PCI Express 3.0 Host Interface

GigaThread Engine



L2 Cache



My research experience



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Reconstruction challenge in CBM at FAIR/GSI



- Future fixed-target heavy-ion experiment
- 10⁷ Au+Au collisions/sec
- ~ 1000 charged particles/collision
- Non-homogeneous magnetic field
- Double-sided strip detectors (85% fake space-points)

Full event reconstruction will be done on-line at the First-Level Event Selection (FLES) and off-line using the same FLES reconstruction package.

Cellular Automaton (CA) Track Finder Kalman Filter (KF) Track Fitter KF short-lived Particle Finder

All reconstruction algorithms are vectorized and parallelized.



Reconstruction challenge in CBM at FAIR/GSI





Many-Core CPU/GPU Architectures



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Multi-/Many-Core CPU systems

SMP (symmetric **multi**-processor) systems:

NUMA (non-uniform memory access) systems:

- Homogeneous
- "Equal-time" access for each processor to any part of the memory
- Heterogeneous
- Non uniform access to different parts of the main memory – different speed, data should be close to the CPU







Many-core system with > 8 cores

Stages of Event Reconstruction



- Hough Transformation
- Elastic Neural Net

(r, C)

K-

 π^+

Stages of Event Reconstruction



- Hough Transformation
- Elastic Neural Net

(r, C)

K-

 π^+

Cellular Automaton - Game "Life"



Each cell has 8 neighboring cells: 4 adjacent orthogonally, 4 adjacent diagonally. The rules are: Survival: Each living cell with 2 or 3 adjacent living cells survives for the next generation.

Death: Each living cell with 4 or more living neighbors dies of overpopulation, with 1 or none neighbor dies of isolation. Birth: Each empty cell adjacent to exactly 3 living neighbors is a birth cell.

All births and deaths (the transition to a new generation) occur for all cells simultaneously, i.e. in discrete steps.



Martin Gardner, Mathematical games: The fantastic combinations of John Conway's new solitaire game "Life", Sci. Amer., 223 (1970) 120-123

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1

Cellular Automaton (CA) Track Finder



Useful for complicated event topologies with heavy combinatorics

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Cellular Automaton (CA) Track Finder



Fast and efficient track finder

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A number of minimum bias events is gathered into a group (super-event), which is then treated by the CA track finder as a single event.



1 mbias event, <N_{reco}> = 109

5 mbias events, <N_{reco}> = 572

100 mbias events, $\langle N_{reco} \rangle = 10340$



Reliable reconstruction efficiency and time as a second order polynomial w.r.t. to the track multiplicity

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Hits at high input rates



Reconstructed tracks clearly represent groups, which correspond to the original events

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Stages of Event Reconstruction



- Hough Transformation
- Elastic Neural Net

(r, C)

K-

 π^+

2

Kalman Filter Algorithm

The Kalman filter is a recursive estimator – only the estimated state from the previous time step and the current measurement are needed to compute the estimate for the current state.



For this work, U.S. President Barack Obama rewarded Rudolf Kálmán with the National Medal of Science on October 7, 2009.





state vector:



covariance matrix:



Apollo Flight Journal

December 21, 1968. The Apollo 8 spacecraft has just been sent on its way to the Moon. **003:46:31 Collins:** Roger. At your convenience, would you please go P00 and Accept? We're going to update to your W-matrix.

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Kalman Filter (KF) based Track Fit

Estimation of the track parameters at one or more hits along the track – Kalman Filter (KF)



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CPU/GPU Programming Frameworks



 \sum

• Cg, OpenGL Shading Language, Direct X

- Designed to write shaders
- Require problem to be expressed graphically

AMD Brook

- Pure stream computing
- No hardware specific
- AMD CAL (Compute Abstraction Layer)
 - Generic usage of hardware on assembler level
- NVIDIA CUDA (Compute Unified Device Architecture)
 - Defines hardware platform
 - Generic programming
 - Extension to the C language
 - Explicit memory management
 - Programming on thread level

