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Xrootd – an outlook

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Outline



- A big question
- What is it, how it evolved
- How it should be used
- How it was used
- What's missing (less)
- Support
- Open items





A big question



WHY

- Is it so special?
- Did it survive ?
- Should I use it?
- Will (not) it solve any problem ?





At the beginning



- Filling a big hole in the market
- Absence of a SCALABLE and ROBUST solution for HEP HPC
- Need to interconnect it with other HEP-related systems
- Need to deploy transparent Hierarchical Storage facilities
- Building an efficient LOCAL storage cluster was nearly impossible
 - Always the same use cases: reco, user analysis, custodial storage etc.
 - HEP has PARTICULAR requirements
 - Still partially covered by modern DFSs + tons of glue code







The not so basic requirements



- Data access as immediate as possible
 - Especially if the files are already where expected
- "Basic primitives" Open, rw, close, stat, prestage
 - Born for analysis-like random access and Hierarchical Storages
- "Easy" Cluster-to-cluster data movements: just a robust cp tool
- Deep-knowledged integration with ROOT
 - Fast file open, efficient random workloads, tough "intelligent" readaheads in multiuser systems... if used!
- Fast filename translations (= global namespace coherency)
 - Shield apps from the local mountpoints of sites
 - Also may create a coherent namespace ACROSS sites
- Quality
 - Always headed to the modern HEP reqs, although very generic
 - Efficient use of the available hw resources
 - Self-contained (avoids messy dependencies)
 - Super-robust communication, fault tolerance







What came next



- Robust communication also means that it worked via WAN, somehow
- The idea was to make it work WELL via WAN
 - Big effort started in ~2005 (async primitives and caching, multistreaming)
 - Exploited the ROOT TTreeCache as it came out (ReadV, then async reads, now async ReadV)
- Natural idea: aggregating WAN-wide storages (~ 2008)
- Next idea: WAN federations of sites helping each other
- Bigger outcome: possibility to simplify the glue code + get more quality

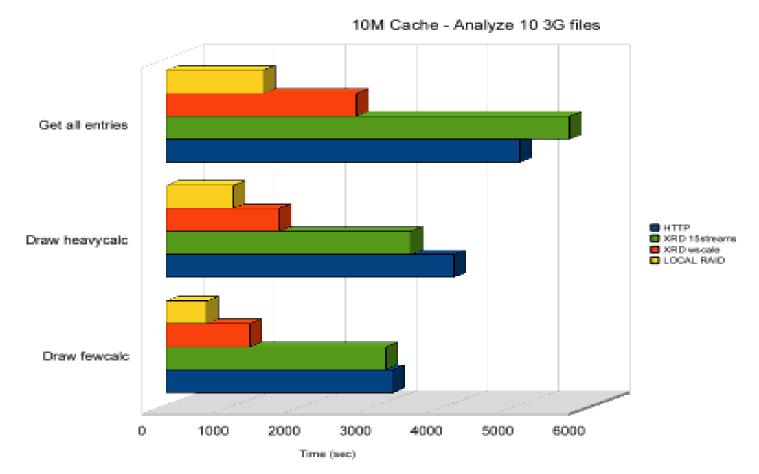




From CHEP09



CALTech 10GB/180msRTT + TTreeCache+Xrd latency hiding More than one year ago, now things should be even better









What does it do



- Hierarchical storage aggregation and access in their strictest meaning
- Plugin-based structure to override the behavior
- From tiny clusters to WAN-wide monster ones with linear scalability
 - They all look the same
- P2P-like internal file discovery, no internal catalogue
 - The files you put are the files you get
 - "You must know the files you've put and their meta-data" ™
 - No danger of misalignment, by definition
- Subset of the POSIX semantics, to support scalability
 - This subset was recently extended (Data Management Client-side functions, Jan 2010)
 - Allows the creation of replicas (replica-agnostic)
- Relies on external systems to do the right things
 - Like accessing the files which are present
 - Bad idea: doing 'find' on a PB repository
 - Another: 'ls' flooding on a PB repository to decide what to do







POSIX vs XROOTD



What becomes possible and what does not

Full File system Semantics Powerful async features
"No-action" server side fault tolerance
Lighweight, speed and scalability
Improved client side fault tolerance

Basic ops:

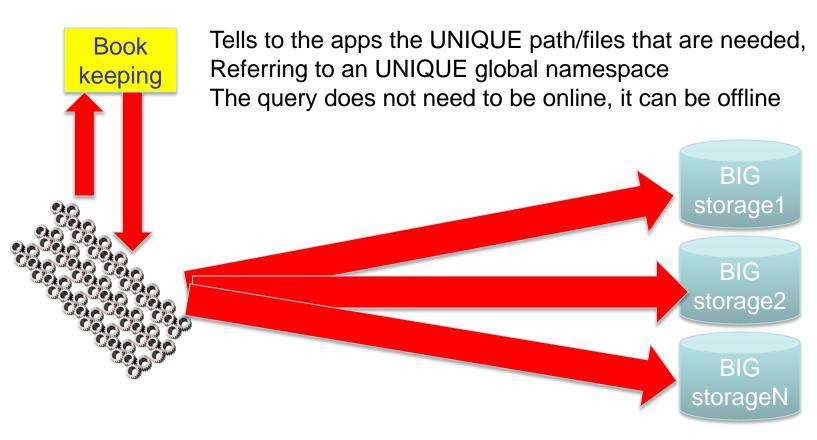
Open, close, read, write





The "expected" architecture





Many variations are possible, sometimes they have a high price.

