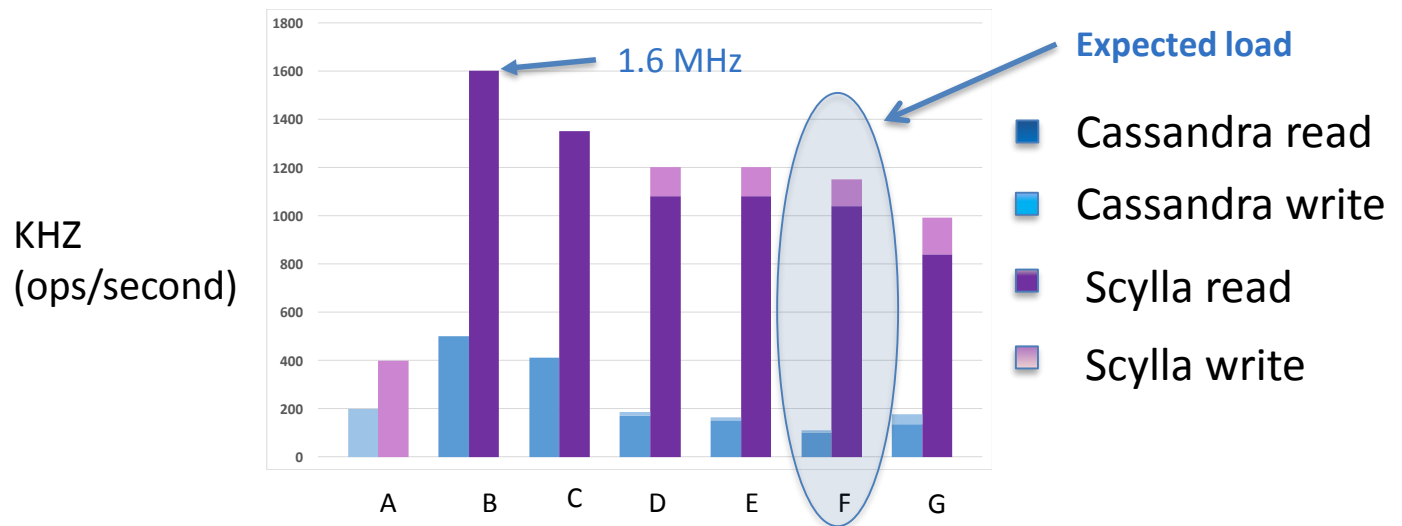
- Apple-ALICE collaborative work on Cassandra
- Answers all general and specific requirements:
 - Horizontal scaling, no single point of failure
 - High query rate, high availability
 - Consistency, easy setup
 - Drawback – SQL to NoSQL operations



Setup and benchmarking

- 6-node test ring, replication factor 3, ~cheap hardware
- Any node (+1) can go away without performance degradation