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• Headers and Vector classes (Vc)

- Overload of C operators with SIMD/SIMT instructions
 - Uniform approach to all CPU/GPU families
 - Uni-Hedeilberg/Uni-Frankfurt/FIAS/GSI

• Intel Ct (C for throughput)

- Extension to the C language
 - Intel CPU/GPU specific
 - SIMD exploitation for automatic parallelism
- OpenCL (Open Computing Language)
 - Open standard for generic programming
 - Extension to the C language
 - Supposed to work on any hardware
 - Usage of specific hardware capabilities by extensions

Cell Processor: Supercomputer on a Chip



External Interconnects:

2

•25.6 GB/sec BW memory interface

- •2 Configurable I/O Interfaces
 - Coherent interface (SMP)
 - Normal I/O interface (I/O & Graphics)
 - •Total BW configurable between interfaces
 - Up to 35 GB/s out
 - •Up to 25 GB/s in



- Sony PlayStation-3 -> cheap
- 32 (8x4) times faster !

Memory Management & Mapping

SPE Local Store aliased into PPE system memory
 MFC/MMU controls SPE DMA accesses

 Compatible with PowerPC Virtual Memory architecture

•S/W controllable from PPE MMIO

- ·Hardware or Software TLB management
- SPE DMA access protected by MFC/MMU

Kalman Filter for Track Fit



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2

"Local" Approximation of the Magnetic Field

Problem:

- Full reconstruction must work within
 256 kB of the Local Store.
- The magnetic field map is too large for that (**70 MB**).
- A position (x,y), to which the track is Stati propagated, is unknown in advance.
- Therefore, access to the magnetic field map is a blocking process.



Solution:

- Use a polynomial approximation (4-th order) of the field in XY planes of the stations.
- 2. Assuming a parabolic behavior of the field between stations calculate the magnetic field **along the track** based on 3 consecutive measurements.



Difference





- The initial track parameters are directly estimated from the input data.
- The propagation step is performed directly from measurement to measurement without intermediate steps.
- Matrix multiplications have been replaced by direct operations on only non-trivial matrix elements.
- Most loops have been unrolled in order to provide additional instructions for interleaving.
- All branches have been eliminated from the algorithm to avoid branch misprediction penalty.
- Calculations have been reordered for better use of the processors pipeline.

Approach:

- 1. Universality (any multi-core architecture)
- 2. Vectorization (SIMDization)
- 3. Run SPEs independently (one collision per SPE)



```
typedef F32vec4 Fvec_t;
                                                                                     SIMD instructions
/* Arithmetic Operators */
friend F32vec4 operator +(const F32vec4 &a, const F32vec4 &b) { return _mm_add_ps(a,b); }
friend F32vec4 operator –(const F32vec4 &a, const F32vec4 &b) { return _mm_sub_ps(a,b); }
friend F32vec4 operator *(const F32vec4 &a, const F32vec4 &b) { return mm mul ps(a,b); }
friend F32vec4 operator /(const F32vec4 &a, const F32vec4 &b) { return mm div ps(a,b); }
/* Functions */
friend F32vec4 min( const F32vec4 &a, const F32vec4 &b){ return mm min ps(a, b); }
friend F32vec4 max( const F32vec4 &a, const F32vec4 &b){ return _mm_max_ps(a, b); }
/* Square Root */
friend F32vec4 sqrt ( const F32vec4 &a ){ return _mm_sqrt_ps (a); }
/* Absolute value */
friend F32vec4 fabs( const F32vec4 &a){ return _mm_and_ps(a, _f32vec4_abs_mask); }
/* Logical */
friend F32vec4 operator&( const F32vec4 &a, const F32vec4 &b ){ // mask returned
  return _mm_and_ps(a, b);
}
friend F32vec4 operator ( const F32vec4 &a, const F32vec4 &b ){ // mask returned
  return _mm_or_ps(a, b);
}
friend F32vec4 operator^( const F32vec4 &a, const F32vec4 &b ){ // mask returned
  return _mm_xor_ps(a, b);
}
friend F32vec4 operator!( const F32vec4 &a ){ // mask returned
  return _mm_xor_ps(a, _f32vec4_true);
}
friend F32vec4 operator || ( const F32vec4 &a, const F32vec4 &b ){ // mask returned
  return _mm_or_ps(a, b);
}
/* Comparison */
friend F32vec4 operator<( const F32vec4 &a, const F32vec4 &b ){ // mask returned
  return mm cmplt ps(a, b);
```

