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GPU-based Clustering Algorithm for the CMS High Granularity Calorimeter

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Luminosity Forecast

Year

[1] https://lhc-commissioning.web.cern.ch/lhc-commissioning/schedule/HL-LHC-plots.htm [2] CMS Collaboration, "Technical Proposal for the Phase-II Upgrade of the Compact Muon Solenoid", Technical Report CERN-LHCC-2015-010, LHCC-P-008, 2015.

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CMS High Granularity Calorimeter

- ❖ Current CMS endcap ECAL and HCAL calorimeters will be replaced (HGCAL), a sampling calorimeter system based on Si sensors and p upgrade.
	- Full system operates at -35° C maintained by a CO_2 cooling system
	- Covers $1.5 < |\eta| < 3.0$ on both left and right sides
	- Total size $z=2m$, $r = 2.3m$. Total weight 215 ton per endcap
	- 620 m² of Silicon sensors (120/200/300 um). 0.5-1.0 cm² \rightarrow 6M channels
	- 400 m² of plastic scintillators with SiPM readout. 4-30 cm² \rightarrow 240k channels.

Challenge of Computing \bullet CMS uses two-level trigger

- (HLT) , to reduce data rate output rate) \rightarrow 1 kHz (HLT
	- **L1 Trigger: based on ASICs and FRG**
	- HLT: based on CPUs. Make

\div HLT in era of HL-LHC expects

- 4x from increased event com
- 7.5x from increased event rate
- ❖ Within 30x increased computer account for only 4x. CPU a challenge.
- $\cdot \cdot$ It is particularly a huge cha the HLT time budget $($20$$
- ❖ GPU could be a solution. (API) for General-Purpose (GPGPU), makes it possible reconstruction with GPUs.

This estimating calculation just sets a scale of upper limit of HGCAL budget in HL⁻ $4(2)$ is factor for optimistic(realistic) increase of # CPU

Clustering in HGCAL

- v Hits of a PU200 event in HGCAL. Color and size represent hit energy. Interaction point is on the left side.
- $\cdot \cdot \cdot$ n~ O(100,000) hits
- \div HGCAL reconstruction starts by reconstructing 2D clusters layer-by-layer. k~O(10,000) clusters.
- ❖ Since cells are small compared to shower lateral size, an "energy density" is defined to better hint regional energy blobs.
- ❖ 3D showers are reconstructed by collecting and associating 2D layer clusters

Features of HGCAL clustering task $n > k \gg \frac{n}{l}$ \boldsymbol{k} in 2D Fast and GPU-friendly

Querying neighborhood N_d is one of the most frequent operations in density-based clustering algorithms. So need fast N_d query.

d-searchBox Ω_d

 $\Omega_d(i) = \{j : j \in \text{tiles touched by square window } (x_i \pm d, y_i \pm d)\}\$

d-neighborhood N_d

 $N_d(i) = \{j : d_{ij} < d\} \subset \Omega_d(i)$

To query N_d , we only need to loop of hits in Ω_d

❖ Build "Grid Spatial Index" for hits on each layer

- Grid tiles are small comparing with the size of HGCal layer
- Each tile in the grid hosts indices of hits inside it and has a fix length of memory to store the hosted indices.
- Points inside a tile can be directly accessed.
- Complexity of query d-neighborhood is $O(1)$, given that d is small.
- v Building spatial index is highly parallelizable on GPU.

v **Step 0: Build Spatial Index**

- **1 CUDA thread for each hit**
- **Register the index of each point to corresponding tile**

❖ Step 1: Calculate Local density

- 1 CUDA thread for each hit
- Density defined on the left

Ex Step 2: Calculate Nearest Higher

- 1 CUDA thread for each hit
- Define $d_m \equiv \max(\delta_s, \delta_o)$, where δ_s, δ_o are algorithm parameters for seed promotion and outlier demotion
- Within $N_{dm}(i)$, find the nearest points with higher density.
- Calculate $\delta_i = dist(i, nh_i)$

\div **Step 3: Promote Seeds and Demote Outliers**

- 1 CUDA thread for each hit
- Promote hit as seed if $\rho_i > \rho_c$, $\delta_i > \delta_s$
- Demote hit as outlier if $\rho_i < \rho_c, \delta_i > \delta_o$

❖ Step 4: Assign Cluster ID

- 1 CUDA thread for each seed
- Push down the cluster ID from seeds through reversed chains of

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CLUE in CMS Software (CMSSW) Framewor

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Average Execution Time of 2D Clustering of PU200 Events

- v CLUE CPU has about 30X speed up over previous clustering in CMSSW.
- ◆ CLUE GPU V3 gives an additional 6X on top of CLUE CPU in CMSSW framework.
- v GPU run time (32 ms) also includes memcpy between host and device (20 ms) and SoA conversion (6 ms). But they can be shared by other GPU reconstruction steps and can be partially hidden if multiple CUDA stream work on different events. 30X speed up over CPU if excluding memcpy and SoA conversion.

Conclusion

- $\cdot \cdot$ We introduce CLUE, an O(n) complexity and GPU-friendly clustering algorithm.
- v CLUE is ideal for heavy clustering task in CMS HGCAL reconstruction during HL-LHC.
- ◆ Thanks to CMS's plan of heterogeneous architecture in HLT and offline reconstruction, HGCal clustering can run on GPU and can provide promising acceleration.
- ◆ In CMS Software (CMSSW) framework, CLUE CPU is 30x faster (203 ms) than previous CPU clustering (6110 ms). CLUE on GPU gives another extra 6x speed up over CLUE on CPU (30x or 6 ms if excluding time of data traffic and SoA conversion).
	- penalty due to data traffic and SoA conversion will be shared with other GPU reconstruction steps.
	- these penalties can also be partially hidden if multiple CUDA streams works on different events.

Backup

CLUE Procedure

Energy Density in the CMS HGCAL Clustering

Definition of energy density

$$
\rho_i = \sum_{j:j \in N_{d_c}(i)} \chi(d_{ij}) w_j
$$

For CMS HGCAL, w_j is energy of hit and $d_c = 13$ mm, which equals to the distance between two adjacent cells. The density kernel used in HGCAL is

$$
\chi(d_{ij}) = \begin{cases} 1 \text{ if } i = j \\ 0.5 \text{ if } 0 < d_{ij} < d_c \end{cases}
$$

Comparing with a single cell of maximum local energy, local energy density is more sensitive to a multi-cell blob of energy.