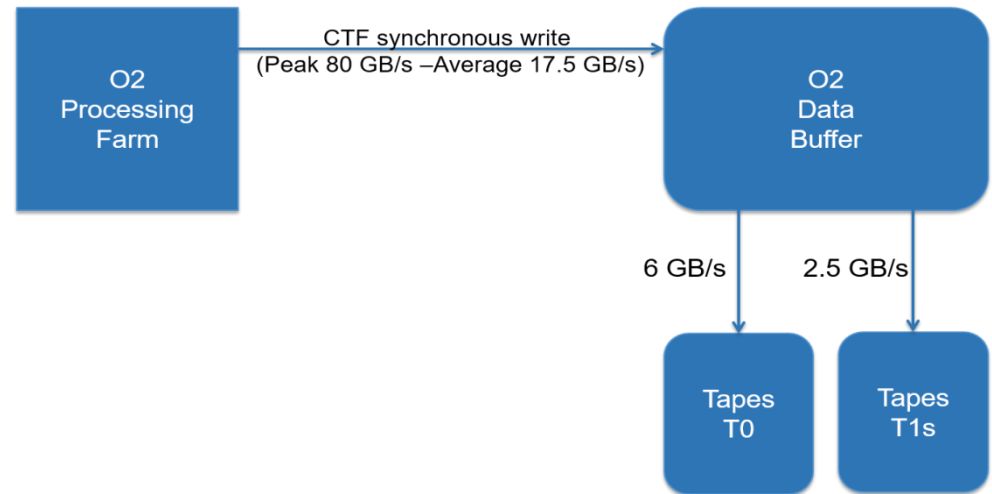
=> **Success!**



	Benchmark (default cassandra-stress)	Cassandra	Scylla	Diff
A	Insert only	18 %util, 2% iowait	100 %util	x2
B	Read-only Gauss(5B,2.5B,10K)	Disk idle, 50% cpu	100 %cpu	x3
C	Read-only Gauss(5B,2.5B,1M)	11 %util, 40% cpu	11 %util, 100 %cpu	x3.28
D	Mixed (10r,1w) Gauss(5B,2.5B,100K)	2 %util, 45% cpu	10 %util, 100 %cpu	x5.8
E	Mixed (10r,1w) Gauss(5B,2.5B,1M)	5 %util, 50% cpu	16 %util, 100 %cpu	x6.4
F	Mixed (10r,1w) Gauss(5B,2.5B,10M)	9% util, 45% cpu	40 %util, 100 %cpu	x6.1
G	Mixed 2K thrd. read, 200 write, G(5B,2.5B,100K)	8 %util, 40% cpu, no iowait	26 %util, 100 %cpu	x5.62

Synchronous processing data path

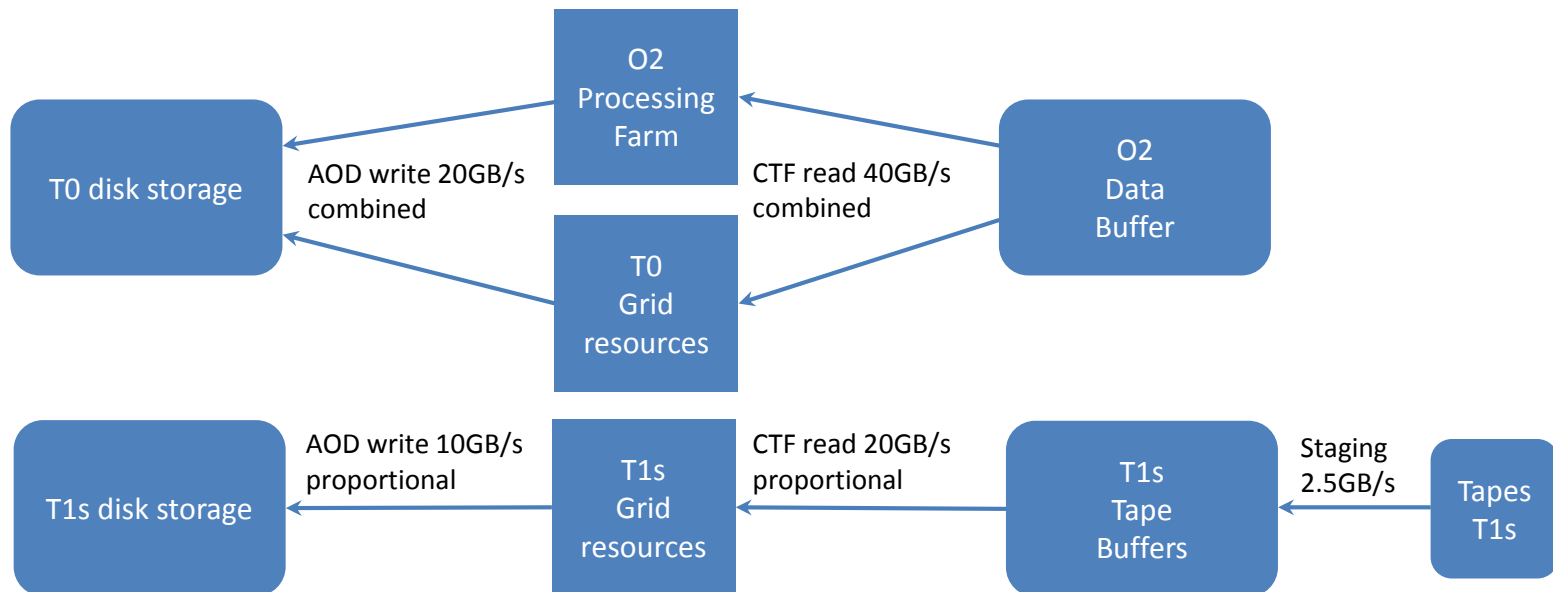
- Simplified schema during data taking periods