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2

Source Code (Part of the Kalman Filter Track Fit)

```
inline void AddMaterial( TrackV & track, Station & st, Fvec t & qp0 )
ł
 cnst mass2 = 0.1396*0.1396;
 Fvec t tx = track.T[2];
 Fvec t ty = track.T[3];
 Fvec t txtx = tx*tx;
 Fvec t tyty = ty*ty;
 Fvec t txtx1 = txtx + ONE;
 Fvec t h = txtx + tyty;
 Fvec t t = sqrt(txtx1 + tyty);
 Fvec t h2 = h^*h;
 Fvec t qp0t = qp0*t;
```



vc = vec_add(va, vb)

cnst c1=0.0136, c2=c1*0.038, c3=c2*0.5, c4=-c3/2.0, c5=c3/3.0, c6=-c3/4.0;

```
Fvec t s0 = (c1+c2*st.logRadThick + c3*h + h2*(c4 + c5*h + c6*h2))*qp0t;
```

Fvec t a = (ONE+mass2*qp0*qp0t)*st.RadThick*s0*s0;

CovV &C = track.C;

```
C.C22 += txtx1*a;
 C.C32 += tx*ty*a; C.C33 += (ONE+tyty)*a;
}
```

Use headers to overload +, -, *, / operators --> the source code is unchanged !

Kalman Filter Track Fit on Cell

	mysim/SPE4: Statistics	
SPU DD3.0		
Total Cycle count Total Instruction count Total CPI ***	335660 643 522.02	72.2%
Performance Cycle count Performance Instruction o Performance CPI	ount 7076 (6638) 1.03 (1.07)	
Branch instructions Branch taken Branch not taken	26 16 10	
Hint instructions Hint hit	7 10	
Contention at LS between	Load/Store and Prefetch 405	
Single cycle Dual cycle Nop cycle Stall due to branch miss Stall due to prefetch mis Stall due to dependency Stall due to fp resource Stall due to waiting for Stall due to dp pipeline Channel stall cycle SPU Initialization cycle Total cycle	s conflict hint target	4440 (62.7%) 1099 (15.5%) 16 (0.2%) 137 (1.9%) 0 (0.0%) 1365 (19.3%) 1 (0.0%) 18 (0.3%) 0 (0.0%) 0 (0.0%) 0 (0.0%) 7076 (100.0%)
stall cycles due to depen FX2 36 (2.6% of SHUF 92 (6.7% of FX3 0 (0.0% of LS 285 (20.9% of BR 0 (0.0% of SPR 0 (0.0% of LNOP 0 (0.0% of FXB 0 (0.0% of FXB 0 (0.0% of FYA 0 (0.0% of FP6 873 (64.0% of FP7 79 (5.8% of FPD 0 (0.0% of	dency on each pipelines all dependency stalls) all dependency stalls) f all dependency stalls) f all dependency stalls) all dependency stalls) all dependency stalls) all dependency stalls) all dependency stalls) f all dependency stalls) f all dependency stalls) all dependency stalls) all dependency stalls) all dependency stalls) all dependency stalls) all dependency stalls) all dependency stalls)	is 100.00
dumped pipeline stats		

Timing profile !

No need to check the assembler code !