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The "expected" architecture



- Design for scalability, this does not come for free.
- Enforces:
 - Separation data/metadata(app-specific)
 - Abstraction from the data location, in the local cluster or ev. In the global domain
 - performance
 - Once a dataset metadata (for an analysis) is ready, no need for the user to requery it
 - Access the storage only when the job has settled down, gives time to the DB part (ev. Quasi-offline)

Hides:

- SE type: tapes/no tapes etc. all look the same
- They only have different characteristics
 - Fast, slow, far, close, big, small
- Tapes can be used transparently. The difference is in knowing how to do it.
- Allows direct access to several SEs at once, via LAN and WAN







Core software process



- HEP-adapted AGILE process
- Small and committed core team
- The basic requirements are well known
 - Leave abundant space for extensions
- Get user's feedbacks as FAST as possible
- Respond as FAST as possible
 - The next beta is now, tagging takes 10 minutes
 - If you don't try it out it means that you are not interested
 - The next stable is "very soon", when the tests succeed
 - But, by definition, the subsequent tag will be better
- Filter requirements and requests
 - Virtuous ones get immediate priority







Security: Authentication



- XrdSec by Gerri Ganis (~2005): one of the best pieces of software around
- Authentication support for virtually any protocol
- The protocols were added one at a time, often as contributions
 - UNIX
 - Simple Shared Secret
 - GSI
 - Kerberos
 - VOMS extensions







Security: Authorization



- Simple plugin-based infrastructure, called XrdAcc
- At least 2 plugins around
 - The default one: simple but effective (and static)
 - AliceTokenAcc: the very efficient ALICE schema





Distributions



- The 'vanilla' one, from the SLAC CVS
 - Requires considerable knowledge and top quality support
 - Full access to any kind of weird configuration/platform
- The 'xrd-installer', from the good old ARDA repository
 - Stays to the 'vanilla' one like Ubuntu stays to the Linux kernel
 - Supports multiple 'bundle flavors'
 - Refurbished several times, very simple and stable. Supports data globalization, Virtual MSS and ML monitoring.
 - Pre-configured, reduces support requests to a minimum
 - Historically used by the ALICE sites (~60 SEs)
- The VDT setup
 - Recently under heavy refurbishment after evaluating the 'xrd-installer'
 - Seems headed towards a middle point between the other two
 - AFAIK Mainly used in the US, mostly ATLAS-driven







Support



- Historical top-level: the usual contributing experts (e.g. Andy, Gerri, Fabrizio, Wilko, Jacek, Andreas, Derek, Artem etc.)
 - Whoever contributes a part typically gives support for that part
 - Often on a best-effort basis but well done so far
- Internally-grown ALICE knowledge, good for first level support
- AFAIK, OSG is setting up a more formal structure







The millionth test



- We have seen so many, sometimes with strange contradictory results
- Large scale tests can be precious
 - Difficult to do
 - Can generate useful feedback, thus improving quality
 - A "secret test" (but publishing the results) does not make life better
 - It declares the unwillingness to have a system which performs well
 - In other words: is your test performing bad? We are the first who want to understand why and make it better.
 - Moreover, can a test of very old stuff give indications for the future? I doubt.
- Not many such tests around now, probably a good sign







Every feature you like



- Xrootd is a fully plugin-based system
 - Developing new plugins can give whatever feature
 - E.g. PROOF, or XrdCASTOR
- Full POSIX-like namespace compliance+additional HEP reqs+unique hexabyte storage is an open research problem
- Weakened push towards reasonable usage patterns and architectures?
 - E.g. 'ls' floodings, namespace-unaware systems, useless file aliasings, apps wasting file resources, ...
 - Difficult to tell.







Open items (1/2)



- Hot: putting clients into servers
 - Or: different criteria to fully differentiate clients in the same app
 - E.g. How to instantiate together:
 - a client tuned for WAN TTreeCache-based random access
 - one optimized for blasting non-TTreeCache LAN traffic ?
 - one optimized for large files xfers
- The Extreme Copy (beta since ~1 year): Torrent-like dynamic multiserver file fetching
 - Needs the previous item to be really strong
- Components for site cooperation
 - Proxies and caching proxies (nothing more than proofs of concept right now)
 - Bandwidth/queuing manager (early alpha)
- 'Personal' persistent caching proxy (or TXNetFile/XrdClient extensions)
- A full-featured command line interface
 - The current one is a quite rough tool







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Open items (2/2)



- Client-side data management funcs (e.g. 'recursive ls' or 'df'): good level but incomplete by now
- WAN performance: huge breakthrough, still to gain
 - Both for file xfer and data access (TTreeCache and not)
- A robust and complete server-to-server file copy
- ROOT integration: very good quality, but still to gain
 - Only partially asynchronous (XrdClient can be fully async instead)
 - Will be more evident with the parallelization of the computing, I/O will likely become the bottleneck again
- "Intelligent" readaheads. After the ATLAS AOD story (Feb) I got no feedbacks...
 - The backport to ROOT 5.22 was done AFAIK
 - Andrei Maslennikov was the only feedback "well, you could try this..." Very partial but still something.
 - Interesting topic, many things could be done.
- An homogeneous, top-class support structure
 - Quality helps support, but vice-versa is true as well





Thank you



Questions?

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