- Defining element for the O2 data buffer
 - Data rates from O2
 - Number of events => buffer size
 - Data rates to T0 tapes and T1s (export)

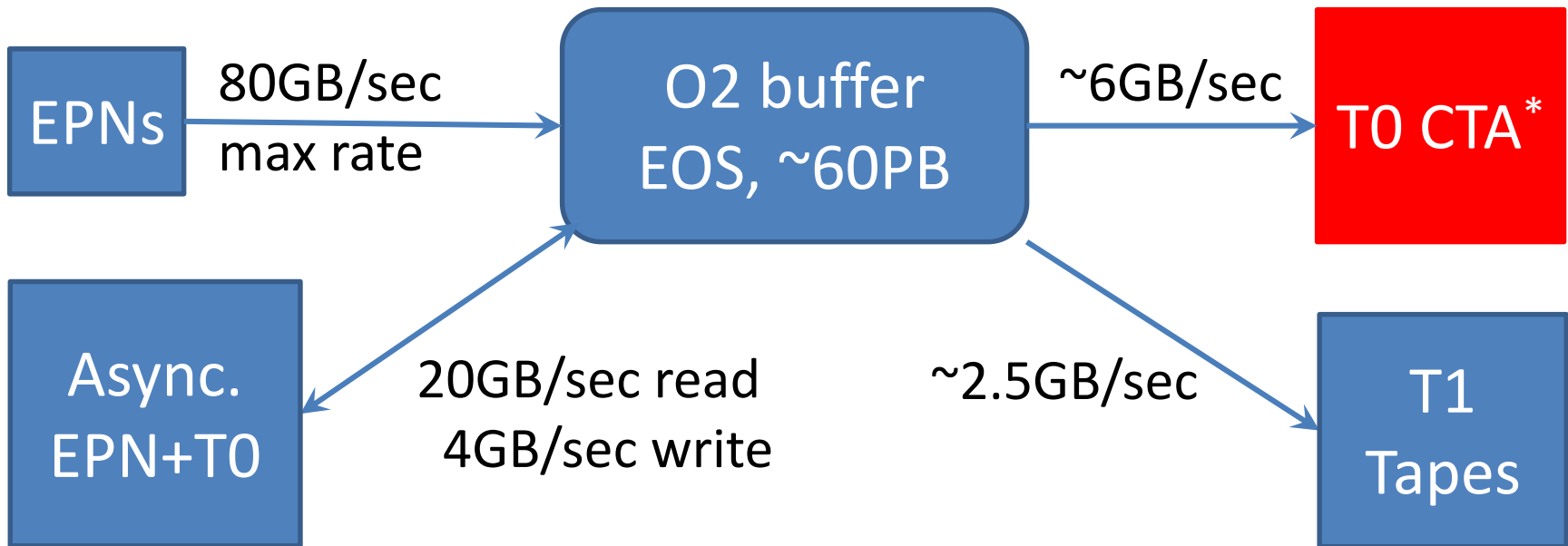
Asynchronous processing

- During beam off periods and after data taking at O2 facility
- All the time at T0/T1s
- Asynchronous processing is similar to synchronous, but with better detector calibration/software
- Output from asynchronous processing is AODs