	Stage	Description	Time/track	Speedup
4		Initial scalar version	12 ms	
	1	Approximation of the magnetic field	$240 \ \mu s$	50
Ĕ	2	Optimization of the algorithm	$7.2 \ \mu s$	35
- L	3	Vectorization	$1.6~\mu { m s}$	4.5
= ∫	4	Porting to SPE	$1.1~\mu { m s}$	1.5
ຶ ງ	5	Parallelization on 16 SPEs	$0.1~\mu{ m s}$	10
		Final simulized version	$0.1 \ \mu s$	120000



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	Stage	Description	Time/track	Speedup	
C		Initial scalar version	12 ms	_	
Intel	1	Approximation of the magnetic field	$240~\mu{\rm s}$	50	10000- 61
	2	Optimization of the algorithm	$7.2~\mu{ m s}$	35 >	on any PC
	3	Vectorization	$1.6~\mu { m s}$	4.5 J	
	4	Porting to SPE	$1.1~\mu{ m s}$	1.5	
ຶ ໂ	5	Parallelization on 16 SPEs	$0.1~\mu{ m s}$	10	
		Final simulized version	$0.1~\mu{ m s}$	120000	

Comp. Phys. Comm. 178 (2008) 374-383



blade11bc4 @IBM, Böblingen: 2 Cell Broadband Engines, 256 kB LS, 2.4 GHz

The KF speed was increased by 5 orders of magnitude



Motivated by, but not restricted to Cell !



Kalman Filter Track Fit Library

Kalman Filter Methods

Kalman Filter Tools:

- KF Track Fitter
- KF Track Smoother
- Deterministic Annealing Filter

Kalman Filter Approaches: • Conventional DP KF

- Conventional SP KF
- Square-Root SP KF
- UD-Filter SP
- Gaussian Sum Filter
- 3D (*x*,*y*,*z*) and 4D (*x*,*y*,*z*,*t*) KF

Track Propagation:

- Runge-Kutta
- Analytic Formula

Detector Types:

- Pixel
- Strip
- Tube
- TPC

Implementations

Vectorization (SIMD):

- Header Files
- Vc Vector Classes
- ArBB Array Building Blocks
- OpenCL

Parallelization (many-cores):

- Open MP
- ITBB
- ArBB
- OpenCL

Precision:

- single precision SP
- double precision DP









Strong many-core scalability of the Kalman filter library

with I. Kulakov, H. Pabst* and M. Zyzak (*Intel)

Prof. Dr. Ivan Kisel, Uni-Frankfurt, FIAS, GSI

1st Real Time Analysis Workshop, Institut Pascal, Université Paris-Saclay, 18.07.2019 29/48

2

Kalman Filter (KF) Track Fit



- Precise estimation of the parameters of particle trajectories is the core of the reconstruction procedure.
- Scalability with respect to the number of logical cores in a CPU is one of the most important parameters of the algorithm.
- The scalability on the Intel Xeon Phi coprocessor is similar to the CPU, but running four threads per core instead of two.
- In case of the graphics cards the set of tasks is divided into working groups of size *local item size* and distributed among compute units (or streaming multiprocessors) and the load of each compute unit is of the particular importance.
- The track fit performance on a single node: **2*CPU+2*GPU = 10⁹ tracks/s** = (100 tracks/event)* 10⁷ events/s = 10⁷ events/s.
- A single compute node is enough to estimate parameters of all particles produced at the maximum 10⁷ interaction rate!

The fastest implementation of the Kalman filter in the world

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Kalman Filter Track Fit on Cell



Available online at www.sciencedirect.com



Computer Physics Communications 178 (2008) 374-383

Computer Physics Communications

www.elsevier.com/locate/cpc

Fast SIMDized Kalman filter based track fit

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Abstract

Modern high energy physics experiments have to process terabytes of input data produced in particle collisions. The core of many data reconstruction algorithms in high energy physics is the Kalman filter. Therefore, the speed of Kalman filter based algorithms is of crucial importance in on-line data processing. This is especially true for the combinatorial track finding stage where the Kalman filter based track fit is used very intensively. Therefore, developing fast reconstruction algorithms, which use maximum available power of processors, is important, in particular for the initial selection of events which carry signals of interesting physics.