Data rates

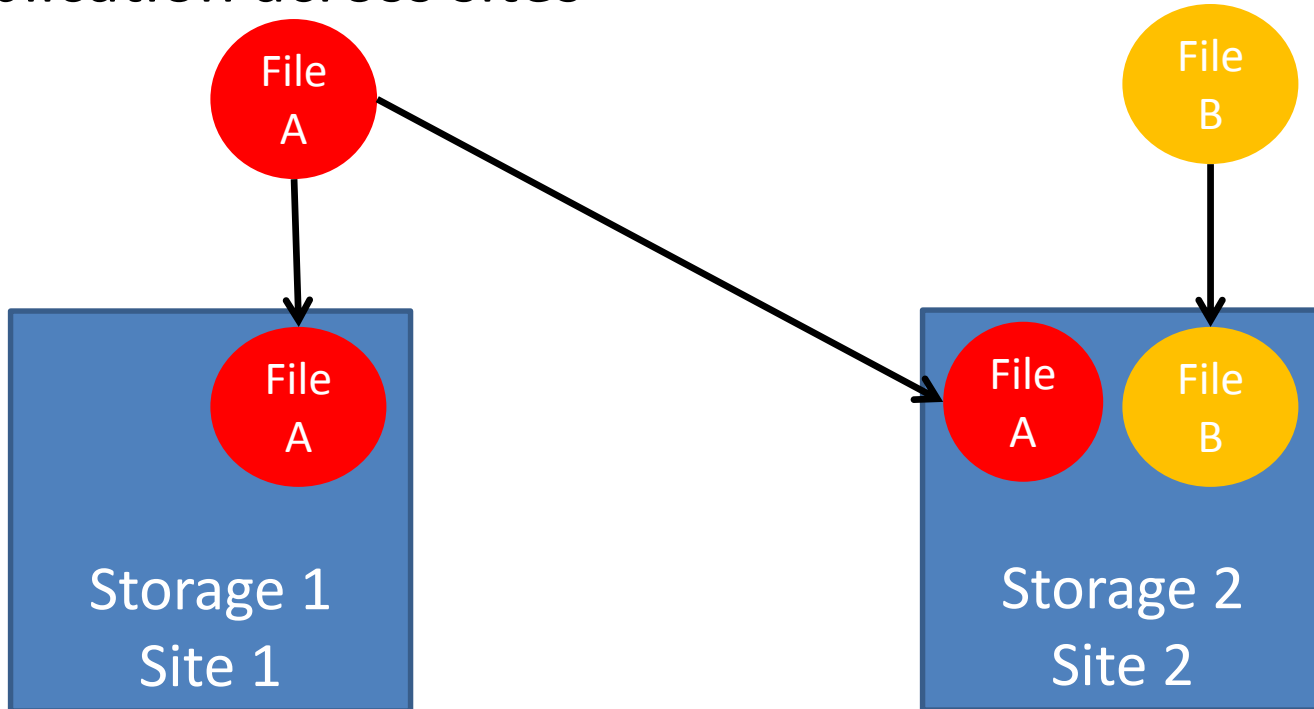
- Final (projected) data rates in Run 3



* CTA – CERN Tape Archive

Data handling on global level

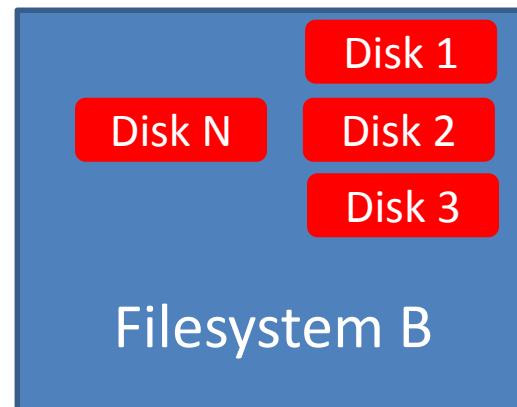
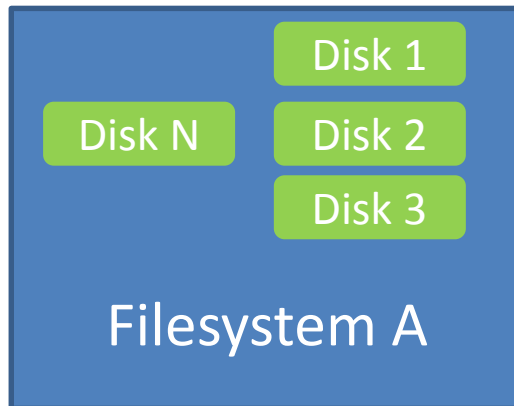
- Replication across sites



- Two (or more) copies – safe, but costly
- One copy of bulk of the data is now a norm
 - With the exception of RAW (2x), smaller analysis objects, frequently accessed and user files

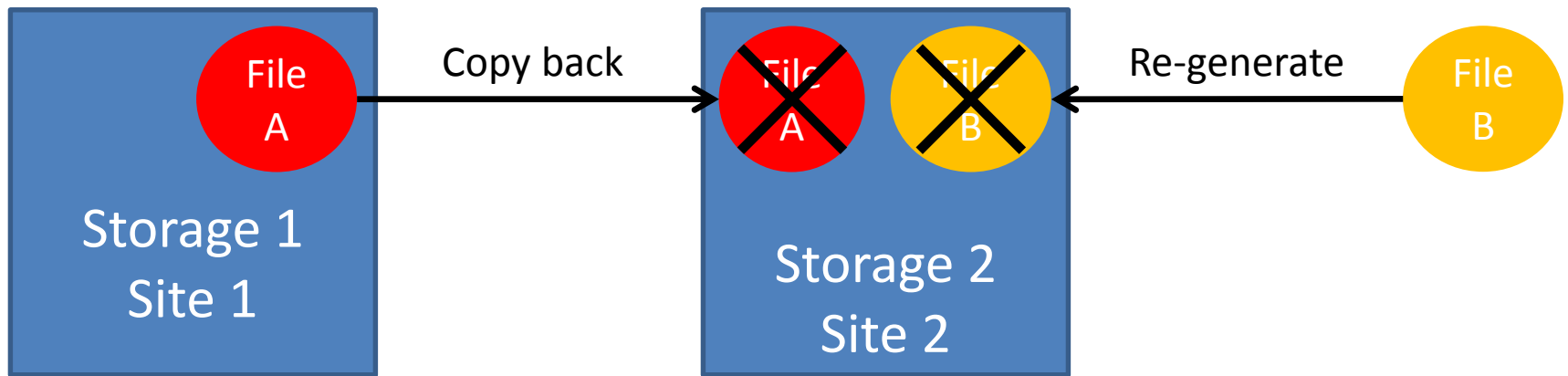
Data handling on site level

- Most common – file servers with some RAID level
 - Usually large (hundred TB) filesystems per server
 - Classical technology, supposedly safe
 - All redundancy within same filesystem



Data loss recovery on global level

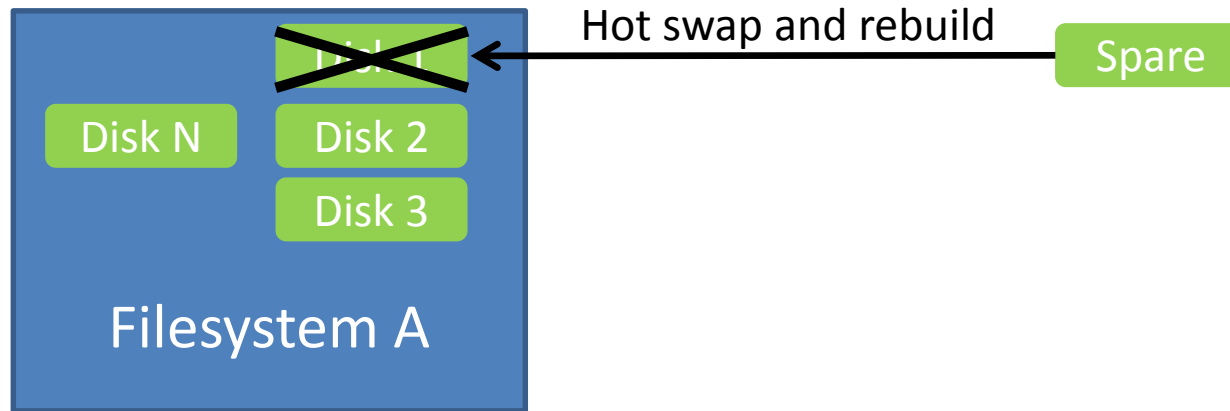
- Re-replication (or re-generation) of lost data



- Both methods external to site and used in case of major or locally unrecoverable data loss