One of such powerful feature supported by almost all up-to-date PC processors is a SIMD instruction set, which allows packing several data items in one register and to operate on all of them, thus achieving more operations per clock cycle. The novel Cell processor extends the parallelization further by combining a general-purpose PowerPC processor core with eight streamlined coprocessing elements which greatly accelerate vector processing applications.

In the investigation described here, after a significant memory optimization and a comprehensive numerical analysis, the Kalman filter based track fitting algorithm of the CBM experiment has been vectorized using inline operator overloading. Thus the algorithm continues to be flexible with respect to any CPU family used for data reconstruction.

Because of all these changes the SIMDized Kalman filter based track fitting algorithm takes 1 µs per track that is 10000 times faster than the initial version. Porting the algorithm to a Cell Blade computer gives another factor of 10 of the speedup.

Finally, we compare performance of the tracking algorithm running on three different CPU architectures: Intel Xeon, AMD Opteron and Cell Broadband Engine.

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PACS: 02.60.Pn; 02.70.-c; 07.05.-t; 07.05.Bx; 07.05.Kf

Keywords: High energy physics; CBM experiment; Data reconstruction; Track fit; Kalman filter; SIMD instruction set; Cell Broadband Engine

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HPC Practical Course at the Goethe University Frankfurt

http://fias.uni-frankfurt.de/de/cs/kisel/lectures/





Stages of Event Reconstruction



- Hough Transformation
- Elastic Neural Net

(r, C)

K-

 π^+



3

Elastic Neural Net (EN) for TSP



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3







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Stages of Event Reconstruction



- Hough Transformation
- Elastic Neural Net

(r, C)

K-

 π^+

KF Particle Reconstruction of short-lived Particles

3 KFParticle: Reconstruction of Vertices and Decayed Particles $r = \{x, y, z, p_x, p_y, p_z, E\}$



 $\overline{\Omega}^+ \rightarrow \overline{\Lambda} \ \mathrm{K}^+$ $\downarrow \overline{p} \pi^+$

KFParticle Lambda(P, Pi);	// c
Lambda.SetMassConstraint(1.1157);	// iı
KFParticle Omega(K, Lambda);	// c
PV -= (P; Pi; K);	// c
PV += Omega;	// a
Omega.SetProductionVertex(PV);	// C
(K; Lambda).SetProductionVertex(Omega);	// K
(P; Pi).SetProductionVertex(Lambda);	// p

/ construct anti Lambda
// improve momentum and mass
/ construct anti Omega
// clean the primary vertex
/ add Omega to the primary vertex
/ Omega is fully fitted
/ K, Lambda are fully fitted
/ p, pi are fully fitted



Concept:

- Mother and daughter particles have the same state vector and are treated in the same way
- · Reconstruction of decay chains
- Kalman filter based the state vector
- · Geometry independent
- Vectorized
- Uncomplicated usage

Functionality:	the KF Particle
 Construction of short-lived particles 	
 Addition and subtraction of particles 	t and vectorised
Transport	
 Calculation of an angle between particles 	ited in the same
 Calculation of distances and deviations 	
Constraints on mass, production point and decay length	3 two reconstruct
KF Particle Finder	

Jent and can be ALICE, STAR).

nformation about

KFParticle provides uncomplicated approach to physics analysis (used in CBM, ALICE and STAR)

V. Akishina, I. Kisel, Uni-Frankfurt, FIAS

MMCP 2017, Dubna, 07.07.2017 11/16

KF Particle provides a simple and very efficient approach to physics analysis

KF Particle Finder for short-lived particles



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4

KF Particle Finder for short-lived particles





Very Clean Probes of Collision Stages



AuAu, 10 AGeV, 3.5M central UrQMD events, MC PID

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Running FLES on HPC Node/Farm







The FLES package is vectorized, parallelized, portable and scalable up to 3 200 CPU cores

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Consolidate Efforts: A Common Reconstruction Package



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