Data loss recovery on site level

- RAID recovery



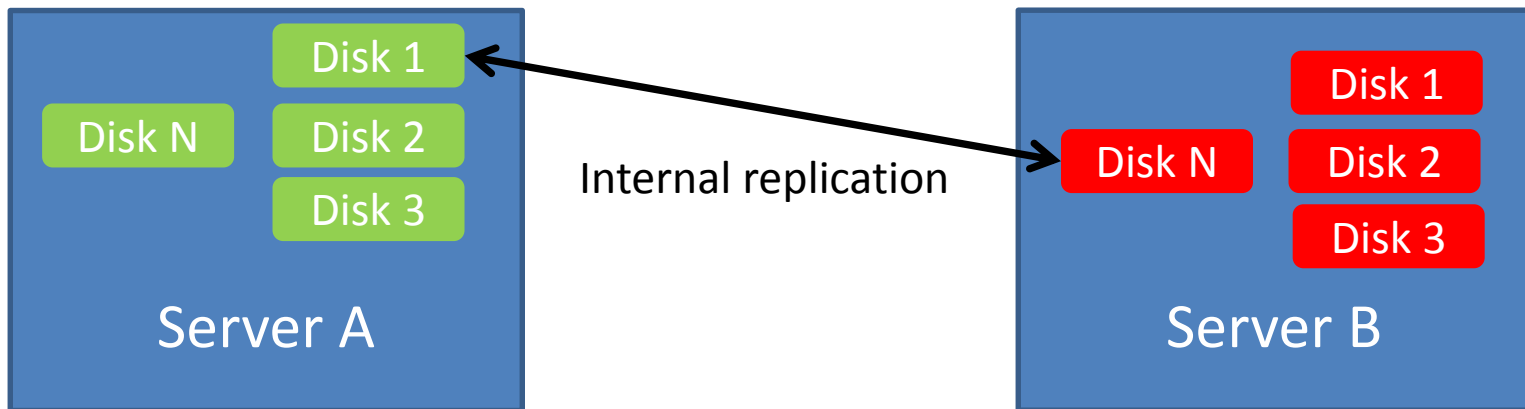
- Long (days!) process with large-volume disks
- Assumes broken disks are less than RAID level – may be worst with power cuts
- If RAID rebuild fails – full filesystem loss
 - Recovery mechanism on previous slide

Data loss recovery on site level (2)

- The recovery of RAID-ed filesystem is heavy and can potentially fail
 - Resulting in major data loss (100s TBs)
 - Source of embarrassment and long recovery procedure, out of site control
- RAIDs are typically expensive... and there is a better solution
 - Coincidentally, it simplifies the setup and operation

Data handling on site level

- (not) New system – RAIN
 - Data is stripped/replicated across physical nodes
 - RAID-x – like
 - Most efficient is RAIN-1, other levels require substantial network traffic between servers
 - RAIN-1 configuration is almost as safe as out of site copy



Summary on data handling

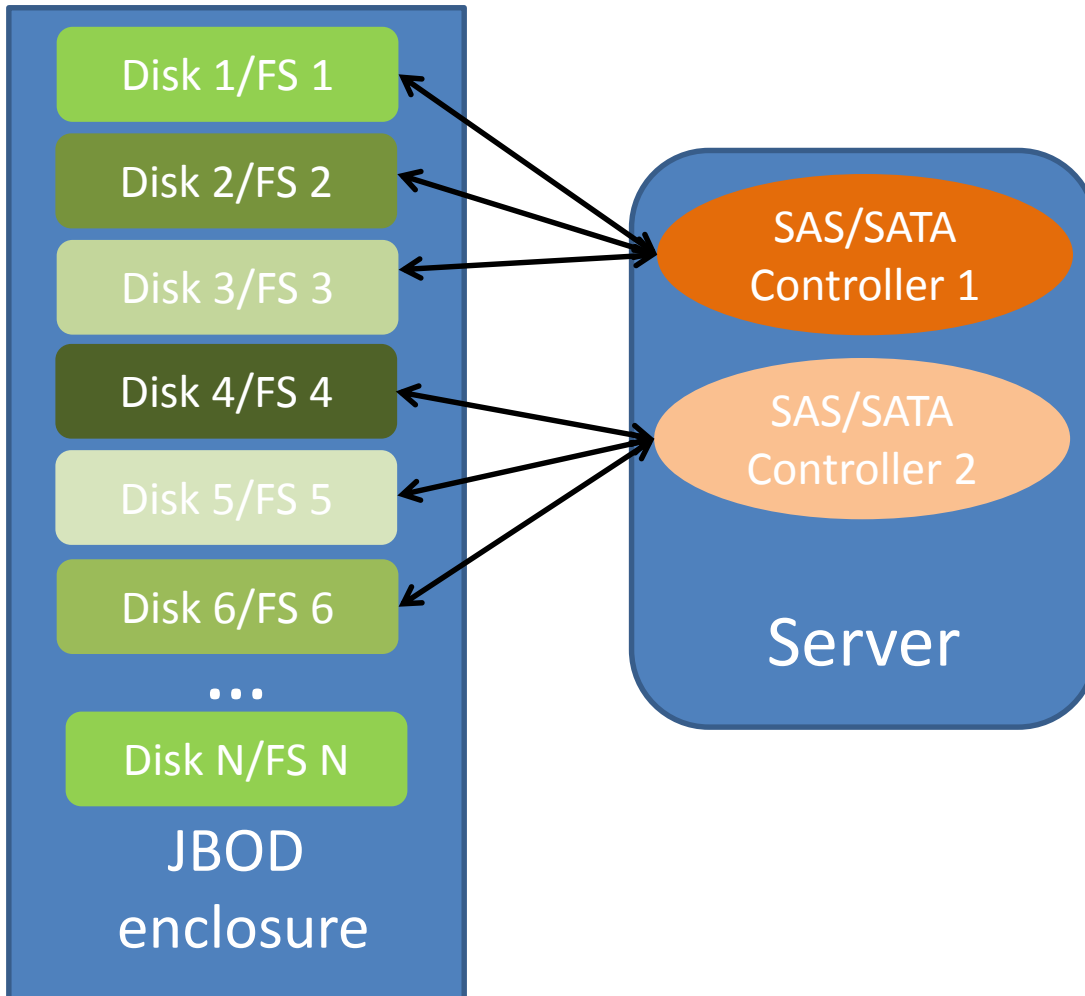
- Sites provide disk storage capacity, usually secured with RAID (think ‘I don’t want to lose data’)
 - In a typical (RAID 6) configuration, the capacity overhead is 15%, the cost of RAID controller is on top
 - Provides some security, but in case of failure there is a massive data loss
- Experiments mostly disregard site redundancy (think ‘Data loss is inevitable’)
 - Precious data is replicated
 - Data is spread around, single site storage failure results only in ‘statistics’ loss
 - Data loss is unpleasant, but tolerated and there are recovery mechanisms in place

Additional considerations - ALICE

- To stay within 'flat budget' on storage ALICE will
 - Have a single copy of RAW data on custodial storage
 - Give up one derived data format (ESD) and remain with a single data format (AOD) for all subsequent data analysis
- Even with the above - at the limit of capacity
 - Provided the growth will not be reduced by unforeseen circumstances

Minimize impact of data loss and tune performance

- Site storage structure



- Each disk is a filesystem
- Max loss of data is one disk
- Number of disk combined under a controller defines the throughput
- Basic configuration – no data redundancy within the storage element
- Software configuration – can define 2 or more copies within the same server for data redundancy

Technical solutions

- Setup of EOS storage at CERN – RAIN-1 for the bulk data
 - Even higher redundancy – directory based
- Setup of EOS storage at Kolkata – RAIN-6 for the bulk data
 - Works, but is not trivial to update
- Setup of EOS storage at ORNL – no redundancy for the bulk data
 - Will serve as a test for the future operation and policy
 - Cheapest storage – RAW=usefull

Technical solutions (2)

- The replication policy will remain the same
 - Most important aspect is the wide distribution of any bulk data – avoid concentration in the same storage
 - Extremely important to limit impact of unavailable (or lost) storage capacity
- Foreseen bulk data with 2 copies will be the AODs
 - Assures continuous availability and quick recovery in case of loss
 - Any other data will ‘suffer’ statistical reduction in case of storage unavailability, potentially can be regenerated

Operational considerations

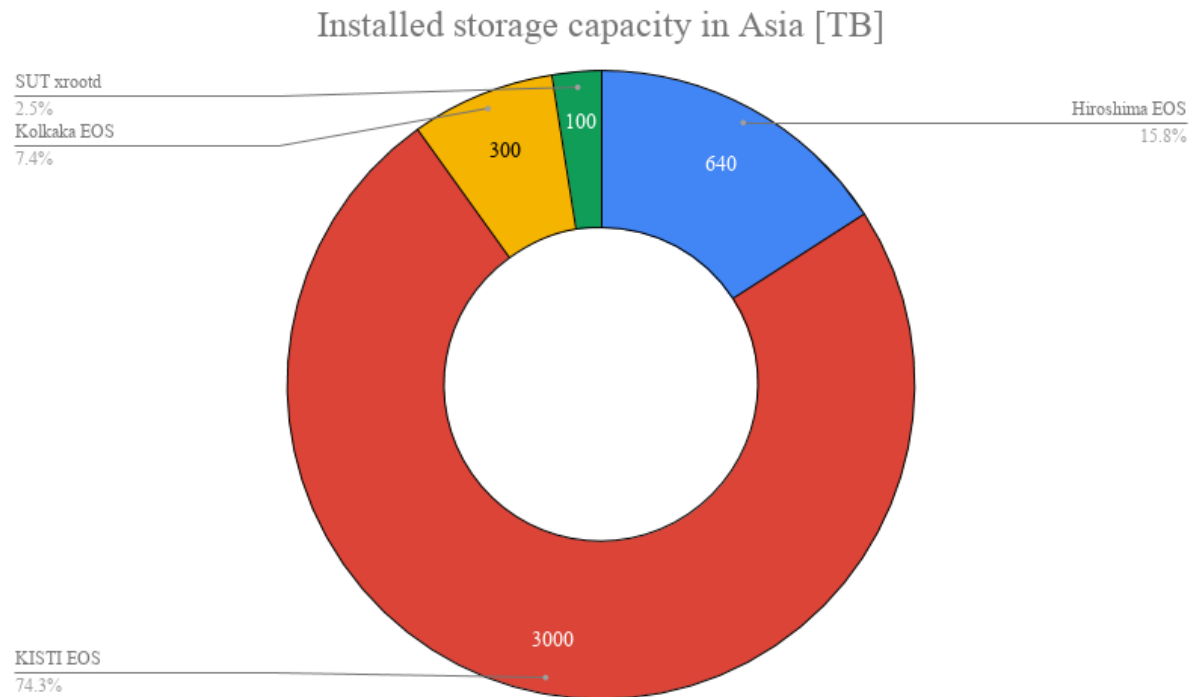
- Reduced local redundancy is taken into account by experiment replication policy
- Local data loss will be limited to single disk capacity
 - Site cannot recuperate data (understood and OK)
 - Should be reported to the experiment quickly to assure proper mitigation (at the discretion of the experiment)
 - Should simplify the local storage operation – replace faulty disk and continue

Operational considerations (2)

- The sites can determine storage performance, as needed by the function
 - Tuned for max throughput per server/network infrastructure – vital for analysis
 - ... or high capacity with low throughput – custodial storage/tape replacement (?)

Storage distribution across Asian sites

- Displayed current capacities (disk only)
- Expected to grow by 20% per year + other sites to join



Future cooperation

- Substantial storage capacity, expected to grow
- Possible areas of cooperation:
 - Selection of common hardware, synchronised purchase (economy of scale)
 - Installation procedure and sharing of experience, perhaps common experts pool
- Single entry point between storage elements with similar/same setup

Summary

- Substantial storage capacity, expected to grow
- Possible areas of cooperation:
 - Selection of common hardware, synchronised purchase (economy of scale)
 - Installation procedure and sharing of experience, perhaps common experts pool
- Single entry point between storage elements with similar/same setup

**Many thanks to the SUT team for
organizing the workshop**

**As with the previous installments – it provides
an opportunity for a focused discussion of a
relevant topic**

**Looking forward to your questions and
discussion**